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Hybrid query plan generation

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Abstract. A hybrid query is a requirement of data produced by data services and a set of QoS preferences w.r.t. the query execution. In this paper we present the problem of the hybrid query optimization and, in particular, the generation of the search space of hybrid query plans. We show how the constraints that characterize the generation of hybrid query plans are modeled and validated by implementing them in an action language. We present graphs with the experiment results that show the complexity of this generation.

Keywords: Hybrid queries, answer set planning, multi-objective optimization, logic programming, non-monotonic reasoning

1 Introduction

A hybrid query [CV11] is a spatio-temporal query over continuous and on-demand data provided by data services and accessible through interfaces (API). In particular, we focus on the characteristics associated to the on-demand data. Data consumers express their data needs along their QoS preferences. For instance, consider the data need "Where are my friends at this moment?" with the QoS preferences "Privilege the execution price over the execution time and this in turn on the battery consumption".

Hybrid queries are evaluated with a service coordination approach [CVVC12] that is represented by a query workflow. A query workflow\footnote{We refer to query workflow and query plan indistinctly.} is a plan that coordinates a set of services with computing capabilities for processing the query. Services have QoS measures associated to their invocation and thus the execution of a query workflow has a cost that is the aggregation of the QoS measures. In centralized data bases the cost of query evaluation is the time, in distributed and parallel data bases the cost is two-folded in execution time and communication cost. The evaluation of hybrid queries in turn has a three dimensional cost: execution time, execution price, and battery consumption.

A finite set of query workflows implements the same hybrid query but there is a subset that satisfies the cost requirements of the user. For example, the
cheaper or the faster. This is a combinatorial optimization problem of minimizing
the cost of a query workflow QW that is a subset of activities A that implement
data operations.

$$\min_{QW \subseteq A} \{ \text{cost}(QW) : QW \text{ is a query workflow} \}$$

The optimization of queries over data services must tackle challenges such
as autonomy, non-durability of data, absence of fine grain data statistics, etc. In
[OB04] the authors discuss some non-centralized query optimization approaches
such as data mediation systems [TRV95, IJKWY97] where the cost model only
considers the execution time. The evaluation of queries over web services pre-
sent in [SMWM06] proposes a service coordination approach. A service co-
ordination is optimized by ordering the service calls in a pipelined fashion and
by tuning the size of data. Nonetheless, the control over data size (i.e., data
chunks) and selectivity statistics is a strong assumption because of the autonomy
of services and the absence of data statistics in a real scenario. Another aspect to
consider during the optimization is the selection of the service instances, which
represents a factor for the service coordination cost. In [CAH05, WCSO08] the
optimal selection of services with a multidimensional cost is proposed for op-
timizing the complete service coordination. This selection is done by solving a
multiobjective assignment problem for a set of abstract services. However, the
possibility to modify the order of service invocations is not considered and this
is a key issue in the hybrid query optimization.

In this paper we focus on the generation of a search space of query work-
flows in order to get a finite set of hybrid query implementations. This work
is the foundation for adding a cost model in the future and be able to choose
the most suitable query workflow w.r.t. user’s preferences. This is important
because a search space without constraints is in the order of n! permutations in
the number of activities required for evaluating the query. We propose a set of
constraints modeled in an action language (e.g., DLV-K) in order to characterize
the construction of query workflows with sequential and parallel compositions,
and to get only the coherent ones and therefore to reduce the size of the search
space.

The paper is organized as follows. Section 2 gives the bases of the environ-
ments where hybrid queries take place and explains how the hybrid query plans
are modeled with the notion of query workflow. In Section 3 we describe a set
of constraints for modeling the query workflow generation using an action lan-
guage. In Section 4 we show experiment results that confirm the complexity of
the problem. Finally, we conclude this paper and present the future work in
Section 5.

2 Querying dynamic environments

A dynamic environments is a set of distributed services that provide on-demand
data, continuous data and computation functionalities. Queries over dynamic en-
vironments, so called Hybrid queries, are expressions of data requirements over
on-demand and continuous data provided by data services and whose execution is constrained by user’s preferences (e.g., price, battery consumption, execution time). For example, the spatio-temporal query presented in the previous section “Where are my friends at this moment?”. In this paper we are interested in the hybrid queries over on-demand data. For example, “What are the interests of my friend Joe”, where only on-demand data are considered.

In [CVVC12] the evaluation of hybrid queries with a service coordination approach is proposed. This approach assumes the existence of computation services with data processing capabilities instead of an off-the-shelf query engine. Service coordination models a plan for evaluating the hybrid query and is represented by a query workflow.

2.1 Data and computing services

A service is a software entity with an abstract functionality described by an Application Program Interface (API) with the specification of methods for accessing data. An operation has a name and a set of typed bound/free parameters [RSU95]. Each parameter has a role identified by a unique name in the method.

For simplicity, we illustrate APIs with the following form

```
service_name:method_name(T: B_1?, . . . , T: B_n?, T: F_1!, . . . , T: F_m!)
```

where the service identified by `service_name` has a method named `method_name` with n bound parameters and m free parameters. Bound and free parameters will be identified by question and exclamation marks respectively. Each parameter is typed by a data type T and has a unique name within in the method header.

From this general notion of service, we distinguish two service specializations. (1) On-demand data service provides data in a request-response fashion through synchronous data operations. (2) Computing service provides a set of methods to perform operations over data (e.g., distance estimation, ordering, correlation).

2.2 Query workflow

A query workflow represents a control flow among a set of activities. Each activity is a program that calls a data service method or a computing service method. The control flow is defined by workflow operators such as `sequence` or `parallel`.

A query workflow QW is modeled as a direct acyclic graph QW=(V, E, init, finish) where:

```
V is a set of vertices
E ⊆ V × V is a set of edges
A ⊆ V is a set of activities named by unique names
init, finish ∈ A are the initial init and finish finish activities of the QW
O ⊆ V is a set of workflow operators of the type op = {seq, par} tagged by unique labels.
```

Data types are defined in the tuple data model in the thesis [CV11].
There are three activity specializations: the activity performs a service method invocation and always has a previous and a next activity, the run activity has no previous activity and its only goal is to launch the first activity of query workflow, the finish activity has no next activity and its only goal is to stop the query workflow execution after the last activity.

The compositions of activities $E \subseteq V \times V$ are defined by two workflow operators: sequential and parallel ones. A sequential composition $\text{seq}$ represents the total order of a pair of activities $a_1, a_2$ such that $e_1 = (a_1, a_2)$. It means that the second activity $a_2$ must be executed after $a_1$. A parallel composition $\text{par}$ represents the partial order of activities $a_1, a_2, a_3, a_4$ such that $e_1 = (a_1, a_2), e_2 = (a_1, a_3), e_3 = (a_2, a_4), e_4 = (a_3, a_4)$. It means that $a_2$ and $a_3$ are independent from each other and are executed in parallel after $a_1$ and before $a_4$.

In the next section we present how we model the query workflow generation by abstracting the semantics of activities [CV11]. This semantics is inherited from the data operations implemented by the services.

3 Query workflow generation

Query workflow generation is the process of enumerating all the query workflows in order to have a search space for choosing a subset that satisfies the user's preferences. The enumeration without constraints is in the order of $n!$, and, in fact, the most of these query workflows would not be coherent w.r.t. the hybrid query.

The generation is done by constructing query workflows whose activities are subject to query workflow constraints. We model these constraints as action rules in the language DLV-K\textsuperscript{6} that allows the sequential or parallel execution of query workflow activities.

In DLV-K, planning problems have a set of facts that represents the problem domain named background knowledge. The facts are predicates of static knowledge and are the input of the planning problem. Planning problems are modeled as state machines described by a set of fluents and a set of actions. A fluent is a property of an object in the world and is part of the states of the world [Bar03]. Fluents may be true, false or unknown. An action is executable if a precondition holds in the current state. Once an action is executed, the fluents and thus the state of the plan are modified. The action rules define the subset of fluents that must be held before the execution of action (i.e., preconditions) and the subset of fluents to be held after the execution (i.e., postconditions). Finally, a goal is a set of fluents that must be reached at the end of plan. A goal is expressed by the conjunction of fluents and by a plan length $l \in \mathbb{Z}_+$.

The mapping from query workflows generation to a planning problem is direct, as shown in Table 1. The APIs and the hybrid query are modeled as facts of the background knowledge. The execution state of a query workflow is modeled as fluents and query workflow activities as actions.

\footnotetext[6]{{http://www.dbai.tuwien.ac.at/proj/dlv/k}}
In the following, we show, along with a simple example, how we represent the background knowledge for query plan generation. Afterwards, we show how the query workflow state and activities are expressed in DLV-K.

3.1 Background knowledge

The background knowledge is the input for query workflow generation and it is represented by a set of facts. The facts are two-folded in (1) service interface representation and (2) hybrid query representation.

Service interface Consider the following three methods of service interfaces:

- \texttt{profile} :: \texttt{profile(string:nickname?, int:age!, string:sex!, string:email!)}
  This method provides the user profile formed by age, sex and email of a given user with a nickname.

- \texttt{interests} :: \texttt{interests(string:nickname?, string:tag!, real:score!)}
  This method provides the interests of a given user identified by a nickname. An interest is represented by a tag and has a score that qualifies the interest over the tag.

The service interface and its methods are described by the facts \texttt{service interface/1} and the \texttt{operation/2} respectively. We distinguish between bound and free parameters with the facts \texttt{bound p/4} and \texttt{free p/4}. The normal form of a parameter is stated by \texttt{parameter/4} and this rule guarantees that the parameter name is unique in the context of the service method.

\begin{verbatim}
    service interface(profile).
    operation(profile, profile).
    bound p(profile, profile, nickname, string).
    free p(profile, profile, age, int).
    free p(profile, profile, sex, string).
    service interface(interests).
    operation(interests, interests).
    bound p(interests, interests, nickname, string).
    free p(interests, interests, tag, string).
    free b(interests, interests, score, real).
    parameter(DSN, ON, PN, T) := bound p(DSN, ON, PN, T).
    parameter(DSN, ON, PN, T) := free p(DSN, ON, PN, T).
\end{verbatim}

Listing 1.1. Service interfaces
Hybrid query Now consider the query "What are the interests of my friend Joe" that is represented by the hybrid query in facts:

\[
\begin{align*}
&\text{project } (p, \text{nickname}, n). \quad 1 \\
&\text{project } (i, \text{score}, s). \quad 2 \\
&\text{project } (i, \text{tag}, t). \quad 3 \\
&\text{retrieve } (\text{profile}, \text{profile}, p). \quad 4 \\
&\text{retrieve } (\text{interests}, \text{interests}, i). \quad 5 \\
&\text{select } (p, \text{nickname}). \quad 6 \\
&\text{join } (p, \text{nickname}, i, \text{nickname}). \quad 7
\end{align*}
\]

Listing 1.2. Hybrid query in facts

This query expresses the need of data over the methods \texttt{profile:profile} and \texttt{interests:interests} methods, the nickname of the profile is filtered and correlated with the nickname of interests. Finally, the parameters nickname, score and tag are projected.

Observe that the selection over the nickname attribute is indicated only in intention because the equality operators are not significant for the query workflow generation.

3.2 Query workflow activities

The query workflow activities are represented as actions in DLV-K. Below we present the intuition of activities, for more details about the semantics of activities please refer to [CV11]. In general, the activities are predicates that hold if their facts from background knowledge are true. There are also activities that are independent from facts.

\textbf{init and finish} These activities have the special purpose to \texttt{init} and \texttt{finish} the query workflow execution accordingly with our model in subsection 2.2. Thus the semantics is not associated with the query processing and there is no dependency with the background knowledge.

\textbf{on-demand} This activity establishes a connection with a data service method and requires a method of a service and the expressed need of the user to be queried.

\[
\text{on-demand(DS) requires operation (DSN,ON), retrieve (DSN,ON,DS)}. \quad 1
\]

Listing 1.3. on-demand activity

\textbf{bind\_selection} This activity invokes and retrieves data from a service method. The invocation is done by a given bound parameter valid in the service interface definition. The hybrid query must express that data are required from this service method respect to a selected bound parameter.

\[
\begin{align*}
&\text{bind\_selection(DS,BP) requires operation (DSN,ON), retrieve (DSN,ON,DS), bound\_p (DSN,ON,BP), select (DS,BP)}. \quad 1
\end{align*}
\]

Listing 1.4. bind-selection activity

\textbf{bind\_join} The required correlation in the hybrid query is performed by this activity. Correlation is binary between data from two service methods on a parameter from each one. The parameter from the outer method must be
bound. This activity is analogous to `bind_selection` but `bind_join` takes the value of a bound parameter from the output of another method (i.e., the inner method).

```plaintext
bind_join(DS1, P1, DS2, BP2)
  requires operation(DSN1, ON1), retrieve_one(DSN1, ON1, DS1),
  parameter(DSN1, ON1, P1, _),
  operation(DSN2, ON2), retrieve_one(DSN2, ON2, DS2),
  bound(DSN2, ON2, BP2, _),
  join(DS1, P1, DS2, BP2).
```

**Listing 1.5. bind-join activity**

`select` The select activity performs the filtering over a valid parameter of a required service method in the hybrid query.

```plaintext
select(DS, P) requires operation(DSN, ON),
  retrieve_one(DSN, ON, DS),
  parameter(DSN, ON, P, _),
  select(DS, P).
```

**Listing 1.6. select activity**

`project` This activity projects a parameter of a service method required in the hybrid query.

```plaintext
project(DS, P) requires project_one(DS, P, _).
```

**Listing 1.7. project activity**

The semantics of activities described above is completed with constraints that define their preconditions and postconditions.

### 3.3 Query workflow constraints

The query workflow constraints define the conditions associated to the execution of activities. A condition is a state of knowledge modifiable by the execution of the activities. Conditions also define when the activities can be executed. Hence, we talk about preconditions and postconditions of activities. A state is composed by a set of fluents that occur during the execution of activities. The rules that define when a fluent occurs can be static or dynamic. Static rules are those that occur given the truth-value of a subset of fluents, and dynamic rules occur given the truth-value of a subset of fluents and after the execution of an activity.

Below we present the constraints that describe how the activities can be performed and how a query plan should be constructed.

- `init` and `finish` The first executable action in a query workflow is `init`. This activity has no previous activity, thus its precondition is that the query workflow is not `initiated` and produces a new state with `initiated`. The last activity is `finish` and there is no other activity executed after this one. The precondition to execute `finish` is that there is not evidence that the query workflow is `finished` and the data is already `delivered` (See output constraint below for details about `delivered`). The postcondition of `finish` is `finished` and this is always the goal for the generation.

123
executable init if ~initiated.  
caused initiated after init.  
executable finish if not finished, delivered.  
caused finished after finish.

Listing 1.8. init and finish activities

- **on-demand** Once initiated the plan, the data services must be connected(DS).
  This fluent is produced by the execution of on-demand(DS) activity.

  executable on-demand(DS) if initiated.  
  caused connected(DS) after on-demand(DS).

Listing 1.9. on-demand activity

During the execution, all data services must be connected. Therefore, there
is a fluent all_connected that is false if there is not evidence that a data service
is connected. Otherwise, it is true.

caused all_connected if not connected(DS).  
caused all_connected if not all_connected.

Listing 1.10. all_connected fluent

- **bind_selection** One of the activities that implements data retrieval is bind
  selection. It is only executable if there is not evidence that the data service
DS has already been retrieved and if there is a connection with DS. Once the
bind selection is executed, the fluent retrieved(DS) is true. retrieved(DS) is
an inertial fluent, thus the re-execution is not possible.

  executable bind_selection(DS,BP)  
  if not retrieved(DS), connected(DS).  
  caused retrieved(DS) after bind_selection(DS,BP).

Listing 1.11. bind_selection activity

- **selection** The filtering of data is done by the selection activity. It is exec-
  utable if there is not evidence that the parameter P of DS has already been
selected. There is also need that the data from DS has been retrieved and
obviously the selection select_(DS,P) must be part of the hybrid query.
  The execution of the selection makes the fluent true selected(DS,P) and it
  is inertial, so the re-execution of the selection select(DS,P) is not possible.

  executable select(DS,P) if not selected(DS,P),
  retrieved(DS), select_(DS,P).  
  caused selected(DS,P) after select(DS,P).

Listing 1.12. select activity

In order to be aware of the state of selection over all the required parameters
of a method DS, the all_selected_from(DSNameName) becomes true if there is
no other parameter pending to be selected.

caused all_selected_from(DS) if not selected(DS,P), select_(DS,P).  
caused all_selected_from(DS) if not all_selected_from(DS),
  retrieved(DS).

Listing 1.13. all_select_from fluent

Analogously, all parameters from all data services must be selected when
they are required. Therefore, there is the fluent all_selected. This fluent is
ture if there is no other method DS pending to be selected.
caused all_selected if all_selected_from(DS), select_(DS,P).
caused all_selected if all_selected_from(DS), not select_(DS,P),
  parameter (DSN,ON,P, _), retrieve_(DSN,ON,DS).
caused all_selected if not all_selected.

Listing 1.14. all_selected fluent

projection This activity is executable if there is not evidence that the parame-
ter P of DSOName has been projected. There is also need that the data
from DSOName be retrieved. The execution of projection makes the fluent
projected(DSOName,P) true and it is inertial, thus the re-execution of the
projection is not possible.

executable project(DS,P) if not projected(DS,P), retrieved(DS),
  project_(DS,P,_).
caused projected(DS,P) after project(DS,P).

Listing 1.15. project activity

The projected fluent is true once the action project(DS,P) is done. During
the query workflow execution, all the required parameters must be projected.
For DS grain the fluent all_projected_from(DS) is true if there is no other
parameter from DS pending to be projected. For the entire query, the fluent
all_projected is true if there is no other DS pending to be projected.

caused all_projected_from(DS) if not projected(DS,P), project_(DS,P,1
  _).
caused all_projected_from(DS) if not all_projected_from(DS) after
  project(DS,P).
caused all_projected if all_projected_from(DS), project_(DS,P,_).
caused all_projected if not all_projected.

Listing 1.16. all_projected_from fluent

output Once the hybrid query is processed, data must be delivered by the ac-
tivity output. In order to model this precondition, the fluent query_processed
is true if all the other possible activities have been processed. Otherwise, the
fluent is query_processed.

caused query_processed if not all_connected.
caused query_processed if not all_retrieved.
caused query_processed if not all_selected.
caused query_processed if not all_projected.
caused query_processed if not query_processed.

Listing 1.17. query_processed fluent

Once the query is processed, the output activity delivers the result and the
fluent delivered becomes true.

executable output if query_processed, not delivered.
caused delivered after output.

Listing 1.18. output fluent

4 Experiments

We performed experiments with a set of hybrid queries described below. The
objective of the experiments is to measure the size of the search space of query
workflows given a fixed number of activities and a goal to define the length of the query workflow.

4.1 Configuration

In order to test the performance of query workflow generation, we have defined four hybrid queries. All are based on the query example of subsection 3.1. In order to make more complex the generation of query plans, we add two data operations. The resulting queries are configured as follows:

<table>
<thead>
<tr>
<th>Query</th>
<th>$\text{union}_p \in P$</th>
<th>$\text{union}_p \in S$</th>
<th>Description</th>
<th># activities</th>
</tr>
</thead>
<tbody>
<tr>
<td>Q1</td>
<td>$\times$</td>
<td>$\times$</td>
<td>no additional operators</td>
<td>11</td>
</tr>
<tr>
<td>Q2</td>
<td>$\checkmark$</td>
<td>$\times$</td>
<td>+-unary data operator</td>
<td>12</td>
</tr>
<tr>
<td>Q3</td>
<td>$\times$</td>
<td>$\checkmark$</td>
<td>+-binary operator</td>
<td>13</td>
</tr>
<tr>
<td>Q4</td>
<td>$\checkmark$</td>
<td>$\checkmark$</td>
<td>+-both unary and binary operator</td>
<td>14</td>
</tr>
</tbody>
</table>

Table 2. Query configuration

Another dimension for the test is the max length of plans. The generated plans depend on a required plan length. In DLV-K, one must specifies the desired length that determines the size of the search space of query workflows and thus the required time for the generation.

In the tests we used lengths from 6 to 14, and we will measure (1) the length of the generated plans, (2) the size of the search space, and (3) the execution time of the generation.

4.2 Results

The data retrieved from the experiments show (1) the behavior of the search space growth during the generation and (2) the time required for generating all the search space.

Each hybrid query requires a fixed number of activities accordingly with the number of data operators required to process the query result. Given that the length $l$ bounds the size of query plans, there are hybrid queries that require a largest $l$ than others. For example, Q1 with 11 activities, requires at least an $l = 7$ to get query workflows with the most possible parallel compositions, and at most $l = 11$ to get completely sequential compositions. For reasons of space we invite the reader to visit the URLs http://goo.gl/XKZuL and http://goo.gl/z4iu3 to appreciate examples of query workflows of 7-length and 11-length respectively. The growth of search space is shown in Table 3 and, as was expected, the size of search space tends to be stable once the maximum length is reached. For example, Q1 reach the maximum length with $l = 11$. This is analogous for Q2, Q3, and Q4. Nevertheless the size of search space for each one is considerably
bigger than the others less complex. In the Figure 1(a), you can see that the search spaces have an exponential growth until the max length.

Besides the size of search space, the time for processing the query workflow generation is also exponential (See 1(b)) and it is not feasible to generate completely the search space in the context of query optimization. Thus, this enumeration must be done implicitly in order to avoid the combinatorial explosion.

5 Conclusion and future work

In this paper we presented the constraints for generating query workflows that model the plans of hybrid queries. The constraints represent the semantics of
query workflow activities that is analogous to the data operation implemented by the associated service method. We used the action language DLV-K for modeling the planning problem of query workflow generation in order to get a first (and naive) search space generation towards the optimization of hybrid queries. The results show that time complexity for query workflow generation is exponential.

The constraints must be extended in order to consider the continuous data operators for completely modeling hybrid queries. These constraints can be taken to model a more sophisticated query workflow generation. This implies the abstraction of data operators types and their associated constraints in a general form.

Theoretically, query workflow generation can be modeled mathematically by taking the constraints presented here and modeling an objective function. Additionally, the equivalence and coherence of query workflows must be proved.

References


