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▶ To cite this version:

Diego Cifuentes. On the degree-chromatic polynomial of a tree. 24th International Conference on Formal Power Series and Algebraic Combinatorics (FPSAC 2012), 2012, Nagoya, Japan. pp.69-72, 10.46298/dmtcs.3020. hal-01283152

HAL Id: hal-01283152 https://hal.science/hal-01283152

Submitted on 5 Mar 2016 $\,$

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On the degree-chromatic polynomial of a tree

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Abstract. The *degree chromatic polynomial* $P_m(G, k)$ of a graph G counts the number of k-colorings in which no vertex has m adjacent vertices of its same color. We prove Humpert and Martin's conjecture on the leading terms of the degree chromatic polynomial of a tree.

Résumé. Le polynôme degré chromatique $P_m(G, k)$ d'un graphe G compte le nombre de k-colorations dans lesquelles aucun sommet n'a m sommets adjacents de sa même couleur. On démontre la conjecture de Humpert et Martin sur les coefficients principaux du polynôme degré chromatique d'un arbre.

Keywords: chromatic polynomial, graph coloring, tree

George David Birkhoff defined the chromatic polynomial of a graph to attack the renowned four color problem. The chromatic polynomial P(G, k) counts the k-colorings of a graph G in which no two adjacent vertices have the same color [Read(1968)].

Given a graph G, Humpert and Martin defined its m-chromatic polynomial $P_m(G, k)$ to be the number of k-colorings of G such that no vertex has m adjacent vertices of its same color. They proved this is indeed a polynomial. When m = 1, we recover the usual chromatic polynomial of the graph.

The chromatic polynomial is of the form

$$P(G,k) = k^{n} - ek^{n-1} + o(k^{n-1}),$$

where n is the number of vertices and e the number of edges of G. For m > 1 the formula is no longer true, but Humpert and Martin conjectured the following formula which we now prove:

Theorem 1 ([Humpert and Martin(2010), Humpert and Martin(2011)], Conjecture) *Let* T *be a tree with* n *vertices and let* m *be an integer with* 1 < m < n, *then the following equation holds.*

$$P_m(T,k) = k^n - \sum_{v \in V(T)} {\binom{d(v)}{m}} k^{n-m} + o(k^{n-m})$$

Proof: For a given coloring of T, say vertices v_1 and v_2 are "friends" if they are adjacent and have the same color. For each v, let A_v be the set of colorings such that v has at least m friends. We want to find the number of colorings which are not in any A_v , and we will use the inclusion-exclusion principle. As the total number of k-colorings is k^n , we have

$$P_m(T,k) = k^n - \sum_{v \in V} |A_v| + \sum_{v_1, v_2 \in V} |A_{v_1} \cap A_{v_2}| - \dots$$

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We first show that $|A_v| = {d(v) \choose m} k^{n-m} + o(k^{n-m})$. Let $A_v^{(l)}$ be the set of k-colorings such that v has exactly l friends. In order to obtain a coloring in $A_v^{(l)}$, we may choose the l friends in ${d(v) \choose l}$ ways, the color of v and its friends in k ways, the color of the remaining adjacent vertices to v in $(k-1)^{d(v)-l}$ ways, and the color of the rest of the vertices in $k^{n-1-d(v)}$ ways. Then

$$|A_{v}| = \sum_{l=m}^{n-1} |A_{v}^{(l)}| = \sum_{l=m}^{n-1} {d(v) \choose l} k^{n-d(v)} (k-1)^{d(v)-l}$$
$$= {d(v) \choose m} k^{n-m} + o(k^{n-m}).$$

To complete the proof, it is sufficient to see that for any set S of at least 2 vertices $|\bigcap_{v \in S} A_v| = o(k^{n-m})$; clearly we may assume $S = \{v_1, v_2\}$. Consider the following cases:

Case 1 v_1 and v_2 are not adjacent.

Split A_{v_1} into equivalence classes with the equivalence relation

$$\sigma_1 \sim \sigma_2 \Leftrightarrow \sigma_1(w) = \sigma_2(w)$$
 for all $w \neq v_2$

Note that each equivalence class C consists of k colorings, which only differ in the color of v_2 . In addition, for each C at most $\frac{d(v_2)}{m}$ of its colorings are in A_{v_2} , as if $\sigma \in A_{v_2}$ there must be m vertices adjacent to v_2 with the color $\sigma(v_2)$. Therefore

$$|A_{v_1} \cap A_{v_2}| = \sum_C |C \cap A_{v_2}| \le \sum_C \frac{d(v_2)}{m} = \frac{|A_{v_1}|}{k} \cdot \frac{d(v_2)}{m}.$$

It follows that $\frac{|A_{v_1} \cap A_{v_2}|}{|A_{v_1}|}$ goes to 0 as k goes to infinity, so $|A_{v_1} \cap A_{v_2}| = o(k^{n-m})$.

Case 2 v_1 and v_2 are adjacent.

Let W be the set of adjacent vertices to v_2 other than v_1 . They are not adjacent to v_1 as T has no cycles. Split A_{v_1} into equivalence classes with the equivalence relation

$$\sigma_1 \sim \sigma_2 \Leftrightarrow \sigma_1(w) = \sigma_2(w)$$
 for all $w \notin W$.

Each equivalence class C consists of $k^{|W|}$ colorings, which may only differ in the colors of the vertices in W. If v_1 and v_2 are friends in the colorings of C, then a coloring in $|C \cap A_{v_2}|$ must contain at least m – 1 vertices in W of the same color as v_2 . Therefore

$$|C \cap A_{v_2}| = \sum_{l=m-1}^{|W|} \binom{|W|}{l} (k-1)^{|W|-l} < \sum_{l=0}^{|W|} \binom{|W|}{l} k^{|W|-1} = 2^{|W|} k^{|W|-1}.$$

Notice that here we are using $m \ge 2$ so that $l \ge 1$. Otherwise, if v_1 and v_2 are not friends in the colorings of C, then

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$$|C \cap A_{v_2}| = \sum_{l=m}^{|W|} {|W| \choose l} (k-1)^{|W|-l} < \sum_{l=0}^{|W|} {|W| \choose l} k^{|W|-1} = 2^{|W|} k^{|W|-1}.$$

Therefore

$$\begin{split} |A_{v_1} \cap A_{v_2}| &= \sum_C |C \cap A_{v_2}| < \sum_C 2^{|W|} k^{|W|-1} \\ &= \frac{|A_{v_1}|}{k^{|W|}} \cdot 2^{|W|} k^{|W|-1} = \frac{|A_{v_1}| \cdot 2^{|W|}}{k} \end{split}$$

and $|A_{v_1} \cap A_{v_2}| = o(k^{n-m})$ follows as in the first case.

This completes the proof of the theorem.

Acknowledgements

I would like to thank Federico Ardila for bringing this problem to my attention, and for helping me improve the presentation of this note. I would also like to acknowledge the support of the SFSU-Colombia Combinatorics Initiative and the Universidad de los Andes.

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