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# On the degree-chromatic polynomial of a tree

## Diego Cifuentes<sup>1</sup>

**Abstract.** The degree chromatic polynomial  $P_m(G,k)$  of a graph G counts the number of k-colorings in which no vertex has m adjacent vertices of its same color. We prove Humpert and Martin's conjecture on the leading terms of the degree chromatic polynomial of a tree.

**Résumé.** Le polynôme degré chromatique  $P_m(G,k)$  d'un graphe G compte le nombre de k-colorations dans lesquelles aucun sommet n'a m sommets adjacents de sa même couleur. On démontre la conjecture de Humpert et Martin sur les coefficients principaux du polynôme degré chromatique d'un arbre.

Keywords: chromatic polynomial, graph coloring, tree

George David Birkhoff defined the chromatic polynomial of a graph to attack the renowned four color problem. The chromatic polynomial P(G, k) counts the k-colorings of a graph G in which no two adjacent vertices have the same color [Read(1968)].

Given a graph G, Humpert and Martin defined its m-chromatic polynomial  $P_m(G,k)$  to be the number of k-colorings of G such that no vertex has m adjacent vertices of its same color. They proved this is indeed a polynomial. When m=1, we recover the usual chromatic polynomial of the graph.

The chromatic polynomial is of the form

$$P(G,k) = k^{n} - ek^{n-1} + o(k^{n-1}),$$

where n is the number of vertices and e the number of edges of G. For m > 1 the formula is no longer true, but Humpert and Martin conjectured the following formula which we now prove:

**Theorem 1** ([Humpert and Martin(2010), Humpert and Martin(2011)], Conjecture) Let T be a tree with n vertices and let m be an integer with 1 < m < n, then the following equation holds.

$$P_m(T, k) = k^n - \sum_{v \in V(T)} {d(v) \choose m} k^{n-m} + o(k^{n-m})$$

**Proof:** For a given coloring of T, say vertices  $v_1$  and  $v_2$  are "friends" if they are adjacent and have the same color. For each v, let  $A_v$  be the set of colorings such that v has at least m friends. We want to find the number of colorings which are not in any  $A_v$ , and we will use the inclusion-exclusion principle. As the total number of k-colorings is  $k^n$ , we have

$$P_m(T,k) = k^n - \sum_{v \in V} |A_v| + \sum_{v_1,v_2 \in V} |A_{v_1} \cap A_{v_2}| - \dots$$

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We first show that  $|A_v| = \binom{d(v)}{m} k^{n-m} + o(k^{n-m})$ . Let  $A_v^{(l)}$  be the set of k-colorings such that v has exactly l friends. In order to obtain a coloring in  $A_v^{(l)}$ , we may choose the l friends in  $\binom{d(v)}{l}$  ways, the color of v and its friends in k ways, the color of the remaining adjacent vertices to v in  $(k-1)^{d(v)-l}$  ways, and the color of the rest of the vertices in  $k^{n-1-d(v)}$  ways. Then

$$|A_v| = \sum_{l=m}^{n-1} |A_v^{(l)}| = \sum_{l=m}^{n-1} {d(v) \choose l} k^{n-d(v)} (k-1)^{d(v)-l}$$
$$= {d(v) \choose m} k^{n-m} + o(k^{n-m}).$$

To complete the proof, it is sufficient to see that for any set S of at least 2 vertices  $|\bigcap_{v \in S} A_v| = o(k^{n-m})$ ; clearly we may assume  $S = \{v_1, v_2\}$ . Consider the following cases:

**Case 1**  $v_1$  and  $v_2$  are not adjacent.

Split  $A_{v_1}$  into equivalence classes with the equivalence relation

$$\sigma_1 \sim \sigma_2 \Leftrightarrow \sigma_1(w) = \sigma_2(w)$$
 for all  $w \neq v_2$ .

Note that each equivalence class C consists of k colorings, which only differ in the color of  $v_2$ . In addition, for each C at most  $\frac{d(v_2)}{m}$  of its colorings are in  $A_{v_2}$ , as if  $\sigma \in A_{v_2}$  there must be m vertices adjacent to  $v_2$  with the color  $\sigma(v_2)$ . Therefore

$$|A_{v_1} \cap A_{v_2}| = \sum_{C} |C \cap A_{v_2}| \le \sum_{C} \frac{d(v_2)}{m} = \frac{|A_{v_1}|}{k} \cdot \frac{d(v_2)}{m}.$$

It follows that  $\frac{|A_{v_1} \cap A_{v_2}|}{|A_{v_1}|}$  goes to 0 as k goes to infinity, so  $|A_{v_1} \cap A_{v_2}| = o(k^{n-m})$ .

**Case 2**  $v_1$  and  $v_2$  are adjacent.

Let W be the set of adjacent vertices to  $v_2$  other than  $v_1$ . They are not adjacent to  $v_1$  as T has no cycles. Split  $A_{v_1}$  into equivalence classes with the equivalence relation

$$\sigma_1 \sim \sigma_2 \Leftrightarrow \sigma_1(w) = \sigma_2(w) \text{ for all } w \notin W.$$

Each equivalence class C consists of  $k^{|W|}$  colorings, which may only differ in the colors of the vertices in W. If  $v_1$  and  $v_2$  are friends in the colorings of C, then a coloring in  $|C \cap A_{v_2}|$  must contain at least m-1 vertices in W of the same color as  $v_2$ . Therefore

$$|C \cap A_{v_2}| = \sum_{l=m-1}^{|W|} \binom{|W|}{l} (k-1)^{|W|-l} < \sum_{l=0}^{|W|} \binom{|W|}{l} k^{|W|-1} = 2^{|W|} k^{|W|-1}.$$

Notice that here we are using  $m \geq 2$ so that  $l \geq 1$ . Otherwise, if  $v_1$  and  $v_2$  are not friends in the colorings of C, then

$$|C \cap A_{v_2}| = \sum_{l=m}^{|W|} {|W| \choose l} (k-1)^{|W|-l} < \sum_{l=0}^{|W|} {|W| \choose l} k^{|W|-1} = 2^{|W|} k^{|W|-1}.$$

Therefore

$$|A_{v_1} \cap A_{v_2}| = \sum_C |C \cap A_{v_2}| < \sum_C 2^{|W|} k^{|W|-1}$$

$$= \frac{|A_{v_1}|}{k^{|W|}} \cdot 2^{|W|} k^{|W|-1} = \frac{|A_{v_1}| \cdot 2^{|W|}}{k}$$

and  $|A_{v_1} \cap A_{v_2}| = o(k^{n-m})$  follows as in the first case.

This completes the proof of the theorem.

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