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Computing periods of rational integrals

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OBJECTIVE

Input $y(t) = \oint R(t, z_0, \dots, z_n) dz_0 \cdots dz_n$

- Multiple integral
- of a rational function
- over a closed path.
- **This is period of rational integral.**

Output $c_m(t)y^{(m)}(t) + \cdots + c_0(t)y(t) = 0$

- Linear differential equation for y
- with polynomial coefficients.
- **This is the Picard-Fuchs equation.**

[Classical theorem – Such a differential equation always exists.]

Computing a differential equation satisfied by a function is computing this function.

A differential equation and initial condition:

- defines uniquely the function,
- allows numerical evaluation, fast and with arbitrary precision,
- allows to compute asymptotics and power series expansions,
- can prove identities.

MOTIVATIONS

• **Multiple binomial sums** can be computed as periods of rational functions through integral representations of generating functions. For example, with $u_n = \sum_{i=0}^n \sum_{j=0}^n \binom{i+j}{j}^2 \binom{4n-2i-2j}{2n-2i}$

$$\sum_{n \geq 0} u_n t^n = \frac{1}{(2i\pi)^2} \oint \frac{(y-1)y^3}{(t-y^2(y-1)^2)((x-1)t-xy^2(y^2+x-1))}$$

The computation of the Picard-Fuchs equation gives a proof of the following identity:

$$\sum_{i=0}^n \sum_{j=0}^n \binom{i+j}{j}^2 \binom{4n-2i-2j}{2n-2i} = (2n+1) \binom{2n}{n}^2$$

This is **end-to-end automated**, cf. Bruno Salvy's talk.

- In the context of **mirror symmetry**, topological invariants of a given variety, like Gromov-Witten invariants, can be computed with a Picard-Fuchs equation related to the mirror family.
- ~> Big integrals, hard to compute.

- And also combinatorics, number theory, physics.

ELEMENTARY RELATIONS

Integral of a derivative $\oint \sum_{i=0}^n \frac{\partial C_i}{\partial z_i} dz_0 \cdots dz_m = 0$

Integration by parts $\oint F \frac{\partial G}{\partial z_i} dz_0 \cdots dz_n = - \oint \frac{\partial F}{\partial z_i} G dz_0 \cdots dz_n$

Differentiation under \oint $\frac{\partial}{\partial t} \oint F dz_0 \cdots dz_m = \oint \frac{\partial F}{\partial t} dz_0 \cdots dz_m$

EXAMPLE

Perimeter of an ellipse as a function of its eccentricity:

$$p(e) = \frac{1}{2i\pi} \oint \frac{dx dy}{1 - \frac{1-e^2 x^2}{(1-x^2)y^2}}$$

Then $(e - e^3)p'' + (1 - e^2)p' + ep = 0$ (Euler, 1733)

Proof.

operator in e and ∂_e only, commutes with \oint the rational function being integrated

$$\left((e - e^3)\partial_e^2 + (1 - e^2)\partial_e + e \right) \cdot \left(\frac{1}{1 - \frac{1-e^2 x^2}{(1-x^2)y^2}} \right) =$$

$$\partial_x \left(\frac{e(-1-x+x^2+x^3)y^2(-3+2x+y^2+x^2(-2+3e^2-y^2))}{(-1+y^2+x^2(e^2-y^2))^2} \right) + \partial_y \left(\frac{2e(-1+e^2)x(1+x^3)y^3}{(-1+y^2+x^2(e^2-y^2))^2} \right)$$

sum of derivatives, its integral over a closed path is zero

QED.

NOTATIONS

- $R(t, z_0, \dots, z_n) = a/f$ a rational function
- $\mathbb{K} = \mathbb{C}(t)$, the base field
- ∂_i denotes the partial differentiation $\partial/\partial z_i$
- $R \mapsto \oint R dz$ denotes really the canonical \mathbb{K} -linear surjection

$$\mathbb{K}[z_0, \dots, z_n, \frac{1}{f}] \rightarrow \frac{\mathbb{K}[z_0, \dots, z_n, \frac{1}{f}]}{\sum_{i=0}^n \partial_i \mathbb{K}[z_0, \dots, z_n, \frac{1}{f}]}$$

HOMOGENEITY

Up to homogenization with a new variable, we can assume that R is homogeneous of degree $-n-1$, that is

$$R(t, \lambda z_0, \dots, \lambda z_n) = \lambda^{-n-1} R(t, z_0, \dots, z_n),$$

or equivalently,

$$R(t, \lambda z_0, \dots, \lambda z_n) d(\lambda z_0) \cdots d(\lambda z_n) = R(t, z_0, \dots, z_n) dz_0 \cdots dz_n.$$

Confinement – For all polynomials f and all integers q , the space of homogeneous fractions of degree $-n-1$ of the form a/f^q is finite dimensional.

CONTRIBUTIONS

- New and fast algorithm.
- Uses only commutative Gröbner basis and linear algebra, rather than \mathcal{D} -module theory.
- Based on an extension of Griffiths-Dwork reduction which handles singular cases.
- Implementation available in Magma. <https://github.com/lairez/periods.git>
- Computes quickly integrals that were previously out-of-reach.

REDUCTION OF POLE ORDER

Principle – Rewrite $\oint \frac{a}{f^q} dz$ as $\oint \frac{a}{f^{q-1}} dz$ using integration by parts.

How? – If $a = b_0 \partial_0 f + \cdots + b_n \partial_n f$, for some polynomials $b_i \in \mathbb{K}[z]$, then

$$\oint \frac{a}{f^q} dz = \underbrace{\oint \sum_{i=0}^n b_i \frac{\partial_i f}{f^q} dz}_{\text{integration by parts}} = \frac{1}{q-1} \oint \sum_{i=0}^n \frac{\partial_i b_i}{f^{q-1}}.$$

Rewriting rule $\frac{\sum_i b_i \partial_i f}{f^q} \rightarrow \frac{1}{q-1} \frac{\sum_i \partial_i b_i}{f^{q-1}}$

By construction, if R rewrites into R' by repeated application of the rewriting rule, then

$$\oint R dz = \oint R' dz.$$

Definition a/f^q is *reducible* if there exists b such that

$$\frac{a}{f^q} \rightarrow \frac{b}{f^{q-1}}.$$

Nota bene a/f^q is reducible if and only if a is in the jacobian ideal of f , which is $\langle \partial_0 f, \dots, \partial_n f \rangle$.

GRIFFITHS-DWORK THEOREM

Assume that f defines a **smooth** hypersurface in the projective space $\mathbb{P}^n(\mathbb{K})$. Let a/f^q a homogeneous fraction of degree $-n-1$.

- If $q > n$, then a/f^q is reducible.
- In particular, there exists a polynomial b such that

$$\oint \frac{a}{f^q} dz = \oint \frac{b}{f^n} dz.$$

- If $\oint \frac{a}{f^q} dz = 0$ then a/f^q is reducible.

ALGORITHM, SMOOTH CASE

1. For each $k \geq 0$, compute b_k such that

$$\frac{\partial^k a}{\partial t^k f} \rightarrow \cdots \rightarrow \frac{b_k}{f^n}.$$

[Always possible, by Griffiths-Dwork theorem.]

2. Find a linear dependency relation over $\mathbb{K} = \mathbb{C}(t)$:

$$c_0(t) \frac{b_0}{f^n} + \cdots + c_m(t) \frac{b_m}{f^n} = 0.$$

[Exists for m large enough, by the confinement property.]

3. When found, return $c_m(t)y^{(m)}(t) + \cdots + c_0(t)y(t) = 0$.

What if f is not smooth?

Generically f is smooth. In practice, most interesting integrals, if not all, have a **singular** f .
generic case \neq usual case

DIMCA'S THEOREM

Without assumptions on f , let a/f^q be a homogeneous fraction of degree $-n-1$. Then there exists a polynomial b such that

$$\oint \frac{a}{f^q} dz = \oint \frac{b}{f^n} dz.$$

But how to compute b/f^n from a/f^q ?

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Reduction rules for general case

EXTENSION OF GRIFFITHS-DWORK REDUCTION

Consider again the rewriting rule

$$\frac{\sum_i b_i \partial_i f}{f^{q+1}} \rightarrow \frac{1}{q} \frac{\sum_i \partial_i b_i}{f^q}.$$

The decomposition of a polynomial as $\sum_i b_i \partial_i f$ is not unique: zero decomposes non trivially. If $\sum_i b_i \partial_i f = 0$, then what about

$$\ll 0 \rightarrow \frac{1}{q} \frac{\sum_i \partial_i b_i}{f^q} \gg ?$$

Not a reduction but indicates that

$$\oint \frac{\sum_i \partial_i b_i}{f^q} dz = 0$$

But it is possible that $(\sum_i \partial_i b_i) / f^q$ is not reducible!

New rule If $\sum_{i=0}^n b_i \partial_i f = 0$, then $\frac{\sum_i \partial_i b_i}{f^q} \rightarrow 0$

Rank 1 rules $\stackrel{\text{def}}{=} \text{Griffiths-Dwork rules}$

Rank 2 rules $\stackrel{\text{def}}{=} \text{Rank 1 rules} + \text{New rules above}$

FURTHER EXTENSION

Sometimes $\frac{a}{f^q} \xrightarrow{\text{rk}_1} \frac{b}{f^{q-1}}$ but b/f^{q-1} is not reducible at rank 2.

New rule If $\frac{a}{f^q} \xrightarrow{\text{rk}_1} \frac{b}{f^{q-1}}$ then $\frac{b}{f^{q-1}} \rightarrow 0$

Rank 3 rules $\stackrel{\text{def}}{=} \text{Rank 2 rules} + \text{New rules above}$

And so on at rank 4, 5, ...

Definition a/f^q is *reducible at rank r* if there exists b such that

$$\frac{a}{f^q} \xrightarrow{\text{rk}_r} \frac{b}{f^{q-1}}.$$

EXAMPLE

What do the new reduction rules bring?

$$f = 2xyz(w-x)(w-y)(w-z) - w^3(w^3 - w^2z + xyz) \quad (n=3)$$

$e(q, r) \stackrel{\text{def}}{=} \#$ of independent rational homogeneous functions a/f^q which are not reducible at rank r

q	0	1	2	3	4	$q > 4$
$e(q, 0)$	0	10	165	680	1771	$\sim 36q^3$
$e(q, 1)$	0	10	86	102	120	$\sim 18q$
$e(q, 2)$	0	10	7	6	6	6
$e(q, 3)$	0	9	1	0	0	0

Interpretation of the reduction rules in terms of a *spectral sequence*.

THEOREM

Without assumptions on f , there exists an integer r depending only on f such that for all homogeneous fraction a/f^q of degree $-n-1$:

- If $q > n$, then a/f^q is reducible at rank r .
- If $\oint \frac{a}{f^q} dz = 0$ then a/f^q is reducible at rank r .

[Compare with Griffiths-Dwork theorem.]

ALGORITHM, GENERAL CASE

1. Set $r = 1$.

2. For each $k \geq 0$, compute b_k , if possible, such that

$$\frac{\partial^k a}{\partial t^k f} \xrightarrow{\text{rk}_r} \cdots \xrightarrow{\text{rk}_r} \frac{b_k}{f^n}.$$

If not possible, increase r and start over.

[Always possible for r large enough, by the theorem above.]

3. Find a linear dependency relation over $\mathbb{K} = \mathbb{C}(t)$:

$$c_0(t) \frac{b_0}{f^n} + \cdots + c_m(t) \frac{b_m}{f^n} = 0.$$

[Exists for m large enough, by the confinement property.]

4. When found, return $c_m(t)y^{(m)}(t) + \cdots + c_0(t)y(t) = 0$.

[No a priori bound on r . Conjecture: r stays $\leq n+1$.]

FAST REDUCTION

Rank 1 – Gröbner basis technology, applied to the ideal $\langle \partial_0 f, \dots, \partial_n f \rangle$
Rank 2 – With Gröbner basis, computation of Koszul homology:

$$\frac{\text{syzygies}}{\text{trivial syzygies}} = \frac{\{ (b_0, \dots, b_n) \in \mathbb{K}[z_0, \dots, z_n]^{n+1} \mid \sum_{i=0}^n b_i \partial_i f = 0 \}}{\mathbb{K}[z_0, \dots, z_n] \left\langle \underbrace{\dots}_{i^{\text{th}} \text{ pos.}}, \dots, \underbrace{-\partial_i f}_{j^{\text{th}} \text{ pos.}}, \dots \right\rangle}$$

Trivial syzygies do not bring new rules.

When f is generic, all syzygies are trivial.

Rank ≥ 3 – Linear algebra