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The canonical crystal isomorphism in Fock spaces

Thomas Gerber *

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Abstract

We introduce an analogue of Schensted's algorithm for affine type A, using its natural interpretation in terms of crystals. In fact, we make explicit a particular crystal graph isomorphism between connected components of Fock spaces representations of $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_e)$. This so-called "canonical" crystal isomorphism turns out to be expressible only in terms of:

- Schensted's bumping procedure,
- the cyclage isomorphism defined in [9],
- a new crystal isomorphism acting on cylindric multipartitions.

1 Introduction

To understand the representations of the quantum algebras of affine type *A*, Jimbo, Misra, Miwa and Okado introduced in [11] Fock spaces of arbitrary level *l*. These are vectors spaces over $\mathbb{C}(q)$ depending on a parameter $\mathbf{s} \in \mathbb{Z}^l$, with basis the set of all *l*-partitions. We denote by $\mathcal{F}_{\mathbf{s}}$ such a space. They made explicit an action of the quantum algebra $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_e)$ which endowed $\mathcal{F}_{\mathbf{s}}$ with the structure of an integrable $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_e)$ -module. In addition, there is a natural subrepresentation $V(\mathbf{s})$ of $\mathcal{F}_{\mathbf{s}}$, the one generated by the empty *l*-partition, that gives a concrete realisation of any abstract highest weight $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_e)$ -module, provided \mathbf{s} is suitably chosen.

According to the works of Kashiwara [13], these abstract modules have a nice underlying combinatorial structure: their crystal. In particular, we can define the so-called crystal operators, crystal graph, crystal basis, and global basis for such modules. The combinatorial nature of Fock spaces provides a convenient framework for the study of the crystal structure of the highest weight $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_e)$ -modules. Actually, the whole space \mathcal{F}_s has a crystal structure, and the features of the highest weight $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_e)$ -representations (most interestingly the global basis) also exist more generally for the whole module \mathcal{F}_s , see e.g. [22], [23] or [17].

There exists a relation between this crystal structure on Fock spaces and the modular representation theory of complex reflection groups. Ariki has proved in [1] the LLT conjecture [15], enabling the computation of the decomposition matrices for Hecke algebras of type G(l, 1, n) (also known as Ariki-Koike algebras) via the matrices of the canonical basis of level *l* Fock spaces. In particular, it is known that the crystal structure on V(s) is isomorphic to the structure induced by the restriction and induction functors on the set of irreducible representations of the corresponding Ariki-Koike algebra \mathcal{H} . According to [21], the set of irreducible representations of any rational Cherednik algebra \mathcal{H} also has a crystal structure, given, again, by the restriction and induction functors, which is isomorphic to the crystal of the whole \mathcal{F}_s (for the appropriate s). As noticed by Gordon and Martino in [6, Section 4.15], the *KZ* functor (which maps a \mathcal{H} -module to a \mathcal{H} -module in a compatible manner) has an image whose labeling set is exactly the set of irreducible \mathcal{H} -modules could be explicited and proved to be identified with that of \mathcal{F}_s . If so, the main result of this paper, namely Theorem 5.26, will make explicit the action of the *KZ* functor on *l*-partitions.

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Indeed, this paper is concerned about the more specific study of crystal graphs within Fock spaces. The notion of crystal graph has been first introduced and studied by Kashiwara in [13], but see also, for instance, [14], [19], [16], [9], or [10]. The whole Fock space has a crystal graph, consisting of several connected components, each of which is entirely determined (and hence parametrised) by its so-called highest weight vertex. In particular, the connected component with highest weight vertex the empty l-partition is the crystal graph of the highest weight representation V(s). What we call a crystal isomorphism within Fock spaces is essentially a graph isomorphism. We demand that it maps the highest weight vertex of some connected component to the highest weight vertex of some other connected component (a priori of a different Fock space), and that it intertwines the structure of oriented colored graph. Because of the ubiquitous combinatorics in Fock spaces, we expect these crystal isomorphisms to have an nice combinatorial description.

In their papers [9] and [10], Jacon and Lecouvey have gathered some information about such mappings. Indeed, in [9], in order to generalise the results of [8] for l = 2, they described the crystal isomorphisms mapping the crystal graph of $V(\mathbf{s})$ to that of $V(\mathbf{r})$, when \mathbf{s} and \mathbf{r} are in the same orbit under the action of $\widehat{\mathfrak{S}}_l$. In [10], they explained how an arbitrary crystal isomorphism should act on a highest weight vertex.

The point of this article is to determine crystal isomorphisms in Fock spaces in a more general setting. By Proposition 2.16, each connected component of the crystal graph of \mathcal{F}_s is isomorphic to the crystal of some $V(\mathbf{r})$, where \mathbf{r} belongs to a particular fundamental domain for the action of $\widehat{\mathfrak{S}}_l$ for which a simple combinatorial description of this crystal is known (its vertices are the "FLOTW" *l*-partitions defined in 2.3). The question is then to determine the associated crystal isomorphism Φ . Precisely, we aim to express explicitly and combinatorially the action of Φ on any $\lambda \in \mathcal{F}_s$. This particular mapping is called the *canonical* $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_e)$ -crystal isomorphism (Definition 2.16). Drawing a parallel with the non-affine case (representations of $\mathcal{U}_q(\mathfrak{sl}_e)$, or $\mathcal{U}_q(\mathfrak{sl}_\infty)$), we can then regard the construction $\lambda \mapsto \Phi(\lambda)$ as an affine analogue of Schensted's bumping algorithm. Indeed, in the case of regular type A, it is known (Theorem 3.1) that the canonical crystal isomorphism is exactly the bumping procedure. Roughly speaking, we replace the Schensted algorithm on semistandard tableaux by a rectification procedure on symbols (also called keys or tabloids in the literature) yielding, in turn, semistandard, cylindric symbols, and FLOTW symbols.

The paper is organized as follows. Section 2 contains the basic notations and definitions needed to handle the combinatorics of crystals in Fock spaces, as well as a quick review on the theory of $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -representations, and more particularly the Fock space representations. We also formally express the problem we wish to solve, which translates to finding the canonical $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -crystal isomorphism.

In Section 3, we first roughly explain how the limit case $e \to \infty$ gives another structure on Fock spaces, namely that of a $\mathcal{U}_q(\mathfrak{sl}_\infty)$ -module. We then recall some classic results about $\mathcal{U}_q(\mathfrak{sl}_\infty)$ -crystals (and their relation to Schensted's bumping procedure) and explain how this solves our problem when Fock spaces are considered as $\mathcal{U}_q(\mathfrak{sl}_\infty)$ -modules. In this perspective, the canonical $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_e)$ -crystal isomorphism we construct in the last section can be seen as an affine version of Schensted's insertion.

Then, we show in Section 4 that both structures are compatible. More precisely, any $\mathcal{U}_q(\mathfrak{sl}_\infty)$ -crystal isomorphism is also a $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_e)$ -crystal isomorphism (Corollary ??). Using this, we determine a way to restrict ourselves only to the study of so-called cylindric multi-partitions. This requires the results about $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_e)$ -crystal isomorphisms obtained in [9].

Section 5 treats the case of cylindric multipartitions. The key ingredient is Theorem 5.20, and the expected canonical crystal isomorphism is described in Theorem 5.26. Together with Section 4, this eventually enables the determination of the canonical crystal isomorphism Φ in full generality.

Finally, Section 6 is an application of these results. We deduce a direct way to compute the highest weight vertex of any connected component of a crystal without using the crystal operators \tilde{e}_i to go back up the crystal graph. Note that this requires the invertibility of the map Φ . This is achieved by adding some "recording data" to the construction of Φ , in a very similar way to the whole (one-to-one) Robinson-Schensted-Knuth correspondence.

2 Formulation of the problem

2.1 Generalities

We recall the usual notations and definitions about the combinatorics of multipartitions.

Fix $n \in \mathbb{N}$ and $l \in \mathbb{Z}_{>0}$. A partition λ of n is a sequence $(\lambda_1, \lambda_2, ...)$ such that $\lambda_a \ge \lambda_{a+1}$ for all a, and $\sum_a \lambda_a = n$. The integer n is then called the *rank* of λ and denoted by $|\lambda|$. Each λ_a is called a part of λ . For convenience, we often consider that a partition has infinitely many parts equal to zero. The set of partitions of n is denoted by $\Pi(n)$. We also set $h(\lambda) = \max_{\lambda_a \neq 0} a$, and call $h(\lambda)$ the *height* of λ . Out of simplicity, we use the multiplicative notation for a partition λ . For instance, $(4, 2, 2, 1, 0, ...) = (4.2^2.1)$. A *l*-partition (or simply multipartition) λ of n is a *l*-tuple of partitions λ^c , for $c \in [\![1, l]\!]$ such that $\sum_{c=1}^l |\lambda^c| = n$. The integer n is again called the *rank* of λ and denoted by $|\lambda|$; and we write $\Pi_l(n)$ for the set of *l*-partitions of n. Moreover, we denote by \emptyset the partition (0, 0, ...) and by \emptyset the multipartition $(\emptyset, \dots, \emptyset)$.

A multipartition λ is often identified with its Young diagram [λ]:

$$[\lambda] := \{ (a, b, c) ; a \ge 1, c \in [[1, l]], b \in [[1, \lambda_a^c]] \}$$

In turn, the Young diagram of λ is itself depicted by a *l*-tuple of array where the *c*-th array is the superposition of λ_a^c boxes, $a \ge 1$. For example, $[(1, 2^2, \emptyset)] = (\Box, \Box, \emptyset)$. Each box in this diagram is then labeled by an element (a, b, c) of $[\lambda]$. These elements are called the *nodes* of λ . A node γ of λ is said to be *removable* (or a node *of type R*) if $[\lambda] \setminus \{\gamma\}$ is still the Young diagram of some *l*-partition μ . In this case, we also call γ an *addable* node (or a node *of type* A) of μ .

A *l*-charge (or simply multicharge) is a *l*-tuple $\mathbf{s} = (s_1, \dots, s_l)$ of integers. A charged multipartition is a formal pair denoted $|\lambda, \mathbf{s}\rangle$, where λ is a *l*-partition and \mathbf{s} is a *l*-charge. We can then define the *content* of a node $\gamma = (a, b, c)$ of λ charged by \mathbf{s} as follows:

$$\operatorname{cont}_{\lambda}(\gamma) = b - a + s_c.$$

Also, given $e \in \mathbb{Z}_{>1}$, we define the *residue* of γ by

$$\operatorname{res}_{\lambda}(\gamma) = \operatorname{cont}_{\lambda}(\gamma) \mod e.$$

Let $i \in [[0, e - 1]]$. A node γ of λ is called an *i*-node if $\operatorname{res}_{\lambda}(\gamma) = i$. If $\gamma = (a, b, c)$ is a node of λ , denote by γ^+ the node (a, b + 1, c), that is the node located on the right of γ . Similarly, denote by γ^- the node (a, b - 1, c), the node located on the left of γ . Of course, $\mu = \gamma^+ \Leftrightarrow \gamma = \mu^-$. Also, if γ is an *i*-node of type R, γ^+ is an (i + 1)-node of type A

We can represent each charged multipartition $|\lambda, s\rangle$ by its Young diagram whose boxes are filled by the associated contents. For instance,

$$|(4.1, 2^2, 3), (0, 3, -2)\rangle = \left(\begin{array}{c} 0 & 1 & 2 & 3 \\ \hline 1 & 1 & 2 & 3 \end{array}, \begin{array}{c} 3 & 4 & 4 \\ \hline 2 & 3 & 3 & 4 \\ \hline 2 & 3 & 3 & 4 \end{array} \right).$$

There is an equivalent way to represent charged multipartitions by another combinatorial object, namely by the so-called *symbols*. Let us recall their definition, following [5]. Having fixed a multipartition λ and a multicharge s, take $p \ge \max_c(1 - s_c + h(\lambda^c))$. We can first define

$$\mathfrak{b}_a^c(\boldsymbol{\lambda}) := \lambda_a^c - a + p + s_c,$$

for $1 \le c \le l$ and $1 \le a \le p + s_c$. We then set $\mathfrak{B}^c_{\mathbf{s}}(\lambda) = (\mathfrak{b}^c_{p+s_c}(\lambda), \dots, \mathfrak{b}^c_1(\lambda))$, and the **s**-symbol of λ of size p is the following l-tuple:

$$\mathfrak{B}_{\mathbf{s}}(\boldsymbol{\lambda}) = (\mathfrak{B}_{\mathbf{s}}^{1}(\boldsymbol{\lambda}), \dots, \mathfrak{B}_{\mathbf{s}}^{l}(\boldsymbol{\lambda})).$$

It is pictured in an array whose *c*-th row, numbered from bottom to top, is $\mathfrak{B}_{s}^{c}(\lambda)$. *Example* 2.1. Take $\lambda = (4.1, 2^{2}, 3)$ and $\mathbf{s} = (0, 3, -2)$. Then we can take p = 4, and we get $\mathfrak{B}_{s}^{1}(\lambda) = (0, 1, 3, 7), \mathfrak{B}_{s}^{2}(\lambda) = (0, 1, 2, 3, 4, 7, 8)$ and $\mathfrak{B}_{s}^{3}(\lambda) = (0, 4)$, which is represented by

$$\mathfrak{B}_{\mathbf{s}}(\boldsymbol{\lambda}) = \begin{pmatrix} 0 & 4 & & \\ 0 & 1 & 2 & 3 & 4 & 7 & 8 \\ 0 & 1 & 3 & 7 & & \end{pmatrix}$$

Note that one recovers the elements $b_a^c(\lambda)$ encoding non-zero parts by translating by p the contents of the rightmost nodes in the Young diagram. In fact, it is easy to see that both objects encode exactly the same data. Throughout this paper, we often switch between the classic approach of "Young diagrams with contents" on the one hand, and the use of symbols on the other hand.

Definition 2.2. Let $\mathbf{s} \in \mathbb{Z}^l$ and $\lambda \in \mathcal{F}_{\mathbf{s}}$. The symbol $\mathfrak{B}_{\mathbf{s}}(\lambda)$ is called *semistandard* if the three following conditions are satisfied:

- $s_1 \leq s_2 \leq \cdots \leq s_l$,
- the entries in each column of $\mathfrak{B}_{s}(\lambda)$ are non-decreasing,
- the entries in each row of $\mathfrak{B}_{s}(\lambda)$ are increasing.

For $e \in \mathbb{Z}_{>1}$, we denote

$$\mathscr{S}_e = \{ \mathbf{s} \in \mathbb{Z}^l \mid 0 \le s_c - s_{c'} < e \text{ for } c < c' \}.$$

We define a family of particular multipartitions, the *FLOTW multipartitions* (for Foda, Leclerc, Okado, Thibon, Welsh [3]). We will see in the Section 2.3 why this is an object of interest in this paper.

Definition 2.3. Let $\mathbf{s} \in \mathscr{S}_e$. A charged multipartition $|\lambda, \mathbf{s}\rangle$ is called FLOTW if:

- For all $1 \le c \le l-1$, $\lambda_a^c \ge \lambda_{a+s_{c+1}-s_c}^{c+1}$ for all $a \ge 1$; and $\lambda_a^l \ge \lambda_{a+e+s_1-s_l}$ for all $a \ge 1$.
- For all α > 0, the residues of the nodes (a, α, c) with λ^c_a = α (i.e. the rightmost nodes of the rows of length α of λ) do not cover [[0, e 1]].

If $|\lambda, s\rangle$ only verifies the first condition, we say that it is *cylindric*.

Remark 2.4. In particular, the symbol of a FLOTW *l*-partition is semistandard.

Multipartitions are the natural objects used to understand combinatorially the representations of the quantum algebras of affine type *A*. In the following section, we recall some definitions and facts about these algebras, their representations, before focusing in Section 2.3 on the crystal structure they are endowed with.

2.2 Fock spaces representations of quantum algebras

In the sequel, we fix $e \in \mathbb{Z}_{>1}$.

For the purpose of this article, it is not really necessary to review completely the theory of quantum algebras. The reader is invited to refer to [2] and [7] for detailed definitions and properties. However, we recall quickly below the theory of highest weight integrable representations of $\mathcal{U}'_{a}(\widehat{\mathfrak{sl}}_{e})$.

Essentially, the quantum algebra $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ is a $\mathbb{C}(q)$ -algebra which is a one-parameter deformation of the universal enveloping algebra of the affine Kac-Moody algebra $\widehat{\mathfrak{sl}_e}$, whose underlying Weyl group is the affine group of type A_{e-1} . It is defined by generators, denoted by $e_i, f_i, t_i^{\pm 1}$ for $i \in [0, e-1]$, and some relations. Besides, it is possible to endow it with a coproduct, which makes it a Hopf algebra.

Denote Λ_i , $i \in [[0, e-1]]$ and δ the fundamental weights for \mathfrak{sl}_e . Recall that the simple roots are then given by $\alpha_i := -\Lambda_{i-1 \mod e} + 2\Lambda_i - \Lambda_{i+1 \mod e}$ for $0 \le i \le e-1$. According to [2, Chapter 6], to each dominant integral weight Λ is associated a particular $\mathcal{U}'_q(\mathfrak{sl}_e)$ -module $V(\Lambda)$, called the *highest weight module of highest weight* Λ , verifying the following property:

Property 2.5. We have $V(\Lambda^{(1)}) \simeq V(\Lambda^{(2)})$ as $\mathcal{U}'_a(\widehat{\mathfrak{sl}_e})$ -modules if and only if $\Lambda^{(1)} - \Lambda^{(2)} \in \mathbb{Z}\delta$.

Moreover, provided we work in the category of so-called "integrable" $\mathcal{U}'_q(\mathfrak{sl}_e)$ -modules, it is known that:

- 1. For all dominant integral weight Λ , the module $V(\Lambda)$ is irreducible,
- 2. Each irreducible $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -module is isomorphic to some $V(\Lambda)$,
- 3. Each $\mathcal{U}'_{a}(\widehat{\mathfrak{sl}_{e}})$ -module is semisimple.

We now construct a practical representation of $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$, the Fock space representation, which turns out to be integrable and has nice combinatorial properties. In particular, this will yield a realisation of the highest weight $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -modules $V(\Lambda)$. In this perspective, we fix $l \in \mathbb{Z}_{>0}$.

Definition 2.6. Let $\mathbf{s} = (s_1, \ldots, s_l) \in \mathbb{Z}^l$. The Fock space associated to \mathbf{s} is the following $\mathbb{C}(q)$ -vector space:

$$\mathcal{F}_{\mathbf{s}} := \bigoplus_{n \in \mathbb{N}} \bigoplus_{\lambda \in \Pi_l(n)} \mathbb{C}(q) | \lambda, \mathbf{s} \rangle.$$

This vector space is made into an integrable $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -module via the action of the generators of $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ detailed in [11, Proposition 3.5]. We do not recall here what theses actions are, since it is not important in our context. However, with this definition of the action, it is easy to see that the element $|\emptyset, \mathbf{s}\rangle$ is a highest weight vector. Denote by $\Lambda(\mathbf{s})$ its weight. By [22, Section 4.2] or [23, Proposition 3.7], we now that for all $|\lambda, \mathbf{s}\rangle$ in $\mathcal{F}_{\mathbf{s}}$,

$$\operatorname{wt}(|\boldsymbol{\lambda}, \mathbf{s}\rangle) = \sum_{c=1}^{l} \Lambda_{s_c \operatorname{mod} e} - \sum_{i=0}^{e-1} M_i(\boldsymbol{\lambda}, \mathbf{s})\alpha_i - \Delta(\mathbf{s})\delta,$$
(1)

where $M_i(\lambda, \mathbf{s})$ is the number of *i*-nodes in λ , and where $\Delta(\mathbf{s})$ is a coefficient depending only on \mathbf{s} and e. In particular, we see that

$$\Lambda(\mathbf{s}) = \operatorname{wt}(|\mathbf{\emptyset}, \mathbf{s}\rangle) = \sum_{c=1}^{l} \Lambda_{s_c \operatorname{mod} e} - \Delta(\mathbf{s})\delta.$$
(2)

Consider the module

$$V(\mathbf{s}) := \mathcal{U}'_q(\widehat{\mathfrak{sl}_e}).|\boldsymbol{\emptyset},\mathbf{s}\rangle.$$

It is an irreducible integrable highest weight $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -module with highest weight $\Lambda(\mathbf{s})$. Hence, by the previous properties of $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -representations, we know that this module $V(\mathbf{s})$ is isomorphic to the "abstract" $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -module $V(\Lambda(\mathbf{s}))$.

Using Property 2.5 and Relation (2), we have a $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -module isomorphism between $V(\mathbf{s})$ and $V(\mathbf{r})$ whenever \mathbf{r} is equal to \mathbf{s} up to a permutation of its components and translations by a multiple of e. In fact, when two multicharges are related in such a way, we regard them as being in the same orbit under some group action. This is why we now introduce the *extended affine symmetric group* $\widehat{\mathfrak{S}}_l$. It is the group with the following presentation:

- Generators: σ_i , $i \in \llbracket 1, l-1 \rrbracket$ and y_i , $i \in \llbracket 1, l \rrbracket$.
- Relations :
 - $σ_i^2 = 1$ for *i* ∈ [[1, *l* − 1]],
 - $\sigma_i \sigma_{i+1} \sigma = \sigma_{i+1} \sigma_i \sigma_{i+1}$ for $i \in \llbracket 1, l-2 \rrbracket$,
 - $\sigma_i \sigma_j = \sigma_j \sigma_i$ if $i j \neq 1 \mod l$,
 - $y_i y_j = y_j y_i$ for $i, j \in [[1, l]]$,
 - $\sigma_i y_j = y_j \sigma_i$ for $i \in \llbracket 1, l-1 \rrbracket$ and $j \in \llbracket 1, l \rrbracket$ such that $j \neq i, i+1 \mod l$,
 - $σ_i y_i σ_i = y_{i+1}$ for *i* ∈ [[1, *l* − 1]].

It can be regarded as the semi-direct product $\mathbb{Z}^l \rtimes \mathfrak{S}_l$, by considering that the elements y_i form the standard basis of \mathbb{Z}^l , and that the elements σ_i are the usual generators of \mathfrak{S}_l (i.e. the transpositions (i, i + 1)).

Now, there is a natural action of $\widehat{\mathfrak{S}}_l$ on \mathbb{Z}^l as follows. Let $\mathbf{s} = (s_1, \dots, s_l) \in \mathbb{Z}^l$. We set

- σ_i .**s** = $(s_1, \ldots, s_{i-1}, s_{i+1}, s_i, \ldots, s_l)$ for all $i \in [[1, l-1]]$, and
- y_i .**s** = $(s_1, \ldots, s_{i-1}, s_i + e, s_{i+1}, \ldots, s_l)$ for all $i \in [[1, l]]$.

It is easy to see that a fundamental domain for this action is given by

$$\mathscr{D}_e = \{ \mathbf{s} \in \mathbb{Z}^l \mid 0 \le s_1 \le \cdots \le s_l < e \}.$$

Remark 2.7. Note that

$$\mathcal{D}_e \subset \mathcal{S}_e.$$

It is important not to be confused between \mathcal{D}_e and the set \mathcal{S}_e defined just before Definition 2.3. Indeed, it is sufficient to work in \mathcal{S}_e to have the explicit combinatorial caracterisation of the crystal of $V(\mathbf{s})$, but it is necessary to work with \mathcal{D}_e if we expect some unicity properties (as in upcoming Proposition 2.16).

With these notations, it is clear that $V(\mathbf{s}) \simeq V(\mathbf{r})$ as $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -modules if and only if \mathbf{r} and \mathbf{s} are in the same orbit under the action of $\widehat{\mathfrak{S}}_l$. Actually, there is a tight connection between these highest weight $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -representations $V(\mathbf{s})$ and the modular representations of Ariki-Koike algebras (the Hecke algebras of the complex reflection groups G(l, 1, n)), which are defined using the parameters e and \mathbf{s} , and which are invariant when \mathbf{s} varies inside a given orbit under the action of $\widehat{\mathfrak{S}}_l$. The main link is known as Ariki's theorem [2, Theorem 12.5],[1], which is a proof of the LLT conjecture [15]. A good review on this subject can also be found in the book of Geck and Jacon [5, Chapters 5 and 6].

From now on, we allow ourselves to write λ instead of $|\lambda, s\rangle$ for an element of a Fock space if there is no possible confusion on the multicharge.

2.3 Crystal isomorphisms and equivalent multipartitions

One of the most important features of Fock spaces is the existence of a *crystal*, in the sense of Kashiwara [13]. This yields a nice combinatorial structure, which is in particular encoded in the *crystal graph* of \mathcal{F}_s . Its definition requires an action of the so-called *crystal operators* \tilde{e}_i and \tilde{f}_i , for $0 \le i \le e - 1$, on the Fock space \mathcal{F}_s . We do not give their original definition, since the only result we need in the sequel is upcoming Theorem 2.9.

In order to determine this action, we introduce an order on the set of nodes of a charged multipartition $|\lambda, \mathbf{s}\rangle$. Let $|\lambda, \mathbf{s}\rangle \in \mathcal{F}_{\mathbf{s}}$ and $\gamma = (a, b, c)$ and $\gamma' = (a', b', c')$ be two removable or addable *i*-nodes of λ . We write

$$\gamma \prec_{\mathbf{s}} \gamma' \text{ if } \begin{cases} b-a+s_c < b'-a'+s_{c'} \text{ or} \\ b-a+s_c = b'-a'+s_{c'} \text{ and } c > c'. \end{cases}$$

Note that this order is needed to define the action of $\mathcal{U}'_{\alpha}(\widehat{\mathfrak{sl}_{e}})$ on \mathcal{F}_{s} in Section 2.2.

For all $i \in [0, e - 1]$, the *i*-word for λ is the sequence obtained by writing the addable and removable *i*-nodes of λ in increasing order with respect to this order \prec_s , each of them being encoded by a letter A if it is addable, and a letter R if it is removable. We denote it by $w_i(\lambda)$. The reduced *i*-word $\hat{w}_i(\lambda)$ is the word in the letters A and R obtained by deleting recursively all occurences RA in $w_i(\lambda)$. Hence $\hat{w}_i(\lambda) = A^{\alpha}R^{\beta}$ for some non-negative integers α and β . If it exists, the rightmost A (respectively the leftmost R) in $\hat{w}_i(\lambda)$ encodes a node which is called the good addable (respectively removable) *i*-node of λ .

Example 2.8. If $|\lambda, \mathbf{s}\rangle = \left(\frac{0}{123}, \frac{34}{23}, \frac{34}{23}, \frac{3}{2}\right), e = 3$, and i = 0, then $w_0(\lambda) = ARARR$ and therefore $\hat{w}_0(\lambda) = ARR$. The good addable 0-node of λ is thus (2, 1, 3), and the good removable 0-node is (2, 2, 2).

Theorem 2.9 ([11]). Let $\lambda \in \mathcal{F}_s$. The crystal operators act as follows:

• $\tilde{e}_i(\lambda) = \begin{cases} \lambda \setminus \{\gamma\} \\ 0 \end{cases}$	if γ is the good removable i-node of λ if λ has no good removable i-node,
• $\tilde{f}_i(\lambda) = \begin{cases} \lambda \cup \{\gamma\} \\ 0 \end{cases}$	if γ is the good addable i-node of λ if λ has no good addable i-node.

Remark 2.10. Clearly, because of Relation (1), we have wt($\tilde{e}_i(\lambda)$) = wt(λ) + α_i and wt($\tilde{f}_i(\lambda)$) = wt(λ) - α_i .

We can now define the *crystal graph* of λ .

Definition 2.11. The crystal graph $B(\lambda)$ of λ is the oriented colored graph with:

vertices : the multipartitions obtained from λ after applying any combination of the operators *e˜_i* and *f˜_i*.

• arrows : $\lambda \xrightarrow{i} \mu$ whenever $\mu = \tilde{f}_i(\lambda)$.

Theorem 2.9 says that when \tilde{e}_i acts non-trivially on λ (i.e. when λ has a good removable *i*-node), then \tilde{e}_i removes a node in λ . Hence, any sequence $\ldots \tilde{e}_{i_k} \tilde{e}_{i_{k-1}} \ldots \tilde{e}_{i_1}(\lambda)$ in $B(\lambda)$ has at most $|\lambda|$ elements. In particular, it is finite, and there is a sequence of maximal length *m* of operators \tilde{e}_{i_k} such that $\tilde{e}_{i_m} \tilde{e}_{i_{m-1}} \ldots \tilde{e}_{i_1}(\lambda)$ is a multipartition and $\tilde{e}_{i_{m+1}} \ldots \tilde{e}_{i_1}(\lambda) = 0$. Consider the multipartition

$$\boldsymbol{\lambda} := \tilde{e}_{i_m} \tilde{e}_{i_{m-1}} \dots \tilde{e}_{i_1}(\boldsymbol{\lambda})$$

It is not complicated to show that it does not depend on the maximal sequence of operators \tilde{e}_{i_k} chosen. In other words, all sequences give the same multipartition λ , which we call the *highest weight vertex* of $B(\lambda)$. Also, every vertex μ in $B(\lambda)$ writes $\mu = \tilde{f}_{i_p} \tilde{f}_{i_{p-1}} \dots \tilde{f}_{i_1}(\lambda)$ for some $p \in \mathbb{N}$. Therefore, any crystal graph is entirely determined by its highest weight vertex, and if we know this highest weight vertex, all other vertices are recursively computable. It turns out that in one particular case, we know an explicit combinatorial description of the vertices of $B(\lambda)$.

Notation: When $\lambda = \emptyset$, we write $B(\emptyset) =: B(\mathbf{s})$.

Theorem 2.12 ([3]). Let $\mathbf{s} \in \mathscr{S}_e$. The vertices of $B(\mathbf{s})$ are the FLOTW multipartitions.

Actually, for all $\mathbf{s} \in \mathbb{Z}^l$, $B(\mathbf{s})$ is the crystal graph of the irreducible highest weight $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_e)$ module $V(\mathbf{s})$ in the sense of Kashiwara [13]. Accordingly, the combinatorics of $B(\lambda)$ has some algebraic interpretation. In particular, the vertices of $B(\lambda)$ (the FLOTW multipartitions if $\mathbf{s} \in \mathscr{S}_e$) label both the so-called crystal and global (or canonical) basis of $V(\mathbf{s})$. We also refer to [5] for details.

We can also define a crystal graph for the whole Fock space \mathcal{F}_s . Its vertices are all the *l*-partitions, and the arrows are defined as in Definition 2.11. We denote it by $B(\mathcal{F}_s)$. It has several connected components, each of them being parametrised by its highest weight vertex. In other terms,

$$B(\mathcal{F}_{\mathbf{s}}) = \bigsqcup_{\lambda} B(\lambda), \tag{3}$$

where the union is taken over all highest weight vertices λ , i.e. vertices without any removable node.

We can now introduce the notion of crystal isomorphism.

Definition 2.13. Let $\lambda \in \mathcal{F}_s$ and $\mu \in \mathcal{F}_r$. A $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -crystal isomorphism is a map $\phi : B(\overset{\bullet}{\lambda}) \longrightarrow B(\overset{\bullet}{\mu})$ verifying:

- 1. $\phi(\lambda) = \dot{\mu}$,
- 2. $\phi \circ \tilde{f}_i = \tilde{f}_i \circ \phi$ whenever \tilde{f}_i acts non trivially.

By (b), the image of $B(\lambda)$ under ϕ is the whole crystal $B(\mu)$. In fact, this definition just says that ϕ intertwines the graph structures of $B(\lambda)$ and $B(\mu)$.

Definition 2.14. Let $|\lambda, \mathbf{s}\rangle \in \mathcal{F}_{\mathbf{s}}$ and $|\mu, \mathbf{r}\rangle \in \mathcal{F}_{\mathbf{r}}$. We say that $|\lambda, \mathbf{s}\rangle$ and $|\mu, \mathbf{r}\rangle$ (or simply λ and μ) are *equivalent* if there is a $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_e)$ -crystal isomorphism ϕ between $B(\lambda)$ and $B(\mu)$ such that $\phi(|\lambda, \mathbf{s}\rangle) = |\mu, \mathbf{r}\rangle$.

Remark 2.15. In other terms, λ and μ are equivalent if they appear at the same place in their respective crystal graphs.

The isomorphisms of $\mathcal{U}'_{q}(\widehat{\mathfrak{sl}}_{e})$ -modules $V(\mathbf{s}) \simeq V(\mathbf{r})$ whenever $\mathbf{r} \in \mathbf{s} \mod \widehat{\mathfrak{S}}_{l}$ (cf. Section 2.2) yield isomorphisms of crystal graphs between $B(\mathbf{s})$ and $B(\mathbf{r})$. There exist also other natural crystal isomorphisms.

Proposition-Definition 2.16. Let $\lambda \in \mathcal{F}_s$. There exists a unique l-charge $\mathbf{r} \in \mathcal{D}_e$ and a unique FLOTW l-partition $\boldsymbol{\mu} \in \mathcal{F}_{\mathbf{r}}$ such that $|\lambda, \mathbf{s}\rangle$ and $|\boldsymbol{\mu}, \mathbf{r}\rangle$ are equivalent. The associated $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -crystal isomorphism is called the canonical crystal isomorphism.

Proof. First of all, if $\lambda \in B(\mathbf{s})$ then, by the remark just above Proposition 2.16, there is a crystal isomorphism between $B(\mathbf{s})$ and $B(\mathbf{r})$ where \mathbf{r} is the representative of \mathbf{s} in the fundamental domain \mathcal{D}_e .

Suppose now that $|\lambda, \mathbf{s}\rangle \in \mathcal{F}_{\mathbf{s}}$ such that $B(\lambda) \neq B(\mathbf{s})$. This means that λ (as a vertex in its crystal graph) is not in the connected component whose highest weight vertex is the empty multipartition. Then there is a sequence (i_1, \ldots, i_m) such that $\tilde{e}_{i_m} \ldots \tilde{e}_{i_1}(\lambda) = \lambda$, the highest weight vertex in $B(\lambda)$. Write wt $(\lambda) = \sum_i a_i \Lambda_i + d\delta$ and define \mathbf{r} to be the increasing *l*-charge containing a_i occurences of *i*. In particular, $\mathbf{r} \in \mathcal{D}_e$. Then we have a natural crystal isomorphism $B(\lambda) \stackrel{\phi}{\simeq} B(\mathbf{r})$, and therefore there is a FLOTW *l*-partition $\mu := \phi(\lambda) = \tilde{f}_{i_1} \ldots \tilde{f}_{i_m}(\mathbf{0})$ in $B(\mathbf{r})$ equivalent to λ .

These elements are clearly unique, since \mathscr{D}_e is a fundamental domain for the action of $\widehat{\mathfrak{S}}_l$.

The goal of this paper is to find a direct and purely combinatorial way to determine this canonical crystal isomorphism, without having to determine the sequence of operators leading to the highest weight vertex and taking the reverse path in $B(|\emptyset, \mathbf{r}\rangle)$ as explained in the previous proof. Note that this question has been answered in [10] in the particular case where λ is a highest weight multipartition. The canonical crystal isomorphism in this case is the so-called "peeling procedure", and is much easier to describe. Of course, in this case, the canonical isomorphism maps λ to the empty *l*-partition.

3 The case $e = \infty$

In this section, we consider the particular (and easier) case where $e = \infty$. This means that we regard Fock spaces as $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ -modules, where $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ is defined as the direct limit of the quantum algebras $\mathcal{U}_q(\mathfrak{sl}_e)$. We refer e.g. to [12] for detailed background on $\mathcal{U}_q(\mathfrak{sl}_{\infty})$. Actually, there is an action of $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ on Fock spaces which generalises the action of $\mathcal{U}'_q(\mathfrak{sl}_e)$ when e tends to ∞ . In particular, \mathcal{F}_s is made into an integrable $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ -module, and all the properties of the $\mathcal{U}'_q(\mathfrak{sl}_e)$ -representation \mathcal{F}_s stated in Section 2.2 still hold for the $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ module structure. With this point of view, the algebra $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ is the natural way to extend $\mathcal{U}'_q(\mathfrak{sl}_e)$ when $e \to \infty$.

In this setting, the notion of being FLOTW for a multipartition simply translates to its symbol being semistandard, with an increasing multicharge. And as a matter of fact, we know a $\mathcal{U}_q(\mathfrak{sl}_\infty)$ -crystal isomorphism which associates to each multipartition a new multipartition whose symbol is semistandard. This is the point of the following section.

3.1 Schensted's bumping algorithm and solution of the problem

Let $\lambda \in \mathcal{F}_{s}$.

We first introduce the reading of the symbol $\mathfrak{B}_{\mathbf{s}}(\lambda)$. It is the word obtained by writing the elements of $\mathfrak{B}_{\mathbf{s}}(\lambda)$ from right to left, starting from the top row. Denote it by read (λ, \mathbf{s}) . The Robinson-Schensted insertion procedure (or simply Schensted procedure, or bumping procedure) enables to construct a semistandard symbol starting from such a word. We only recall it on an example (3.3) below, see also Example 6.1 in Section 6. For proper background, the reader can refer to e.g. [4] or [20]. Denote by $\mathscr{P}(\operatorname{read}(\lambda, \mathbf{s}))$ the semistandard symbol obtained from $\operatorname{read}(\lambda, \mathbf{s})$ applying this insertion procedure. Finally, we set $\mathbf{RS}(\mathbf{s})$ and $\mathbf{RS}(\lambda)$ to be the FLOTW multicharge and multipartition determined by $\mathfrak{B}_{\mathbf{RS}(\mathbf{s})}(\mathbf{RS}(\lambda)) = \mathscr{P}(\operatorname{read}(\lambda, \mathbf{s}))$

We further write **RS** for the map

$$\mathbf{RS} : \begin{array}{ccc} B(\lambda) & \longrightarrow & B(\mathbf{RS}(\lambda)) \\ |\lambda, \mathbf{s}\rangle & \longmapsto & |\mathbf{RS}(\lambda), \mathbf{RS}(\mathbf{s})\rangle. \end{array}$$

Theorem 3.1. $|\lambda, \mathbf{s}\rangle$ and $|\mu, \mathbf{r}\rangle$ are equivalent if and only if $\mathscr{P}(read(\lambda, \mathbf{s})) = \mathscr{P}(read(\mu, \mathbf{r}))$.

For a proof of this statement, see for instance [16, Section 3] or [18], which state the result for $\mathcal{U}_q(\mathfrak{sl}_e)$ -crystals, relying on the original arguments of Kashiwara in [13] and [14]. Moreover, since the symbol associated to $|\mathbf{RS}(\lambda), \mathbf{RS}(\mathbf{s})\rangle$ is semistandard, we have the following result. **Corollary 3.2.** *RS* is the canonical $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ -crystal isomorphism.

Example 3.3. $\mathbf{s} = (0, 2, -1)$ and $\lambda = (2.1, 3, 4.1^2)$.

Then $\mathfrak{B}_{\mathbf{s}}(\lambda) = \begin{pmatrix} 0 & 2 & 3 & 7 & \\ 0 & 1 & 2 & 3 & 4 & 5 & 9 \\ 0 & 1 & 2 & 4 & 6 & \end{pmatrix}$.

The associated reading is read(λ) = 7320954321064210, and the Schensted algorithm yields

 $\mathscr{P}(\operatorname{read}(\lambda)) = \begin{pmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 7 \\ 0 & 1 & 2 & 3 & 6 & 9 \\ 0 & 2 & 4 & & & \end{pmatrix}.$

Hence = **RS**(s) = (-2, 1, 2) and = **RS**(λ) = (2.1, 4.2, 1).

Property 3.4. Suppose **s** is such that $s_1 \leq \cdots \leq s_l$. Take $\lambda \in \mathcal{F}_s$. Then $|\mathbf{RS}(\lambda)| \leq |\lambda|$. Moreover, $|\mathbf{RS}(\lambda)| = |\lambda|$ if and only if $\mathbf{RS}(\lambda) = \lambda$.

Proof. Because of Corollary 3.2, we know in particular that $\mathbf{RS}(\lambda)$ is in the connected component of $B(\mathcal{F}_{\mathbf{RS}(s)})$ whose highest weight vertex is \emptyset . Hence, if we write $\emptyset = \tilde{e}_{i_m} \dots \tilde{e}_{i_1}(\mathbf{RS}(\lambda))$,

we have $|\mathbf{RS}(\lambda)| = m$. Now, because **RS** is a crystal isomorphism, we have $\lambda = \tilde{e}_{i_m} \dots \tilde{e}_{i_1}(\lambda)$. If $\lambda \neq \mathbf{RS}(\lambda)$, then the symbol of $|\lambda, \mathbf{s}\rangle$ is not semistandard, and the fact that $s_1 \leq \dots \leq s_l$ ensures that $\lambda \neq \emptyset$. Therefore, we have $|\lambda| = |\lambda| + m > m = |\mathbf{RS}(\lambda)|$.

3.2 Another $\mathcal{U}_q(\mathfrak{sl}_\infty)$ -crystal isomorphism

Let $\sigma \in \mathfrak{S}_l$, and for $\mathbf{s} = (s_1, \ldots, s_l) \in \mathbb{Z}^l$ denote $\mathbf{s}^{\sigma} = (s_{\sigma(1)}, \ldots, s_{\sigma(l)})$.

According to [9, Section 2], we know an explicit combinatorial description of the following $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ -crystal isomorphism:

$$\begin{array}{cccc} \chi_{\sigma} : & B(\mathbf{s}) & \longrightarrow & B(\mathbf{s}^{\sigma}) \\ & \lambda & \longmapsto & \chi_{\sigma}(\lambda) \end{array}$$

In fact, in [9, Corollary 2.3.3], the map χ_{σ} is described in the case where σ is a transposition (c, c + 1). We do not recall here the combinatorial construction of $\chi_{\sigma}(\lambda)$, since it is not really important for our purpose. However, we notice the following property. It will be used in the proof that the algorithm we construct in Section 4 terminates.

Property 3.5. *For all* $\sigma \in \mathfrak{S}_l$ *, and for all* $\lambda \in V(\mathbf{s})$ *,*

 $|\chi_{\sigma}(\lambda)| = |\lambda|.$

We also denote simply by χ the isomorphism corresponding to a permutation σ verifying $s_{\sigma(1)} \leq s_{\sigma(2)} \leq \cdots \leq s_{\sigma(l)}$ (i.e. the reordering of **s**).

4 General case and reduction to cylindric multipartitions

4.1 Compatibility between $\mathcal{U}'_{q}(\widehat{\mathfrak{sl}_{e}})$ -crystals and $\mathcal{U}_{q}(\mathfrak{sl}_{\infty})$ -crystals

In this section, we use the subscript or superscript e or ∞ to specify which module structure we are interested in, in particular for the (reduced) *i*-word, crystal operators, crystal graph.

The aim is to show that any $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ -crystal isomorphism is also a $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -crystal isomorphism. This comes as a natural consequence of the existence of an embedding of $B_e(\lambda)$ in $B_{\infty}(\lambda)$, as explained in [9, Section 4]. Note that the embedding in our case will just map any $\lambda \in B(\lambda)$ onto itself, unlike in [9].

Lemma 4.1. Let $i \in [[0, e-1]]$. Suppose there is an arrow $\lambda \xrightarrow{i} \mu$ in the crystal graph $B_e(\lambda)$, and denote $\gamma := [\mu] \setminus [\lambda]$. Then there is an arrow $\lambda \xrightarrow{j} \mu$ in the crystal graph $B_{\infty}(\lambda)$, where $j = cont(\gamma)$.

Proof. Denote by $w_i^e(\lambda)$ (resp. $w_i^{\infty}(\lambda)$) the *i*-word for λ with respect to the $\mathcal{U}'_q(\mathfrak{sl}_e)$ -crystal (resp. $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ -crystal) structure. Then $w_i^e(\lambda)$ is the concatenation of *i*-words for the $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ -crystal structure, precisely

$$w_i^e(\lambda) = \prod_{k \in \mathbb{Z}} w_{i+ke}^{\infty}(\lambda).$$
(4)

We further denote $\hat{w}_i^e(\lambda)$ and $\hat{w}_i^{\infty}(\lambda)$ the reduced *i*-words (that is, after recursive cancellation of the factors *RA*). The node γ is encoded in both $w_i^e(\lambda)$ and $w_j^{\infty}(\lambda)$ by a letter *A*. Now if this letter *A* does not appear in $\hat{w}_j^{\infty}(\lambda)$, this means that there is a letter *R* in $w_j^{\infty}(\lambda)$ which simplifies with this *A*. Hence, because of (4), this letter *R* also appears in $w_i^e(\lambda)$ and simplifies with the *A* encoding γ , and γ cannot be the good addable *i*-node of λ for the $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_e)$ -crystal structure, whence a contradiction. Thus γ produces a letter *A* in $\hat{w}_j^{\infty}(\lambda)$.

In fact, this letter A is the rightmost one in $\hat{w}_j^{\infty}(\lambda)$. Indeed, suppose there is another letter A in $\hat{w}_j^{\infty}(\lambda)$ to the right of the A encoding γ . Then it also appears in $\hat{w}_i^e(\lambda)$ at the same place (again because of Relation (4)). This contradicts the fact that γ is the good addable *i*-node for the $\mathcal{U}_a(\widehat{\mathfrak{sl}}_e)$ -crystal structure.

Therefore, γ is the good addable *j*-node of λ for the $\mathcal{U}_q(\mathfrak{sl}_\infty)$ -crystal structure.

Lemma 4.2. Let $i \in [[0, e-1]]$, and let φ be a $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ -crystal isomorphism. Suppose there is an arrow $\lambda \xrightarrow{i} \mu$ in the crystal graph $B_e(\lambda)$, and denote $\gamma := [\mu] \setminus [\lambda]$. Then

- 1. there is an arrow $\varphi(\lambda) \xrightarrow{i} v$ in the crystal graph $B_e(\varphi(\lambda))$,
- 2. $\mathbf{v} = \tilde{f}_i^{\infty}(\varphi(\lambda))$, where $j = cont(\gamma)$.

Proof. First, for all *k*, we have the following relation:

$$\hat{w}_{k}^{\infty}(\lambda) = \hat{w}_{k}^{\infty}(\varphi(\lambda)). \tag{5}$$

Indeed, if $\hat{w}_k^{\infty}(\lambda) = A^{\alpha} R^{\beta}$, then α can be seen as the number of consecutive arrows labeled by k in $B_{\infty}(\lambda)$ starting from λ , and β as the number of consecutive arrows labeled by k leading to λ . Subsequently, the integers α and β are invariant by φ , and the relation (5) is verified. Hence, by concatenating, we get

$$\prod_{k\in\mathbb{Z}}\hat{w}_{i+ke}^{\infty}(\lambda) = \prod_{k\in\mathbb{Z}}\hat{w}_{i+ke}^{\infty}(\varphi(\lambda)),\tag{6}$$

and therefore

$$\hat{w}_i^e(\lambda) = \hat{w}_i^e(\varphi(\lambda)),\tag{7}$$

which proves the first point.

Besides, we know by Lemma 4.1 that \tilde{f}_i^e acts like \tilde{f}_j^∞ on λ . Together with (6), this implies that \tilde{f}_i^e acts like \tilde{f}_j^∞ on $\varphi(\lambda)$. In other terms, $\nu = \tilde{f}_i^\infty(\varphi(\lambda))$, and the second point is proved. \Box

Proposition 4.3. Every $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ -crystal isomorphism is also a $\mathcal{U}'_q(\widehat{\mathfrak{sl}}_{e})$ -crystal isomorphism.

Proof. The fact that φ is a $\mathcal{U}_q(\mathfrak{sl}_\infty)$ -crystal isomorphism is encoded in the following diagram:



The first point of Lemma 4.2 tells us that we have:

Besides,

$$\begin{aligned} \mathbf{v} &= \tilde{f}_j^{\infty}(\varphi(\lambda)) & \text{by Point 2 of Lemma 4.2} \\ &= \varphi(\tilde{f}_j^{\infty}(\lambda)) & \text{because } \varphi \text{ is a } \mathcal{U}_q(\mathfrak{sl}_{\infty})\text{-crystal isomorphism} \\ &= \varphi(\boldsymbol{\mu}). \end{aligned}$$

Hence we can complete the previous diagram in

$$\begin{array}{c|c} \lambda & \xrightarrow{\varphi} & \varphi(\lambda) \\ & & & & \downarrow \\ \tilde{f}_i^e & & & & \downarrow \\ \mu & \xrightarrow{\varphi} & \varphi(\mu) \end{array}$$

which illustrates the commutation between \tilde{f}_i^e and φ .

As a consequence, the two particular $\mathcal{U}_q(\mathfrak{sl}_{\infty})$ -crystal isomorphisms **RS** and χ_{σ} , defined respectively in Section 3.1 and 3.2, are $\mathcal{U}'_q(\mathfrak{sl}_e)$ -crystal isomorphisms (for all values of $e \in \mathbb{N}_{>1}$).

4.2 The cyclage isomorphism

One of the most natural $\mathcal{U}'_q(\widehat{\mathfrak{sl}_e})$ -crystal isomorphisms to determine is the so-called *cyclage* isomorphism. For $\mathbf{s} = (s_1, \ldots, s_l)$, let $\mathbf{s}' := (s_l - e, s_1, \ldots, s_{l-1})$. Then the following result is easy to show (see for instance [9, Proposition 5.2.1], or [8, Proposition 3.1] for the simpler case l = 2):

Proposition 4.4. The map

$$\begin{aligned} \xi : & \overrightarrow{B(\lambda)} & \longrightarrow & \overrightarrow{B(\mathcal{F}_{s'})} \\ & (\lambda^1, \dots, \lambda^l) & \longmapsto & (\lambda^l, \lambda^1, \dots, \lambda^{l-1}) \end{aligned}$$

is a $\mathcal{U}'_{a}(\widehat{\mathfrak{sl}_{e}})$ -crystal isomorphism. It is called the cyclage isomorphism.

Therefore, in the sequel, we denote $\xi(\mathbf{s}) := (s_l - e, s_1, \dots, s_{l-1})$ and $\xi(\lambda) := (\lambda^l, \lambda^1, \dots, \lambda^{l-1})$.

Remark 4.5. Actually, we have more than this. Indeed, the map ξ is clearly invertible. Hence, because $\tilde{e}_i \circ \tilde{f}_i = \tilde{f}_i \circ \tilde{e}_i = \text{Id}$ whenever they act non trivially, we have

$$\hat{f}_i \circ \xi = \xi \circ \hat{f}_i$$

i.e.
$$\xi = \tilde{e}_i \circ \xi \circ \tilde{f}_i$$

i.e.
$$\xi \circ \tilde{e}_i = \tilde{e}_i \circ \xi.$$

Hence, \tilde{e}_i and ξ also commute.

Remark 4.6. Note that in [9], the cyclage is defined slightly differently, namely by

$$\zeta: \begin{array}{ccc} B(\lambda) & \longrightarrow & B(\mathcal{F}_{(s_2,s_3,\ldots,s_l,s_{1+e})}) \\ & (\lambda^1,\ldots,\lambda^l) & \longmapsto & (\lambda^2,\ldots,\lambda^l,\lambda^1). \end{array}$$

It is easy to see that both definitions are equivalent. Indeed, one recovers ξ by:

- 1. applying l 1 times ζ ,
- 2. translating all components of the multicharge $\zeta^{l-1}(\mathbf{s})$ by -e (which is a transformation that has clearly no effect on the multipartition).

Remark 4.7. Note that the cylindricity condition defined in Definition 2.3 is conveniently expressible in terms of symbols using the cyclage ξ . Precisely, $|\lambda, \mathbf{s}\rangle$ is cylindric if and only if the three following conditions are verified:

- 1. $\mathbf{s} \in \mathscr{S}_e$,
- 2. $\mathfrak{B}_{\mathbf{s}}(\lambda)$ is semistandard,
- 3. $\mathfrak{B}_{\xi(\mathbf{s})}(\xi(\lambda))$ is semistandard.

Remark 4.8. We will see (Remark 5.23) that the canonical isomorphism we aim to determine can be naturally regarded as a generalisation of this cyclage isomorphism...

Finally, it is straightforward from the definition of ξ that the following property holds:

Property 4.9. *For all* $\lambda \in \mathcal{F}_s$ *, we have* $|\xi(\lambda)| = |\lambda|$ *.*

4.3 Finding a cylindric equivalent multipartition

In this section, we make use of the $\mathcal{U}'_{q}(\widehat{\mathfrak{sl}_{e}})$ -crystal isomorphisms **RS** (see Section 3.1) and ξ (defined in Proposition 4.4) to construct an algorithm which associates to any charged multipartition $|\lambda, \mathbf{s}\rangle$ an equivalent charged multipartition $|\mu, \mathbf{r}\rangle$ which is cylindric (see Definition 2.3). In the sequel, we will denote by $\mathscr{C}_{\mathbf{s}}$ the subset of $\mathcal{F}_{\mathbf{s}}$ of cylindric *l*-partitions. In particular, this implies that $\mathbf{s} \in \mathscr{S}_{e}$. First of all, let us explain why restricting ourselves to cylindric multipartitions is relevant.

Proposition 4.10. Let $s \in \mathcal{S}_e$. The set \mathcal{C}_s is stable under the action of the crystal operators.

Proof. Let $\lambda \in \mathscr{C}_s$. By Remark 4.7, we know that $\mathfrak{B}_s(\lambda)$ is semistandard and $\mathfrak{B}_{\xi(s)}(\xi(\lambda))$ is semistandard.

It is easy to see that $\mathfrak{B}_{s}(\tilde{f}_{i}(\lambda))$ (resp. $\mathfrak{B}_{s}(\tilde{e}_{i}(\lambda))$) is still semistandard, whenever \tilde{f}_{i} (resp. \tilde{e}_{i}) acts non trivially on λ . Indeed, denote γ the good addable *i*-node and let $j = \operatorname{cont}(\gamma)$. It is encoded by an entry j + p, where p is the size of the symbol, see Section 2.1. By definition of being a good node is the leftmost *i*-node amongst all *i*-node of content j. Hence, there is no other no entry below the entry j + p encoding γ . Since \tilde{e}_{i} just turns this j + p into j + p + 1, the symbol of $\tilde{f}_{i}(\lambda)$ is still semistandard. The similar argument applies to $\tilde{e}_{i}(\lambda)$.

Since the symbol of $\xi(\lambda)$ is semistandard, by the same argument as above, we deduce that the symbol of $\tilde{f}_i(\xi(\lambda))$ is still semistandard, i.e. the symbol of $\xi(\tilde{f}_i(\lambda))$ is semistandard, since ξ commutes with \tilde{f}_i (by Proposition 4.4). This result also holds for $\xi(\tilde{e}_i(\lambda))$ because ξ also commutes with \tilde{e}_i (see Remark 4.5).

By Remark 4.7, this means that $\tilde{f}_i(\lambda)$ and $\tilde{e}_i(\lambda)$ are both cylindric.

The algorithm expected can now be stated. Firstly, if $\mathfrak{B}_{\mathbf{s}}(\lambda)$ is not semistandard, then we can apply **RS** to get a charged multipartition whose symbol is semistandard. Hence we can assume that $\mathfrak{B}_{\mathbf{s}}(\lambda)$ is semistandard. In particular, this implies that $s_c \leq s_{c+1}$ for all $c \in \llbracket 1, l-1 \rrbracket$. Then,

- 1. If $s_l s_1 < e$, then:
 - (a) if $\mathfrak{B}_{\xi(s)}(\xi(\lambda))$ is semistandard, then $|\lambda, s\rangle$ is cylindric, hence we stop and take $\mu = \lambda$ and $\mathbf{r} = \mathbf{s}$.
 - (b) if $\mathfrak{B}_{\xi(s)}(\xi(\lambda))$ is not semistandard, then put $\lambda \leftarrow \mathbf{RS}(\xi(\lambda))$ and $\mathbf{s} \leftarrow \mathbf{RS}(\xi(\mathbf{s}))$ and start again.
- 2. If $s_l s_1 \ge e$, then put $\lambda \leftarrow \mathbf{RS}(\xi(\lambda))$ and $\mathbf{s} \leftarrow \mathbf{RS}(\xi(\mathbf{s}))$ and start again.

Proposition 4.11. The algorithm above terminates.

Proof. For a multicharge $\mathbf{s} = (s_1, \dots, s_l)$, we denote $||\mathbf{s}|| := \sum_{k=2}^l (s_k - s_1)$. Hence, if $\mathfrak{B}_{\mathbf{s}}(\lambda)$ is semistandard, we have $||\mathbf{s}|| \ge 0$. In particular, at each step in the algorithm, this statistic is always non-negative, since we replace \mathbf{s} by $\mathbf{RS}(\xi(\mathbf{s}))$.

Suppose we are in case 1.(b). Since $s_l - s_1 < e$ and $s_1 \le s_2 \le \cdots \le s_l$, we have $(s_l - e) < s_1 \le s_2 \le \cdots \le s_{l-1}$. In other terms, the multicharge $\xi(\mathbf{s}) = (s_l - e, s_1, \dots, s_{l-1})$ is an increasing sequence. Hence Property 3.4 applies, and we have $|\mathbf{RS}(\xi(\lambda))| < |\lambda|$.

Suppose we are in case 2. The first thing to understand is that we get the same multipartition and multicharge applying ξ and **RS**, or applying ξ , then χ (see Section 3.2) and **RS**. Indeed, χ just reorders the multicharge and gives the associated multipartition, which is a transformation already included in **RS**, (which gives a multipartition whose symbol is semistandard). Hence, we consider that **RS**($\xi(\lambda)$) is obtained by applying successively ξ , then χ , and finally **RS**, to λ . In this procedure, it is possible that **RS** acts trivially (i.e. that $\chi(\xi(\lambda))$) is already semistandard). In fact,

If **RS** acts non trivially, then on the one hand χ(ξ(λ)) is non semistandard; and on the other hand χ(ξ(s)) is an increasing sequence (by definition of χ). Thus, we have

$$\begin{aligned} |\mathbf{RS}(\chi(\xi(\lambda)))| &< |\chi(\xi(\lambda))| & \text{applying Property 3.4} \\ &= |\lambda| & \text{by Properties 3.5 and 4.9} . \end{aligned}$$

Hence in this case, $|\mathbf{RS}(\xi(\lambda))| < |\lambda|$.

• If **RS** acts trivially, then this argument no longer applies. However, we have $||\mathbf{RS}(\chi(\xi(\mathbf{s})))|| = ||\chi(\xi(\mathbf{s})))|| < ||\mathbf{s}||$. Indeed, denote $\mathbf{s}' = \chi(\xi(\mathbf{s}))$. Since $s_l - s_1 \ge e$, we have $s_l - e \ge s_1$. This implies that the smallest element of $\xi(\mathbf{s}) = (s_l - e, s_1, \dots, s_{l-1})$ is again s_1 , and that $s'_1 = s_1$. Hence

$$||\mathbf{s}'|| = \sum_{k=2}^{l} (s'_k - s'_1)$$

=
$$\sum_{k=2}^{l} (s'_k - s_1)$$

=
$$\sum_{k=2}^{l-1} (s_k - s_1) + (s_l - e) - s_1$$

=
$$\sum_{k=2}^{l} (s_k - s_1) - e$$

=
$$||\mathbf{s}|| - e$$

<
$$||\mathbf{s}||$$

Note also that in this case, $|\mathbf{RS}(\xi(\lambda))| = |\chi(\xi(\lambda))| = |\lambda|$ by Properties 3.5 and 4.9.

We see that at each step, the rank |.| can never increase. In fact, since it is always nonnegative, there is necessarily a finite number of steps at which this statistic decreases. Moreover, when the rank does not increase, then the second statistic ||.|| decreases. Since it can never be negative (as noted in the beginning of the proof), there is also finite number of such steps. In conclusion, there is a finite number of steps in the algorithm, which means that it terminates.

Remark 4.12. This algorithm can also be stated in the simpler following way:

- 1. If $|\lambda, s\rangle$ is cylindric, then stop and take $\mu = \lambda$ and $\mathbf{r} = \mathbf{s}$.
- 2. Else, put $\lambda \leftarrow \mathbf{RS}(\xi(\lambda))$ and $\mathbf{s} \leftarrow \mathbf{RS}(\xi(\mathbf{s}))$ and start again.

Example 4.13. Set e = 4, l = 3, s = (0, 9, 5), and $\lambda = (4.2^2.1^3, 5.2^3.1^4, 7.6.4^2.2^2.1^3) \in \mathcal{F}_s$. Firstly, we see that

$$\mathfrak{B}_{\mathbf{s}}(\boldsymbol{\lambda}) = \begin{pmatrix} 0 & 1 & 3 & 4 & 5 & 7 & 8 & 11 & 12 & 15 & 17 & 18 \\ 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 9 & 10 & 11 & 12 & 14 & 15 & 16 & 20 \\ 0 & 2 & 3 & 4 & 6 & 7 & 10 & & & & & \\ \end{pmatrix}$$

is not semistandard. Thus we first compute $\tilde{\lambda} := \mathbf{RS}(\lambda)$ and $\tilde{\mathbf{s}} := \mathbf{RS}(\mathbf{s})$. We obtain

$$\mathfrak{B}_{\tilde{\mathbf{s}}}(\tilde{\boldsymbol{\lambda}}) = \begin{pmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 10 & 11 & 12 & 14 & 15 & 16 & 17 & 18 \\ 0 & 1 & 3 & 4 & 5 & 7 & 9 & 11 & 12 & 15 & 20 \\ 0 & 2 & 3 & 4 & 6 & 7 & 10 \end{pmatrix},$$

i.e. $\tilde{\lambda} = (4.2^2.1^3, 10.6.4^2.3.2.1^3, 2^5.1^3)$ and $\tilde{\mathbf{s}} = (0, 4, 10)$. We see that $|\tilde{\lambda}, \tilde{\mathbf{s}}\rangle$ is not cylindric. Hence we compute $\lambda^{(1)} := (\mathbf{RS} \circ \xi)(\lambda)$ and $\mathbf{s}^{(1)} := (\mathbf{RS} \circ \xi)(\lambda)$. We get

$$\mathfrak{B}_{\mathbf{s}^{(1)}}(\boldsymbol{\lambda}^{(1)}) = \begin{pmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 7 & 8 & 9 & 11 & 12 & 13 & 14 & 15 & 20 \\ 0 & 1 & 3 & 4 & 6 & 7 & 10 & 11 & 12 \\ 0 & 2 & 3 & 4 & 6 & 7 & 10 & & & \end{pmatrix}$$

i.e. $\lambda^{(1)} = (4.2^2.1^3, 4^3.2^2.1^2, 6.2^5.1^3)$ and $s^{(1)} = (0, 2, 8)$.

We keep on applying **RS** $\circ \xi$ until ending up with a cylindric multipartition. In fact, if we denote $\lambda^{(k)} := (\mathbf{RS} \circ \xi)^k(\lambda)$ and $\mathbf{s}^{(k)} := (\mathbf{RS} \circ \xi)^k(\mathbf{s})$, we can compute $|\lambda^{(2)}, \mathbf{s}^{(2)}\rangle$, $|\lambda^{(3)}, \mathbf{s}^{(3)}\rangle$, $|\lambda^{(4)}, \mathbf{s}^{(4)}\rangle$, and we finally have

$$\mathfrak{B}_{\mathbf{s}^{(5)}}(\boldsymbol{\lambda}^{(5)}) = \begin{pmatrix} 0 & 1 & 2 & 3 & 4 & 6 & 7 & 9 & 11 & 12 \\ 0 & 1 & 3 & 4 & 5 & 7 & 8 & 10 & 11 & 16 \\ 0 & 2 & 3 & 4 & 6 & 7 & 10 & & & \\ \end{pmatrix}$$

i.e. $\lambda^{(5)} = (4.2^2.1^3, 7.3^2.2^2.1^3, 3^2.2.1^2)$ and $\mathbf{s}^{(5)} = (-4, -1, -1)$. We see that $|\lambda^{(5)}, \mathbf{s}^{(5)}\rangle$ is cylindric.

This charged multipartition has the following Young diagram with contents:



With this representation, we see that $|\lambda^{(5)}, \mathbf{s}^{(5)}\rangle$ is not FLOTW (cf. Definition 2.3). Therefore, It remains to understand how to obtain an equivalent FLOTW multipartition from a cylindric multipartition. This is the point of the next section, which contains the main result of this paper (Theorem 5.20).

5 The case of cylindric multipartitions

Recall that for $\mathbf{s} \in \mathscr{S}_e$, we have denoted $\mathscr{C}_{\mathbf{s}}$ the set of cylindric *l*-partitions.

5.1 Pseudoperiods in a cylindric multipartition

Let $|\lambda, \mathbf{s}\rangle \in \mathcal{C}_{\mathbf{s}}$ such that λ is not FLOTW. Then there is a set of parts of the same size, say α , such that the residues at the end of these parts cover [[0, e - 1]]. This is formalised in the following definition.

Definition 5.1.

• The *first pseudoperiod* of λ is the sequence $P(\lambda)$ of its rightmost nodes

$$\gamma_1 = (a_1, \alpha, c_1), \ldots, \gamma_e = (a_e, \alpha, c_e)$$

verifying:

- 1. there is a set of parts of the same size $\alpha \ge 1$ such that the residues at the rightmost nodes of these parts cover [[0, e 1]].
- 2. α is the maximal integer verifying 1. 3. $\operatorname{cont}(\gamma_1) = \max_{\substack{c \in \llbracket 1, I \\ a \in \llbracket 1, h(\lambda^c) \rrbracket}} \operatorname{cont}(a, \alpha, c) \text{ and } c_1 = \min_{\substack{a \in \llbracket 1, h(\lambda^c) \rrbracket}} c,$ 4. for all $i \in \llbracket 2, e \rrbracket$, $\operatorname{cont}(\gamma_i) = \max_{\substack{c \in \llbracket 1, I \\ a \in \llbracket 1, h(\lambda^c) \rrbracket} \operatorname{cont}(a, \alpha, c) < \operatorname{cont}(a, \alpha, c)} \operatorname{cont}(a, \alpha, c)$ and $c_i = \min_{\substack{a \in \llbracket 1, h(\lambda^c) \rrbracket} \operatorname{cont}(\alpha, \alpha, c) = \operatorname{cont}(\gamma_i)} c.$

In this case, $P(\lambda)$ is also called a α -pseudoperiod of λ , and α is called the width of $P(\lambda)$.

• Denote $\lambda^{[1]} := \lambda \setminus P(\lambda)$, that is the multipartition obtained by forgetting ¹ in λ the parts α whose rightmost node belongs to $P(\lambda)$. Let $k \ge 2$. Then the *k*-th pseudoperiod of λ is defined recursively as being the first pseudoperiod of $\lambda^{[k]}$, if it exists, where $\lambda^{[k]} := \lambda^{[k-1]} \setminus P(\lambda^{[k-1]})$.

Any *k*-th pseudoperiod of λ is called a *pseudoperiod* of λ .

Example 5.2. Let e = 3, $\mathbf{s} = (2, 3, 4)$, and $\lambda = (2.1^2, 2.1^3, 2.1^4)$. One checks that λ is cylindric for e but not FLOTW. Then λ has the following diagram with contents:



Then λ has one 2-pseudoperiod and two 1-pseudoperiods. Its first pseudoperiod consists of $\gamma_1 = (1, 2, 3), \gamma_2 = (1, 2, 2)$ and $\gamma_3 = (1, 2, 1)$, with respective contents 5, 4 and 3, colored in blue. The second pseudoperiod is ($\gamma_1 = (2, 1, 3), \gamma_2 = (2, 1, 2), \gamma_3 = (2, 1, 1)$), with red contents; and the third (and last) pseudoperiod is ($\gamma_1 = (3, 1, 3), \gamma_2 = (3, 1, 2), \gamma_3 = (3, 1, 1)$), with green contents.

- **Lemma 5.3.** 1. $cont(\gamma_i) = cont(\gamma_{i-1}) 1$ for all $i \in [\![2, e]\!]$. In other terms, the contents of the elements of the pseudoperiod are consecutive.
 - 2. $c_i \le c_{i-1}$ for all $i \in [[2, e]]$.
- *Proof.* 1. Suppose there is a gap in the sequence of these contents. Then the pseudoperiod must spread over e + 1 columns in the symbol $\mathfrak{B}_{s}(\lambda)$. Denote by b the integer of $\mathfrak{B}_{s}(\lambda)$ corresponding to the last element of $P(\lambda)$, and k the column where it appears. The integer of $\mathfrak{B}_{s}(\lambda)$ corresponding to the gap must be in column k + 1, and since λ is cylindric, it must be greater than or equal to b + e. In fact, it cannot be greater than b + e since the corresponding part is below a part of size α , and it has to correspond to a part of size α , and there cannot be a gap, whence a contradiction.
 - 2. Since the nodes of $P(\lambda)$ are the rightmost nodes of parts of the same size α , together with the fact $s_1 \leq \cdots \leq s_l$, and point 1., γ_i is necessarily either to the left of γ_{i-1} or in the same component.

Remark 5.4. If $\alpha = \max_{i,j} \lambda_j^i$, then the first pseudoperiod corresponds to a "period" in $\mathfrak{B}_{\mathbf{s}}(\lambda)$, accordingly to [10, Definition 2.2] This is the case in Example 5.2. In the case where each pseudoperiod corresponds to a period in the symbol associated to $\lambda^{[k]}$, one can directly recover the empty *l*-partition and the corresponding multicharge using the "peeling procedure" explained in [10]. However, in general, $\mathfrak{B}_{\mathbf{s}}(\lambda)$ might not have a period, as shown in the following example.

Example 5.5. e = 4, $\mathbf{s} = (5, 6, 8)$ and $\lambda = (6^2 \cdot 2 \cdot 1, 3 \cdot 2^3 \cdot 1^2, 6 \cdot 2^2 \cdot 1^3)$. Then $\lambda \in \mathscr{C}_{\mathbf{s}}$ but is not FLOTW for *e*. It has the following Young diagram with residues:



Then there is a 2-pseudoperiod and a 1-pseudoperiod. The first pseudoperiod of λ consists of the nodes $\gamma_1 = (2, 2, 3)$, $\gamma_2 = (3, 2, 3)$, $\gamma_3 = (2, 2, 2)$ and $\gamma_4 = (3, 2, 2)$, with respective contents 8, 7, 6 and 5. The 1-pseudoperiod is ((4, 1, 3), (5, 1, 3), (6, 1, 3), (4, 1, 1)).

¹This means that one considers only the nodes of λ that are in parts whose rightmost node is not in $P(\lambda)$, but without changing the indexation nor the contents of these nodes.

Of course, one could also describe pseudoperiods on the **s**-symbol associated to λ . However, this approach is not that convenient, and in the setting of cylindric multipartitions, we favour the "Young diagram with contents" approach, which encodes the same information. Nevertheless, we notice this property, which will be used in the proof of Lemma 5.16:

Proposition 5.6. Let $P(\lambda)$ be a pseudoperiod of λ . Denote by B the set of entries of $\mathfrak{B}_{s}(\lambda)$ corresponding to the nodes of $P(\lambda)$. Then each column of $\mathfrak{B}_{s}(\lambda)$ contains at most one element of B. Moreover, the elements of B appear in consecutive columns of $\mathfrak{B}_{s}(\lambda)$.

Proof. This is direct from the fact that the nodes of $P(\lambda)$ are all rightmost nodes of parts of the same size α , together with Lemma 5.3.

We will now determine the canonical $\mathcal{U}'_{q}(\widehat{\mathfrak{sl}_{e}})$ -crystal isomorphism for cylindric multipartitions. In the following section, we only determine the suitable multicharge. In Section 5.3, we explain how to construct the actual corresponding FLOTW multipartition.

5.2 Determining the multicharge

In [10], Jacon and Lecouvey have proved that when $\lambda = \lambda$ is a highest vertex, then it suffices to "peel" the symbol of λ in order to get an empty equivalent multipartition. We do not recall here this procedure in detail, but it basically consists in removing all periods in the symbol of λ (see Example 5.10).

When we start from a multipartition λ which is no longer a highest weight vertex, we can, in general, no longer peel the symbol, for it does not necessarily contain a period anymore (see Remark 5.4). However, the multicharge we look for is constant along the crystal, hence entirely determined by the highest weight vertex $\hat{\lambda}$. Therefore, it is the representative in \mathcal{D}_e of the multicharge associated to the empty multipartition obtained after peeling $\hat{\lambda}$. We denote it by $\varphi(\mathbf{s})$.

Now, since $|\hat{\lambda}, \mathbf{s}\rangle$ is already cylindric, it turns out that the period can be easily read in the symbol of $\hat{\lambda}$. In fact, we have the following property:

Proposition 5.7. Let $|\lambda, \mathbf{s}\rangle \in \mathcal{C}_{\mathbf{s}}$ such that $\mathfrak{B}_{\mathbf{s}}(\lambda)$ has a period *P*. Then each of the rightmost *e* columns of $\mathfrak{B}_{\mathbf{s}}(\lambda)$ contains a unique element of *P*.

Proof. Since a period is nothing but a pseudoperiod whose width is the largest part in λ (cf. Remark 5.4), this result is just a particular case of Proposition 5.6.

Hence, deleting the first period (first step of the peeling) $\dot{\lambda}$, we get a multicharge $s^{(1)}$ verifying

$$s_{l}^{(1)} = s_{l} - (s_{l} - s_{l-1}) = s_{l-1}$$

$$s_{l-1}^{(1)} = s_{l-1} - (s_{l-1} - s_{l-2}) = s_{l-2}$$

$$\vdots$$

$$s_{2}^{(1)} = s_{2} - (s_{2} - s_{1}) = s_{1}$$

$$s_{1}^{(1)} = s_{1} - (e - (s_{l} - s_{1})) = s_{l} - e.$$

In other terms, we have $s^{(1)} = \xi(s)$, where ξ is the cyclage operator defined in Section 4.2.

Applying this recursively, we obtain

Proposition 5.8. For all $k \ge 1$, $\mathbf{s}^{(k)} = \xi^k(\mathbf{s})$, where $\mathbf{s}^{(k)}$ denotes the multicharge associated with the peeled symbol after k steps (with $\mathbf{s}^{(0)} = \mathbf{s}$).

Remark 5.9. It is possible that, at some step, a multicharge $\mathbf{s}^{(k)}$ will not be in \mathscr{S}_e anymore. However, for any k, one always has $s_i^{(k)} \le s_j^{(k)}$ for i < j, and $s_l^{(k)} - s_1^{(k)} \le e$. Moreover, if $s_l^{(k)} - s_1^{(k)} = e$, then $s_l^{(k+m)} - s_1^{(k+m)} < e$, where $m \ge 1$ is the number of components of $s^{(k)}$ equal to $s_l^{(k)}$.

Note that this never happens if $s_i < s_j$ for all i < j.

Example 5.10. e = 3, $\mathbf{s} = (3, 3, 4)$, $\lambda = (3^2.2, 2^2.1, 3.1^2) = \lambda$. One checks that $|\lambda, \mathbf{s}\rangle$ is cylindric but not FLOTW. The associated symbol is

$$\begin{pmatrix} 0 & 1 & 3 & 4 & 7 \\ 0 & 2 & 4 & 5 \\ 0 & 3 & 5 & 6 \end{pmatrix}$$

Peeling this symbol, we get successively

$$\begin{pmatrix} 0 & 1 & 3 & 4 \\ 0 & 2 & 4 & 5 \\ 0 & 3 & & \end{pmatrix} \text{ and } \mathbf{s}^{(1)} = (1, 3, 3),$$
$$\begin{pmatrix} 0 & 1 & 3 & 4 \\ 0 & 2 & & \\ 0 & & & \end{pmatrix} \text{ and } \mathbf{s}^{(2)} = (0, 1, 3),$$
$$\begin{pmatrix} 0 & 1 \\ 0 & \\ 0 & & \end{pmatrix} \text{ and } \mathbf{s}^{(3)} = (0, 0, 1).$$

Note that $\mathbf{s}^{(2)} \notin \mathscr{S}_e$.

The following proposition is now easy to prove.

Proposition 5.11. Let $\mathbf{s} \in \mathscr{S}_e$. There exists $k \in \mathbb{Z}$ such that $\xi^k(\mathbf{s}) \in \mathscr{D}_e$. In fact, we have $\varphi(\mathbf{s}) = \xi^k(\mathbf{s})$.

Proof. Recall that $\varphi(\mathbf{s})$ is the representative of the multicharge \mathbf{s}' associated to the peeled symbol of $\hat{\boldsymbol{\lambda}}$. Because of Consequence 5.8, \mathbf{s}' is obtained from \mathbf{s} by applying several times (say *t* times) $\boldsymbol{\xi}$. In other terms, $\mathbf{s}' = \boldsymbol{\xi}^t(\mathbf{s})$. Now,

- if $\mathbf{s}' \in \mathcal{D}_e$, then $\varphi(\mathbf{s}) = \mathbf{s}'$ and k = t.
- if $\mathbf{s}' \in \mathscr{S}_e \setminus \mathscr{D}_e$, then the representative of \mathbf{s}' in \mathscr{D}_e is of the form $\xi^v(\mathbf{s}')$ for some $v \in \mathbb{Z}$. Hence, $\varphi(\mathbf{s}) = \xi^k(\mathbf{s})$ with k = t + v.
- if s' ∉ S_e, we are however ensured (see Remark 5.9) that there exists m ∈ [[1, e − 1]] such that ξ^m(s') ∈ S_e. We are then in the previous situation, i.e. there exists v ∈ Z such that ξ^{m+v}(s') ∈ D_e, and therefore φ(s) = ξ^k(s) with k = t + m + v.

Example 5.12. As in Example 5.10, take e = 3, $\mathbf{s} = (3, 3, 4)$ and $\lambda = (3^2.2, 2^2.1, 3.1^2)$. Then $\xi^3(\mathbf{s}) = (0, 0, 1) \in \mathcal{D}_e$. Hence $\varphi(\mathbf{s}) = (0, 0, 1)$.

5.3 Determining the FLOTW multipartition

In order to compute the multipartition $\varphi(\lambda)$, we need to introduce a new crystal isomorphism, which acts on cylindric multipartitions. In fact, the only difference between a cylindric multipartition and a FLOTW multipartition is the possible presence of pseudoperiods. Therefore, we want to determine an isomorphism which maps a cylindric multipartition to another cylindric multipartition with one less pseudoperiod. Applying this recursively, we will eventually end up with a FLOTW multipartition equivalent to λ .

Take $\mathbf{s} \in \mathscr{S}_e$ and $\lambda \in \mathscr{C}_s$. Let α be the width of $P(\lambda)$, the first pseudoperiod of λ . Denote by $\psi(\lambda)$ the multipartition μ charged by $\xi(\mathbf{s})$ defined as follows:

- μ^c contains all parts λ_a^c of λ^c such that $1 \le \lambda_a^c < \alpha$, for $c \in [[1, l]]$,
- μ^c contains all parts λ_a^{c-1} of λ^{c-1} such that $\lambda_a^{c-1} > \alpha$ for c > 1, and μ^1 contains all parts λ_a^l of λ^l such that $\lambda_a^l > \alpha$,
- μ^c contains all parts $\lambda_a^c = \alpha$ whose rightmost node does not belong to $P(\lambda)$, for $c \in [1, l]$.

This naturally defines a mapping $\gamma \leftrightarrow \Gamma$ from the set of nodes of $\lambda | P(\lambda)$ (see Definition 5.1) onto the set of nodes of $\psi(\lambda)$. We then say that Γ is *canonically associated to* γ , and conversely.

Example 5.13. Let us go back to Example 5.5. We had e = 4, s = (5, 6, 8) and



where the bold contents correspond to the first pseudoperiod (whose width is $\alpha = 2$). Then $\xi(\mathbf{s}) = (4, 5, 6)$, and



We observe that the contents of canonically associated nodes are unchanged, except for the nodes of λ that lie in part of λ^l greater than α , whose content is translated by -e. More formally, this writes:

Proposition 5.14. *Denote* $\psi(\lambda) =: \mu$.

- 1. Let $\Gamma = (A, B, C)$ be a node of μ with $\mu_A^C > \alpha$. Denote by γ the node of λ canonically associated to γ .
 - If c > 1, then $cont_{\mu}(\Gamma) = cont_{\lambda}(\gamma)$.
 - If c = 1, then $cont_{\mu}(\Gamma) = cont_{\lambda}(\gamma) e$.
- 2. Let $\Gamma = (A, B, C)$ be a node of μ with $\mu_A^C \leq \alpha$. Denote by γ its canonically associated node in λ . Then $cont_{\mu}(\Gamma) = cont_{\lambda}(\gamma)$.

Notation 5.15. Let s be a *l*-charge in \mathscr{S}_e , $\lambda \in \mathscr{C}_s$ non FLOTW, and $c \in [[1, l]]$. Set α to be the width of the first pseudoperiod of λ . We denote:

- $N_c^{>\alpha}$ the number of parts greater than α in λ^c ,
- N_c^{α} the number of parts equal to α in λ^c that are deleted in λ to get $\psi(\lambda)$ (i.e. parts whose rightmost node belongs to the pseudoperiod).

Proof. Of course, it is sufficient to prove this for only one node in each part considered, since the contents of all other nodes of the part is then determined. We prove it only for the leftmost nodes, i.e. the ones of the form (A, 1, C).

- 1. This is clear since the multicharge associated to μ is simply the cyclage of **s** (that shifts **s** "to the right" and maps s_1 to $s_l e$), and the parts greater than α are similarly shifted in λ to get μ .
- 2. Let $\Gamma = (A, 1, C)$ be a node in μ such that $\mu_A^C \le \alpha$, so that $\gamma = (a, 1, c)$ with c = C. First, assume C > 1. Then

$$\operatorname{cont}_{\mu}(\Gamma) = s_{C-1} - N_{C-1}^{>\alpha}$$
.

On the other hand,

$$\operatorname{cont}_{\lambda}(\gamma) = s_C - N_C^{>\alpha} - N_C^{\alpha}.$$

Now by definition of $P(\lambda)$, which is charged by $\xi(s)$, we have

$$s_C - s_{C-1} = N_C^{>\alpha} + N_C^{\alpha} - N_{c-1}^{>\alpha},$$

which is equivalent to

$$N_C^{\alpha} = s_C - N_C^{>\alpha} - s_{C-1} - N_{C-1}^{>\alpha}.$$

Hence we have

$$\operatorname{cont}_{\lambda}(\gamma) = s_{C} - N_{C}^{>\alpha} - N_{C}^{\alpha}$$

= $s_{C} - N_{C}^{>\alpha} - (s_{C} - N_{C}^{>\alpha} - (s_{C-1} - N_{C-1}^{>\alpha}))$
= $s_{C-1} - N_{C-1}^{>\alpha}$
= $\operatorname{cont}_{\mu}(\Gamma).$

Now, assume C = 1. The argument is the same:

$$\operatorname{cont}_{\mu}(\Gamma) = s_l - e - N_l^{>\alpha}$$
, and
 $\operatorname{cont}_{\lambda}(\gamma) = s_1 - N_1^{>\alpha} - N_1^{\alpha}$.

Moreover, we have:

$$N_{1}^{\alpha} = e - \sum_{d=2}^{l} N_{d}^{\alpha}$$

= $e - \sum_{d=2}^{l} (r_{d} - r_{d-1})$
= $e - r_{l} + r_{1}$
= $e - (s_{l} - N_{l}^{>\alpha}) + (s_{1} - N_{1}^{>\alpha})$ using the above case,

which implies that

$$\operatorname{cont}_{\mu}(\Gamma) = s_{l} - e - N_{l}^{>\alpha}$$

= $s_{1} - N_{1}^{>\alpha} - e + s_{l} - N_{l}^{>\alpha} - s_{1} + N_{1}^{>\alpha}$
= $s_{l} - e - N_{l}^{>\alpha}$
= $\operatorname{cont}_{\lambda}(\gamma).$

We now aim to prove that the map ψ we have just defined is in fact a crystal isomorphism between connected components of Fock spaces crystals (this is upcoming Theorem 5.20). In order to do that, we will need the following three lemmas, in which we investigate the compatibility between ψ and the possible actions of the crystal operators \tilde{f}_i .

Lemma 5.16. Suppose that $\gamma^+ = (a, \alpha + 1, c)$ is the good addable *i*-node of λ , with $\gamma \in P(\lambda)$. Then

- $\Delta^+ = (a, \alpha + 1, c + 1)$ is the good addable *i*-node of $\psi(\lambda)$ if $1 \le c < l$,
- $\Delta^+ = (a, \alpha + 1, 1)$ is the good addable *i*-node of $\psi(\lambda)$ if c = l.

Proof. The proof first splits in three cases (which, in turn split in subcases), even though ultimately, the argument is the same. In each case (and subcase), we determine a certain node δ of λ which is not in $P(\lambda)$. Then, we show that the node of $\psi(\lambda)$ canonically associated to δ is the node Δ we expect.

- 1. Assume first that γ is the first element of $P(\lambda)$. Denote by $\gamma_e = (a_e, \alpha, c_e)$ the last node of $P(\lambda)$. Then $\operatorname{cont}(\gamma_e) = \operatorname{cont}(\gamma^+) e$, and since γ^+ is an *i*-node, then so is γ_e .
 - (a) Suppose that γ_e is removable (cf Example 5.17, 1.). Since γ^+ is the good addable *i*-node, the letter *R* produced by γ_e in the *i*-word must not simplify with the *A* produced by γ^+ . This means that there exists an addable node $\tilde{\gamma}^+$ in $\lambda^{\tilde{c}}$ with either
 - $\tilde{c} > c$ and $\operatorname{cont}(\tilde{\gamma}^+) = \operatorname{cont}(\gamma^+)$, or
 - $\tilde{c} < c_e$ and $\operatorname{cont}(\tilde{\gamma}^+) = \operatorname{cont}(\gamma^+) e$ (one cannot have $\tilde{c} = c_e$ since otherwise γ_e would not be removable).

In the first case, this means that there is an integer β in the \tilde{c} -th row of $\mathfrak{B}_{\mathbf{s}}(\lambda)$ which is also in the *c*-th row and the same column. By the semistandardness of $\mathfrak{B}_{\mathbf{s}}(\lambda)$ (because $|\lambda, \mathbf{s}\rangle$ is cylindric), we are ensured that β is also present in the row *d* and the same column for all $d \in [\![c, \tilde{c}]\!]$. This is equivalent to saying that there is a part of size α in each component λ^d with $d \in [\![c, \tilde{c}]\!]$ whose rightmost node $\hat{\gamma}$ verifies $\operatorname{cont}(\hat{\gamma}) = \operatorname{cont}(\gamma)$. We denote by δ the one located in the component λ^{c+1} . In particular, δ^+ is an *i*-node. Moreover, δ^+ is an addable node of λ . Indeed, if it is not, then there is a part of size α just above the part whose rightmost node is $\tilde{\gamma}$. Denote by $\overline{\gamma}$ its rightmost node. Then $\operatorname{cont}(\overline{\gamma}) = \operatorname{cont}(\delta) + 1 = \operatorname{cont}(\gamma) + 1$. This implies that $\overline{\gamma}$ must be in $P(\lambda)$, which contradicts the fact that γ is the first element of $P(\lambda)$.

In the second case, there is an integer β in the \tilde{c} -th row of $\mathfrak{B}_{\mathbf{s}}(\lambda)$ which is also in the c_e -th row and the same column, say column k. Again, since the symbol is semistandard, the elements $\mathfrak{b}_{d,k}$ appearing in row d, with $d \in \llbracket 1, \tilde{c} \rrbracket$, and in column k verify $\mathfrak{b}_d \geq \mathfrak{b}_{d'} \geq \beta$ for $1 \leq d' < d < c_e$. But the cylindricity property also implies that $\mathfrak{b}_{l,k+e} \geq \mathfrak{b}_{1,k} + e$. Actually, the (k + e)-th column of the symbol is also the column that contains the integer corresponding to the first node of $P(\lambda)$, since pseudoperiods have length e. Moreover, this element is equal to $\beta + e$. In other terms, $\mathfrak{b}_{c,k+e} = \beta + e$. By semistandardness again, one must have $\beta + e \geq \mathfrak{b}_{d,k+e} \geq$ $\mathfrak{b}_{d',k+e}$ for all $c \leq d < d' \leq l$. To sum up, we have

$$\beta + e = \mathfrak{b}_{c,k+e} \ge \mathfrak{b}_{c+1,k+e} \ge \cdots \ge \mathfrak{b}_{l,k+e} \ge \mathfrak{b}_{1,k} + e \ge \cdots \ge \mathfrak{b}_{c_e-1,k} + e \ge \mathfrak{b}_{c_e,k} + e = \beta + e,$$

thus all inequalities are in fact equalities. As in the first case, this means in particular that there is a part of size α in λ^{c+1} (if c < l) or in λ^1 (if c = l) whose rightmost node δ verifies cont(δ) = cont(γ) – e. In particular, δ^+ is an *i*-node. Moreover, δ^+ is an addable node of λ . Indeed, if it is not, then there is a part of size α just above the part whose rightmost node is $\tilde{\gamma}$. Denote by $\overline{\gamma}$ its rightmost node. Then cont($\overline{\gamma}$) = cont(δ) + 1 = cont(γ) – e + 1 = cont(γ_e). This implies that $\overline{\gamma}$ must be the last element of $P(\lambda)$ instead of γ_e , which is a contradiction.

- (b) Suppose that γ_e is not removable (cf Example 5.17, 2.). This means that there exists a part of size α below the part whose rightmost node is γ_e. Note that the rightmost node ỹ of this part has content cont(ỹ) = cont(γ_e) − 1. By the same cylindricity argument used in 1.(a), this part of size α spreads in all components of λ. This means that there exists a part α in λ^{c+1} (if c < l) or in λ¹ (if c = l) with rightmost node δ verifying cont(δ) = cont(γ) (if c < l) and cont(δ) = cont(γ) − e (if c = l). In particular, δ⁺ is an *i*-node. Moreover, if δ⁺ is addable, unless, of course, it is in the component λ^{c_e}. This can be seen using the exact same argument as in 1.(a).
- 2. Assume now that γ is the last element of $P(\lambda)$. First of all, note that if l > 1, one can never have c = l. Indeed, in this case γ^+ would not be an addable node.
 - (a) Suppose that $s_{c+1} > s_c$ (cf Example 5.17, 3.). Then, using Proposition 5.6, we can claim that there exists a part of size α in the component λ^{c+1} . Denote $\tilde{\gamma}$ its rightmost node. By Lemma 5.3, $\operatorname{cont}(\tilde{\gamma}) = \operatorname{cont}(\gamma) + 1 = \operatorname{cont}(\gamma^+)$, and $\tilde{\gamma}$ is an *i*-node. Now, if $\tilde{\gamma}$ is removable, then $\tilde{\gamma}$ and γ^+ yield an occurence of *RA* which contradicts the fact that γ^+ is the good *i*-node of type *A*. Hence, there is necessarily a part α below the part whose rightmost node is $\tilde{\gamma}$. Denote by δ its rightmost node, so that δ^+ is an *i*-node.
 - (b) Suppose now that $s_{c+1} = s_c$ (cf Example 5.17, 4.). Consider the previous node in $P(\lambda)$, denote it by γ_{e-1} . By Lemma 5.3 again, $\operatorname{cont}(\gamma_{e-1}) = \operatorname{cont}(\gamma) + 1 =$ $\operatorname{cont}(\gamma^+)$, so that γ_{e-1} is an *i*-node. Besides, since γ^+ is addable, γ_{e-1} is removable unless there is a part α below the part whose rightmost node is γ_{e-1} . But if γ_{e-1} is removable, then there exists an addable *i*-node $\tilde{\gamma}$ in $\lambda^{\tilde{c}}$ with $\tilde{c} \in [c+1, \ldots, c_1-1]$, where c_1 is the component of λ which contains the first node of $P(\lambda)$ (otherwise γ^+ and γ_{e-1} yield and occurence RA and γ^+ cannot be the good addable *i*-node). Then, by the cylindricity argument again, there is a part α in each component λ^d with $d \in [c+1, \ldots, \tilde{c}]$, whose rightmost node has content $\operatorname{cont}(\gamma)$. In particular, this is true for d = c + 1. We denote δ the one located in the component λ^{c+1} . Now if γ_{e-1} is not removable, i.e. if there exists a part α below the part whose rightmost node is γ_{e-1} , then again the cylindricity implies that this parts spreads to all components λ^d with $d \in [c+1, \ldots, c_{e-1} - 1]$, where c_{e-1} is the component of λ which contains γ_{e-1} , and have the same contents. Again, we denote δ the rightmost node of the part α of λ^{c+1} which has content $\operatorname{cont}(\delta) = \operatorname{cont}(\gamma)$.
- 3. Assume finally that γ is neither the first nor the last node of $P(\lambda)$. Again, *c* cannot be equal to *l* because then it would be the first node of $P(\lambda)$.

- (a) Suppose that $s_{c+1} > s_c$ (cf Example 5.17, 5.). We can use the same arguments as in 2.(a), and define δ in the exact same way.
- (b) Suppose that s_c = s_{c+1} (cf Example 5.17, 6 .). Then since γ⁺ is not the first node of P(λ), there exists a part α in a component λ^c with c̃ > c whose rightmost node γ̃ has content cont(γ̃) = cont(γ) + 1. If c̃ = c + 1, then we can use the previous case 3.(a). If c̃ > c + 1, then:
 - if $\tilde{\gamma}$ is removable, then there exists a part α in a component \bar{c} with $c < \bar{c} < \tilde{c}$ whose rightmost node has content cont(γ), otherwise $\tilde{\gamma}$ produces, together with γ^+ , an occurence *RA*, whence the usual contradiction. By cylindricity, such a part α also exists in the (c + 1)-th component of λ . We denote δ its rightmost node.
 - if γ̃ is not removable, then there exists a part α below the part whose rightmost node is γ̃, with rightmost content equal to cont(γ). Again, by cylindricity, it also exists in the (c + 1)-th component of λ, and we denote δ its rightmost node.

It is obvious, but important to notice, that δ is not a node of $P(\lambda)$. Hence, there is a node Δ of $\psi(\lambda)$ which is canonically associated to δ . By Proposition 5.14, $\operatorname{cont}_{\psi(\lambda)}(\Delta) = \operatorname{cont}_{\lambda}(\delta)$ and in fact $\Delta = (a, \alpha, c + 1)$ if c < l and $\Delta = (a, \alpha, 1)$ if c = l. In particular Δ^+ is an *i*-node. Moreover, it is addable in $\psi(\lambda)$ since δ^+ is either addable in λ , or the rightmost node of a part above which sit parts that are deleted after applying ψ .

In fact, Δ^+ is the good addable *i*-node of $\psi(\lambda)$. Indeed, consider the *i*-word for $\psi(\lambda)$. Denote it by w_i^{ψ} , and denote w_i the *i*-word for λ . By construction of $\psi(\lambda)$, the subword of w_i corresponding to the rightmost nodes of the parts that are either greater than α (respectively smaller than or equal to α but not in $P(\lambda)$) is also a subword of w_i^{ψ} . The only differences that are likely to appear are the following:

- The letters *R* and *A* that correspond to nodes in $P(\lambda)$ vanish. Note that we have assumed that there is always such a letter *A* in w_i (since γ is in $P(\lambda)$).
- The parts of size α that are below a part whose rightmost node is in P(λ) give a new letter A in ψ(λ).

Now by construction of δ :

- 1. If δ^+ gives a letter *A* in w_i , then it is adjacent to the letter *A* encoding γ^+ , to its left. Hence, since the *A* encoding γ^+ is no longer in w_i^{ψ} , the letter *A* corresponding to Δ^+ in w_i^{ψ} plays the same role as the one corresponding to γ^+ in λ : it is the rightmost *A* in the reduced *i*-word of $\psi(\lambda)$. In other terms, Δ^+ is the good addable *i*-node of $\psi(\lambda)$.
- 2. If δ^+ does not give a letter *A* in w_i , that is if there is an element of $P(\lambda)$ just above δ , then it is clear that, again, the *A* encoding Δ^+ in w_i^{ψ} plays the same role as the *A* encoding γ^+ in w_i , and that Δ^+ is the good addable *i*-node of $\psi(\lambda)$.

Example 5.17.

- 1. $\lambda = (4.2, 2^2, 5.2)$, $\mathbf{s} = (2, 3, 4)$, e = 3 and i = 2.
- 2. $\lambda = (6^2 \cdot 2 \cdot 1^3, 4 \cdot 2^3 \cdot 1^2, 6 \cdot 2^2 \cdot 1^4)$, $\mathbf{s} = (5, 6, 8)$, e = 4 and i = 1.
- 3. $\lambda = (2.1, 1^3, 1)$, $\mathbf{s} = (3, 4, 5)$, e = 4 and i = 3.
- 4. $\lambda = (1, 1, 1^3)$, $\mathbf{s} = (4, 4, 7)$, e = 4 and i = 1.
- 5. $\lambda = (3.2.1^2, 4.2.1, 2^3)$, $\mathbf{s} = (2, 3, 4)$, e = 4 and i = 0.
- 6. $\lambda = (2^2, 3.2, 2)$, $\mathbf{s} = (3, 4, 4)$, e = 3 and i = 0.

Lemma 5.18. Suppose that $\gamma^+ = (a, \lambda_a^c + 1, c)$ is the good addable *i*-node of λ , with $\lambda_a^c < \alpha$ or $[\lambda_a^c = \alpha \text{ and } \gamma \notin P(\lambda)]$. Then $\Gamma^+ = (a - D, \lambda_a^c + 1, c)$ is the good addable *i*-node of $\psi(\lambda)$, where

- $D = N_c^{>\alpha} N_{c-1}^{>\alpha} + N_c^{\alpha} N_{c-1}^{\alpha}$ if $c \ge 1$,
- $D = N_1^{>\alpha} N_l^{>\alpha} + N_1^{\alpha} N_l^{\alpha}$ if c = 1 (see Notation 5.15).

Proof. First, note that Γ is nothing but the node of $\psi(\lambda)$ canonically associated to γ in the definition of ψ . Besides, by definition of $\psi(\lambda)$, together with Proposition 5.14, we are ensured that Γ^+ is in fact an addable *i*-node of $\psi(\lambda)$. Moreover, w_i^{ψ} is likely to contain new letters A, namely the one corresponding to nodes η^+ with $\eta = (b, \alpha, d) \notin P(\lambda)$ lying just below a node of $P(\lambda)$. Denote by H the node of $\psi(\lambda)$ corresponding to η . Note that $\operatorname{cont}(\eta) \neq \max_{\tilde{\gamma} \in P(\lambda)}(\operatorname{cont}(\tilde{\gamma}))$. Indeed, since there is a node of $P(\lambda)$ just above η , it is not the first node of $P(\lambda)$ and hence it does not have maximal content. Since $\eta \notin P(\lambda)$, there exists a node $\tilde{\gamma} = (\tilde{a}, \alpha, \tilde{c}) \in P(\lambda)$ such that $\operatorname{cont}(\tilde{\gamma}) = \operatorname{cont}(\eta)$, $\tilde{c} < d$, and $\tilde{\gamma}^+$ is addable. Hence H plays the same role in $\psi(\lambda)$ as $\tilde{\gamma}$ in λ . In particular, since the good addable *i*-node of λ is in a part of size smaller than α , there is necessarily a letter R, encoding a node ρ , that simplifies with the letter A encoding $\tilde{\gamma}^+$. Now:

- If ρ is the rightmost node of a part of size different than α , then it is obviously not in $P(\lambda)$.
- If ρ is the rightmost node of a part of size α and cont(ρ) = cont(η^+), then it is not in $P(\lambda)$ either. Indeed, there is a node of $P(\lambda)$ just above η whose content is cont(η^+), which is therefore not of type R, and hence different from ρ , and since all nodes of $P(\lambda)$ have different contents, ρ is not in $P(\lambda)$.
- If ρ is the rightmost node of a part of size α and cont(ρ) < cont(η^+), then it is not in $P(\lambda)$ either, because the contents of the nodes of $P(\lambda)$ are consecutive (cf. Lemma 5.3), and because ρ does not have maximal content.

Therefore, the letter *R* encoding ρ is also present in w_i^{ψ} , and simplifies with the *A* encoding H. This implies that Γ^+ is the good addable *i*-node of $\psi(\lambda)$.

Lemma 5.19. Suppose that $\gamma^+ = (a, \lambda_a^c + 1, c)$ is the good addable *i*-node of λ , with $\lambda_a^c > \alpha$. Then

- $\Gamma^+ = (a, \lambda_a^c + 1, c + 1)$ is the good addable *i*-node of $\psi(\lambda)$ if $1 \le c < l$,
- $\Gamma^+ = (a, \lambda_a^c + 1, 1)$ is the good addable *i*-node of $\psi(\lambda)$ if c = l.

Proof. As in Lemma 5.18, Γ is the node of $\psi(\lambda)$ canonically associated to γ . Consider the letter *A* encoding γ^+ . It is the rightmost letter *A* in w_i and does not simplify. By Proposition 5.14, Γ^+ is also encoded by a letter *A* which the rightmost letter *A* in w_i^{ψ} . It remains to show that it does not simplify with any letter *R* either. In fact, as noticed in the previous proofs, the deletion of the pseudoperiod, in the construction of $\psi(\lambda)$, cannot yield any new letter *R*. However, some letters *A* encoding nodes of $P(\lambda)$ can vanish. Denote by η^+ such a node.

Suppose first that η is not the first node of $P(\lambda)$. Since η^+ is addable, there cannot be another element of $P(\lambda)$ above η . Then denote by η_1 the node of $P(\lambda)$ which has content $\operatorname{cont}(\eta_1) = \operatorname{cont}(\eta) + 1$ (i.e. the previous node of $P(\lambda)$). Note that if η is in λ^d , then η_1 is in λ^{d+1} . Suppose now that η is the first node of $P(\lambda)$. Then, similarly, consider the last node of $P(\lambda)$ and denote it by $P(\lambda)$. In each case, either:

- there is a part α just below the part whose rightmost node is η₁, in which case the node just below η₁ is not in P(λ) and yields an addable node in ψ(λ) which plays exactly the same role in ψ(λ) as η in λ;
- or there is no node just below η₁, in which case η₁ is encoded by a letter R which simplifies with the letter A encoding η⁺.

As a consequence, we are ensured that the letter A encoding Γ^+ does not simplify in w_i^{ψ} , and hence Γ^+ is the good addable *i*-node of $\psi(\lambda)$.

We are now ready to prove the following key result.

Theorem 5.20. The map

$$\begin{array}{ccccc} \psi: & \mathscr{C}_{\mathbf{s}} & \longrightarrow & \mathcal{F}_{\xi(\mathbf{s})} \\ & \lambda & \longmapsto & \psi(\lambda) \end{array}$$

is a crystal isomorphism. We call it the reduction isomorphism for cylindric multipartitions.

Proof. We need to prove that for all $i \in [[0, e-1]]$,

$$\tilde{f}_i(\psi(\lambda)) = \psi(\tilde{f}_i(\lambda)). \tag{8}$$

Thanks to the previous lemmas, we know precisely what $\tilde{f}_i(\psi(\lambda))$ is. It remains to understand the right hand side of (8), by looking at the pseudoperiod of $\tilde{f}_i(\lambda)$. Let $P(\lambda) =$ $(\gamma_1, \gamma_2, \ldots, \gamma_e).$

The operator \tilde{f}_i acts on λ either by:

1. Adding a node to a part α whose rightmost node belongs to $P(\lambda)$. This is the setting of Lemma 5.16. Let $\gamma = \gamma_k = (a, \alpha, c)$ be the node of $P(\lambda)$ such that γ^+ is the good addable *i*-node of λ . In this case, we have

$$P(\psi(\lambda)) = (\gamma_1, \gamma_2, \dots, \gamma_{k-1}, \delta, \gamma_{k+1}, \dots, \gamma_e),$$

where:

- $\delta = (b, \alpha, c + 1)$, with $b = a + N_{c+1}^{>\alpha} + N_{c+1}^{\alpha} N_c^{>\alpha}$, if c < l, and $\delta = (b, \alpha, 1)$, with $b = a + N_1^{>\alpha} + N_1^{\alpha} N_l^{>\alpha}$, if c = l.

Indeed, this node δ is the same as the one determined in the proof of Lemma 5.16 (and whose canonically associated node is Δ). The value of the row b is simply computed using the fact that:

- (a) there is no part α above the part of rightmost node γ ,
- (b) all parts α above the part of rightmost node δ in $\tilde{f}_i(\lambda)$ have an element of $P(\lambda)$ as rightmost node.

But the part of $\tilde{f}_i(\lambda)$ whose rightmost node is γ^+ is a part of size greater than α , and is is therefore shifted to the (c + 1)-th component (if c < l), or the first component (if c = l) when building $\psi(\tilde{f}_i(\lambda))$. Hence, by deleting the elements of $P(\psi(\lambda))$ and shifting the parts greater than α , we end up with the same multipartition as $\tilde{f}_i(\psi(\lambda))$, whence the identity $\tilde{f}_i(\psi(\lambda)) = \psi(\tilde{f}_i(\lambda))$. This is illustrated in Example 5.21 below.

- 2. Adding a node to a part $\leq \alpha$ whose rightmost node does not belong to $P(\lambda)$. This is the setting of Lemma 5.18. In this case, we have $P(\psi(\lambda)) = P(\lambda)$. It is then straightforward that $f_i(\psi(\lambda)) = \psi(f_i(\lambda))$.
- 3. Adding a node to a part > α (whose rightmost node necessarily does not belong to $P(\lambda)$). This is the setting of Lemma 5.19. Here, we also have $P(\psi(\lambda)) = P(\lambda)$, as in the previous point.

Example 5.21. We take the same example as 5.17, 5., namely $\lambda = (3.2.1^2, 4.2.1, 2^3)$, **s** = (2, 3, 4), e = 4 and i = 0. Then we have the following constructions:



The bold contents represent the pseudoperiods. This illustrates the commutation between the operators ψ and f_i .

Remark 5.22. Note that the charged multipartition $|\psi(\lambda), \xi(\mathbf{s})\rangle$ that we get is not cylindric anymore in general, because $\xi(\mathbf{s})$ might not be in \mathscr{S}_e . If it is, then it is clear that $\psi(\lambda)$ has one pseudoperiod less that λ .

Remark 5.23. Interestingly, this isomorphism ψ can be seen as a generalisation of the cyclage isomorphism ξ . Indeed, ξ would be the version of ψ for pseudoperiods of width 0 (which can be found in any multipartition, considering that they have infinitely many parts of size 0).

By a simple use of the cyclage isomorphism, we can now easily determine a refinement Ψ of the reduction isomorphism ψ which maps a cylindric multipartition to another cylindric multipartition with one pseudoperiod less.

Let $\lambda \in \mathcal{C}_{\mathbf{s}}$. If $\xi(\mathbf{s}) \notin \mathcal{S}_{e}$, denote by *m* the number of components of $\xi(\mathbf{s})$ equal to $\xi(\mathbf{s})_{l}$. By Remark 5.9 and the proof of Proposition 5.11, $\xi^{1+m}(\mathbf{s}) \in \mathcal{S}_{e}$.

Define $\Psi(\lambda)$ and $\Psi(s)$ in the following way:

- If $\xi(\mathbf{s}) \in \mathscr{S}_e$, then $\Psi(\lambda) := \psi(\lambda)$ and $\Psi(\mathbf{s}) := \xi(\mathbf{s})$
- If $\xi(\mathbf{s}) \notin \mathscr{S}_e$, then $\Psi(\lambda) := (\xi^m \circ \psi)(\lambda)$ and $\Psi(\mathbf{s}) := \xi^{1+m}(\mathbf{s})$.

We denote $\Psi : |\lambda, s\rangle \longmapsto |\Psi(\lambda), \Psi(s)\rangle$. Then by construction, the following result holds:

Proposition 5.24. For all $\lambda \in \mathcal{C}_s$, we have $\Psi(\lambda) \in \mathcal{C}_{\Psi(s)}$. Moreover, Ψ is a $\mathcal{U}'_q(\mathfrak{sl}_e)$ -crystal isomorphism, and $\Psi(\lambda)$ has one pseudoperiod less than λ .

We can now determine the canonical crystal isomorphism φ for cylindric multipartitions. Recall that we have already determined $\varphi(\mathbf{s})$ in Proposition 5.11. It writes $\varphi(\mathbf{s}) = \xi^k(\mathbf{s})$ for some *k* explicitly determined.

Remark 5.25. The integer *t* defined in the proof of Proposition 5.11 is simply the number of pseudoperiods in λ .

Denote *t* the number of pseudoperiods in λ . Applying *t* times Ψ to λ , we end up with a FLOTW multipartition, but charged by an element $\Psi^t(\mathbf{s})$ which might not be in \mathcal{D}_e . We now simply need to adjust it by some iterations of the cyclage isomorphism ξ . Since $\Psi^t(\mathbf{s}) \in \mathcal{S}_e$, we are ensured that $\varphi(\mathbf{s}) = (\xi^u \circ \Psi^t)(\mathbf{s})$ for some $u \in \mathbb{Z}$ easily computable.

Hence, we set

$$\varphi(\lambda) := (\xi^u \circ \Psi^t)(\lambda),$$

and the following theorem is straightforward.

Theorem 5.26. The map

$$egin{array}{cccc} arphi & \colon & \mathscr{C}_{\mathbf{s}} & \longrightarrow & \mathscr{C}_{arphi(\mathbf{s})} \ & |oldsymbol{\lambda}, \mathbf{s}
angle & \longmapsto & |arphi(oldsymbol{\lambda}), arphi(\mathbf{s})
angle \end{array}$$

is the canonical $\mathcal{U}'_{a}(\widehat{\mathfrak{sl}_{e}})$ -crystal isomorphism for cylindric multipartitions.

Remark 5.27. Note that to determine $\varphi(\mathbf{s})$, we could also have built first $|\Psi^t(\lambda), \Psi^t(\mathbf{s})\rangle$, and found *u* such that $\xi^u(\Psi^t(\mathbf{s})) \in \mathcal{D}_e$. Then, we would have set $\varphi(\mathbf{s}) = \xi^u(\Psi^t(\mathbf{s}))$. Clearly, this construction would give the same multicharge as the construction of $\varphi(\mathbf{s})$ in Section 5.2. The point of Section 5.2 is to show that the suitable multicharge is directly computable using only cyclages of \mathbf{s} .

We can therefore express the canonical crystal isomorphism Φ in full generality. Starting from any charged multipartition $|\lambda, \mathbf{s}\rangle$, we first apply **RS**, we get $\lambda' = \mathbf{RS}(\lambda)$. Then, according to Section 4.3, there is an integer *m* such that $(\mathbf{RS} \circ \xi)^m(\lambda')$ is cylindric. Finally, we use ψ^t to delete all pseudoperiods and adjust everything using ξ^u so that we end up in the fundamental domain \mathcal{D}_e (Theorem 5.26).

Hence, we have

$$\Phi = \xi^{u} \circ \psi^{t} \circ (\mathbf{RS} \circ \xi)^{m} \circ \mathbf{RS}$$
$$= \varphi \circ (\mathbf{RS} \circ \xi)^{m} \circ \mathbf{RS}.$$

6 Application to determining the highest weight vertex

Let $\mathbf{s} \in \mathbb{Z}^l$ and $\lambda \in \mathcal{F}_{\mathbf{s}}$.

In this short section, we explain how to determine the highest weight vertex λ in $B(\lambda)$ using the canonical crystal isomorphism, and without having to find a path in the crystal graph. Recall that we have denoted by $\Phi : |\lambda, s\rangle \mapsto |\mu, r\rangle$ the canonical crystal isomorphism. In Section 4.3, we have seen how to find a cylindric multipartition equivalent to λ . In fact, this uses only the crystal isomorphisms **RS** (Section 3.1) and ξ (Proposition 4.4). Then, after defining the reduction isomorphism ψ in Section 5.3, we proved in Theorem 5.26 that the canonical isomorphism φ for cylindric multipartitions consists only of iterations of ψ and ξ .

Hence, in the general case, there are only three different crystal isomorphisms needed to construct Φ , namely **RS**, ξ , and ψ . Amongst them, ξ is the only map which is clearly invertible. However, it is possible, keeping extra information, to make **RS** and ψ invertible.

6.1 Invertibility of the crystal isomorphism RS

First, it is well known (e.g. [4]) that the correspondence read(λ , **s**) $\mapsto \mathscr{P}(\text{read}(\lambda, \mathbf{s}))$ becomes a bijection if we also associate to read(λ , **s**) its "recording symbol", i.e. the symbol with the same shape as $\mathscr{P}(\text{read}(\lambda, \mathbf{s})) =: \mathscr{P}$ in which we put the entry *k* in the spot where a letter appears at the *k*-th step. We denote by $\mathscr{Q}(\text{read}(\lambda, \mathbf{s}))$ or simply \mathscr{Q} this symbol.

Example 6.1. Take s = (2, 0) and $\lambda = (3.2, 3^2)$. Then

$$\mathfrak{B}_{\mathbf{s}}(\boldsymbol{\lambda}) = \begin{pmatrix} 0 & 4 & 5 \\ 0 & 1 & 2 & 5 & 7 \end{pmatrix}.$$

We get read(λ , s) = 54075210. We can thus give the sequence of symbols leading to \mathcal{P} , and, on the right, the corresponding recording symbols, leading to \mathcal{Q} .

	(5)							(1)				
	(4	5)						(1	2)			
	(0	4	5)					(1	2	3)		
	$\begin{pmatrix} 0\\7 \end{pmatrix}$	4	5)					$\begin{pmatrix} 1\\ 4 \end{pmatrix}$	2	3)		
	$\begin{pmatrix} 0\\5 \end{pmatrix}$	4 7	5)					$\begin{pmatrix} 1\\ 4 \end{pmatrix}$	2 5	3)		
	$\begin{pmatrix} 0\\ 2 \end{pmatrix}$	4 5	$\begin{pmatrix} 5\\7 \end{pmatrix}$					$\begin{pmatrix} 1\\ 4 \end{pmatrix}$	2 5	$\begin{pmatrix} 3 \\ 6 \end{pmatrix}$		
	$\begin{pmatrix} 0\\1 \end{pmatrix}$	2 5	4 7	5)				$\begin{pmatrix} 1\\ 4 \end{pmatrix}$	2 5	3 6	7)	
=	$\begin{pmatrix} 0\\ 0 \end{pmatrix}$	1 5	2 7	4	5)		<i>Q</i> =	$\begin{pmatrix} 1\\ 4 \end{pmatrix}$	2 5	3 6	7	8)

Therefore, this extra data \mathcal{Q} turns **RS** into a bijection.

P

6.2 Invertibility of the reduction isomorphism

Recall that the multicharge charging $\psi(\lambda)$ is nothing but $\xi(\mathbf{s})$. Hence, the computation of \mathbf{s} from $\psi(\mathbf{s}) = \xi(\mathbf{s})$ is straightforward. Moreover, starting from $\psi(\lambda)$, it is easy to recover almost the *l*-partitions λ provided we know the width α of the pseudoperiod that has been deleted. In fact:

- 1. This data determines which parts will stay in the same component (namely the ones smaller than or equal to α), and which will be shifted "to the left" (namely the ones greater than α). Moreover, the property on the contents (Proposition 5.14), which says that all nodes must keep the same content, ensures that we can keep the boxes filled in with the same integers.
- 2. It remains to insert the *e* parts of the α -pseudoperiod at the right locations, i.e. so that the diagram obtained is in fact the Young diagram of a charged multipartition (which is possible thank, again, to Proposition 5.14).

Example 6.2. We take, as in Example 5.13, e = 4 and



and we suppose that we know the width of $P(\lambda)$, namely $\alpha = 2$. Then the two steps above give:

1. Shifting the parts greater that α and keeping the same filling:



2. Inserting coherently the 4 missing parts of size 2 (represented in bold type):

(~	6	-	0	0	1.0		6	7	8		8	9	10111213	١
	5	6	1	8	9	10		5	6	-		7	8		
	4	5	6	1	8	9		4	5			6	7		
	3	4					,	3	4		,	5			ļ,
	2							2				4			
l	1							1	1			3		,)

which is indeed equal to λ .

Hence, this extra data turns ψ into an invertible map.

6.3 Finding the highest weight vertex

Both **RS** and ψ being turned into bijections with the appropriate extra information, we can turn the canonical crystal isomorphism $\Phi : B(\lambda) \longrightarrow B(\mathbf{r})$ into an invertible map. We write $\Phi^{-1} : B(\mathbf{r}) \longrightarrow B(\lambda)$ for the inverse map.

Theorem 6.3. Let $\mathbf{s} \in \mathbb{Z}^l$ and $\lambda \in \mathcal{F}_s$. Denote λ the highest weight vertex in $B(\lambda)$. Then

$$\dot{\boldsymbol{\lambda}} = \boldsymbol{\Phi}^{-1}(\boldsymbol{\emptyset}).$$

Proof. Write $\lambda = \tilde{e}_{i_m} \dots \tilde{e}_{i_1}(\lambda)$. Then

$$\Phi(\boldsymbol{\lambda}) = \Phi(\tilde{e}_{i_m} \dots \tilde{e}_{i_1}(\boldsymbol{\lambda})) = \tilde{e}_{i_m} \dots \tilde{e}_{i_1}(\Phi(\boldsymbol{\lambda})) = \boldsymbol{\emptyset},$$

and since Φ is invertible, this gives $\lambda = \Phi^{-1}(\emptyset)$.

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