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Ballistic motion planning

Mylène Campana^{1,2} Jean-Paul Laumond^{1,2}

Abstract— This paper addresses the motion planning problem of a jumping point-robot. Each jump consists in a ballistic motion linking two positions in contact with obstacle surfaces. A solution path is thus a sequence of parabola arcs. The originality of the approach is to consider non-sliding constraints at contact points: slipping avoidance is handled by constraining takeoff and landing velocity vectors to 3D friction cones. Furthermore the magnitude of these velocities is bounded. The ballistic motion lying in a vertical plane, we transform the 3D problem into a 2D one. We then solve the motion equations. The solution gives rise to an exact steering method computing a jump path between two contact points while respecting all constraints. The method is integrated into a standard probabilistic roadmap planner. Probabilistic completeness is proven. Simulations illustrate the performance of the approach.

I. INTRODUCTION

Geometrical motion planning is a well-known problem [1]. Today, most motion planners are inspired by the random sampling seminal approaches [2], [3]. They can also be adapted for systems that handle special types of paths (e.g. spline-based approach [4]), or which embed special constraints (e.g. constrained optimization [5]).

In this paper we consider the ballistic motion planning for a jumping robot in environments containing slippery surfaces. It is well-known that ballistic motion results in a parabola trajectory. According to the Coulomb friction law, a condition for the robot not to slide during its takeoff is that the contact force belongs to a so-called friction cone. This latter property extends to the landing phase. We consider a simple point mass robot with simplified contact dynamics: we assume that the robot is submitted to an impulse force as soon as it lands, so that the transition between landing and takeoff is instantaneous. This gives rise to a discontinuity between the contact forces and the contact velocities. Moreover we assume that the robot has limited energy resources, which limit the velocity at takeoff. The landing velocity is also constrained to avoid requiring to dissipate too much energy (and damaging the robot). These energy restrictions are realized by limiting the magnitude of the velocity vectors during the takeoff and landing phases. Constraints on the velocity vector magnitudes are named velocity constraints.

How to plan a collision-free path satisfying both sliding and velocity constraints in such a context? This is the question addressed by this paper.

Ballistic motion planning has been relatively little addressed. [6] shows a one-legged robot hopping while keeping balance. [7] makes a miniature robot jump to climb horizontal stairs. In [8], a multi-articulated gymnast robot jumps over obstacles on a horizontal surface, while taking into account the whole-body angular momentum. In [9]–[13] ballistic motion is considered from a character animation viewpoint. Focus is done on jump motion preparation and adaptation. The authors of [14] present a simulated humanoid robot that runs and jumps on a horizontal platform. The takeoff leg angle and intensity are computed to cross the large gap. [15] and [16] produce near optimal hopping motions online for quadrupeds from reinforcement based learning techniques. However, as they simplify the location of the contacts, the method cannot deal with arbitrary 3D environments. [17] and [18] consider how existing motions may be edited of re-timed to consider friction. Motion planning and obstacle avoidance are considered in [19]. The method computes a sequence of parabolas. It consists of tuning the parabola heights in order to reach different levels while avoiding obstacles.

This paper does not consider the full dynamics of articulated avatars, but is restricted to point-robots. With respect to the state of the art, the contribution is to account for slipping prevention as well as takeoff and landing velocity limitations. Furthermore, the proposed approach applies on 3D environments and rough terrains without any restriction.

We formally state our problem in Section II. In Section III, we detail the unconstrained parabola trajectory equations. Then, we present and solve the non-sliding constraints as well as the velocity limitations in Section IV. Finally, this resolution is integrated into a motion planner (Section V) whose probabilistic completeness is proven. Simulation results are provided in Section VI.

II. PROBLEM STATEMENT

Let us consider a point-robot moving in a 3D environment. The robot begins from a starting position \mathbf{c}_s and wants to reach a goal position \mathbf{c}_g , only by performing jumps from one contact to another. Both \mathbf{c}_s and \mathbf{c}_g are assumed to be in contact with the environment. There is no distinction between ground and obstacles. The purpose of this paper is to determine a sequence of jumps, under the following assumptions:

- The robot is modeled by a point mass m of position \mathbf{c} with respect to the origin.
- The only force that applies to the robot during a jump is mg .
- Contact phases are instantaneous, so that the velocity at a contact point is discontinuous, i.e. transition from

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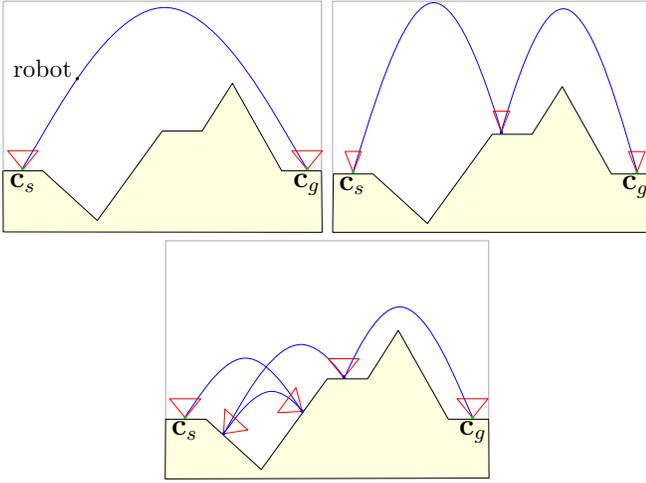


Fig. 1. Three sequences of parabolas between \mathbf{c}_s and \mathbf{c}_g positions for different constraints. In the case illustrated on the right, friction cones are narrower than those on the left, so that the parabola on top is not admissible anymore, and a waypoint has to be used. In the bottommost case, the initial velocity has been limited compared to the left case, resulting in a sequence with numerous jumps.

landing to takeoff results from an impulsion.

- The non-sliding constraint is modeled by a friction cone. According to the contact model, the robot is not sliding as soon as its impulsion vector belongs to the friction cone. The direction of the impulsion vector is the same as the difference between the takeoff and the landing velocity vectors. Then, if both velocity vectors belong to the convex friction cone, so does the impulsion direction. We detail the benefit of this sufficient condition in the motion planning section. Moreover, the surface material is uniform in the environment, i.e. the non sliding constraints can be modeled everywhere by a friction cone with a constant coefficient. We denote by μ the tangent of the cone half-angle.
- Takeoff and landing velocity magnitudes are bounded by the same value, so that an admissible jump path can be followed in a reversed direction.
- There is no constraint on the energy balance between one jump and the next.
- The robot cannot collide with the obstacles.

Fig. 1 illustrates the effects of the friction and velocity constraints on the existence of parabola sequences. The following section reminds the basics of ballistic motion and details the equations of parabolas linking two points.

III. UNCONSTRAINED BALLISTIC MOTION

A. Accessible space of ballistic motion

We denote the global frame basis by $(\mathbf{e}_x, \mathbf{e}_y, \mathbf{e}_z)$. When Newton's second law of motion is integrated with respect to time for a ballistic shot from the \mathbf{c}_s position with a $\dot{\mathbf{c}}_s$ initial velocity, the following robot trajectory is obtained:

$$\mathbf{c}(t) = -\frac{g}{2} t^2 \mathbf{e}_z + \dot{\mathbf{c}}_s t + \mathbf{c}_s \quad (1)$$

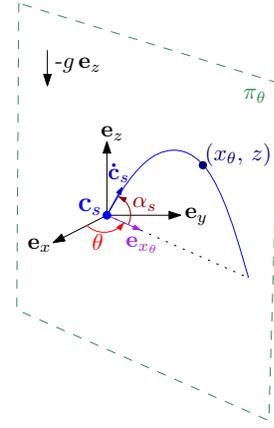


Fig. 2. The parabola always belongs to the plane π_θ defined by $(\mathbf{c}_s; \mathbf{e}_{x_\theta}; \mathbf{e}_z)$, where $\mathbf{e}_{x_\theta} = \cos(\theta)\mathbf{e}_x + \sin(\theta)\mathbf{e}_y$.

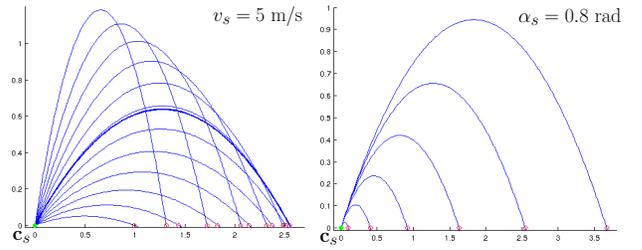


Fig. 3. Accessible space from $\mathbf{c}_s = (0, 0)^T$ when varying takeoff angle α_s (left), or when varying initial velocity v_s (right). On left, the bold parabola leads to the maximal range at given initial velocity, and is obtained for $\alpha_s = \frac{\pi}{4}$.

Let $(x, y, z)^T$ be the coordinates of \mathbf{c} . The ballistic motion belongs to a vertical plane denoted by π_θ . The orientation of the plane is given by the initial velocity components as follows:

$$\theta = \text{atan2}(\dot{y}_s, \dot{x}_s) \in [-\pi; \pi]$$

Considering $\Theta = [\cos(\theta) \sin(\theta) 0]^T$, we introduce the following variable changes involving the scalar product:

$$\begin{aligned} x_\theta &= \mathbf{c} \cdot \Theta, & x_{\theta_s} &= \mathbf{c}_s \cdot \Theta \\ x_{\theta_g} &= \mathbf{c}_g \cdot \Theta, & \dot{x}_{\theta_s} &= \dot{\mathbf{c}}_s \cdot \Theta \end{aligned}$$

Thus from (1), one can rewrite the main equations of motion determining the robot coordinates $(x_\theta, z)^T$ in π_θ (see Fig. 2):

$$z = -\frac{g}{2} \frac{(x_\theta - x_{\theta_s})^2}{\dot{x}_{\theta_s}^2} + \frac{\dot{z}_s}{\dot{x}_{\theta_s}} (x_\theta - x_{\theta_s}) + z_s \quad (2)$$

$$\frac{\dot{z}}{\dot{x}_{\theta_s}} = -g \frac{x_\theta - x_{\theta_s}}{\dot{x}_{\theta_s}^2} + \frac{\dot{z}_s}{\dot{x}_{\theta_s}} \quad (3)$$

Let us denote the takeoff angle by $\alpha_s = \text{atan2}(\dot{z}_s, \dot{x}_{\theta_s})$ and the velocity value $\|\dot{\mathbf{c}}_s\|$ by v_s . Equations (2)-(3) highlight the two parameters α_s and \dot{x}_{θ_s} that determine a parabola in π_θ . For instance, Fig. 3 presents the parabola beams when v_s (resp. α_s) is fixed. This can also be viewed as the accessible space of the robot performing ballistic motions.

B. Goal-oriented ballistic motion

Now we want our robot to reach the goal position \mathbf{c}_g with a jump starting from \mathbf{c}_s . Therefore, the value $\theta = \text{atan2}(y_g - y_s, x_g - x_s)$ is now known. Let X_θ equal $x_{\theta_g} - x_{\theta_s}$ and Z equal $z_g - z_s$. Since Z is fixed, it appears that from (2), α_s is the only remaining variable to compute the parabola beam leading to \mathbf{c}_g . In fact, the initial velocity \dot{x}_{θ_s} can be obtained with the following equation:

$$\dot{x}_{\theta_s} = \sqrt{\frac{gX_\theta^2}{2(X_\theta \tan(\alpha_s) - Z)}} \quad (4)$$

Equation (4) implies that \dot{x}_{θ_s} is only defined for top-curved parabolas, which is consistent with gravity. Non-physically-feasible parabolas such as down-curved ones are not considered. Therefore we impose:

$$\text{atan2}(Z, X_\theta) < \alpha_s < \frac{\pi}{2} \quad (5)$$

An example of a goal-oriented parabola beam is presented Fig. 4. Finally, we denote a parabola starting from \mathbf{c}_s and its parameters by $\mathcal{P}_s(\theta, \alpha, v)$.

IV. BALLISTIC MOTION WITH CONSTRAINTS

So far we have defined a beam of feasible parabolas to connect two positions. In this section, the non-sliding and velocity constraints are introduced, and the resulting reduction of the space of admissible parabola beams is detailed.

A. Non-sliding constraints

1) *2D reduction*: Let us consider two points \mathbf{c}_s and \mathbf{c}_g at the contact of environment surfaces. Non-sliding constraints impose the robot to land and take off along velocity vectors that belong to 3D friction cones of apexes based on \mathbf{c}_s and \mathbf{c}_g .

Since the motion has to lie in a vertical plane π_θ , the problem of computing a parabola between the two 3D friction cones is reduced to a 2D problem. Corresponding 2D cones result from the intersections of the 3D cones with the plane π_θ (see Fig. 5 top). If one of both intersection sets is reduced to a point, there is no possible jump between

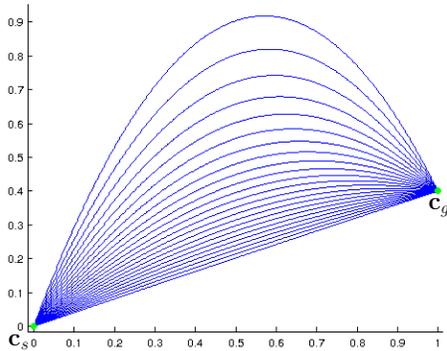


Fig. 4. Physically-feasible parabolas linking \mathbf{c}_s and \mathbf{c}_g , for multiple values of α_s in $[0.91; 1.27]$ rad.

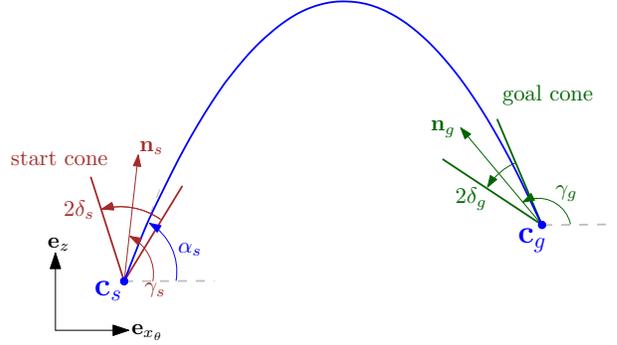
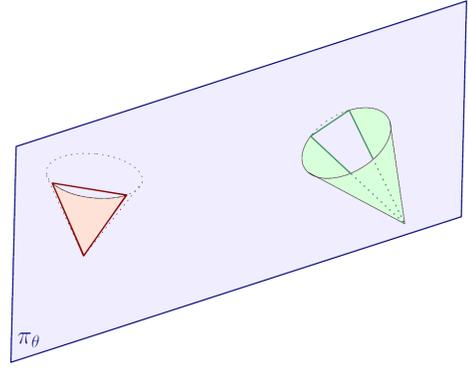


Fig. 5. (Top) Representation of the intersection between π_θ and two 3D cones. (Bottom) 2D cones resulting of the intersection and an example of parabola belonging to both cones.

\mathbf{c}_s and \mathbf{c}_g . Otherwise, let us denote their half-apex angles¹ respectively by δ_s and δ_g , and their directions projected in π_θ respectively by \mathbf{n}_s and \mathbf{n}_g . Note that δ_s and δ_g may be smaller than $\arctan(\mu)$. For the 2D start cone of direction $\mathbf{n}_s = (n_{x_s}, n_{y_s}, n_{z_s})^T$, let us denote by γ_s the angle between the cone direction and the horizontal line:

$$\gamma_s = \text{atan2}(n_{z_s}, n_{x_s} \cos(\theta) + n_{y_s} \sin(\theta))$$

γ_g is similarly defined, respectively to the 2D goal cone. Thus the problem of slippage avoidance is reduced to the problem of finding a parabola going through the 2D cones (see Fig. 5 bottom).

2) *Constraint formulation*: Since the equation of a parabola starting at \mathbf{c}_s and ending at \mathbf{c}_g only depends on α_s , the four constraints are just formulated relatively to α_s .

For the non-sliding takeoff constraint, the inequalities on α_s are immediate:

$$\alpha_1^- \leq \alpha_s \leq \alpha_1^+ \quad \text{with} \quad \begin{cases} \alpha_1^- = \gamma_s - \delta_s \\ \alpha_1^+ = \gamma_s + \delta_s \end{cases}$$

To express the three remaining constraints according to α_s , (2)-(3) are brought back to the parabola origin \mathbf{c}_s . Thus constraints are still expressed as inequalities:

$$\alpha_i^- \leq \alpha_s \leq \alpha_i^+, \quad i \in \{2..4\}$$

For the landing cone constraint, different cases appear, depending on the accessibility of the cone. They are tackled by

¹Detailed computation of δ_s and δ_g is presented in the appendix file.

Algorithm 1 Resolution of the landing cone constraint.

Output: Defined constraint bounds α_2^-, α_2^+

if $\gamma_g > 0$ **then**

$$\alpha_g^- = \gamma_g - \pi - \delta_g$$

$$\alpha_g^+ = \gamma_g - \pi + \delta_g$$

if $\alpha_g^+ < -\frac{\pi}{2}$ **then**

No solution

else

$$\alpha_2^- = \arctan\left(\frac{2Z}{X_\theta} - \tan(\alpha_g^+)\right)$$

if $\alpha_2^- > -\frac{\pi}{2}$ **then**

$$\alpha_2^+ = \arctan\left(\frac{2Z}{X_\theta} - \tan(\alpha_g^-)\right)$$

else

α_2^+ not defined

else

$$\alpha_g^- = \gamma_g + \pi - \delta_g$$

$$\alpha_g^+ = \gamma_g + \pi + \delta_g$$

if $\alpha_g^+ > \frac{\pi}{2}$ **then**

No solution

else

$$\alpha_2^+ = \arctan\left(\frac{2Z}{X_\theta} - \tan(\alpha_g^-)\right)$$

if $\alpha_2^+ < \frac{\pi}{2}$ **then**

$$\alpha_2^- = \arctan\left(\frac{2Z}{X_\theta} - \tan(\alpha_g^+)\right)$$

else

α_2^- not defined

Algorithm 1, which returns the constraint bounds (α_2^-, α_2^+) . Note that one of the bounds may not exist, and that the constraint may also not be satisfied.

B. Velocity constraints

Takeoff velocity limitation is expressed as $v_s \leq V_{max}$. Equation (2) leads to:

$$gX_\theta^2 \tan(\alpha_s)^2 - 2X_\theta V_{max}^2 \tan(\alpha_s) + gX_\theta^2 + 2ZV_{max}^2 \leq 0 \quad (6)$$

$$\Delta = V_{max}^4 - 2gZV_{max}^2 - g^2X_\theta^2$$

If $\Delta < 0$, (6) has no solution. In other words, it means that the goal position is not reachable with an initial velocity satisfying the limitation. In the case where the constraint is solvable, we write:

$$\begin{cases} \alpha_3^- = (V_{max}^2 - \sqrt{\Delta})/gX_\theta \\ \alpha_3^+ = (V_{max}^2 + \sqrt{\Delta})/gX_\theta \end{cases}$$

The same argument can be applied for the landing velocity limitation (see Fig. 6). Using (3)-(4), we can rewrite the constraint equation $v_f \leq V_{max}$ as:

$$gX_\theta^2 \tan(\alpha_s)^2 - (4X_\theta Zg + 2X_\theta V_{max}^2) \tan(\alpha_s) + gX_\theta^2 + 2ZV_{max}^2 + 4gZ^2 \leq 0 \quad (7)$$

$$\Lambda = V_{max}^4 + 2gZV_{max}^2 - g^2X_\theta^2$$

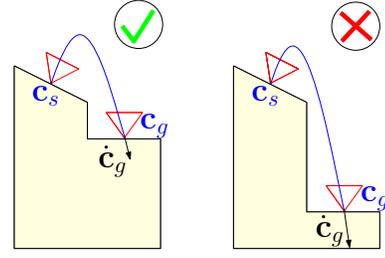


Fig. 6. The landing velocity limitation allows to reject parabolas that have too important an impact velocity magnitude v_g (right).

If $\Lambda < 0$, (7) has no solution. Otherwise, we write:

$$\begin{cases} \alpha_4^- = (V_{max}^2 + 2gZ - \sqrt{\Lambda})/gX_\theta \\ \alpha_4^+ = (V_{max}^2 + 2gZ + \sqrt{\Lambda})/gX_\theta \end{cases}$$

A symmetry property of the parabola implies that, given one parabola from c_s to c_g determined by \dot{c}_s , the same parabola can be obtained from c_g to c_s with $-\dot{c}_g$ as initial velocity. In the latter case, the contact velocity becomes $-\dot{c}_s$. We use this property to apply the same bound V_{max} for the takeoff and landing velocities. Thus the parabola can be traveled both ways without violating the velocity limitation constraints.

C. Constraints collection and solution existence

The domains where constraints are satisfied are convex. Thus, we intersect these domains to determine if an interval $]\alpha_s^-; \alpha_s^+[$ of α_s values complying with all the constraints exists. The interval bounds are given by:

$$\begin{cases} \alpha_s^- = \max(\alpha_1^-, \alpha_2^-, \alpha_3^-, \alpha_4^-) \\ \alpha_s^+ = \min(\alpha_1^+, \alpha_2^+, \alpha_3^+, \alpha_4^+) \end{cases}$$

Note that (5) has to be simultaneously satisfied to consider an admissible parabola. Fig. 7 presents an illustration of this constraint intersection. Constraint bounds $(\alpha_i^-, \alpha_i^+)_{i \in \{1..4\}}$ are used to plot parabolas, representing the domains where constraints are satisfied. The intersection of these domains leads to the set of possible solutions.

Finally, the existence of an admissible jump between two points is guaranteed as soon as:

- Neither of the intersections between both friction cone and π_θ is reduced to a point.
- (α_s^-, α_s^+) are defined and $\alpha_s^- \leq \alpha_s^+$.

The two conditions are necessary and sufficient. The interval $]\alpha_s^-; \alpha_s^+[$ gives a simple parametrization of the solution beam. Choosing α_s as the average $0.5(\alpha_s^- + \alpha_s^+)$ allows to optimize the distance to the constraints, e.g. to be far from the limits of the friction cones, and so far from sliding. Fig. 8 illustrates the constraint effects on a simple example.

Topological Property: Let us consider an admissible parabola $\mathcal{P}_s(\theta, \alpha, v)$ starting at c_s and ending at c_g . There exists a neighborhood \mathcal{N}_s (resp. \mathcal{N}_g) of c_s (resp. c_g) such that any pair of points (c_s^*, c_g^*) belonging to $\mathcal{N}_s \times \mathcal{N}_g$ can be linked by an admissible parabola.

1. Takeoff from initial cone
2. Landing in final cone
3. Takeoff velocity limitation
4. Landing velocity limitation

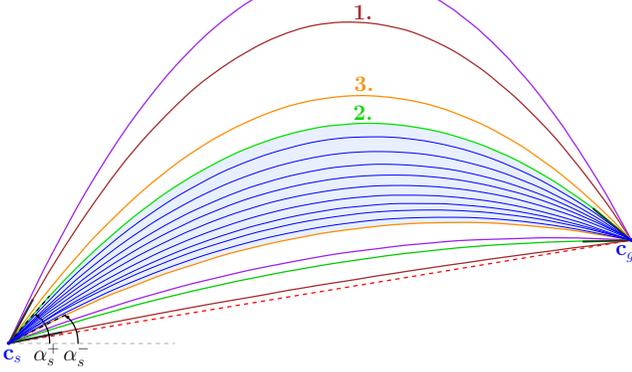


Fig. 7. Illustration of the constraints on a practical example: each constraint bound is used as α_s and represents a bold parabola. Between these bounds, the constraint is satisfied, out of them not. The constraint intersection is given by the bounds (α_s^-, α_s^+) and illustrated by the gray zone: blue parabolas belonging to it are admissible solutions to the problem.

Proof: Let us consider the parabola family $\{\mathcal{P}_s(\theta + e_1, \alpha + e_2, v + e_3), e_i \in] - \epsilon, \epsilon[\}$ starting at \mathbf{c}_s . The function giving the three parabola parameters from the three coordinates of \mathbf{c}_g is a homeomorphism. Therefore such a family spans a neighborhood of \mathbf{c}_g . Let \mathbf{c}_g^* be a point of this neighborhood and $\mathcal{P}_s(\theta + e_1^*, \alpha + e_2^*, v + e_3^*)$ the parabola from \mathbf{c}_s to \mathbf{c}_g^* . \mathbf{c}_g^* may be chosen close enough from \mathbf{c}_g to guarantee that e^* is small enough, and then $\mathcal{P}_s(\theta + e_1^*, \alpha + e_2^*, v + e_3^*)$ is admissible. Because the construction is symmetric, let us consider the same parabola as starting at \mathbf{c}_g^* and ending at \mathbf{c}_s . We get a new parametrization of the same parabola, i.e. $\mathcal{P}_g(\theta^*, \alpha^*, v^*)$. By using the same argument as above, the parabola family $\{\mathcal{P}_g(\theta^* + e_1, \alpha^* + e_2, v^* + e_3), e \in] - \epsilon, \epsilon[\}$ starting at \mathbf{c}_g^* spans a neighborhood of \mathbf{c}_s . The property holds for any point \mathbf{c}_g^* sufficiently close to \mathbf{c}_g . Therefore, there exists a neighborhood \mathcal{N}_s (resp. \mathcal{N}_g) of \mathbf{c}_s (resp. \mathbf{c}_g) such that any pair of points $(\mathbf{c}_s^*, \mathbf{c}_g^*)$ belonging to $\mathcal{N}_s \times \mathcal{N}_g$ can be linked by an admissible parabola. ■

V. MOTION PLANNING

To find a sequence of parabola arcs between an initial position \mathbf{c}_s and a final one \mathbf{c}_g , we use a simple PRM-based probabilistic roadmap planner [2] (see Algorithm 2). Note that the roadmap may contain cycles. The planner builds a roadmap in the 3D space by randomly sampling contact positions similarly to [20] (RANDOMSAMPLE), and by linking them (STEER) with admissible collision-free parabola arcs. The roadmap construction is over as soon as either a path linking \mathbf{c}_s and \mathbf{c}_g is found (ARECONNECTED), or computation time is over. Then, the function FINDSHORTESTPATH explores the roadmap to return the shortest path sequence, in terms of sum of parabola lengths.

Note that the sufficient non-slipping condition reduces the dimensionality of the motion planning problem, because it removes the relationship between the entering velocity and

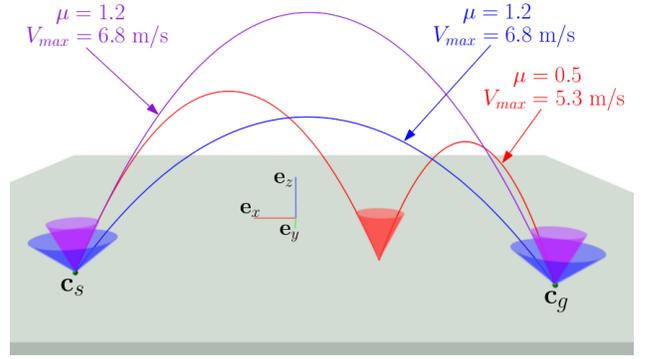


Fig. 8. Three parabola examples linking \mathbf{c}_s to \mathbf{c}_g with different constraints. Large (blue) and narrow (violet) friction cones are considered, forcing the solution parabola to be adapted. With a large velocity limitation, \mathbf{c}_g can be directly reached (blue). Otherwise, an intermediate position has to be considered (red).

the exiting velocity of a node. With a classical kynodynamic planner [21], to verify whether a trajectory can be connected with another one, it is required to extend the state space with the velocities and so it doubles the dimensionality of the problem. In our case, this is only required to verify whether there exists a velocity vector belonging to the cone each time we want to add a new path.

The function BEAM is described in Algorithm 3. It computes the interval of takeoff angles that generate constrained parabolas to link \mathbf{c}_s and \mathbf{c}_g . The algorithm starts by calculating each cone and plane π_θ intersection, and continues only if both intersections are not reduced to the cone apexes. Then, takeoff angle bounds related to constraints are computed as in Section IV. A boolean *fail* conveys the feasibility of constraints, i.e. if one constraint cannot be satisfied, *fail* is set to *true*. At this stage, an admissible parabola exists if the global constraint bounds verify $\alpha_s^- \leq \alpha_s^+$.

Then the steering method STEER detailed in Algorithm 4 selects a takeoff angle α_s and tests the corresponding parabola for collisions. If the parabola is not collision-free (HASCOLLISIONS), then we select a new parabola α_s by dichotomy on the interval $]\alpha_s^-; \alpha_s^+[$ until the resolution

Algorithm 2 Probabilistic roadmap planner for ballistic motion planning.

Input: *environment*, \mathbf{c}_s , \mathbf{c}_g , μ , V_{max}

Output: Collision-free solution *sequence* to problem

$path \leftarrow \text{STEER}(\mathbf{c}_s, \mathbf{c}_g)$

$finished \leftarrow \text{ARECONNECTED}(\mathbf{c}_s, \mathbf{c}_g)$

while not(*finished*) **do**

$\mathbf{c}_{random} \leftarrow \text{RANDOMSAMPLE}()$

$\text{ADDTOROADMAP}(\mathbf{c}_{random})$

for $\mathbf{c}_{node} \in \text{Roadmap}$ **do**

$path \leftarrow \text{STEER}(\mathbf{c}_{node}, \mathbf{c}_{random})$

$\text{ADDTOROADMAP}(path)$

end for

$finished \leftarrow \text{ARECONNECTED}(\mathbf{c}_s, \mathbf{c}_g)$

return *sequence* $\leftarrow \text{FINDSHORTESTPATH}()$

Algorithm 3 BEAM($\mathbf{c}_s, \mathbf{c}_g$): Computes the parabola beam represented by the takeoff angle interval $]\alpha_s^-; \alpha_s^+]$.

Input: $\mathbf{c}_s, \mathbf{c}_g, \mu, V_{max}$

Output: Interval of takeoff angles I_{beam}

```

 $cone_s^{2D} \leftarrow \text{COMPUTEINTERSECTION}(cone_s^{3D}, \pi_\theta)$ 
 $cone_g^{2D} \leftarrow \text{COMPUTEINTERSECTION}(cone_g^{3D}, \pi_\theta)$ 
if ISREDUCED( $cone_s^{2D}$ ) or ISREDUCED( $cone_g^{2D}$ ) then
  return  $I_{beam} \leftarrow \emptyset$ 
 $(\alpha_i^-, \alpha_i^+, fail)_{i \in \{1..4\}} \leftarrow \text{COMPUTECONSTRAINTS}()$ 
if  $fail = true$  then return  $I_{beam} \leftarrow \emptyset$ 
 $\alpha_s^- \leftarrow \max(\alpha_1^-, \alpha_2^-, \alpha_3^-, \alpha_4^-)$ 
 $\alpha_s^+ \leftarrow \min(\alpha_1^+, \alpha_2^+, \alpha_3^+, \alpha_4^+)$ 
return  $I_{beam} \leftarrow ]\alpha_s^-; \alpha_s^+[$ 

```

Algorithm 4 STEER($\mathbf{c}_s, \mathbf{c}_g$): Steering method based on a constrained parabola. Returns a collision-free path linking \mathbf{c}_s and \mathbf{c}_g . Otherwise, returns an empty path.

Input: $\mathbf{c}_s, \mathbf{c}_g, \mu, V_{max}, n_{limit}$

Output: Collision-free parabola path $path$

```

 $]\alpha_s^-; \alpha_s^+[\leftarrow I_{beam}$ 
 $n \leftarrow 1$ 
 $I_{beam} \leftarrow \text{BEAM}(\mathbf{c}_s, \mathbf{c}_g)$ 
if ISEMPY( $I_{beam}$ ) then return  $path \leftarrow emptyPath$ 
else
 $\alpha_s \leftarrow 0.5(\alpha_s^- + \alpha_s^+)$ 
 $path \leftarrow \text{COMPUTEPARABOLA}(\mathbf{c}_s, \mathbf{c}_g, \alpha_s)$ 
while HASCOLLISIONS( $path$ ) and  $n < n_{limit}$  do
   $\alpha_s \leftarrow \text{DICHOTOMY}(]\alpha_s^-; \alpha_s^+[ , n)$ 
   $path \leftarrow \text{COMPUTEPARABOLA}(\mathbf{c}_s, \mathbf{c}_g, \alpha_s)$ 
   $n \leftarrow n + 1$ 
if HASCOLLISIONS( $path$ ) then  $path \leftarrow emptyPath$ 
return  $path$ 

```

threshold n_{limit} is reached. In the worst case, the DICHOTOMY function allows to span the almost entire parabola beam. Doing so, the algorithm is probabilistically complete as proven by the following property.

Convergence Property: Let us consider a sequence of collision-free ballistic jumps between two points \mathbf{c}_s and \mathbf{c}_g in a given environment. Let us assume that the entire path is at a distance of about ϵ from the obstacles. Then the probability for Algorithm 2 to find a sequence of collision-free ballistic jumps between \mathbf{c}_s and \mathbf{c}_g converges to 1 when running time tends to infinity.

Proof: The property is a direct consequence of the topological property exposed in Section IV. Indeed, let us consider a sequence of collision-free ballistic jumps between two points \mathbf{c}_s and \mathbf{c}_g . Let \mathbf{c}_i and \mathbf{c}_{i+1} two consecutive points in the sequence. \mathbf{c}_i and \mathbf{c}_{i+1} are linked by a collision-free parabola \mathcal{P}_i . From the topological property, there are two neighborhoods \mathcal{N}_i (resp. \mathcal{N}_{i+1}) of \mathbf{c}_i (resp. \mathbf{c}_{i+1}) such that any pair of points $(\mathbf{c}_i^*, \mathbf{c}_{i+1}^*)$ belonging to $\mathcal{N}_i \times \mathcal{N}_{i+1}$ can be linked by an admissible parabola \mathcal{P}_i^* . Because Algorithm 2 tends to sample the environment uniformly, the probability of sampling two points in \mathcal{N}_i and \mathcal{N}_{i+1} respectively tends to

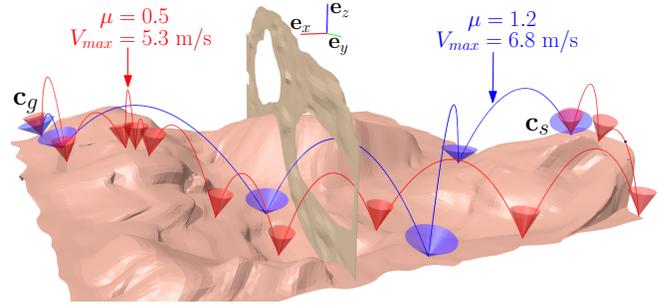


Fig. 9. Full path planning results in an environment containing two windows to cross from right to left. The red solution is more constrained, so it results in a longer sequence of parabolas. Number of triangles: 47 733.

Parameters		Computation time (s)	Collision found	Roadmap nodes	Path length (m)
μ	V_{max}				
0.5	6.5 m/s	9.89	3270	1995	39.9
0.5	7 m/s	9.99	4002	1835	37.4
1.2	6.5 m/s	1.04	601	282	28.1
1.2	7 m/s	0.909	540	237	27.0

TABLE I

AVERAGES OF 40 BALLISTIC PLANNING OF THE EXAMPLE FIG. 10, FOR FOUR COMBINATIONS OF THE PARAMETERS.

1 when time tends to infinity. \mathcal{N}_i and \mathcal{N}_{i+1} can be arbitrarily small. As a consequence, \mathcal{P}_i^* can be arbitrarily close to \mathcal{P}_i . Because \mathcal{P}_i is away about ϵ from the obstacles, \mathcal{P}_i^* is guaranteed to be collision-free. ■

VI. RESULTS

The ballistic motion planner was tested in 3D environments containing slippery surfaces, using the software Humanoid Path Planner². Graphical renderings were done using Blender 2.7. In all described examples, the parameter n_{limit} from Algorithm 4 was set to 6.

We planned sequences of parabolas for a point-robot in three environments. For each example, we considered weak and strong constraints. The results are shown in Fig. 9, 10 and 11. Movies of the trajectories are available in the companion video³. Solutions under strong constraints tend to increase the number of waypoints. It is not only a consequence of the velocity limitation that forces to reach closer positions (see also Fig. 8 bottom), but it is also a result of the cone narrowness. In fact, as it is shown in Fig. 8 top, narrow cones provide parabolas with greater heights, more likely to produce collisions or to exceed the environment bounds.

Table I presents the average performance results of the ballistic motion planner run on the Fig. 10 benchmark. The velocity limitation is less restrictive in terms of computation time than the cone coefficient. However, the velocity limitation cannot be reduced without endangering the existence of a solution. In fact, the robot has to reach other platforms in order to find a solution path sequence. Benchmarking was

²<http://humanoid-path-planner.github.io/hpp-doc/index.html>

³<https://youtu.be/vv.K7HqANmk>

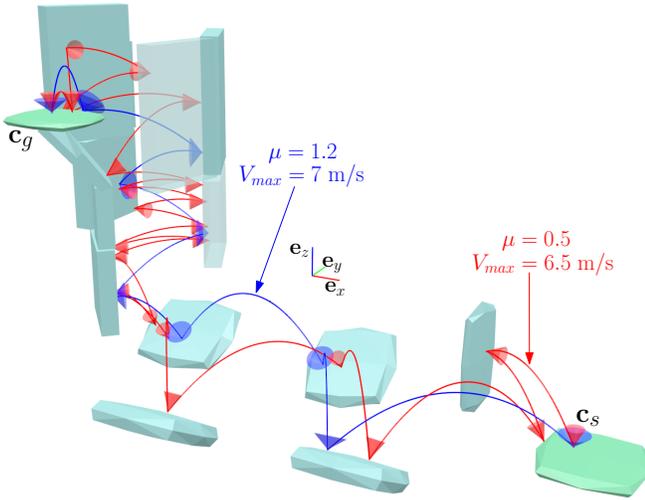


Fig. 10. Full path planning results in an environment containing platforms and a chimney. The blue path is constrained by large cones, the red path by narrow cones. Reducing μ prevents the creation of a parabola linking two low platforms with the same inclination. Number of triangles: 696.

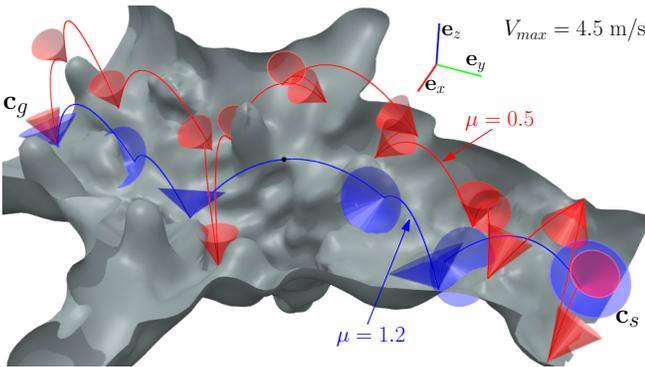


Fig. 11. Path planning results in a cave. The robot has to avoid the numerous stalactites, stalagmites and holes. Number of triangles: 33 513.

done on a PC with 15.8 GB of main memory and using one core of an Intel Xeon E3-1240 processor running at 3.4 GHz

VII. CONCLUSION

We presented a method that analytically computes a non-sliding jump for a point-robot, resulting in a parabola going from one friction cone to another. The method has been implemented as a steering method in a probabilistic-roadmap motion planner in order to determine a sequence of jumps between given start and goal positions.

We can easily extend the parabola-based steering method to devise a diffusing type of probabilistic planner, such as RRT [3]. Besides, our algorithm is the first stage of a more ambitious challenge. Our final purpose is to address dynamic motion planning for digital artifacts. The solution which we provide can be used to compute the center of mass path when the artifact is jumping. Now, it remains to consider more realistic models of contacts (e.g. multiple contacts involving feet and hands) and impacts (e.g. including energy balance). Moreover, the steering method we consider in the motion

planner is assumed to be symmetric. This assumption is not realistic. Indeed, for a given parabola, the energy required to overcome the gravity effect from a position is greater than the energy to dissipate when landing at the same position. The extension of the motion planner to more realistic energetic models is the purpose of future developments.

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