# THE RIGHT ANGLE TO LOOK AT ORTHOGONAL SETS

#### FRANK O. WAGNER

ABSTRACT. If X and Y are orthogonal hyperdefinable sets such that X is simple, then any group G interpretable in  $X \cup Y$  has a normal hyperdefinable X-internal subgroup N such that G/N is Y-internal; N is unique up to commensurability. In order to make sense of this statement, local simplicity theory for hyperdefinable sets is developed.

# INTRODUCTION

Two definable sets X and Y in some structure are said to be *orthogonal* if every definable subset of  $X \times Y$  is a finite union of *rectangles*, i.e. of subsets of the forme  $U \times V$  with  $U \subseteq X$  and  $V \subseteq Y$  definable. It easily follows that if X and Y are orthogonal groups, every definable subgroup of  $X \times Y$  is of the form  $U \times V$  with  $U \leq X$  and  $V \leq Y$  subgroups. However, the situation is considerably more complicated for a group G definable, or more generally interpretable, in  $X \cup Y$ , as it need not be the sum of a group interpretable in X and a group interpretable in Y. In fact, an example by Berarducci and Mamino [2, Example 1.2] shows that G need not have any subgroup interpretable in either X or Y. However, they prove [2, Theorem 7.1] that if X is superstable of finite and definable Lascar rank, then any group G interpretable in  $X \cup Y$  has a normal subgroup N interpretable in X, such that G/N is interpretable in Y.

In this paper we shall generalize their result to the case where X is merely simple. In this context, definability has to be replaced by type-definability, as even for a definable group the tools of simplicity theory in general only yield type-definable subgroups. In fact, we even have to study *hyperdefinable* groups, since the quotient G/N, for N

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type-definable, will be of that form. We therefore put ourselves in the hyperdefinable context and assume right from the start that our orthogonal sets X and Y are merely hyperdefinable. To this end, we shall include a quick development of local simplicity theory for hyperdefinable sets in section 4.

Another problem is that of parameters. The usual hypothesis would be that of *stable embedding*, i.e. that every hyperdefinable subset of X is hyperdefinable with parameters in X. We shall circumvent this issue by only ever considering parameters from  $X \cup Y$ , as orthogonality automatically yields stable embeddedness of X and of Y in  $X \cup Y$ .

We shall work in a big  $\kappa$ -saturated and strongly  $\kappa$ -homogeneous monster model  $\mathfrak{M}$ , where  $\kappa$  is bigger than any cardinality we wish to consider. We shall not usually distinguish between elements and tuples.

# 1. Orthogonality

**Definition 1.1.** A set X is hyperdefinable over some parameters A if it is of the form Y/E, where Y is a type-definable set in (at most) countably many variables and E a type-definable equivalence relation on Y, both with parameters in A. For an element  $y \in Y$  we denote the class of y modulo E by  $y_E$ . If  $y = (y_0, \ldots, y_n)$  is a tuple, we put  $y_E = ((y_0)_E, \ldots, (y_n)_E)$ ; similarly, we write yEz if  $y_iEz_i$  for all  $i \leq n$ .

For a tuple  $y_E$  in X and some parameters  $B \supseteq A$ , the type  $\operatorname{tp}(y_E/B)$  is given by all partial types over B of the form

$$\exists z \in Y \ [zEy \land \varphi(z)]$$

true of y. It is easy to see that (in the monster model) two tuples in X have the same type over B if and only if they are conjugate by an automorphism fixing B. We similarly define the type over a hyperimaginary set containing A.

Note that we can also consider the type of a hyperimaginary over B as a partial real type over B. If B is hyperimaginary, this corresponds to a partial type over a representative of B, and any two such types (for different representatives) are equivalent. We shall say that a partial type  $\pi(y)$  is a (partial) X-type if  $\pi(y) \vdash y_E \in X$  and  $\pi(y)$  is E-invariant.

If a is hyperimaginary, a *representative* for a is any real (or imaginary) tuple  $\bar{a}$  with  $a \in dcl^{heq}(\bar{a})$ .

From now on, all tuples and parameter sets are hyperimaginary, unless stated otherwise.

**Definition 1.2.** Let X, Y be A-hyperdefinable sets in some structure  $\mathfrak{M}$ . We say that X and Y are *orthogonal* over A, denoted  $X \perp_A Y$ , if for any tuples  $a \in X$  and  $b \in Y$ , the partial type  $\operatorname{tp}(a/A) \cup \operatorname{tp}(b/A)$  determines  $\operatorname{tp}(ab/A)$ .

**Example 1.3.** If  $\mathfrak{M}_1$  and  $\mathfrak{M}_2$  are two structures and  $\mathfrak{N} = \mathfrak{M}_1 \times \mathfrak{M}_2$  with a predicate X for  $\mathfrak{M}$  and a predicate Y for  $\mathfrak{N}$ , then X and Y are orthogonal in  $\mathfrak{N}$ .

**Remark 1.4.** If X and Y are orthogonal type-definable sets and  $Z \subseteq X \times Y$  is relatively definable, then Z is a finite union of rectangles  $A_i \times B_i$ , where  $A_i \subseteq X$  and  $B_i \subseteq Y$  are relatively definable.

Proof: For any  $z = (x, y) \in Z$  we have that  $\operatorname{tp}(x) \cup \operatorname{tp}(y) \vdash (x, y) \in Z$ . By compactness there are relatively definable subsets  $A_z \subseteq X$  in  $\operatorname{tp}(x)$ and  $B_z \subseteq Y$  in  $\operatorname{tp}(y)$  with  $A_z \times B_z \subseteq Z$ . Again by compactness, finitely many of these rectangles suffice to cover Z.

For the rest of this section, X and Y will be orthogonal  $\emptyset$ -hyper-definable sets.

**Proposition 1.5.** If  $X' \subseteq X^{heq}$  and  $Y' \subseteq Y^{heq}$  are hyperdefinable over some parameters  $A \subseteq X^{heq} \cup Y^{heq}$ , then  $X' \perp_A Y'$ .

*Proof:* Suppose A = (a, b) with  $a \in X^{heq}$  and  $b \in Y^{heq}$ , and consider tuples  $a' \in X'$  and  $b' \in Y'$ . Choose representatives  $\bar{a}, \bar{a}' \in X$  of a, a' and  $\bar{b}, \bar{b}' \in Y$  of b, b'. Then  $\operatorname{tp}(\bar{a}\bar{a}') \cup \operatorname{tp}(\bar{b}\bar{b}') \vdash \operatorname{tp}(\bar{a}\bar{a}'\bar{b}\bar{b}')$ .

Now if  $a'' \equiv_A a'$  and  $b'' \equiv_A b'$ , we can find A-conjugates  $\tilde{a}\bar{a}''$  of  $\bar{a}\bar{a}'$  and  $\tilde{b}\bar{b}''$  of  $\bar{b}\bar{b}'$  such that  $a''\tilde{a}\bar{a}'' \equiv_A a'\bar{a}\bar{a}'$  and  $b''\tilde{b}\bar{b}'' \equiv_A b'\bar{b}\bar{b}'$ . By orthogonality of X and Y, we obtain  $\bar{a}''\tilde{a}\bar{b}''\tilde{b} \equiv \bar{a}'\bar{a}\bar{b}'\bar{b}$ , whence  $a''ab''b \equiv a'ab'b$ , and thus  $a''b'' \equiv_A a'b'$ .

Note that this in particular shows that orthogonality is preserved under adding parameters from  $X^{heq} \cup Y^{heq}$ .

**Proposition 1.6.** X is stably embedded in  $X \cup Y$ : For tuples  $a \in X^{heq}$ and  $b \in Y^{heq}$ , every ab-hyperdefinable subset X' of  $X^{heq}$  is hyperdefinable over a.

*Proof:* If  $\Phi(x, a, b)$  hyperdefines X and  $\Psi(y) = \operatorname{tp}(b)$ , put

 $\Phi'(x,a) = \exists y \, [\Psi(y) \land \Phi(x,a,y)].$ 

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Clearly  $\Phi(x, a, b) \vdash \Phi'(x, a)$ . Conversely, suppose  $a' \models \Phi'(x, a)$ , and choose  $b' \models \Psi$  with  $a' \models \Phi(x, a, b')$ . By orthogonality  $a'ab \equiv a'ab'$ , whence  $a' \models \Phi(x, a, b)$ , and  $\Phi'(x, a)$  hyperdefines X'.  $\Box$ 

We put  $\operatorname{dcl}_X^{heq}(A) = \operatorname{dcl}^{heq}(A) \cap X^{heq}$  and  $\operatorname{bdd}_X(A) = \operatorname{bdd}(A) \cap X^{heq}$ .

**Corollary 1.7.** Suppose  $a \in X^{heq}$  and  $b \in Y^{heq}$ . Then

$$\operatorname{dcl}_X^{heq}(a,b) = \operatorname{dcl}_X^{heq}(a)$$
 and  $\operatorname{bdd}_X(a,b) = \operatorname{bdd}_X(a)$ .

*Proof:* Immediate from Proposition 1.6.

# 2. Weak elimination of hyperimaginaries

In this section, X and Y will be  $\emptyset$ -hyperdefinable sets.

**Definition 2.1.** Let Z be  $\emptyset$ -hyperdefinable. We say that Z has weak elimination of hyperimaginaries with respect to  $X^{heq}$  and  $Y^{heq}$  if for every  $z \in (X \cup Y)^{heq}$  there is some  $x \in X^{heq}$  and  $y \in Y^{heq}$  with  $xy \in bdd(z)$  and  $z \in dcl^{heq}(xy)$ .

**Theorem 2.2.** If X and Y are orthogonal, then  $X \cup Y$  has weak elimination of hyperimaginaries with respect to  $X^{heq}$  and  $Y^{heq}$ .

*Proof:* Consider  $z \in (X \cup Y)^{heq}$ , say  $z = (z_X, z_Y)_E$  for some tuples  $z_X \in X, z_Y \in Y$  and type-definable equivalence relation E. For  $x \equiv z_X$  consider the hyperdefinable equivalence relation  $E_x$  on  $tp(z_Y)$  given by

$$yE_xy' \quad \Leftrightarrow \quad (x,y)E(x,y').$$

Then  $E_x$  is  $\emptyset$ -hyperdefinable by Proposition 1.6, and we may replace  $z_Y$  by  $z_Y/E_x$ . Similarly, we may replace  $z_X$  by  $z_X/E_y$ . We claim that now E has bounded classes on  $\operatorname{tp}(z_X) \times \operatorname{tp}(z_Y)$ . If not, there is a nonconstant indiscernible sequence  $(x_i, y_i : i < \omega)$  in some E-class. By orthogonality, for i < j

$$\operatorname{tp}(x_i, x_j) \cup \operatorname{tp}(y_i, y_j) \vdash \operatorname{tp}((x_i, y_i), (x_j, y_j)).$$

But  $tp(y_i, y_j) = tp(y_i, y_k)$  for i < j < k, whence

$$tp((x_i, y_i), (x_j, y_j)) = tp((x_i, y_i), (x_j, y_k))$$

and  $(x_j, y_k)E(x_i, y_i)E(x_j, y_j)$  holds. By triviality of  $E_x$  we get  $y_j = y_k$ , whence  $x_j = x_k$  by triviality of  $E_y$ , a contradiction.

Thus 
$$(z_X, z_Y) \in bdd(z)$$
, and trivially  $z \in dcl^{heq}(z_X, z_Y)$ .

We shall put  $bdd_{XY}(A) = bdd(A) \cap (X^{heq} \cup Y^{heq}).$ 

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**Corollary 2.3.** If  $X \perp Y$ , then  $\operatorname{bdd}_{XY}(A)$  and  $\operatorname{bdd}(A) \cap (X \cup Y)^{heq}$ are interdefinable for any  $A \subset X \cup Y$ . Moreover, for  $a \in (X \cup Y)^{heq}$ we have  $\operatorname{tp}(a/\operatorname{bdd}_{XY}(A)) \vdash \operatorname{tp}(a/\operatorname{bdd}(A))$ .

*Proof:* The first assertion is immediate from weak elimination of hyperimaginaries. For the second assertion, note that equality of type over bdd(A) is a bounded type-definable equivalence relation, and the class of a is definable over  $bdd_{XY}(A)$ .

### 3. INTERNALITY AND ANALYSABILITY

Again, X and Y will be  $\emptyset$ -hyperdefinable sets.

**Definition 3.1.** We say that X is (almost) Y-internal if there is some parameter set A such that for every  $a \in X$  there is a tuple  $b \in Y$  with  $a \in dcl^{heq}(Ab)$  (or  $a \in bdd(Ab)$ , respectively).<sup>1</sup>

If the parameters A can be chosen in some set Z, we say that X is Y-internal within Z.

We say that X is Y-analysable (within Z) if there is a sequence  $(a_i : i < \alpha)$  such that  $\operatorname{tp}(a_i/a_j : j < i)$  is Y-internal (within Z) for every  $i < \alpha$ , and  $a \in \operatorname{bdd}(a_i : i < \alpha)$ .

**Proposition 3.2.** Suppose that X and Y are orthogonal. If an  $\emptyset$ -hyperdefinable set X' is X-analysable within  $X \cup Y$ , then X' is almost X-internal within  $\operatorname{bdd}_Y(\emptyset)$ ; if X' is X-internal within  $X \cup Y$ , then X' is X-internal within  $\operatorname{bdd}_Y(\emptyset)$ .

*Proof:* We first show the assertion for (almost) X-internal X'. So suppose  $\bar{a} \in X$  and  $\bar{b} \in Y$  are such that for every  $x \in X'$  there is a tuple  $a \in X$  with  $x \in bdd(\bar{a}\bar{b}a)$ . Let  $\bar{x}$  be the set of  $\bar{a}\bar{b}a$ -conjugates of x, and E the type-definable equivalence relation on  $tp(\bar{a}\bar{b}a)$  given by

$$(\bar{a}'\bar{b}'a')E(\bar{a}''\bar{b}''a'') \Leftrightarrow \exists \bar{x}' \ \bar{x}'\bar{a}'\bar{b}'a' \equiv \bar{x}'\bar{a}''\bar{b}''a'' \equiv \bar{x}\bar{a}\bar{b}a.$$

Clearly,  $\bar{x}$  is interdefinable with  $(\bar{a}ba)_E$ . By weak elimination of hyperimaginaries, there is  $\tilde{a} \in X^{heq}$  and  $\tilde{b} \in Y^{heq}$  with  $\tilde{a}\tilde{b} \in \text{bdd}(\bar{x})$  and  $\bar{x} \in \text{dcl}^{heq}(\tilde{a}\tilde{b})$ . Note that if X' is X-internal and  $x \in \text{dcl}^{heq}(\bar{a}\bar{b}a)$ , then  $x = \bar{x}$ .

Suppose  $\tilde{b} \notin \text{bdd}_Y(\emptyset)$ . Then there is an  $\emptyset$ -conjugate  $\tilde{b}'$  of  $\tilde{b}$  outside  $\text{bdd}(\bar{a}\bar{b})$ ; if  $\sigma$  is an automorphism mapping  $\tilde{b}'$  to  $\tilde{b}$ , put  $\bar{a}'\bar{b}' = \sigma(\bar{a}\bar{b})$ .

<sup>&</sup>lt;sup>1</sup>In simplicity theory, this is called *finite generation*; for *internality* we would require for every  $a \in X$  the existence of some  $A \perp a$  and tuple  $b \in Y$  with  $a \in dcl^{heq}(Ab)$ .

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Then  $\tilde{b} \notin \text{bdd}(\bar{a}'\bar{b}')$ . On the other hand, since  $\bar{a}'\bar{b}' \equiv \bar{a}\bar{b}$  and  $\bar{x} \subset X'$ , there is  $a' \in X$  with  $x' \in \text{bdd}(\bar{a}'\bar{b}'a')$  for all  $x' \in \bar{x}$ . Then in particular  $\tilde{b} \in \text{bdd}(\bar{a}'\bar{b}'a')$ , whence  $\tilde{b} \in \text{bdd}(\bar{b}')$  by Corollary 1.7, a contradiction.

Now assume that  $x \in X'$  and  $(x_i : i < \alpha)$  is an X-analysis of x within  $X \cup Y$ . We show inductively on i that  $\operatorname{tp}(x_j : j < i)$  is X-internal within  $\operatorname{bdd}_Y(\emptyset)$ . So suppose  $\operatorname{tp}(x_j : j < k)$  is X-internal within  $\operatorname{bdd}_Y(\emptyset)$  for all k < i. If i is limit, then clearly  $\operatorname{tp}(x_j : j < i)$  is X-internal within  $\operatorname{bdd}_Y(\emptyset)$ . If i = k + 1, then  $\operatorname{tp}(x_k/x_j : j < k)$  is X-internal within  $\operatorname{bdd}_Y(x_j : j < k)$ , and there is  $a \in X$  with

$$x_k \in \operatorname{dcl}^{heq}(a, \operatorname{bdd}_Y(x_j : j < k), x_j : j < k).$$

Or, by X-internality of  $\operatorname{tp}(x_j : j < k)$  within  $\operatorname{bdd}_Y(\emptyset)$  there is  $a' \in X$  with  $(x_j : j < k) \in \operatorname{dcl}^{heq}(a', \operatorname{bdd}_Y(\emptyset))$ . Then

 $\operatorname{bdd}_Y(x_j : j < k) \subseteq \operatorname{bdd}_Y(a', \operatorname{bdd}_Y(\emptyset)) = \operatorname{bdd}_Y(\emptyset)$ 

by Corollary 1.7, and  $x_k \in bdd(a, a', bdd_Y(\emptyset))$ . So  $tp(x_j : j < i)$  is X-internal within  $bdd_Y(\emptyset)$ , and tp(x) is almost X-internal within  $bdd_Y(\emptyset)$ .

**Corollary 3.3.** Suppose that X and Y are orthogonal. Let X' and Y' be  $\emptyset$ -hyperdefinable. If X' is X-internal within  $X \cup Y$  and Y' is Y-internal within  $X \cup Y$ , then  $X' \perp_{\text{bdd}_{XY}(\emptyset)} Y'$ .

*Proof:* This is immediate from Propositon 3.2 and 1.5.

**Corollary 3.4.** Suppose that X and Y are orthogonal. If an  $(X \cup Y)$ -hyperdefinable set Z is almost X- and almost Y-internal within  $X \cup Y$ , then it is bounded.

Proof: We may assume that Z is hyperdefinable over  $\emptyset$ . Let  $z \in Z$ . Since Z is almost X-internal, there is  $b \in \text{bdd}_Y(\emptyset)$  and  $x \in X$  such that  $z \in \text{bdd}(xb)$ . Let  $\bar{z}$  be the set of xb-conjugates of z. Then  $\text{tp}(\bar{z})$  is X-internal within  $\text{bdd}_Y(\emptyset)$ , and still almost Y-internal within  $X \cup Y$ . So there is  $a \in \text{bdd}_X(\emptyset)$  and  $y \in Y$  such that  $\bar{z} \in \text{bdd}(ya)$ . Let  $\bar{z}$  be the set of ya-conjugates of  $\bar{z}$ . Then  $\text{tp}(\bar{z})$  is Y-internal within  $\text{bdd}_X(\emptyset)$ , and still X-internal within  $\text{bdd}_Y(\emptyset)$ . Hence  $\text{tp}(\bar{z}) \perp_{\text{bdd}_{XY}(\emptyset)} \text{tp}(\bar{z})$ , so  $\text{tp}(\bar{z})$  is bounded. Thus Z is bounded as well.  $\Box$ 

**Corollary 3.5.** Suppose that X and Y are orthogonal. If  $Z \subseteq (X \cup Y)^{heq}$  is  $\emptyset$ -hyperdefinable and almost X-internal within  $X \cup Y$ , then it is X-internal within  $\operatorname{bdd}_Y(\emptyset)$ .

*Proof:* By weak elimination of hyperimaginaries there is  $x \in X^{heq}$  and  $y \in Y^{heq}$  with  $z \in dcl^{heq}(xy)$  and  $xy \in bdd(z)$ . So tp(y) is Y-internal and almost X-internal, whence bounded.

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#### 4. Local simplicity

**Definition 4.1.** Let  $A \subseteq B$ , and  $\pi(x, B)$  be a partial type over B. We say that  $\pi(x, B)$  does not divide over A if for any indiscernible sequence  $(B_i : i < \omega)$  in  $\operatorname{tp}(B/A)$  the partial type

$$\bigcup_{i<\omega}\pi(x,B_i)$$

is consistent. Clearly, tp(a/B) divides over A if and only if  $tp(a_0/B)$  does so for some finite subtuple  $a_0 \subseteq a$ .

**Example 4.2.** If  $tp(a) \perp tp(b)$ , then tp(a/b) does not divide over  $\emptyset$ .

We now define the appropriate version of local rank. We follow Ben Yaacov's terminology [1, Definition 1.4], more general than [5, Definiton 4.3.5].

**Definition 4.3.** Let  $\pi(x)$ ,  $\Phi(x, y)$  and  $\Psi(y_1, \ldots, y_k)$  be partial types in (at most) countably many variables.

(1)  $\Psi$  is a *k*-inconsistency witness for  $\Phi$  if

$$\models \forall y_1 \dots y_k \neg \exists x \left[ \Psi(y_1, \dots, y_k) \land \bigwedge_{i=1}^k \Phi(x, y_i) \right].$$

- (2) The local  $(\Phi, \Psi)$ -rank  $D(., \Phi, \Psi)$  is defined on partial types in x as follows:
  - $D(\pi(x), \Phi, \Psi) \ge 0$  if  $\pi(x)$  is consistent.
  - $D(\pi(x), \Phi, \Psi) \ge n + 1$  if there is a sequence  $(a_i : i < \omega)$ such that  $\models \Psi(\bar{a})$  for any k-tuple  $\bar{a} \subset (a_i : i < \omega)$ , and  $D(\pi(x) \land \Phi(x, a_i), \Phi, \Psi) \ge n$  for all  $i < \omega$ .
  - If  $D(\pi, \Phi, \Psi) \ge n$  for all  $n < \omega$ , we put  $D(\pi, \Phi, \Psi) = \infty$ .

Note that in Definition 4.3(2) we may require  $(a_i : i < \omega)$  to be indiscernible. Moreover,  $D(\pi(x, a), \Phi, \Psi) \ge n$  is a closed condition on a, and  $D(\operatorname{tp}(x/a), \Phi, \Psi) \ge n$  is a closed condition on x over a.

**Definition 4.4.** Let I be an ordered set. A sequence  $I = (a_i : i \in I)$  is *independent over* A, or A-*independent*, if  $tp(a_i/A, a_j : j < i)$  does not divide over A for all  $i \in I$ . If  $A \subseteq B$  and  $p \in S(B)$ , the sequence  $(a_i : i \in I)$  is a Morley sequence in p over A if it is B-indiscernible,  $a_i \models p$  and  $tp(a_i/B, a_j : j < i)$  does not divide over A for all  $i \in I$ . If  $A \subseteq B$  are  $p \in S(B)$ , the sequence  $a_i \models p$  and  $tp(a_i/B, a_j : j < i)$  does not divide over A for all  $i \in I$ . If A = B, we simply call it a Morley sequence in p.

**Fact 4.5.** [5, Corollary 3.2.5] or [4, Proposition 16.12] If tp(b/cd) does not divide over d and tp(a/cbd) does not divide over bd, then tp(ab/cd) does not divide over d.

**Theorem 4.6.** Let X be a hyperdefinable set over  $\emptyset$ . The following are equivalent:

- (1) Symmetry holds on X: For all  $a, b, c \in X$ , tp(a/bc) does not divide over b if and only if tp(c/ab) does not divide over b.
- (2) Transitivity holds on X: If  $a, b, c, d \in X$ , then tp(a/bcd) does not divide over b if and only if tp(a/bc) does not divide over b and tp(a/bcd) does not divide over bc.
- (3) Local character holds on X: There is  $\kappa$  such that for all countable  $a \in X$  and  $A \subset X$  there is  $A_0 \subseteq A$  with  $|A_0| \leq \kappa$  such that  $\operatorname{tp}(a/A)$  does not divide over  $A_0$ . In fact, we can take  $\kappa = 2^{|T|}$ .
- (4)  $D(., \Phi, \Psi) < \infty$  for any partial X-type  $\Phi(x, y)$  and inconsistency witness  $\Psi$  for  $\Phi$ .
- (5) For any  $A \subseteq B \subset X$ , a partial X-type  $\pi(x, B)$  does not divide over A if and only if there is a Morley sequence I in  $\operatorname{tp}(B/A)$ such that  $\{\pi(x, B') : B' \in I\}$  is consistent.

If any of these conditions is satisfied, then for all  $A \subseteq B \subset X$  and  $a \in X$  the type  $\operatorname{tp}(a/B)$  does not divide over A if and only if

$$D(\operatorname{tp}(a/B), \Phi, \Psi) = D(\operatorname{tp}(a/A), \Phi, \Psi)$$

for all  $(\Phi, \Psi)$ . Moreover, Extension holds on X: For any partial Xtype  $\pi(x)$  over B, if  $\pi$  does not divide over A then it has a completion which does not divide over A.

Proof: (1)  $\Rightarrow$  (2) Clearly, if tp(a/bcd) does not divide over b, it does not divide over bc and tp(a/bc) does not divide over b. Conversely, suppose that tp(a/bcd) does not divide over bc and tp(a/bc) does not divide over b. By symmetry, tp(d/abc) does not divide over bc and tp(c/ab) does not divide over b. By Fact 4.5 tp(cd/ab) does not divide over b, so again by symmetry tp(a/bcd) does not divide over b.

(2)  $\Rightarrow$  (4) Suppose there is a partial X-type  $\Phi$  and an inconsistency witness  $\Psi$  for  $\Phi$  such that  $D(x = x, \Phi, \Psi) = \infty$ . This means that for all  $n < \omega$  there are

$$(b_i, a_i^j : i < n, j < \omega)$$

such that  $\models \Psi(\bar{a})$  for all i < n and  $\bar{a} \subset (a_i^j : j < \omega)$  of the right length,  $b_i \models \bigwedge_{k \le i} \Phi(x, a_k^0)$ , and  $(a_i^j : j < \omega)$  is indiscernible over  $\{b_k a_k^0 : k < i\}$ .

By Ramsey's Theorem and compactness we may assume that i is indexed by  $\{\pm(1+\frac{1}{n}): n > 0\}$  and that the sequence  $(b_i a_i^0)_i$  is indiscernible. Put  $I = \{b_i a_i^0: i < -1\}$  and  $J = \{b_i a_i^0: i > 1\}$ . Then  $\operatorname{tp}(b_1/IJ)$  does not divide over I and  $\operatorname{tp}(b_1/IJa_{-1})$  does not divide over IJ, since the former is finitely satisfiable in I and the latter in J. However,  $(a_{-1}^j: j < \omega)$  witnesses that  $\Phi(x, a_{-1})$ , and hence  $\operatorname{tp}(b_1/IJa_{-1})$ , divides over I, contradicting transitivity.

(4) 
$$\Rightarrow$$
 (3) Assume (4). First, we note that for  $A \subseteq B \subset X$ , if  
 $D(\operatorname{tp}(a/B), \Phi, \Psi) = D(\operatorname{tp}(a/A), \Phi, \Psi)$ 

for all  $(\Phi, \Psi)$ , then  $\operatorname{tp}(a/B)$  does not divide over A. This is obvious, as if some A-indiscernible sequence  $(B_i : i < \omega)$  in  $\operatorname{tp}(B/A)$  witnesses dividing, we can take  $\Phi(x, y) = \operatorname{tp}(a, B)$  and  $\Psi = \operatorname{tp}(B_1, B_2, \ldots, B_n)$ for  $n < \omega$  sufficiently large. Then  $\Psi$  is an *n*-inconsistency witness (clearly, we may restrict to countable B), and

$$D(\operatorname{tp}(a/B), \Phi, \Psi) < D(\operatorname{tp}(a/A), \Phi, \Psi).$$

Given  $\operatorname{tp}(a/A)$  it is hence enough to take  $A_0 \subseteq A$  big enough such that

$$D(\operatorname{tp}(a/A), \Phi, \Psi) = D(\operatorname{tp}(a/A_0), \Phi, \Psi)$$

for all  $(\Phi, \Psi)$ . There are only  $2^{|T|}$  such pairs, so we need at most that many parameters.

(3)  $\Rightarrow$  (4) Suppose  $D(x = x, \Phi, \Psi) = \infty$ . Then for any cardinal  $\kappa$  we can find an indiscernible sequence  $(a_i b_i : i \leq \kappa^+)$  as in (2)  $\Rightarrow$  (4), such that  $b_{\kappa^+} \models \bigwedge_{i \leq \kappa^+} \Phi(x, a_i)$  and  $\Phi(x, a_i)$  divides over  $\{a_j : j < i\}$  for all  $i \leq \kappa^+$ . Then  $\operatorname{tp}(b_{\kappa^+}/a_i : i < \kappa^+)$  divides over any subset of its domain of cardinality  $\leq \kappa$ .

 $(4) \Rightarrow (5)$ . Assume (4). Given  $a \in X$  and  $A \subseteq B \subset X$ , for any pair  $(\Phi, \Psi)$  and any formula  $\varphi(y)$  we can adjoin either  $\exists y [xEy \land \varphi(y)]$  or  $\exists y [xEy \land \neg \varphi(y)]$  and preserve  $D(., \Phi, \Psi)$ -rank. By compactness we can thus complete  $\operatorname{tp}(a/A)$  to a type p over B of the same  $D(., \Phi, \Psi)$ -rank. In particular, no  $\Phi$ -instance in p divides over A with  $\Psi$  as inconsistency witness. Coding finitely many pairs  $(\Phi_i, \Psi_i : i < n)$  in a single one, one obtains an extension such that no  $\Phi_i$ -instance  $\Psi_i$ -divides for any i < n; by compactness we can do this for all pairs  $(\Phi, \Psi)$  simultaneously and obtain an extension which does not divide over A. Taking for B a sufficiently saturated model, a sequence  $(a_i : i < \omega) \subset B$  such that  $a_i \models p \upharpoonright_{(A,a_j:j < i)}$  is a Morley sequence in  $\operatorname{tp}(a/A)$ .

This shows in particular that if  $\pi(x, B)$  does not divide over A, then there is a Morley sequence I in  $\operatorname{tp}(B/A)$  such that  $\{\pi(x, B') : B' \in I\}$ is consistent.

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Conversely, suppose that  $\pi(x, B)$  divides over A, as witnessed by an A-indiscernible sequence  $(B_i : i < \omega)$  in  $\operatorname{tp}(B/A)$  with  $\bigcup_{i < \omega} \pi(x, B_i)$  inconsistent. Take any Morley sequence I in  $\operatorname{tp}(B/A)$ . By [5, Corollary 2.2.8] (which is shown there for real tuples, but transfers easily to hyperimaginaries) we may assume that  $B_i \cap I$  is A-conjugate to I for all  $i \in I$  and that  $(B_i : i < \omega)$  is indiscernible over AI. If  $\overline{\pi}(x) = \bigcup_{B' \in I} \pi(x, B')$  were consistent, then  $(B_i : i < \omega)$  would witness that

$$D(\bar{\pi}(x) \wedge \pi(x, B_0), \pi(x, y), \Psi) < D(\bar{\pi}(x), \pi(x, y), \Psi)$$

for some inconsistency witness  $\Psi$ . But by A-conjugacy the two ranks must be equal, a contradiction.

 $(5) \Rightarrow (1)$  Let us first show *Extension*. If  $A \subseteq B \subset X$  and  $\pi(x, B)$  is a partial X-type which does not divide over A, let  $(B_i : i < \alpha)$  be a very long Morley sequence in  $\operatorname{tp}(B/A)$ . Consider any realization  $a \models \bigwedge_{i < \alpha} \pi(x, B_i)$ . Since  $\alpha$  is large, there is an infinite subset  $J \subset \alpha$  such that  $p(aAX) = \operatorname{tp}(aAB_i)$  is constant for  $i \in J$ . Then p(xAB) is a completion of  $\pi$  which does not divide over A, as witnessed by  $(B_i : i \in J)$ .

Now given  $a, b, c \in X$  such that  $\operatorname{tp}(a/bc)$  does not divide over b, let  $B \ni bc$  be a sufficiently saturated model, and p an extension of  $\operatorname{tp}(a/bc)$  to B which does not divide over b. Choose a sequence  $(a_i : i < \omega) \subset B$  such that  $a_i \models p \upharpoonright_{(bc,a_j:j < i)}$ . This is a Morley sequence in  $\operatorname{tp}(a/bc)$  over b. Then  $(a_i : i < \omega)$  is a Morley sequence in  $\operatorname{tp}(a/bc)$ , and  $a_i \models \operatorname{tp}(a/bc)$  for all  $i < \omega$ . Hence  $\operatorname{tp}(c/ba)$  does not divide over b, and symmetry holds.

Finally we show that if (1) - (5) hold and  $\operatorname{tp}(a/B)$  does not divide over A for  $A \subseteq B \subset X$  and  $a \in X$ , then  $D(\operatorname{tp}(a/A), \Phi, \Psi) \ge n$  implies  $D(\operatorname{tp}(a/B), \Phi, \Psi) \ge n$  for all  $(\Phi, \Psi)$ . For n = 0 this is obvious. So suppose  $D(\operatorname{tp}(a/A), \Phi, \Psi) \ge n + 1$ . Then there is  $(d_i : i < \omega)$  indiscernible over A such that  $\overline{d} \models \Psi$  for all  $\overline{d} \subset (d_i : i < \omega)$  of the right length, and  $D(\operatorname{tp}(a/A) \land \Phi(x, d_i), \Phi, \Psi) \ge n$  for all  $i < \omega$ . Let q be a completion of  $\operatorname{tp}(a/A) \land \Phi(x, d_0)$  with  $D(q, \Phi, \Psi) \ge n$ . Clearly, we may assume  $a \models q$ , and that  $\operatorname{tp}(d_0/aB)$  does not divide over aA. As  $\operatorname{tp}(a/B)$  does not divide over A, by symmetry and transitivity  $\operatorname{tp}(ad_0/B)$  does not divide over A, and  $\operatorname{tp}(a/d_0B)$  does not divide over  $d_0A$ . By induction hypothesis,

$$D(\operatorname{tp}(a/d_0A), \Phi, \Psi) \ge n$$
 implies  $D(\operatorname{tp}(a/d_0B), \Phi, \Psi) \ge n$ .

As  $\operatorname{tp}(d_0/B)$  does not divide over A and  $(d_i : i < \omega)$  is A-indiscernible, we may assume that it remains indiscernible over B. But then it witnesses

$$D(\operatorname{tp}(a/B), \Phi, \Psi) \ge D(\operatorname{tp}(a/d_0B), \Phi, \Psi) + 1 \ge n + 1.$$

**Definition 4.7.** An A-hyperdefinable set X is simple (over A) if it satisfies any of the conditions of Theorem 4.6 when we adjoin A to the language. If X is simple over A and  $a, b, c \in X$ , we shall say that a and c are independent over Ab, written  $a op_{Ab} c$ , if tp(a/Abc) does not divide over Ab.

Note that we only allow tuples and parameters from  $A \cup X$ . If X is stably embedded, we can of course allow parameters from anywhere.

**Remark 4.8.** If X is merely hyperdefinable, it may be simple although no definable or even type-definable imaginary set in the ambient structure is simple.

If X is simple, it is now standard to extend the notions of dividing and independence to hyperimaginaries in  $X^{heq}$ . Moreover, we can develop basic simplicity theory (canonical bases, the independence theorem, stratified ranks, generic types, stabilizers, see [5]) within  $X^{heq}$ , replacing the notion of a model by a boundedly closed subset C of  $X^{heq}$ such that  $X^{heq}$ -types over C are finitely satisfiable in C.

**Proposition 4.9.** Let X and Y be orthogonal  $\emptyset$ -hyperdefinable sets such that X is simple over  $\emptyset$ . If  $A \subset Y$  is a set of parameters, then X is simple over A, and over  $bdd_Y(A)$ . In particular, let Z be a set hyperdefinable over some parameters  $A \subset X \cup Y$ . If Z is X-internal within  $X \cup Y$ , then Z is simple over A; if  $Z \subseteq (X \cup Y)^{heq}$  is almost X-internal, then Z is simple as well.

*Proof:* Simplicity over A is obvious from orthogonality; simplicity over  $\operatorname{bdd}_Y(A)$  follows. Now if Z is X-internal within  $X \cup Y$ , it lives in  $\operatorname{dcl}^{eq}(X, \operatorname{bdd}_Y(A))$  by Corollary 3.2, and must be simple as well; if  $Z \subseteq (X \cup Y)^{heq}$  is almost X-internal within  $X \cup Y$ , it is X-internal within  $\operatorname{bdd}_Y(\emptyset)$  by Corollary 3.5.

# 5. Groups interpretable in orthogonal sets

Recall that two hyperdefinable subgroups  $H_1$  and  $H_2$  of some group G are *commensurable* if  $H_1 \cap H_2$  has bounded index both in  $H_1$  and in  $H_2$ .

**Theorem 5.1.** Suppose X and Y are orthogonal  $\emptyset$ -hyperdefinable sets in a structure  $\mathfrak{M}$ , and G is an  $\emptyset$ -hyperdefinable group in  $(X \cup Y)^{heq}$ . If X is simple over  $\emptyset$ , there is an  $\emptyset$ -hyperdefinable normal X-internal subgroup N of G such that the quotient G/N is Y-internal. N is unique up to commensurability.

*Proof:* Let us first show uniqueness: If N' is a second such group, then  $N/(N \cap N')$  and  $N'/(N \cap N')$  are X-internal and Y-internal, and hence bounded by orthogonality of X and Y. Thus N and N' are commensurable.

By Theorem 2.2 every element  $g \in G$  is of the form  $(g_X, g_Y)_E$  for some  $g_X \in X^{heq}$  and  $g_Y \in Y^{heq}$ , both bounded over g, and some typedefinable equivalence relation E with bounded classes, depending on  $\operatorname{tp}(g)$ . Hence  $\operatorname{tp}(g/g_Y)$  is X-internal and  $\operatorname{tp}(g/g_X)$  is Y-internal. Now if  $h = (h_X, h_Y)_E$  and  $gh = ((gh)_X, (gh)_Y)_E$ , then

$$(gh)_X \in \mathrm{bdd}(g_X, g_Y, h_X, h_Y),$$

whence  $(gh)_X \in \text{bdd}(g_X, h_X)$  by Corollary 1.7. Similarly  $(gh)_Y \in \text{bdd}(g_Y, h_Y)$ .

Now X is simple, as is  $\operatorname{tp}(g, h/g_Y, h_Y)$  for any  $g, h \in G$ . Hence we can consider  $g, h \in G$  such that  $g \bigcup_{g_Y, h_Y} h$ . Then for any stratified local rank D

(†) 
$$D(gh/(gh)_Y) \ge D(gh/(gh)_Y, g_Y, h_Y, g) = D(h/g_Y, h_Y, g) = D(h/g_Y, h_Y, g) = D(h/g_Y, h_Y) = D(h/h_Y),$$

as also  $h 
ightarrow g_Y$  by orthogonality (Example 4.2). Similarly

$$D(gh/(gh)_Y) \ge D(g/g_Y).$$

Now suppose G is a subset of  $(X^m \times Y^n)/E$ , where m, n are at most countable. Then if  $g = (\bar{x}_g, \bar{y}_g)_E$  we have  $g 
eq g_Y \bar{y}_g$ , and  $g_Y \in \text{bdd}(\bar{y}_g)$ , whence  $D(g/g_Y) = D(g/\bar{y}_g)$ . By compactness, there is a G-type  $p((\bar{x}, \bar{y})_E)$ implying that  $D((\bar{x}, \bar{y})_E/\bar{y})$  is maximal for all local stratified ranks. But if  $g, h \models p$  with  $g 
eq g_{Y,h_Y} h$ , then we must have equality in  $(\dagger)$ , which implies that g, h and gh are pairwise independent over  $g_Y, h_Y, (gh)_Y$ . Put

$$S = \operatorname{stab}(g/g_Y, h_Y, (gh)_Y).$$

By [3, Lemme 1.2 and Remarque 1.3] the element g is generic in the coset Sg, and this coset is hyperdefinable over  $bdd(g_Y, h_Y, (gh)_Y)$ . Moreover, S is X-internal since  $tp(g/g_Y)$  is.

Suppose that S is not commensurable with  $S^h$  for some  $h \in G$ . Then  $SS^h$  is still X-internal, with

$$D(SS^h) \ge D(S) = D(p)$$

for every stratified local rank D, and for at least one such rank  $D_0$  we have  $D_0(SS^h) > D_0(p)$ . Choose  $g' \in SS^h$  with  $D_0(g'/h) = D_0(SS^h)$ . By pre-multiplying with a generic element of S and post-multiplying with a generic element of  $S^h$ , the inequality ( $\dagger$ ) implies that we may assume  $D(g'/h) \ge D(p)$  for every stratified local rank D. However,  $\operatorname{tp}(g'/h)$  is X-internal, so  $g'_Y \in \operatorname{bdd}(g')$  implies  $g'_Y \in \operatorname{bdd}(h)$  by Lemma 1.7 and Proposition 3.2. Thus

$$D(g'/g'_Y) \ge D(g'/h) \ge D(p)$$
 and  $D_0(g'/g'_Y) \ge D_0(g'/h) > D_0(p)$ ,

contradicting our choice of p. Hence S is commensurable with all its conjugates, and the locally connected component N of S is normal in G [5, Definition 4.5.15 and Corollary 4.5.16]. In fact the same proof, with NZ instead of  $SS^h$ , shows that if Z is an X-internal hyperdefinable subset of G, then Z is covered by boundedly many cosets of N. In particular for any  $g' \in G$  the type  $\operatorname{tp}(g'/g'_Y)$  is covered by boundedly many cosets of N. But then  $g'N \in \operatorname{bdd}(g'_Y)$ , and G/N is almost Y-internal, whence Y-internal by Corollary 3.5.

Question 5.2. If X and Y are orthogonal type-definable sets and G is a relatively definable group in  $(X \cup Y)^{eq}$ , can we find a relatively definable normal X-internal subgroup N such that G/N is Y-internal?

Question 5.3. What can we say if neither X nor Y is simple? Is it true that in every hyperdefinable subgroup of  $((X \cup Y)^{heq})$  there is a maximal normal hyperdefinable X-internal subgroup  $N_X$ , a maximal normal hyperdefinable Y-internal subgroup  $N_Y$ , an X-internal hyperdefinable local group  $G_X$ , a Y-internal hyperdefinable local group  $G_Y$  and a hyperdefinable locally bounded equivalence relation E on  $G_X \times G_Y$  such that  $G/(N_X N_Y)$  is isogenous, or even isomorphic, to  $(G_X \times G_Y)/E$ ?

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Université de Lyon; CNRS; Université Lyon 1; Institut Camille Jordan UMR5208, 43 bd du 11 novembre 1918, 69622 Villeurbanne Cedex, France

E-mail address: wagner@math.univ-lyon1.fr