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CODING MULTITYPE BRANCHING FORESTS: APPLICATION TO THE LAW OF THE TOTAL PROGENY

LOIĆ CHAUMONT AND RONGLI LIU

ABSTRACT. By extending the breadth first search algorithm to any d -type critical or subcritical irreducible branching tree, we show that such trees may be encoded through d independent, integer valued, d -dimensional random walks. An application of this coding together with a multivariate extension of the Ballot Theorem allow us to give an explicit form of the law of the total progeny, jointly with the number of subtrees of each type, in terms of the offspring distribution of the branching process.

1. INTRODUCTION

Let $u_1, u_2 \dots$ be the enumeration in the breadth first search order of the vertices of a critical or subcritical branching forest with progeny distribution ν . Call $p(u_i)$, the size of the progeny of the i -th vertex, then the stochastic process $(X_n)_{n \geq 0}$ defined by,

$$X_0 = 0 \quad \text{and} \quad X_{n+1} - X_n = p(u_{n+1}) - 1, \quad n \geq 0$$

is a downward skip free random walk with step distribution $\mathbb{P}(X_1 = n) = \nu(n+1)$, from which the entire structure of the original branching forest can be recovered. We will refer to this random walk as the *Lukasiewicz-Harris coding path* of the branching forest, see Section 6 of [7], Section 1.1 of [3] or Section 6.2 of [13]. A nice example of application of this coding is that the total progeny of the k first trees $\mathbf{t}_1, \mathbf{t}_2, \dots, \mathbf{t}_k$ of the forest, see Figure 1, may be expressed as the first passage time of $(X_n)_{n \geq 0}$ at level $-k$, that is,

$$T_k = \inf\{i : X_i = -k\}.$$

This result combined with the following Kemperman's identity (also known as the Ballot Theorem, see Lemma 5 in [2] or Section 6.2 in [13]):

$$P(T_k = n) = \frac{k}{n} P(X_n = -k),$$

allows us to compute the law of the total progeny of $\mathbf{t}_1, \mathbf{t}_2, \dots, \mathbf{t}_k$ in terms of the progeny distribution ν . Note that the total progeny is actually a functional of the associated branching process, $(Z_n, n \geq 0)$, since the random variable Z_n represents the number of individuals at the n -th generation in the forest. The expression of this law was first obtained by Dwass [4].

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Theorem 1.1 (Dwass, 69). *Let $Z = (Z_n)$ be a critical or subcritical branching process. Let \mathbb{P}_k be its law when it starts from $Z_0 = k \geq 1$ and denote by ν its progeny law. Let O be the total progeny of Z , that is $O = \sum_{n \geq 0} Z_n$. Then for any $n \geq k$,*

$$(1.1) \quad \mathbb{P}_k(O = n) = \frac{k}{n} \nu^{*n}(n - k),$$

where ν^{*n} is the n -th iteration of the convolution product of the probability ν by itself.

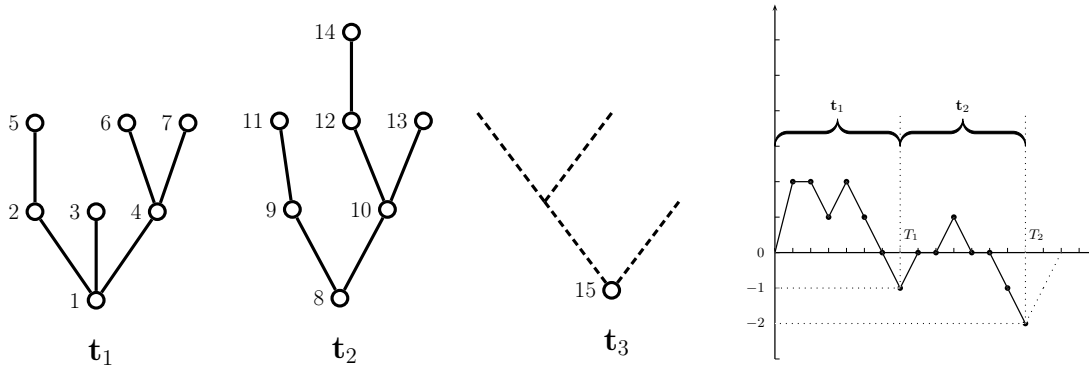


FIGURE 1. A forest labeled according to the breath first search order and the associated Lukasiewicz-Harris coding path.

More generally, whenever a functional of the branching forest admits a nice expression in terms of the Lukasiewicz-Harris coding path, we may expect to obtain an explicit form of its law. For instance, the law of the number of individuals with a given degree in the k first trees can be obtained in this way. We refer to Proposition 1.6 in [8] where the law of the number of leaves, first obtained in [10], is derived from the Lukasiewicz-Harris coding. The aim of this paper is to extend the coding that is recalled above to any irreducible, critical or subcritical multitype branching forest, and to derive the joint law of the total progeny of each type. Up to now, most of the results in this direction concern non irreducible, 2-type branching processes, as we will recall after the following definitions and notations.

We set $\mathbb{Z}_+ = \{0, 1, 2, \dots\}$ and $\mathbb{N} = \{1, 2, \dots\}$, and in all the sequel of this paper, d will be an integer such that $d \geq 2$. On a probability space (Ω, \mathcal{G}, P) , we define a d -type branching process $\mathbf{Z} := \{(Z_n^{(1)}, \dots, Z_n^{(d)}), n \geq 0\}$, as a \mathbb{Z}_+^d valued Markov chain with transition probabilities:

$$P(\mathbf{Z}_{n+1} = (k_1, \dots, k_d) \mid \mathbf{Z}_n = (r_1, \dots, r_d)) = \nu_1^{*r_1} * \dots * \nu_d^{*r_d}(k_1, \dots, k_d),$$

where ν_i are distributions on \mathbb{Z}_+^d and ν_i^{*r} is the r -th iteration of the convolution product of ν_i by itself, with $\nu_i^{*0} = \delta_0$. For $\mathbf{r} = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$, we will denote by $\mathbb{P}_{\mathbf{r}}$ the probability law $P(\cdot \mid \mathbf{Z}_0 = \mathbf{r})$. The vector $\nu = (\nu_1, \dots, \nu_d)$ will be called the progeny distribution of \mathbf{Z} . According to this process, each individual of type i gives birth to a random number of children with law ν_i , independently of the other individuals of its generation. The integer

valued random variable $Z_n^{(i)}$ is the total number of individuals of type i , at generation n . For $i, j \in \{1, \dots, d\}$, let us define the rate

$$m_{ij} = \sum_{\mathbf{z} \in \mathbb{Z}_+^d} z_j \nu_i(\mathbf{z}),$$

that corresponds to the mean number of children of type j , given by an individual of type i and let

$$\mathbf{M} := (m_{ij})_{1 \leq i, j \leq d}$$

be the mean matrix of \mathbf{Z} . Suppose that the extinction time T is a.s. finite, that is

$$(1.2) \quad T = \inf\{n : \mathbf{Z}_n = 0\} < \infty, \quad \text{a.s.}$$

Then let O_i , be the total number of offspring of type i , i.e. the total number of individuals of type i which are born up to time T :

$$O_i = \sum_{n=0}^T Z_n^{(i)} = \sum_{n \geq 0} Z_n^{(i)}.$$

The vector (O_1, \dots, O_d) will be called the *total progeny* of the multitype branching process. In the case where $d = 2$ and when $m_{12} > 0$ and $0 < m_{11} \leq 1$ but $m_{22} = m_{21} = 0$, it may be derived from Theorem 1 (ii) in [2], that the distribution of the total progeny of \mathbf{Z} is given by

$$(1.3) \quad \mathbb{P}_{(r_1, 0)}(O_1 = n_1, O_2 = n_2) = \frac{r_1}{n_1} \nu_1^{*n_1}(n_1 - r_1, n_2), \quad 1 \leq r_1 \leq n_1.$$

When $m_{12} > 0$ and $0 < m_{11}, m_{22} \leq 1$ but $m_{21} = 0$, after some elementary computation, combining the identities in (1.1) and (1.3), we obtain that for $n_2 \geq 1$,

$$(1.4) \quad \mathbb{P}_{(r_1, 0)}(O_1 = n_1, O_2 = n_2) = \frac{r_1}{n_1 n_2} \sum_{j=0}^{n_2} j \nu_1^{*n_1}(n_1 - r_1, j) \nu_2^{*n_2}(0, n_2 - j).$$

Note that (1.3) and (1.4) concern only the reducible case, when $d = 2$ and $T < \infty$, a.s. As far as we know, those are the only situations where the law of the total progeny of multitype branching processes is known explicitly.

Recall that if \mathbf{M} is irreducible, then according to Perron-Frobenius Theorem, it admits a unique eigenvalue ρ which is simple, positive and with maximal modulus. In this case, we will also say that \mathbf{Z} is irreducible. If moreover, \mathbf{Z} is non-degenerated, that is, if individuals have exactly one offspring with probability different from 1, then extinction, that is (1.2), holds if and only if $\rho \leq 1$, see [6], [11] and Chapter V of [1]. If $\rho = 1$, we say that \mathbf{Z} is critical and if $\rho < 1$, we say that \mathbf{Z} is subcritical. The results of this paper will be concerned by the case where \mathbf{Z} is *irreducible, non-degenerated, and critical or subcritical* so that (1.2) holds, that is the multitype branching process \mathbf{Z} becomes extinct with probability 1.

The next result gives the joint law of the total progeny together with the number of connected components of each type in the forest, which can also be defined as the total number of individuals of type j , whose parent is of type i , $i \neq j$, up to time T . Let us

denote by A_{ij} this random variable. Then Theorem 1.1 and identity (1.4) are extended as follows:

Theorem 1.2. *Assume that the d -type branching process \mathbf{Z} is irreducible, non-degenerated and critical or subcritical. For $i = 1, \dots, d$, let O_i be the total offspring of type i , up to the extinction time T and for $i \neq j$, let A_{ij} be the total number of individuals of type j , whose parent is of type i , up to time T .*

Then for all integers $r_i, n_i, k_{ij}, i, j = 1, \dots, d$, such that $r_i \geq 0, r_1 + \dots + r_d \geq 1, k_{ij} \geq 0$, for $i \neq j, -k_{jj} = r_j + \sum_{i \neq j} k_{ij}$, and $n_i \geq -k_{ii}$,

$$\begin{aligned} & \mathbb{P}_{\mathbf{r}}(O_1 = n_1, \dots, O_d = n_d, A_{ij} = k_{ij}, 1 \leq i, j \leq d, i \neq j) \\ &= \frac{\det(K)}{\bar{n}_1 \bar{n}_2 \dots \bar{n}_d} \prod_{i=1}^d \nu_i^{*n_i}(k_{i1}, \dots, k_{i,i-1}, n_i + k_{ii}, k_{i,i+1}, \dots, k_{id}), \end{aligned}$$

where $\mathbf{r} = (r_1, \dots, r_d)$, $\nu_i^{*0} = \delta_0$, $\bar{n}_i = n_i \vee 1$ and K is the matrix $(-k_{ij})_{1 \leq i, j \leq d}$ to which we removed the line i and the column i , for all i such that $n_i = 0$.

As in the single type case, it is theoretically possible to obtain the law of the total progeny (O_1, \dots, O_d) of a critical or subcritical multitype branching process by applying the multivariate Lagrange inversion formula, to its generating function. An attempt in this direction has been done in [5], but the results are not directly explicit in terms of the progeny distribution. Moreover, we emphasize that the random variables A_{ij} are not functionals of the multitype branching process \mathbf{Z} . So, their formal definition and the computation of their law require a more complete information provided by the forest.

Our proof of Theorem 1.2 uses a bijection, displayed in Theorem 2.7, between multitype forests and a particular set of multidimensional, integer valued sequences. A consequence of this result is that any critical or subcritical irreducible multitype branching forest is encoded by d independent, d -dimensional random walks, see Theorem 3.1. Then, in a similar way to the single type case, the total progeny, jointly with the number of subtrees of each type in the forest, is expressed as the first passage time of this multivariate process in some domain. An extension of the Ballot Theorem that can be considered as one of the main result of this paper, see Theorem 3.4, allows us to conclude, as in the single type case. We emphasize that this coding should allow us to compute the law of other functionals of branching forests such as the number of vertices of a given degree.

This paper is organized as follows. Section 2 is devoted to deterministic multitype forests. In Subsection 2.1, we define the space of these forests. Then Subsection 2.2 is devoted to the definition of the space of the coding sequences and to obtain the bijection between this space and the space of multitype forests. This result is stated in Theorem 2.7. Then in Section 3, we define the probability space of multitype branching forests, we display their multitype Lukasiewicz-Harris coding in Theorem 3.1 and we prove its application to the total progeny that is stated in Theorem 1.2. This result requires a multivariate extension of the Ballot Theorem, see Theorem 3.4, whose proof bears on the crucial combinatorial Lemma 3.3. The latter is proved in Section 4.

2. MULTITYPE FORESTS

2.1. **The space of multitype forests.** A *forest*, is a directed planar graph $\mathbf{f} \subset \mathbf{v} \times \mathbf{v}$, with a finite or infinite set of vertices $\mathbf{v} = \mathbf{v}(\mathbf{f})$, such that the outer degree of each vertex is equal to 0 or 1 and whose connected components, which are called the *trees*, are finite. A forest consisting of a single connected component is also called a tree. In a tree \mathbf{t} , the only vertex with outer degree equal to 0 is called the *root* of \mathbf{t} . It will be denoted by $r(\mathbf{t})$. The roots of the connected components of a forest \mathbf{f} are called the roots of \mathbf{f} . For two vertices u and v of a forest \mathbf{f} , if (u, v) is a directed edge of \mathbf{f} , then we say that u is a *child* of v , or that v is the *parent* of u .

A label will be given to each vertex of a forest \mathbf{f} in the following way. We first give an order to the trees of \mathbf{f} and denote them by $\mathbf{t}_1(\mathbf{f}), \mathbf{t}_2(\mathbf{f}), \dots, \mathbf{t}_k(\mathbf{f}), \dots$ (we will usually write $\mathbf{t}_1, \mathbf{t}_2, \dots, \mathbf{t}_k, \dots$ if no confusion is possible). Then each tree is labelled according to the *breadth first search algorithm*: we read the tree from its root to its last generation by running along each generation from the left to the right. This definition should be obvious from the example of Figure 1. If a forest \mathbf{f} contains at least i vertices, then the i -th vertex of \mathbf{f} is denoted by $u_i(\mathbf{f})$. When no confusion is possible, we will simply denote the i -th vertex by u_i . The set of plane forests will be denoted by \mathcal{F} . The elements of \mathcal{F} will simply be called forests.

Recall that d is an integer such that $d \geq 2$. To each forest $\mathbf{f} \in \mathcal{F}$, we associate an application $c_{\mathbf{f}} : \mathbf{v}(\mathbf{f}) \rightarrow \{1, 2, \dots, d\}$ such that if $u_i, u_{i+1}, \dots, u_{i+j} \in \mathbf{v}(\mathbf{f})$ have the same parent, then $c_{\mathbf{f}}(u_i) \leq c_{\mathbf{f}}(u_{i+1}) \leq \dots \leq c_{\mathbf{f}}(u_{i+j})$. For $v \in \mathbf{v}(\mathbf{f})$, the integer $c_{\mathbf{f}}(v)$ is called the *type* (or the *color*) of v . The couple $(\mathbf{f}, c_{\mathbf{f}})$ is called a d -*type forest*. When no confusion is possible, we will simple write \mathbf{f} . The set of d -type forests will be denoted by \mathcal{F}_d . A 2-type forest is represented on Figure 2 below.

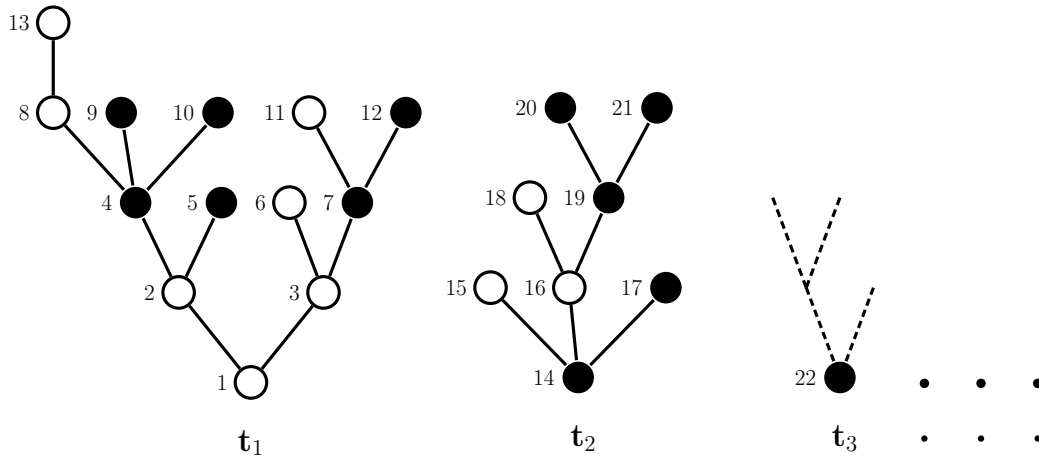


FIGURE 2. A two type forest labeled according to the breath first search order. Vertices of type 1 (resp. 2) are represented in white (resp. black).

A *subtree of type* $i \in \{1, \dots, d\}$ of a d -type forest $(\mathbf{f}, c_{\mathbf{f}}) \in \mathcal{F}_d$ is a maximal connected subgraph of $(\mathbf{f}, c_{\mathbf{f}})$ whose all vertices are of type i . Formally, \mathbf{t} is a subtree of type i of $(\mathbf{f}, c_{\mathbf{f}})$, if it is a connected subgraph whose all vertices are of type i and such that either $r(\mathbf{t})$ has no parent or the type of its parent is different from i . Moreover, if the parent of a vertex $v \in \mathbf{v}(\mathbf{t})^c$ belongs to $\mathbf{v}(\mathbf{t})$, then $c_{\mathbf{f}}(v) \neq i$. Subtrees of type i of $(\mathbf{f}, c_{\mathbf{f}})$ are ranked according to the order of their roots in \mathbf{f} and are denoted by $\mathbf{t}_1^{(i)}, \mathbf{t}_2^{(i)}, \dots, \mathbf{t}_k^{(i)}, \dots$. The forest $\mathbf{f}^{(i)} := \{\mathbf{t}_1^{(i)}, \mathbf{t}_2^{(i)}, \dots, \mathbf{t}_k^{(i)}, \dots\}$ is called *the subforest of type* i of $(\mathbf{f}, c_{\mathbf{f}})$. It may be considered as an element of \mathcal{F} . We denote by $u_1^{(i)}, u_2^{(i)}, \dots$ the elements of $\mathbf{v}(\mathbf{f}^{(i)})$, ranked in the breadth first search order of $\mathbf{f}^{(i)}$. The subforests of type 1 and 2 of a 2-type forest are represented in Figure 3.

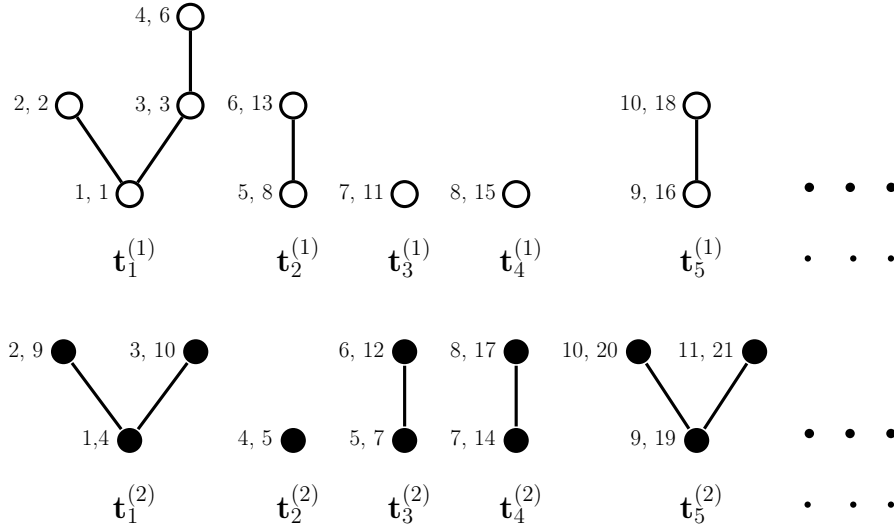


FIGURE 3. Labeled subforests $\mathbf{f}^{(1)}$ and $\mathbf{f}^{(2)}$ associated with Figure 2. Beside each vertex, the first number corresponds to its rank in $\mathbf{f}^{(i)}$, $i = 1, 2$ and the second one is its rank in the original forest \mathbf{f} .

To any forest $(\mathbf{f}, c_{\mathbf{f}}) \in \mathcal{F}_d$, we associate the *reduced forest*, denoted by $(\mathbf{f}_r, c_{\mathbf{f}_r}) \in \mathcal{F}_d$, which is the forest of \mathcal{F}_d obtained by aggregating all the vertices of each subtree of $(\mathbf{f}, c_{\mathbf{f}})$ with a given type, in a single vertex with the same type, and preserving an edge between each pair of connected subtrees. An example is given in Figure 4.

2.2. Coding multitype forests. For a forest $(\mathbf{f}, c_{\mathbf{f}}) \in \mathcal{F}_d$ and $u \in \mathbf{v}(\mathbf{f})$, when no confusion is possible, we denote by $p_i(u)$ the number of children of type i of u . For each $i \in \{1, \dots, d\}$, let $n_i \geq 0$ be the number of vertices in the subforest $\mathbf{f}^{(i)}$ of $(\mathbf{f}, c_{\mathbf{f}})$. Then let us define the d -dimensional chain $x^{(i)} = (x^{i,1}, \dots, x^{i,d})$, with length n_i and whose values belong to the set \mathbb{Z}^d , by $x_0^{(i)} = 0$ and if $n_i \geq 1$,

$$(2.5) \quad x_{n+1}^{i,j} - x_n^{i,j} = p_j(u_{n+1}^{(i)}), \quad \text{if } i \neq j \quad \text{and} \quad x_{n+1}^{i,i} - x_n^{i,i} = p_i(u_{n+1}^{(i)}) - 1, \quad 0 \leq n \leq n_i - 1,$$

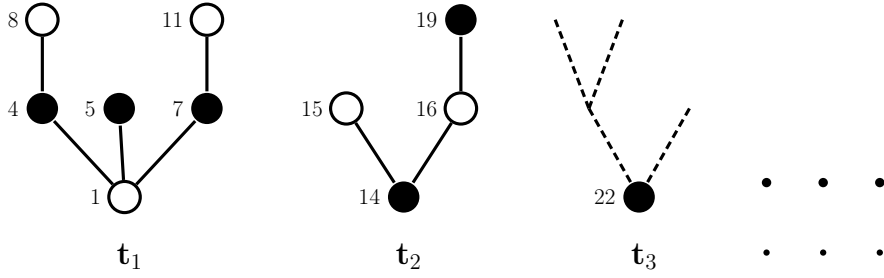


FIGURE 4. Reduced forest associated with the example of Figure 2. Beside each vertex is the rank of the root of the corresponding subtree in the original forest \mathbf{f} .

where we recall that $(u_n^{(i)})_{n \geq 1}$ is the enumeration of the subforest $\mathbf{f}^{(i)}$ in its own breadth first search order. Note that the chains $(x_n^{i,j})$, for $i \neq j$ are nondecreasing whereas $(x_n^{i,i})$ is a downward skip free chain, i.e. $x_{n+1}^{i,i} - x_n^{i,i} \geq -1$, for $0 \leq n \leq n_i - 1$. The chain $(x_n^{i,i})$ corresponds to the Lukasiewicz-Harris coding walk of the subforest $\mathbf{f}^{(i)}$, as defined in the introduction, see also Section 6.2 in [13] for a proper definition. In particular, if n_i is finite, then $n_i = \min\{n : x_n^{i,i} = \min_{0 \leq k \leq n_i} x_k^{i,i}\}$. These properties of the chains $x^{(i)}$ lead us to the following definition.

Definition 2.1. Let S_d be the set of $[\mathbb{Z}^d]^d$ -valued sequences, $x = (x^{(1)}, x^{(2)}, \dots, x^{(d)})$, such that for all $i = 1, \dots, d$, $x^{(i)} = (x^{i,1}, \dots, x^{i,d})$, is a \mathbb{Z}^d -valued sequence defined on some interval of integers, $\{0, \dots, n_i\}$, $0 \leq n_i \leq \infty$, which satisfies $x_0^{(i)} = 0$ and if $n_i \geq 1$ then

- (i) for $i \neq j$, the sequence $(x_n^{i,j})_{0 \leq n \leq n_i}$ is nondecreasing,
- (ii) for all i , $x_{n+1}^{i,i} - x_n^{i,i} \geq -1$, $0 \leq n \leq n_i - 1$.

We will also denote sequences in S_d by $(x_k^{i,j}, 0 \leq k \leq n_i, i, j = 1, \dots, d)$. The vector $\mathbf{n} = (n_1, \dots, n_d) \in \overline{\mathbb{Z}}_+^d$, where $\overline{\mathbb{Z}}_+ = \mathbb{Z}_+ \cup \{+\infty\}$ will be called the length of x .

Relation (2.5) defines an application from the set \mathcal{F}_d to the set S_d . Let us denote by Ψ this application, that is

$$(2.6) \quad \begin{aligned} \Psi : \mathcal{F}_d &\rightarrow S_d \\ (\mathbf{f}, c_{\mathbf{f}}) &\mapsto \Psi((\mathbf{f}, c_{\mathbf{f}})) = x. \end{aligned}$$

For $x \in S_d$, set $k_i = -\inf_{0 \leq n \leq n_i} x_n^{i,i}$ and define the *first passage time process* of the chain $(x_n^{i,i})$ as follows:

$$(2.7) \quad \tau_k^{(i)} = \min\{n \geq 0 : x_n^{i,i} = -k\}, \quad 0 \leq k \leq k_i,$$

where $\tau_{k_i}^{(i)} = \infty$, if $k_i = \infty$. If x is the image by Ψ of a forest $(\mathbf{f}, c_{\mathbf{f}}) \in \mathcal{F}_d$, i.e. $x = \Psi((\mathbf{f}, c_{\mathbf{f}}))$, then k_i is the (finite or infinite) number of trees in the subforest $\mathbf{f}^{(i)}$ and for $k < \infty$, the time $\tau_k^{(i)}$ is the total number of vertices which are contained in the k first trees of $\mathbf{f}^{(i)}$, i.e. $\mathbf{t}_1^{(i)}, \mathbf{t}_2^{(i)}, \dots, \mathbf{t}_k^{(i)}$. This fact is well known and easily follows from the Lukasiewicz-Harris

coding of the single type forest $\mathbf{f}^{(i)}$, see the introduction and Lemma 6.3 in [13]. Then for $i, j = 1, \dots, d$, define the integer valued sequence

$$(2.8) \quad \bar{x}_k^{i,j} = x_k^{i,j}(\tau_k^{(i)}), \quad 0 \leq k \leq k_i.$$

If $x = \Psi((\mathbf{f}, c_{\mathbf{f}}))$, then we may check that when $i \neq j$, $\bar{x}_k^{i,j}$ is the number of subtrees of type j whose root is the child of a vertex in $\mathbf{t}_1^{(i)}, \mathbf{t}_2^{(i)}, \dots, \mathbf{t}_{k_i}^{(i)}$. Or equivalently, it is the number of vertices of type j whose parent is a vertex of $\mathbf{t}_1^{(i)}, \mathbf{t}_2^{(i)}, \dots, \mathbf{t}_{k_i}^{(i)}$. For each $i = 1, \dots, d$, we set

$$\bar{x}^{(i)} = (\bar{x}^{i,1}, \dots, \bar{x}^{i,d}) \quad \text{and} \quad \bar{x} = (\bar{x}^{(1)}, \bar{x}^{(2)}, \dots, \bar{x}^{(d)}).$$

Clearly for $i \neq j$, the sequence $(\bar{x}_k^{i,j})_{0 \leq k \leq k_i}$ is increasing and $\bar{x}_k^{i,i} = -k$, for all $i = 1, \dots, d$ and $0 \leq k \leq k_i$, so that $\bar{x} \in S_d$ and recalling the definition of the reduced forest, $(\mathbf{f}_{\mathbf{r}}, c_{\mathbf{f}_{\mathbf{r}}})$, see the end of Section 2.1, we may check that:

$$(2.9) \quad \Psi((\mathbf{f}_{\mathbf{r}}, c_{\mathbf{f}_{\mathbf{r}}})) = \bar{x}.$$

For a forest $(\mathbf{f}, c_{\mathbf{f}}) \in \mathcal{F}_d$ with trees $\mathbf{t}_1, \mathbf{t}_2, \dots$, we will denote by $c_{(\mathbf{f}, c_{\mathbf{f}})}$ the sequence of types of the roots of $\mathbf{t}_1, \mathbf{t}_2, \dots$, i.e.

$$c_{(\mathbf{f}, c_{\mathbf{f}})} := (c_{\mathbf{f}}(r(\mathbf{t}_1)), c_{\mathbf{f}}(r(\mathbf{t}_2)), \dots).$$

Note that $c_{(\mathbf{f}, c_{\mathbf{f}})} \in \cup_{1 \leq r \leq \infty} \{1, \dots, d\}^r$ and that $c_{(\mathbf{f}, c_{\mathbf{f}})} = c_{(\mathbf{f}_{\mathbf{r}}, c_{\mathbf{f}_{\mathbf{r}}})}$. When no confusion is possible, $c_{(\mathbf{f}, c_{\mathbf{f}})}$ will simply be denoted by $c = (c_1, c_2, \dots)$ and we will call it the *root type sequence* of the forest.

Then before we state the general result on the coding of multitype forests in Theorem 2.7, we first need to show that the sequences $(\bar{x}^{i,j})_{i \neq j}$ together with $c = (c_1, c_2, \dots)$ allow us to encode the reduced forest $(\mathbf{f}_{\mathbf{r}}, c_{\mathbf{f}_{\mathbf{r}}})$, i.e. this forest can be reconstructed from (\bar{x}, c) . This claim is stated in Lemma 2.5 below. In order to prove it, we first need to describe the set of sequences which encode reduced forests and to state the preliminary Lemma 2.2 regarding these sequences.

Recall that $\bar{\mathbb{Z}}_+ = \mathbb{Z}_+ \cup \{+\infty\}$ and let us define the following (non total) order in $\bar{\mathbb{Z}}_+^d$: for two elements $\mathbf{q} = (q_1, \dots, q_d)$ and $\mathbf{q}' = (q'_1, \dots, q'_d)$ of $\bar{\mathbb{Z}}_+^d$ we write $\mathbf{q} \leq \mathbf{q}'$ if $q_i \leq q'_i$ for all $i = 1, \dots, d$. Moreover we write $\mathbf{q} < \mathbf{q}'$ if $\mathbf{q} \leq \mathbf{q}'$ and if there is $i = 1, \dots, d$ such that $q_i < q'_i$. For an element $x = (x_k^{i,j}, 0 \leq k \leq q_i, i, j = 1, \dots, d)$ of S_d with length $\mathbf{q} = (q_1, \dots, q_d)$, and $\mathbf{r} = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$, we say that the system of equations (\mathbf{r}, x) admits a solution if there exists $\mathbf{s} = (s_1, \dots, s_d) \in \mathbb{Z}_+^d$, such that $\mathbf{s} \leq \mathbf{q}$ and

$$(2.10) \quad r_j + \sum_{i=1}^d x_k^{i,j}(s_i) = 0, \quad j = 1, \dots, d.$$

Lemma 2.2. *Let $x \in S_d$ and $\mathbf{r} = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$. Assume that the system (\mathbf{r}, x) admits a solution, then*

- (i) *there exists a unique solution $\mathbf{n} = (n_1, \dots, n_d)$ of the system (\mathbf{r}, x) such that if \mathbf{n}' is any solution of (\mathbf{r}, x) , then $\mathbf{n} \leq \mathbf{n}'$. Moreover we have $n_i = \min\{n : x_n^{i,i} = \min_{0 \leq k \leq n_i} x_k^{i,i}\}$, for all $i = 1, \dots, d$. A solution such as \mathbf{n} will be called the smallest solution of the system (\mathbf{r}, x) .*

- (ii) Let $\mathbf{r}' \in \mathbb{Z}_+^d$ be such that $\mathbf{r}' \leq \mathbf{r}$. Then the system (\mathbf{r}', x) admits a solution. Let us denote its smallest solution by \mathbf{n}' . Then the system $(\mathbf{r} - \mathbf{r}', \tilde{x})$, where $\tilde{x}^{i,j}(k) = x^{i,j}(n'_i + k) - x^{i,j}(n'_i)$, $0 \leq k \leq n_i - n'_i$, admits a solution, and its smallest solution is $\mathbf{n} - \mathbf{n}'$.

A proof of this lemma is given in Section 4. For $\mathbf{r} = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$, with $r = r_1 + \dots + r_d \geq 1$, that is $\mathbf{r} > 0$, we define,

$$C_d^{\mathbf{r}} = \{c \in \{1, \dots, d\}^r : \text{Card}\{1 \leq j \leq r : c_j = i\} = r_i, i = 1, \dots, d\}.$$

We emphasize that the root type sequence of a forest $(\mathbf{f}, c_{\mathbf{f}}) \in \mathcal{F}_d$ with $r = r_1 + \dots + r_d$ trees amongst which exactly r_i trees have a root of type i is an element $c \in C_d^{\mathbf{r}}$. Now we define the subsets of forests and reduced forests whose root type sequence is in $C_d^{\mathbf{r}}$ and that contain at least one vertex of each type.

Definition 2.3. Let $\mathbf{r} = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$, such that $\mathbf{r} > 0$.

- (i) We denote by $\mathcal{F}_d^{\mathbf{r}}$, the subset of \mathcal{F}_d of forests $(\mathbf{f}, c_{\mathbf{f}})$ with $r_1 + \dots + r_d$ trees, which contain at least one vertex of each type, and such that $c_{(\mathbf{f}, c_{\mathbf{f}})} \in C_d^{\mathbf{r}}$.
- (ii) We denote by $\bar{\mathcal{F}}_d^{\mathbf{r}}$ the subset of $\mathcal{F}_d^{\mathbf{r}}$, of reduced forests, more specifically, $(\mathbf{f}, c_{\mathbf{f}}) \in \bar{\mathcal{F}}_d^{\mathbf{r}}$ if $(\mathbf{f}, c_{\mathbf{f}}) \in \mathcal{F}_d^{\mathbf{r}}$ and if for each i , vertices of type $i = 1, \dots, d$ in $\mathbf{v}(\mathbf{f})$ have no child of type i .

Then we define the sets of coding sequences related to $\mathcal{F}_d^{\mathbf{r}}$ and $\bar{\mathcal{F}}_d^{\mathbf{r}}$.

Definition 2.4. Let $\mathbf{r} = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$, such that $\mathbf{r} > 0$.

- (i) We denote by $S_d^{\mathbf{r}}$ the subset of S_d of sequences x whose length belongs to \mathbb{N}^d and corresponds to the smallest solution of the system (\mathbf{r}, x) defined in (2.10).
- (i) We denote by $\bar{S}_d^{\mathbf{r}}$ the subset of $S_d^{\mathbf{r}}$ consisting in sequences, such that $x_k^{i,i} = -k$, for all k and i .

Then we first establish a bijection between the sets $\bar{\mathcal{F}}_d^{\mathbf{r}}$ and $\bar{S}_d^{\mathbf{r}} \times C_d^{\mathbf{r}}$. Recall the definition of Ψ in (2.6).

Lemma 2.5. Let $\mathbf{r} = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$, be such that $\mathbf{r} > 0$, then the mapping

$$\begin{aligned} \Phi : \bar{\mathcal{F}}_d^{\mathbf{r}} &\rightarrow \bar{S}_d^{\mathbf{r}} \times C_d^{\mathbf{r}} \\ (\mathbf{f}, c_{\mathbf{f}}) &\mapsto (\Psi((\mathbf{f}, c_{\mathbf{f}})), c_{(\mathbf{f}, c_{\mathbf{f}})}) \end{aligned}$$

is a bijection.

Proof. Let $(\mathbf{f}, c_{\mathbf{f}}) \in \bar{\mathcal{F}}_d^{\mathbf{r}}$ and let k_i be the total number of subtrees of type i which are contained in $(\mathbf{f}, c_{\mathbf{f}})$. (Note that since $(\mathbf{f}, c_{\mathbf{f}})$ is a reduced forest, its subtrees are actually single vertices.) By definition, $k_i \geq 1$, for each i . The fact that, $c_{(\mathbf{f}, c_{\mathbf{f}})} \in C_d^{\mathbf{r}}$ follows from Definition 2.3 (ii). Then let us show that $x = \Psi((\mathbf{f}, c_{\mathbf{f}})) \in \bar{S}_d^{\mathbf{r}}$. Since $(\mathbf{f}, c_{\mathbf{f}})$ is a reduced forest, then $x = \bar{x}$. Besides, from (2.8), x has length $\mathbf{k} = (k_1, \dots, k_d)$ and for $j \neq i$, $x^{i,j}(k_i)$ is the number of subtrees of type j whose root is the child of a vertex of $\{\mathbf{t}_1^{(i)}, \mathbf{t}_2^{(i)}, \dots, \mathbf{t}_{k_i}^{(i)}\}$, i.e. of any subtree of type i in $(\mathbf{f}, c_{\mathbf{f}})$. Hence for $j = 1, \dots, d$, $\sum_{i \neq j} x^{i,j}(k_i)$ is the total number of subtrees of type j in $(\mathbf{f}, c_{\mathbf{f}})$, whose root is the child of a vertex of type $i = 1, \dots, d, i \neq j$. Then we obtain k_j by adding to $\sum_{i \neq j} x^{i,j}(k_i)$, the

r_j subtrees of type j whose root is one of the roots of $\mathbf{t}_1, \dots, \mathbf{t}_r$, where $r = r_1 + \dots + r_d$. That is $k_j = r_j + \sum_{i \neq j} x_k^{i,j}(k_i)$. Since moreover from (2.8), $x_k^{i,i} = -k$, for all $0 \leq k \leq k_i$, we have proved that \mathbf{k} is a solution of the system (\mathbf{r}, x) . It remains to prove that it is the smallest solution.

Let us first assume that $\mathbf{r} = (1, 0, \dots, 0)$, so that $(\mathbf{f}, c_{\mathbf{f}})$ consists in a single tree \mathbf{t}_1 whose root has color $c_{\mathbf{f}}(r(\mathbf{t}_1)) = 1$. Then we can reconstruct, this tree from the d sequences $(x_k^{(i)}, 0 \leq k \leq k_i)$, $i = 1, \dots, d$ by inverting the procedure defined in (2.5) and this *reconstruction procedure* gives a unique tree. Indeed, by definition of the application Ψ , each sequence $(x_k^{(i)}, 0 \leq k \leq k_i)$, is associated to a unique 'marked subforest', say $\tilde{\mathbf{f}}^{(i)}$, of type i whose vertices kept the memory of their progeny. More specifically, for $k = 1, \dots, k_i$, the increment $x_k^{(i)} - x_{k-1}^{(i)}$ gives the progeny of the k -th vertex of the subforest $\tilde{\mathbf{f}}^{(i)}$. This connection between marked subforests $\tilde{\mathbf{f}}^{(i)}$ and sequences $(x_k^{(i)}, 0 \leq k \leq k_i)$ is illustrated on Figure 5. Now let $\mathbf{k}' \leq \mathbf{k}$ be the smallest solution of the system (\mathbf{r}, x) . Let

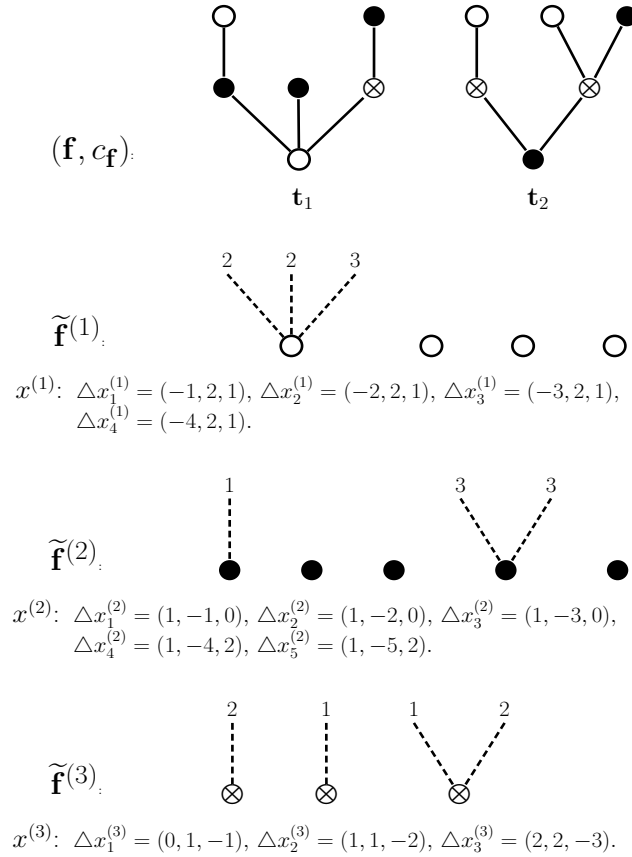


FIGURE 5. A three types (reduced) forest $(\mathbf{f}, c_{\mathbf{f}})$, its three marked subforests $\tilde{\mathbf{f}}^{(1)}$, $\tilde{\mathbf{f}}^{(2)}$, $\tilde{\mathbf{f}}^{(3)}$, and its coding sequences $x^{(1)}$, $x^{(2)}$, $x^{(3)}$. Here we have set $\Delta x_k^{(i)} = x_k^{(i)} - x_{k-1}^{(i)}$.

$\mathbf{q} = (q_1, \dots, q_d) < \mathbf{k}'$ and suppose that we have been able to perform the reconstruction

procedure until q , that is from the sequences $(x_k^{(i)}, 0 \leq k \leq q_i)$, $i = 1, \dots, d$. Then since q is not a solution of (r, x) , we see from what has been proved just above that the tree that is obtained is 'not complete'. That is, at least one of its leaves (say of type j) is marked, so that this leaf should still get children whose types and numbers are given by the next jump $x_{q_j+1}^{(j)} - x_{q_j}^{(j)}$, for $q_j < k'_j$, according to the reconstruction procedure. Thus, doing so, we necessarily end up with a tree from the sequences $(x_k^{(i)}, 0 \leq k \leq k'_i)$, $i = 1, \dots, d$, and this tree is complete, that is none of its leaves is marked. Then since the reconstruction procedure obtained by inverting (2.5), gives a unique tree, we necessarily have $k' = k$.

Then let $r = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$. Assume with no loss of generality that the root of the first tree \mathbf{t}_1 of $(\mathbf{f}, c_{\mathbf{f}})$ has color 1. Let k_i^1 be the number of subtrees of type i in \mathbf{t}_1 . From Lemma 2.2, the system (r^1, x) , where $r^1 := (1, 0, \dots, 0)$, admits a smallest solution. Moreover from the reconstruction procedure which is described above, this solution is $k^1 = (k_1^1, \dots, k_d^1)$. Suppose now with no loss of generality that the second tree, \mathbf{t}_2 in $(\mathbf{f}, c_{\mathbf{f}})$ has color 2. Let k_i^2 be the number of subtrees of type i in \mathbf{t}_2 . Then from the same arguments as for the reconstruction of the first tree, \mathbf{t}_2 may be reconstructed from the system (r^2, y) , where $r^2 := (0, 1, 0, \dots, 0)$ and $y^{i,j}(k) = x^{i,j}(k_i^1 + k) - x^{i,j}(k_i^1)$, $k \geq 0$. Moreover (r^2, y) admits $k^2 = (k_1^2, \dots, k_d^2)$ as a smallest solution. Then from part (ii) of Lemma 2.2, $k^1 + k^2$ is the smallest solution of the system $(r^1 + r^2, x)$. So we have proved the result for the forest consisting in the trees \mathbf{t}_1 and \mathbf{t}_2 . Then by iterating these arguments for each tree of $(\mathbf{f}, c_{\mathbf{f}})$, we obtain that $x \in \bar{S}_d^r$.

Conversely, let $c = (c_1, \dots, c_r) \in C_d^r$, $x \in \bar{S}_d^r$ and let $k = (k_1, \dots, k_d)$ be the smallest solution of the system (r, x) . Then let us show that there is a forest $(\mathbf{f}, c_{\mathbf{f}}) \in \bar{\mathcal{F}}_d^r$ such that $\Psi((\mathbf{f}, c_{\mathbf{f}})) = x$ and $c_{(\mathbf{f}, c_{\mathbf{f}})} = c$. Assume, without loss of generality that $c_1 = 1$. From Lemma 2.2 (ii), there is a smallest solution, say $k^1 = (k_1^1, \dots, k_d^1)$, to the system (r^1, x) , where $r^1 := (1, 0, \dots, 0)$. Then we may reconstruct a unique forest $(\mathbf{t}_1, c_{\mathbf{t}_1}) \in \bar{\mathcal{F}}_d^1$ (consisting in a single tree) such that $\Psi((\mathbf{t}_1, c_{\mathbf{t}_1})) = (x_k^{i,j}, 0 \leq k \leq k_i^1, i, j = 1, \dots, d)$ and $c_{(\mathbf{t}_1, c_{\mathbf{t}_1})} = 1$ by inverting the procedure that is described in (2.5). Assume for instance that $c_2 = 2$ and set $r^2 := (1, 1, 0, \dots, 0) \leq r$, then from Lemma 2.2 (ii), there is a smallest solution, say k^2 , to the system (r^2, x) . Moreover, $k^3 = k^2 - k^1$ is the smallest solution of the system $(r^2 - r^1, y)$, where $y^{i,j}(k) = x^{i,j}(k_i^1 + k) - x^{i,j}(k_i^1)$, $k \geq 0$. Then as before, we can reconstruct a unique tree $(\mathbf{t}_2, c_{\mathbf{t}_2})$ such that $\Psi((\mathbf{t}_2, c_{\mathbf{t}_2})) = y$ and such that the forest $\hat{\mathbf{f}} = \{\mathbf{t}_1, \mathbf{t}_2\}$ satisfies $\Psi((\hat{\mathbf{f}}, c_{\hat{\mathbf{f}}})) = (x_k^{i,j}, 0 \leq k \leq k_i^2, i, j = 1, \dots, d)$ and $c_{(\hat{\mathbf{f}}, c_{\hat{\mathbf{f}}})} = (1, 1, 0, \dots, 0)$. Then iterating these arguments, we may reconstruct a unique forest $(\mathbf{f}, c_{\mathbf{f}}) \in \bar{\mathcal{F}}_d^r$ such that $\Psi((\mathbf{f}, c_{\mathbf{f}})) = x$ and $c_{(\mathbf{f}, c_{\mathbf{f}})} = c$. \square

Let $x \in S_d$ with length $\mathbf{n} = (n_1, \dots, n_d)$ and recall from (2.8), the definition of the associated sequence \bar{x} , with length $\mathbf{k} = (k_1, \dots, k_d)$, such that $k_i = -\min_{0 \leq n \leq n_i} x_n^{i,i}$.

Lemma 2.6. *Let $r \in \mathbb{Z}_+^d$, such that $r > 0$ and $x \in S_d$, with length $\mathbf{n} \in \mathbb{N}^d$ and set $k_i = -\min_{0 \leq n \leq n_i} x_n^{i,i}$, $i = 1, \dots, d$. If \mathbf{n} is the smallest solution of the system (r, x) (i.e. $x \in S_d^r$), then $\mathbf{k} = (k_1, \dots, k_d)$ is the smallest solution of the system (r, \bar{x}) . Conversely, if $n_i = \tau_{k_i}^{(i)}$, for all $i = 1, \dots, d$ and if \mathbf{k} is the smallest solution of (r, \bar{x}) , (i.e. $\bar{x} \in \bar{S}_d^r$), then \mathbf{n} is the smallest solution of (r, x) .*

Proof. Assume that n is the smallest solution of the system (r, x) . Then from part (i) of Lemma 2.2, $n_i = \tau_{k_i}^{(i)}$, hence k is a solution of (r, \bar{x}) . Let $k' \leq k$ be such that

$$r_j + \sum_{i=1}^d \bar{x}^{i,j}(k'_i) = 0, \quad j = 1, \dots, d.$$

Then by definition of \bar{x} there is $n' \leq n$ such that $n'_i = \tau_{k'_i}^{(i)}$ and

$$r_j + \sum_{i=1}^d x^{i,j}(n'_i) = 0, \quad j = 1, \dots, d.$$

So $n' = n$ and hence $k' = k$.

The converse is proved in the same way. Suppose that $n_i = \tau_{k_i}^{(i)}$, $i = 1, \dots, d$, and that k is the smallest solution of (r, \bar{x}) . Then clearly, n is a solution of (r, x) . Let n' be the smallest solution of (r, x) . Then from Lemma 2.2, there is k' such that $n'_i = \tau_{k'_i}^{(i)}$, hence k' is a solution of (r, \bar{x}) . This implies that $k \leq k'$, so that $n \leq n'$, hence $n = n'$. \square

Now we extend the application Φ defined in Lemma 2.5 to the set \mathcal{F}_d^r . Here is the main result of this section, that can be considered as an extension of Proposition 1.1 in [3].

Theorem 2.7. *Let $r = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$, be such that $r > 0$, then the mapping*

$$\begin{aligned} \Phi : \mathcal{F}_d^r &\rightarrow S_d^r \times C_d^r \\ (\mathbf{f}, c_{\mathbf{f}}) &\mapsto (\Psi((\mathbf{f}, c_{\mathbf{f}})), c_{(\mathbf{f}, c_{\mathbf{f}})}) \end{aligned}$$

is a bijection.

Proof. Let us first check that for any $(\mathbf{f}, c_{\mathbf{f}}) \in \mathcal{F}_d^r$, we have $\Phi((\mathbf{f}, c_{\mathbf{f}})) \in S_d^r \times C_d^r$. By definition 2.3, (i), $c_{(\mathbf{f}, c_{\mathbf{f}})} \in C_d^r$. Now set $x = \Psi((\mathbf{f}, c_{\mathbf{f}}))$ and let $(\mathbf{f}_r, c_{\mathbf{f}_r}) \in \bar{\mathcal{F}}_d^r$ be the forest, $(\mathbf{f}, c_{\mathbf{f}})$ once reduced. Then from (2.9) and Lemma 2.5, this reduced forest is encoded by $(\bar{x}, c_{\mathbf{f}})$. Let $k = (k_1, \dots, k_d)$ be the number of subtrees of type i in this forest (this is actually the number of vertices of type i), then k is the length of \bar{x} and it is the smallest solution of (r, \bar{x}) , i.e. $\bar{x} \in \bar{S}_d^r$. Moreover $n = (n_1, \dots, n_d)$, where $n_i = \tau_{k_i}^{(i)}$ is the length of x , and from Lemma 2.6, it is the smallest solution of (r, x) . So, we have proved that $\Psi((\mathbf{f}, c_{\mathbf{f}})) \in S_d^r$.

Conversely let $(x, c) \in S_d^r \times C_d^r$. From Lemma 2.5, to (\bar{x}, c) , we may associate a unique forest $(\mathbf{f}_r, c_{\mathbf{f}_r}) \in \bar{\mathcal{F}}_d^r$. Then let k_i be the number of vertices of type i in this forest. For $1 \leq k \leq k_i$, let $u_k^{(i)}$ be the k -th vertex of type i in the breadth first search order of $(\mathbf{f}_r, c_{\mathbf{f}_r})$. Then in $(\mathbf{f}_r, c_{\mathbf{f}_r})$, we replace the vertex $u_k^{(i)}$ by the subtree of type i which is encoded by the Lukasiewicz-Harris path $(x^{i,i}(\tau_{k-1}^{(i)} + l) + k - 1, 0 \leq l \leq \tau_k^{(i)} - \tau_{k-1}^{(i)})$. We know about the progeny of each vertex of this subtree, thanks to the chains $(x^{i,j}(\tau_{k-1}^{(i)} + l) - x^{i,j}(\tau_{k-1}^{(i)}), 0 \leq l \leq \tau_k^{(i)} - \tau_{k-1}^{(i)})$, so that we can graft at the proper place, on this subtree, all the corresponding subtrees of the other types which have been constructed from the same procedure. Proceeding this way, we construct a unique forest $(\mathbf{f}, c_{\mathbf{f}}) \in \mathcal{F}_d^r$ and we easily check that $\Psi((\mathbf{f}, c_{\mathbf{f}})) = x$. \square

3. MULTITYPE BRANCHING TREES AND FORESTS

Let ν_i , $i = 1, \dots, d$ be distributions on \mathbb{Z}_+^d , such that $\nu = (\nu_1, \dots, \nu_d)$ is the progeny law of an irreducible, critical or subcritical, non-degenerated branching process, as defined in Section 1. Assume that we can define on the reference probability space (Ω, \mathcal{G}, P) introduced in Section 1, a family $(\mathbf{P}_c)_{c \in \{1, \dots, d\}^\infty}$ of probability measures and an infinite sequence $F = (T_k)_{k \geq 1}$ of independent random trees, such that for each $c = (c_1, c_2, \dots) \in \{1, \dots, d\}^\infty$ and $k \geq 1$, under \mathbf{P}_c , T_k is a branching tree, with progeny law ν , whose root has type $r(T_k) = c_k$. In particular, for any random time $\alpha : (\Omega, \mathcal{G}) \rightarrow \mathbb{N} \cup \{+\infty\}$, the sequence $\{T_1, \dots, T_\alpha\}$ is an element of \mathcal{F}_d . The infinite sequence F will be called a d -type branching forest with progeny law ν .

Let us denote by $F^{(i)}$ the subforest of type i of F , as it is defined in subsection 2.1. From the properties of ν , it follows that for each $i = 1, \dots, d$, the subforest $F^{(i)}$ is a.s. infinite, so that we may define a \mathbb{Z}^d valued infinite random sequence $X^{(i)} = (X^{i,1}, \dots, X^{i,d})$, for $i = 1, \dots, d$, in the same way as in (2.5), that is $X_0^{(i)} = 0$ and

$$(3.11) \quad X_{n+1}^{i,j} - X_n^{i,j} = p_j(u_{n+1}^{(i)}), \text{ if } i \neq j \text{ and } X_{n+1}^{i,i} - X_n^{i,i} = p_i(u_{n+1}^{(i)}) - 1, \quad n \geq 0,$$

where $(u_n^{(i)})_{n \geq 1}$ is the enumeration of $F^{(i)}$ in its breadth first search order.

Theorem 3.1. *Let F be a d -type branching forest with progeny law ν .*

1. *Then for any $c = (c_1, c_2, \dots) \in \{1, \dots, d\}^\infty$, under \mathbf{P}_c , the chains*

$$(3.12) \quad X^{(i)} = (X^{i,1}, \dots, X^{i,d}), \quad i = 1, \dots, d$$

are independent random walks with step distribution

$$\mathbf{P}_c(X_1^{(i)} = (q_1, \dots, q_d)) = \nu_i(q_1, \dots, q_{i-1}, q_i + 1, q_{i+1}, \dots, q_d).$$

In particular, their laws do not depend on c . For each $i = 1, \dots, d$, $X^{i,i}$ is a downward skip free random walk such that $\liminf_{n \rightarrow +\infty} X_n^{i,i} = -\infty$, a.s. and for $j \neq i$, $X^{i,j}$ is a renewal process.

2. *For all integer $r \geq 1$, almost surely there is $\mathbf{r} \in \mathbb{Z}_+^d$, with $r = r_1 + \dots + r_d$ and such that there is a smallest solution \mathbf{n} to the system (\mathbf{r}, X) .*
3. *Conversely, let Y be a copy of X and $c = (c_1, c_2, \dots) \in \{1, \dots, d\}^\infty$. Then to Y and c , we may associate a unique d -type branching forest, with progeny law ν and root type sequence c , whose coding random walk is Y .*

Proof. Part 1. just follows from the construction (3.11) of X . Since the order on the subforests $F^{(i)}$ does not depend on the particular topology of F , from the branching property, it is clear that the chains $X^{(i)}$, $i = 1, \dots, d$ are independent random walks. Then the expression of the law of $X^{(i)}$ is a direct consequence of (3.11). Recall that $X^{i,i}$ is the Lukasiewicz-Harris path of the subforest $F^{(i)}$, see section 6.2 of [13]. Moreover, since from the properties of ν , each subforest $F^{(i)}$ is a.s. infinite, the random walk $X^{i,i}$ satisfies $\liminf_{n \rightarrow \infty} X_n^{i,i} = -\infty$, a.s. The fact that $X^{i,j}$, for $i \neq j$ is a renewal process is obvious.

Then part 2 is a direct consequence of the construction of X and Theorem 2.7. Let $r \geq 1$ and first assume that the finite forest $\{T_1, \dots, T_r\}$ consisting in the r first trees of

F contains at least one vertex of each type. Let r be the unique element of \mathbb{Z}^d such that $\{T_1, \dots, T_r\} \in \mathcal{F}_d^r$. Then, by coding the forest $\{T_1, \dots, T_r\}$ and by applying Theorem 2.7, we obtain that there is $n \in \mathbb{N}^d$ which is the smallest solution of the system (r, X) . Now if for instance $n = 1, \dots, d - 1$ types are missing in the r first trees of F , then we can apply the same arguments by replacing d by $d - n$ in Theorem 2.7.

Then part 3. is a consequence of part 2. For each $r \geq 1$, we may associate a unique forest to Y with r trees. Since r can be arbitrarily large, the result is proved. \square

In the same spirit as in [9], the Lukasiewicz-Harris type coding that is displayed in Theorem 3.1 might be used to obtain invariance principles, for any functional that can be encoded simply enough. Besides this result should provide a way to obtain a proper definition of continuous multitype branching trees and forests. Actually, it is natural to think that the latter objects are coded by d independent, d -dimensional Lévy processes, with $d - 1$ increasing coordinates and a spectrally positive coordinate.

Now we are going to apply our coding of multitype branching forests to the law of their total progeny and give a proof of Theorem 1.2. To that aim, we first need to establish the crucial combinatorial Lemma 3.3. Let E be \mathbb{Z}_+ or a finite integer interval of the type $\{0, 1, \dots, m\}$, with $m \geq 1$ and let $g : E \rightarrow \mathbb{Z}^d$, be any application such that $g(0) = 0$. For $n \in E$ such that $n \geq 1$, the n -cyclical permutations of g are the n applications $g_{q,n}$, $q = 0, \dots, n - 1$ which are defined on E by:

$$(3.13) \quad g_{q,n}(h) \stackrel{(\text{def})}{=} \begin{cases} g(q+h) - g(q) & \text{if } 0 \leq h \leq n - q, \\ g(h - (n - q)) + g(n) - g(q) & \text{if } n - q \leq h \leq n \\ g(h) & \text{if } h \geq n. \end{cases}$$

Note that $g_{0,n} \equiv g$. The transformation $g \mapsto g_{q,n}$ consists in inverting the parts $\{g(h), 0 \leq h \leq q\}$ and $\{g(h), q \leq h \leq n\}$ in such a way that the new application, $g_{q,n}$, has the same values as g at 0 and n , i.e. $g_{q,n}(0) = 0$ and $g_{q,n}(n) = g(n)$.

Let $x \in S_d$, with finite length $n = (n_1, \dots, n_d) \in \mathbb{N}^d$ and recall the notation $x^{(i)} = (x^{i,1}, \dots, x^{i,d})$ from Definition 2.1. Then we define the n -cyclical permutations of x by

$$(3.14) \quad x_{q,n} := (x_{q_1, n_1}^{(1)}, \dots, x_{q_d, n_d}^{(d)}), \quad \text{for all } q = (q_1, \dots, q_d) \leq n - 1_d,$$

where we have set $1_d = (1, 1, \dots, 1)$. Each sequence $x_{q,n}$ will simply be called a cyclical permutation of x . Note that there are $n_1 n_2 \dots n_d$, cyclical permutations of x . Let $r = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$ be such that $r > 0$ and assume that n is a solution of the system (r, x) . Then note that n is also a solution of the system $(r, x_{q,n})$, for all $q \leq n - 1_d$, that is,

$$r_j + \sum_{i=1}^d x_{q,n}^{i,j}(n_i) = 0, \quad j = 1, \dots, d.$$

This remark raises the question of the number of cyclical permutations $x_{q,n}$ of x , such that n is the smallest solution of the system $(r, x_{q,n})$.

Definition 3.2. Let $x \in S_d$, with finite length $\mathbf{n} = (n_1, \dots, n_d) \in \mathbb{N}^d$. Let $\mathbf{r} = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$ be such that $\mathbf{r} > 0$ and assume that \mathbf{n} is a solution of the system (\mathbf{r}, x) . For $\mathbf{q} \leq \mathbf{n} - \mathbf{1}_d$, we say that $x_{\mathbf{q}, \mathbf{n}}$ is a good \mathbf{n} -cyclical permutation of x with respect to \mathbf{r} , if \mathbf{n} is the smallest solution of the system $(\mathbf{r}, x_{\mathbf{q}, \mathbf{n}})$, that is $x_{\mathbf{q}, \mathbf{n}} \in S_{\mathbf{q}, \mathbf{n}}^{\mathbf{r}}$. When no confusion is possible, we will simply say that $x_{\mathbf{q}, \mathbf{n}}$ is a good cyclical permutation of x .

Here is our crucial combinatorial cyclic lemma.

Lemma 3.3. Let $x \in S_d$, with length $\mathbf{n} = (n_1, \dots, n_d) \in \mathbb{N}^d$ and let $\mathbf{r} = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$ be such that $\mathbf{r} > 0$. Assume that \mathbf{n} is a solution of the system (\mathbf{r}, x) such that $x^{i,i}(n_i) \neq 0$, for all $i = 1, \dots, d$. Then the number of good cyclical permutations of x is $\det((-x^{i,j}(n_i))_{1 \leq i, j \leq d})$.

This lemma will be proved in Section 4. It is used for the proof of the following extension of the Ballot Theorem.

Theorem 3.4 (Multivariate Ballot Theorem). Let $Y = (Y^{(1)}, \dots, Y^{(d)})$ be a stochastic process defined on (Ω, \mathcal{G}, P) , with $Y^{(i)} = (Y^{i,j}(h), 1 \leq j \leq d, h \geq 0)$, $i = 1, \dots, d$ and $Y_0 = 0$. We assume that the coordinates $Y^{i,j}$, for $i \neq j$ are \mathbb{Z}_+ valued, nondecreasing and that the coordinates $Y^{i,i}$ are \mathbb{Z} valued and downward skip free. Fix $\mathbf{n} = (n_1, \dots, n_d) \in \mathbb{N}^d$, then we assume further that the process Y is \mathbf{n} -cyclically exchangeable, that is for any $\mathbf{q} = (q_1, \dots, q_d) \in \mathbb{Z}_+^d$ such that $\mathbf{q} \leq \mathbf{n} - \mathbf{1}_d$,

$$Y_{\mathbf{q}, \mathbf{n}} \stackrel{(law)}{=} Y,$$

where $Y_{\mathbf{q}, \mathbf{n}}$ is defined as in (3.14) for deterministic functions. Then for any $\mathbf{r} = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$ such that $\mathbf{r} > 0$ and k_{ij} , $1 \leq i, j \leq d$, such that $k_{ij} \in \mathbb{Z}_+$, for $i \neq j$ and $-k_{jj} = r_j + \sum_{i \neq j} k_{ij}$,

$$(3.15) \quad \begin{aligned} & P(Y_{n_i}^{i,j} = k_{ij}, 1 \leq i, j \leq d \text{ and } \mathbf{n} \text{ is the smallest solution of } (\mathbf{r}, Y)) \\ &= \frac{\det(-k_{ij})}{n_1 n_2 \dots n_d} P(Y_{n_i}^{i,i} = k_{ii}, Y_{n_i}^{i,j} = k_{ij}, i \neq j, 1 \leq i, j \leq d). \end{aligned}$$

Proof. If $P(Y_{n_i}^{i,j} = k_{ij}, i \neq j) = 0$, then the result is clearly true. Suppose that it is not the case and let $y = (y^{(1)}, \dots, y^{(d)})$ be a deterministic function such that for all $1 \leq i, j \leq d$, $y^{i,j}$ is defined on \mathbb{Z}_+ , $y^{i,j}(0) = 0$, $y^{i,j}(n_i) = k_{ij}$ and

$$(3.16) \quad P((Y^{(i)}(h), 0 \leq h \leq n_i)) = (y^{(i)}(h), 0 \leq h \leq n_i) > 0.$$

For $\mathbf{h} = (h_1, \dots, h_d) \in \mathbb{Z}_+^d$, we set $Y(\mathbf{h}) := (Y^{(1)}(h_1), \dots, Y^{(d)}(h_d))$ and for $0 \leq \mathbf{q} \leq \mathbf{n} - \mathbf{1}_d$, using the notation of (3.14), we set $y_{\mathbf{q}, \mathbf{n}}(\mathbf{h}) := (y_{q_1, n_1}^{(1)}(h_1), \dots, y_{q_d, n_d}^{(d)}(h_d))$. Let us consider the set

$$E_{y, \mathbf{n}} = \{(y_{\mathbf{q}, \mathbf{n}}(\mathbf{h}), 0 \leq \mathbf{h} \leq \mathbf{n}) : 0 \leq \mathbf{q} \leq \mathbf{n} - \mathbf{1}\},$$

of \mathbf{n} -cyclical permutations of y over the multidimensional interval $[0, \mathbf{n}]$. Then $\text{Card}(E_{y, \mathbf{n}}) = n_1 n_2 \dots n_d$ and since $Y = (Y^{(1)}, \dots, Y^{(d)})$ is a cyclically exchangeable chain, the law of $(Y(\mathbf{h}), 0 \leq \mathbf{h} \leq \mathbf{n})$, conditionally to the set $\{(Y(\mathbf{h}), 0 \leq \mathbf{h} \leq \mathbf{n}) \in E_{y, \mathbf{n}}\}$ is the uniform law in the set $E_{y, \mathbf{n}}$. Moreover, assume that $k_{ii} \neq 0$ for all $i = 1, \dots, d$, then conditionally

to the set $\{(Y(h), 0 \leq h \leq n) \in E_{y,n}\}$, from Lemma 3.3, the number of good cyclical permutations of $(Y(h), 0 \leq h \leq n)$ is $\det(-k_{ij})$. Therefore,

$$P(Y_{n_i}^{i,j} = k_{ij}, 1 \leq i, j \leq d \text{ and } n \text{ is the smallest solution of } (r, Y) \mid (Y(h), 0 \leq h \leq n) \in E_{y,n}) = \frac{\det(-k_{ij})}{n_1 n_2 \dots n_d}.$$

Then we obtain the result by summing the identity

$$\begin{aligned} & P(Y_{n_i}^{i,j} = k_{ij}, 1 \leq i, j \leq d \text{ and } n \text{ is the smallest solution of } (r, Y), (Y(h), 0 \leq h \leq n) \in E_{y,n}) \\ &= \frac{\det(-k_{ij})}{n_1 n_2 \dots n_d} P((Y(h), 0 \leq h \leq n) \in E_{y,n}), \end{aligned}$$

over all functions y satisfying (3.16) and with different sets $E_{y,n}$ of cyclical permutations. Finally, if $k_{ii} = 0$, for some $i = 1, \dots, d$, then since $n_i \geq 1$, we can see that both members of identity (3.15) are equal to 0. \square

Proof of Theorem 1.2: Let r , n and k_{ij} be as in the statement. Let F be a d -type branching forest with progeny law ν , as defined at the beginning of this section and such that the r first trees have root type sequence $(c_1, \dots, c_r) \in C_d^r$. Let X be the coding random walk of F , as defined in (3.11). Recall the notation of Theorem 1.2, then from the coding of Subsection 2.2 and Theorem 3.1, we may check that

$$(3.17) \quad \begin{aligned} & \mathbb{P}_r(O_1 = n_1, \dots, O_d = n_d, A_{ij} = k_{ij}, 1 \leq i, j \leq d, i \neq j) \\ &= \mathbf{P}_c(X_{n_i}^{i,j} = k_{ij}, 1 \leq i, j \leq d \text{ and } n \text{ is the smallest solution of } (r, X)). \end{aligned}$$

Assume first that $n \in \mathbb{N}^d$, then since X is clearly cyclically exchangeable in the sense of Theorem 3.4, we obtain by applying this theorem,

$$\begin{aligned} & \mathbf{P}_c(X_{n_i}^{i,j} = k_{ij}, 1 \leq i, j \leq d \text{ and } n \text{ is the smallest solution of } (r, X)) \\ &= \frac{\det(-k_{ij})}{n_1 n_2 \dots n_d} P(X_{n_i}^{i,i} = k_{ii}, X_{n_i}^{i,j} = k_{ij}, i \neq j, 1 \leq i, j \leq d). \end{aligned}$$

On the other hand, since from Theorem 3.1, the random walks $X^{(i)}$, $i = 1, \dots, d$ are independent, we have

$$\begin{aligned} & \mathbf{P}_c(X_{n_i}^{i,j} = k_{ij}, 1 \leq i, j \leq d \text{ and } n \text{ is the smallest solution of } (r, X)) \\ &= \frac{\det(-k_{ij})}{n_1 n_2 \dots n_d} \prod_{i=1}^d P(X_{n_i}^{i,i} = k_{ii}, X_{n_i}^{i,j} = k_{ij}, i \neq j, 1 \leq j \leq d). \end{aligned}$$

Then from the expression of the law of X given in Theorem 3.1, we obtain

$$(3.18) \quad \begin{aligned} & \mathbf{P}_c(X_{n_i}^{i,j} = k_{ij}, 1 \leq i, j \leq d \text{ and } n \text{ is the smallest solution of } (r, X)) \\ &= \frac{\det(-k_{ij})}{n_1 n_2 \dots n_d} \prod_{i=1}^d \nu_i^{*n_i}(k_{i,1}, \dots, k_{i,i-1}, n_i + k_{ii}, k_{i,i+1}, \dots, k_{i,d}), \end{aligned}$$

and the result is proved in this case.

Now with no loss of generality, let us assume that for some $0 < d' < d$, we have $n_1, \dots, n_{d'} \in \mathbb{N}$ and $n_{d'+1} = \dots = n_d = 0$. We point out that from the assumption $n_j \geq -k_{jj} = r_j + \sum_{i \neq j} k_{ij}$, in this case we necessarily have $r_j = 0$ and $k_{ij} = 0$, for all $i = 1, \dots, d$ and $j = d' + 1, \dots, d$. Then provided we also have $k_{ij} = 0$, for all $i = d' + 1, \dots, d$ and $j = 1, \dots, d, i \neq j$,

$$\begin{aligned} & \mathbf{P}_c(X_{n_i}^{i,j} = k_{ij}, d' + 1 \leq i \leq d, 1 \leq j \leq d) \\ &= \prod_{i=d'+1}^d \nu_i^{*n_i}(k_{i1}, \dots, k_{i,i-1}, n_i + k_{ii}, k_{i,i+1}, \dots, k_{i,d}) = 1. \end{aligned}$$

Define the chain X restricted to $\mathbb{Z}^{d'}$, by $X' = (X'^{(1)}, \dots, X'^{(d')})$, where $X'^{(i)} = (X^{i,1}, \dots, X^{i,d'})$, $i = 1, \dots, d'$. Set also $\mathbf{n}' = (n_1, \dots, n_{d'})$ and $\mathbf{r}' = (r_1, \dots, r_{d'})$. Then under our assumption on the integers k_{ij} , the following identity is satisfied,

$$\begin{aligned} & \{X_{n_i}^{i,j} = k_{ij}, 1 \leq i, j \leq d \text{ and } \mathbf{n} \text{ is the smallest solution of } (\mathbf{r}, X)\} \\ &= \{X_{n_i}^{i,j} = k_{ij}, 1 \leq i, j \leq d' \text{ and } \mathbf{n}' \text{ is the smallest solution of } (\mathbf{r}', X') \text{ and} \\ & \quad X_{n_i}^{i,j} = 0, 1 \leq i \leq d', d' + 1 \leq j \leq d\}, \end{aligned}$$

so that identity (3.17) can be rewritten as,

$$\begin{aligned} & \mathbb{P}_r(O_1 = n_1, \dots, O_d = n_d, A_{ij} = k_{ij}, 1 \leq i, j \leq d, i \neq j) \\ &= \mathbf{P}_c(X_{n_i}^{i,j} = k_{ij}, 1 \leq i, j \leq d' \text{ and } \mathbf{n}' \text{ is the smallest solution of } (\mathbf{r}', X') \text{ and} \\ & \quad X_{n_i}^{i,j} = 0, 1 \leq i \leq d', d' + 1 \leq j \leq d). \end{aligned}$$

Moreover, conditionally on the set

$$\{X_{n_i}^{i,j'} = k_{ij'}, X_{n_i}^{i,j} = 0, 1 \leq i \leq d', 1 \leq j' \leq d', d' + 1 \leq j \leq d\},$$

the chain X' is cyclically exchangeable, so that we can conclude in the same way as above that

$$\begin{aligned} & \mathbb{P}_r(O_1 = n_1, \dots, O_d = n_d, A_{ij} = k_{ij}, 1 \leq i, j \leq d, i \neq j) \\ &= \frac{\det((-k_{ij})_{1 \leq i, j \leq d'})}{n_1 n_2 \dots n_{d'}} \prod_{i=1}^{d'} \nu_i^{*n_i}(k_{i1}, \dots, k_{i,i-1}, n_i + k_{ii}, k_{i,i+1}, \dots, k_{i,d}) \\ &= \frac{\det((-k_{ij})_{1 \leq i, j \leq d'})}{n_1 n_2 \dots n_{d'}} \prod_{i=1}^d \nu_i^{*n_i}(k_{i1}, \dots, k_{i,i-1}, n_i + k_{ii}, k_{i,i+1}, \dots, k_{i,d}). \end{aligned}$$

Finally, if $k_{ij} \neq 0$, for some $i = d' + 1, \dots, d$ and $j = 1, \dots, d$, then the first and the third members of the above equality are equal to 0. So the proof is complete. \square

4. PROOF OF LEMMAS 2.2 AND 3.3

Proof of Lemma 2.2. Assume that there is a solution $\mathbf{s} = (s_1, \dots, s_d)$ to the system (\mathbf{r}, \mathbf{x}) , that is: $r_j + \sum_{i=1}^d x^{i,j}(s_i) = 0$, $j = 1, \dots, d$. Let us write this equation in the following form:

$$r_j + \sum_{i \neq j} x^{i,j}(s_i) + x^{j,j}(s_j) = 0, \quad j = 1, \dots, d.$$

Then recall that for fixed j , when the k_i 's increase, the term $\sum_{i \neq j} x^{i,j}(k_i)$ increases and when k_j increases, the term $x^{j,j}(k_j)$ may decrease only by jumps of amplitude -1 .

For $\mathbf{k} = (k_1, \dots, k_d) = 0$, we have $r_j + \sum_{i \neq j} x^{i,j}(k_i) + x^{j,j}(k_j) = r_j$, $j = 1, \dots, d$. So for the left hand side of the later equation to reach 0, each k_i has to be at least $\tau^{(i)}(r_i)$, where $\tau^{(i)}$ has been defined in (2.7). (In this proof, we found more convenient to use the notation $\tau^{(i)}(k)$ for $\tau_k^{(i)}$.) Then either

$$r_j + \sum_{i \neq j} x^{i,j}(\tau^{(i)}(r_i)) + x^{j,j}(\tau^{(j)}(r_j)) = 0, \quad j = 1, \dots, d,$$

or all of the terms $r_j + \sum_{i \neq j} x^{i,j}(\tau^{(i)}(r_i)) + x^{j,j}(\tau^{(j)}(r_j))$, $j = 1, \dots, d$ are greater or equal than 0, at least one of them being strictly greater than 0.

Then in the latter case, for $r_j + \sum_{i \neq j} x^{i,j}(k_i) + x^{j,j}(k_j)$ to attain 0, each of the k_j 's has to be at least $\tau^{(j)}(r_j + \sum_{i \neq j} x^{i,j}(\tau^{(i)}(r_i)))$. This argument can be repeated until all of the terms $r_j + \sum_{i \neq j} x^{i,j}(k_i) + x^{j,j}(k_j)$ attain 0. More specifically, set $v_1^{(j)} = r_j$ and for $n \geq 1$,

$$v_{n+1}^{(j)} = r_j + \sum_{i \neq j} x^{i,j}(\tau^{(i)}(v_n^{(i)})),$$

and set $k_j^{(n)} = \tau^{(j)}(v_n^{(j)})$. For $n \geq 1$, either

$$r_j + \sum_{i \neq j} x^{i,j}(k_i^{(n)}) + x^{j,j}(k_j^{(n)}) = 0, \quad j = 1, \dots, d,$$

or all of the terms $r_j + \sum_{i \neq j} x^{i,j}(k_i^{(n)}) + x^{j,j}(k_j^{(n)})$, $j = 1, \dots, d$ are greater or equal than 0, at least one of them being strictly greater than 0. In the later case, for all of the terms $r_j + \sum_{i \neq j} x^{i,j}(k_i) + x^{j,j}(k_j)$, $j = 1, \dots, d$ to vanish, the index \mathbf{k} has to be at least $\mathbf{k}^{(n+1)} = (k_1^{(n+1)}, \dots, k_d^{(n+1)})$. But since there is a solution \mathbf{s} to the equation (r, x) , there is necessarily a finite index n_0 such that $\mathbf{k}^{(n_0)} \leq \mathbf{s}$ and

$$r_j + \sum_{i \neq j} x^{i,j}(k_i^{(n_0)}) + x^{j,j}(k_j^{(n_0)}) = 0.$$

That is, for all $j = 1, \dots, d$, $k_j^{(n_0)} = \tau^{(j)}(r_j + \sum_{i \neq j} x^{i,j}(k_i^{(n_0)}))$. Hence $\mathbf{k}^{(n_0)}$ is the smallest solution of the system (r, x) . Moreover by definition, $k_i^{(n_0)} = \min\{n : x_n^{i,i} = \min_{0 \leq k \leq k_i^{(n_0)}} x_k^{i,i}\}$. This proves the first part of the lemma.

Let $\mathbf{r}' = (r'_1, \dots, r'_d) \in \mathbb{Z}_+^d$ be such that $\mathbf{r}' \leq \mathbf{r}$. Then we prove that there is a smallest solution to the system (\mathbf{r}', x) similarly. More specifically, since there is a smallest solution to the equation

$$r'_j + \sum_{i \neq j} x^{i,j}(s_i) + x^{j,j}(s_j) = r'_j - r_j \leq 0, \quad j = 1, \dots, d,$$

then by the same arguments as in the first part, we prove that there is a smallest solution to the equation

$$r'_j + \sum_{i \neq j} x^{i,j}(s_i) + x^{j,j}(s_j) = 0, \quad j = 1, \dots, d.$$

Let \mathbf{k} and \mathbf{k}' be respectively the smallest solutions of (\mathbf{r}, x) and (\mathbf{r}', x) . Then we have,

$$r_j - r'_j + r'_j + \sum_{i=1}^d x^{i,j}(k'_i) + \sum_{i=1}^d \tilde{x}^{i,j}(k_i - k'_i) = 0.$$

Since $r'_j + \sum_{i=1}^d x^{i,j}(k'_i) = 0$, the above equation shows that $\mathbf{k} - \mathbf{k}'$ is a solution of the system $(\mathbf{r} - \mathbf{r}', \tilde{x})$. Moreover, if \mathbf{k}'' was a strictly smaller solution of $(\mathbf{r} - \mathbf{r}', \tilde{x})$, than $\mathbf{k} - \mathbf{k}'$ (i.e. $\mathbf{k}'' < \mathbf{k} - \mathbf{k}'$) then we would have

$$\begin{aligned} r_j - r'_j + r'_j + \sum_{i=1}^d x^{i,j}(k'_i) + \sum_{i=1}^d \tilde{x}^{i,j}(k''_i) &= r_j + \sum_{i=1}^d x^{i,j}(k'_i + k''_i) \\ &= 0, \end{aligned}$$

so that $\mathbf{k}' + \mathbf{k}''$ would be a solution of (\mathbf{r}, x) , strictly smaller than \mathbf{k} , which is a contradiction. \square

Proof of Lemma 3.3. Recall from Lemma 2.5 that to each forest $(\mathbf{f}, c_{\mathbf{f}}) \in \bar{\mathcal{F}}_d^{\mathbf{r}}$, we can associate a system (\mathbf{r}, x) , with smallest solution \mathbf{k} , where $x \in \bar{S}_d^{\mathbf{r}}$. We then define the *Laplacian matrix* of each forest $(\mathbf{f}, c_{\mathbf{f}}) \in \bar{\mathcal{F}}_d^{\mathbf{r}}$ by $(k_{ij}) = (x^{i,j}(k_i))$. We will also sometimes say that (k_{ij}) is the *Laplacian matrix* of the associated coding sequence x . (Actually, in the conventional terminology of graph theory, (k_{ij}) would rather be the Laplacian matrix of a directed graph on the set of vertices $\{1, \dots, d\}$, with incidence matrix $U = (u_{ij})_{1 \leq i, j \leq d}$, where $u_{ij} = k_{ij}$, if $i \neq j$ and $u_{ii} = 0$, $i = 1, \dots, d$.) Recall that $-k_{jj} = r_j + \sum_{i \neq j} k_{ij}$ and that for $i \neq j$, k_{ij} is the total number of vertices in $\mathbf{v}(\mathbf{f})$ with type j , whose parent has type i and that $-k_{ii}$ is the total number of vertices of type i in $\mathbf{v}(\mathbf{f})$.

Definition 4.1. Let $\mathbf{r} \in \mathbb{Z}_+^d$ be such that $\mathbf{r} > 0$.

1. An element $(\mathbf{f}, c_{\mathbf{f}})$ of $\bar{\mathcal{F}}_d^{\mathbf{r}}$ is said to be a *simple forest* if for each type $i = 1, \dots, d$, at most one vertex of type i in $\mathbf{v}(\mathbf{f})$ has children (the others are leaves). We assume moreover that the types of the roots of simple forests are ranked in the increasing order.
2. For $x \in \bar{S}_d^{\mathbf{r}}$, the system (\mathbf{r}, x) , with smallest solution $\mathbf{k} = (k_1, \dots, k_d)$ is called a *simple system* if for all $i = 1, \dots, d$, there is $k'_i \leq k_i - 1$, such that for all $k = 0, \dots, k'_i$ and $i \neq j$, $x^{i,j}(k) = 0$ and for all $k = (k'_i + 1) \vee k_i, \dots, k_i$ and $i \neq j$, $x^{i,j}(k) = x^{i,j}(k_i)$.

In other words, simple systems are elements $x \in \bar{S}_d^{\mathbf{r}}$ such that for all $i = 1, \dots, d$, the sequences $x^{i,j}$, $j = 1, \dots, d$ have at most one positive jump that occurs at the same point in $\{0, \dots, k_i\}$. Then we have the following straightforward result.

Proposition 4.2. A forest of $\bar{\mathcal{F}}_d^{\mathbf{r}}$ is simple if and only if its associated system (as defined in Lemma 2.5) is simple.

Definition 4.3. An elementary forest is a forest of \mathcal{F}_d that contains exactly one vertex of each type. In particular, each elementary forest contains exactly d vertices and is coded by the d couples (j_i, i) , $i = 1, \dots, d$, where j_i is the type of the parent of the vertex of type i . If the vertex of type i is a root, then we set $j_i = 0$. We define the set D of vectors (j_1, \dots, j_d) , $0 \leq j_i \leq d$ such that (j_i, i) , $i = 1, \dots, d$ codes an elementary forest.

Lemma 4.4. Let $\mathbf{r} = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$ such that $\mathbf{r} > \mathbf{0}$ and let (\mathbf{r}, x) be a simple system with smallest solution $\mathbf{k} = (k_1, \dots, k_d) \in \mathbb{N}^d$, then the number of good cyclical permutations of x is

$$\sum_{(j_1, \dots, j_d) \in D} \prod_{i=1}^d k_{j_i, i},$$

where for $i, j = 1, \dots, d$, $k_{ij} = x^{i,j}(k_i)$ and where we have set $k_{0,i} = r_i$.

Proof. The proof will be performed by reasoning on forests. Let (\mathbf{r}, x) be as in the statement. We first make the additional assumption that for all $i = 1, \dots, d$, there is $j \neq i$ such that $k_{ij} > 0$. Then observe that under this assumption, all the cyclical permutations of x are different from each other, since each sequence $x^{i,j}$, for $i \neq j$ has exactly one (strictly positive) jump.

Then note that for each good cyclical permutation $x_{\mathbf{q}, \mathbf{k}}$ of x , the system $(\mathbf{r}, x_{\mathbf{q}, \mathbf{k}})$ is simple itself. Thus from Lemma 2.5 and Proposition 4.2, to each good cyclical permutation $x_{\mathbf{q}, \mathbf{k}}$ of x , we may associate a unique simple forest of $\bar{\mathcal{F}}_d^{\mathbf{r}}$, since, by definition, root type sequences of simple forests are ranked in the increasing order. Moreover if $\mathbf{q} \neq \mathbf{q}'$ and $\mathbf{q}, \mathbf{q}' < \mathbf{k}$, then from our additional assumption, $x_{\mathbf{q}, \mathbf{k}} \neq x_{\mathbf{q}', \mathbf{k}}$, so that the associated simple forests are different. Conversely the simple system (\mathbf{r}, y) which is associated to each simple forest with Laplacian matrix (k_{ij}) , through Lemma 2.5 is necessarily obtained from a good cyclical permutation of x . Indeed the corresponding sequences $y^{i,j}$, for $i \neq j$, have exactly one jump and are such that $y^{i,j}(k_i) = k_{ij}$. So y is nothing but a cyclical permutation of x . These arguments prove that the number of good cyclical permutations of x is equal to the number of simple forests with Laplacian matrix (k_{ij}) .

Then let us prove that the number of simple forests with Laplacian matrix (k_{ij}) is $\sum_{(j_1, \dots, j_d) \in D} \prod_{i=1}^d k_{j_i, i}$. First observe that to each simple forest, we can associate a unique elementary forest in the following way: the vertex of type j in the elementary forest is the parent of the vertex of type i if, in the simple forest the parent of the vertex of type i who has children has type j (recall that $j = 0$ if the vertex of type i is a root). An example of an elementary forest associated to a simple forest is given in Figure 6. Then let $(j_1, \dots, j_d) \in D$. We easily see that the monomial $\prod_{i=1}^d k_{j_i, i}$ is the number of simple forests such that for each i , the parent of the vertex of type i who has children has type j_i . Indeed, there are $k_{j_i, i}$ possibilities to choose the vertex of type i who has children. In other words, $\prod_{i=1}^d k_{j_i, i}$ is the number of all possible simple forests to which we can associate the same elementary forest which is coded by (j_i, i) , $i = 1, \dots, d$. Then in order to obtain the total number of simple forests with Laplacian matrix (k_{ij}) , it remains to perform the summation of these monomials over all the possible elementary forests. So we obtained the formula of the statement, under the additional assumption made at the beginning of this proof.

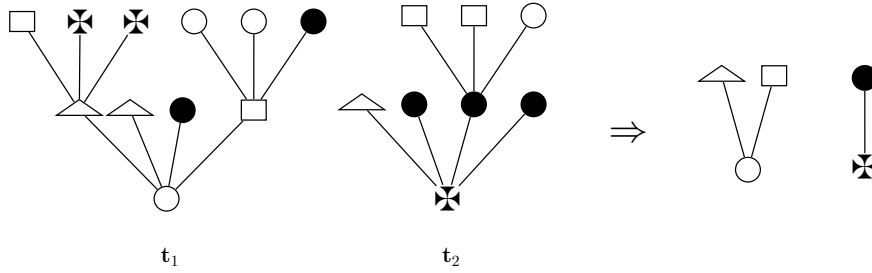


FIGURE 6. An elementary forest and its associated simple forest.

Now assume that for all $i = 1, 2, \dots, d-1$, there is $j \neq i$ such that $k_{ij} > 0$ and that $k_{dj} = 0$, for all $j = 1, \dots, d$. Then in any forest with Laplacian matrix (k_{ij}) , all vertices of type d are leaves. Since these leaves are undistinguished, such a simple forest is determined only by the relative positions of the vertices of type $1, 2, \dots, d-1$. So we derive from what has just been proved, that the number of simple forests with Laplacian matrix (k_{ij}) is $h_{d-1} := \sum_{(j_1, \dots, j_{d-1}) \in D'} \prod_{i=1}^{d-1} k_{j_i, i}$ where D' is the set that is defined in Definition 4.3 and where we replaced d by $d-1$. To these h_{d-1} forests correspond h_{d-1} simple systems to which correspond h_{d-1} good cyclical permutations of x . Then in order to obtain all the good cyclical permutations of x it remains to consider the k_d cyclical permutations of the sequence $x^{(d)}$. Since the latter are all identical, we have actually $k_d \times h_{d-1}$ good cyclical permutations of x . Then we conclude by noticing that $k_d \times h_{d-1} = \sum_{(j_1, \dots, j_d) \in D} \prod_{i=1}^d k_{j_i, i}$, since $k_d = r_d + \sum_{i \neq d} k_{id} = \sum_{i=0, i \neq d}^d k_{id}$. The general case where $k_{ij} = 0$, for all $j = 1, \dots, d$, for other types i is obtained in the same manner. \square

Lemma 4.5. For any $r \in \mathbb{Z}_+^d$ and any integer valued matrix $(k_{ij})_{1 \leq i, j \leq d}$, such that k_{ij} is nonnegative for $i \neq j$ and $-k_{jj} = r_j + \sum_{i \neq j} k_{ij}$, $j = 1, \dots, d$,

$$\det(-k_{ij}) = \sum_{(j_1, \dots, j_d) \in D} \prod_{i=1}^d k_{j_i, i},$$

where the set D is defined in Definition 4.3 and $k_{0,i} = r_i$.

Proof. The proof is a direct consequence of the matrix tree theorem for directed graphs, due to Tutte [14], see Section 3.6, page 470 therein. However, Theorem 3.1 in [12] that implies Tutte's theorem is actually easier to apply, since it uses a setting which is closer to ours. Let us consider a set $\{v_0, v_1, \dots, v_d\}$ of $d+1$ vertices and in the notations of [12], set $W = L = \{v_0\}$. Then the family $\mathcal{F}_{W,L}$ that is described in Theorem 3.1 of [12] is in bijection with the set of elementary forests, or equivalently with the set D , and identity (3.2) in this theorem is exactly $\det(-k_{ij}) = \sum_{(j_1, \dots, j_d) \in D} \prod_{i=1}^d k_{j_i, i}$. \square

Lemma 4.6. *Let $x \in S_d$, with length $k = (k_1, \dots, k_d) \in \mathbb{N}^d$. Let $r = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$ be such that $r > 0$ and assume that k is a solution of the system (r, x) . Assume moreover that $x_k^{i,i} = -k$, for all $k = 0, \dots, k_i$ and $i = 1, \dots, d$. Then the number of good cyclical permutations of x is $\det((x^{i,j}(k_i))_{1 \leq i, j \leq d})$.*

Proof. Let x be as in the statement. If for all $i = 1, \dots, d$, the sequences $x^{i,j}$, $j = 1, \dots, d$ have at most one positive jump that occurs at the same point in $\{0, \dots, k_i\}$, then from Definition 4.1, Lemma 4.4 and Lemma 4.5, there is a cyclical permutation x' of x such that (r, x') is a simple system and the result follows in this case. In general, let us prove that there is a simple system, with Laplacian matrix $(x^{i,j}(k_i))$, and whose number of good cyclical permutations is the same as this of x .

Fix any index $m = 1, \dots, d$ and assume without loss of generality, that $m \neq 1$, and

$$x^{m,1}(k_m) - x^{m,1}(k_m - 1) > 0.$$

From x , we define a new sequence $\tilde{x} \in S_d$ with length k as follows:

$$(4.19) \quad \begin{cases} \tilde{x}^{m,j}(k) = x^{m,j}(k), & j = 2, \dots, d, k = 1, \dots, k_m, \\ \tilde{x}^{m,1}(k) = x^{m,1}(k), & k = 1, \dots, k_m - 2, k_m, \\ \tilde{x}^{m,1}(k_m - 1) = x^{m,1}(k_m - 1) + 1. \end{cases}$$

All the other coordinates remain unchanged, that is,

$$\tilde{x}^{i,j}(k) = x^{i,j}(k), \quad j = 1, \dots, d, \quad k = 1, \dots, k_i, \quad i = 1, \dots, d, \quad i \neq m.$$

The sequence \tilde{x} is obtained from x by decreasing by one unit, the last jump of the coordinate $x^{m,1}$, that is $x^{m,1}(k_m) - x^{m,1}(k_m - 1)$. (Therefore the jump $x^{m,1}(k_m - 1) - x^{m,1}(k_m - 2)$ is increased by one unit.) Denote by $N_{r,x}$ the number of good cyclical permutations of x . We claim that

$$(4.20) \quad N_{r,\tilde{x}} \geq N_{r,x}.$$

To achieve this aim, first observe that k is a solution of the system (r, \tilde{x}) , that is

$$r_j + \sum_{i=1}^d \tilde{x}^{i,j}(k_i) = 0, \quad j = 1, \dots, d,$$

and note the straightforward inequality,

$$\tilde{x}^{i,j}(k) \geq x^{i,j}(k), \quad i, j = 1, \dots, d, \quad k = 1, \dots, k_i.$$

Therefore, if k is the smallest solution of (r, x) then it is also the smallest solution of (r, \tilde{x}) . Moreover, for $q = (q_1, \dots, q_d) \in \mathbb{N}^d$ satisfying $q_m < k_m - 1$, one easily checks that the same inequality holds, that is

$$\tilde{x}_{q,k}^{i,j}(k) \geq x_{q,k}^{i,j}(k), \quad i, j = 1, \dots, d, \quad k = 1, \dots, k_i,$$

so that if k is the smallest solution of $(r, x_{q,k})$ (i.e. $x_{q,k}$ is a good cyclical permutation of x) then it is also the smallest solution of $(r, \tilde{x}_{q,k})$.

Now it remains to study the case where $q_m = k_m - 1$. Assume that $x_{q,k}$ is a good cyclical permutation of x , but that $\tilde{x}_{q,k}$ is not a good cyclical permutation of \tilde{x} . Then in order to obtain the inequality (4.20), we have to find a good cyclical permutation $\tilde{x}_{1,k}$ of \tilde{x}

such that $x_{1,k}$ is not a good cyclical permutation of x . Let us define the sequence $\hat{x} \in S_d$ with length k , which is obtained by decreasing by one unit the first coordinate of $x_{q,k}^{(m)}$, that is

$$\begin{cases} \hat{x}^{(i)} \equiv x_{q,k}^{(i)}, & i \neq m, \\ \hat{x}^{(m)}(k) = x_{q,k}^{(m)}(k) - \mathbf{e}_1, & k \geq 1, \end{cases}$$

where $\mathbf{e}_1 = (1, 0, \dots, 0)$, is the d dimensional unit vector. Since k is the smallest solution of $(r, x_{q,k})$, then k is the smallest solution of $(r + \mathbf{e}_1, \hat{x})$ by the definition of \hat{x} . Moreover, from Lemma 2.2 (ii), the system (\mathbf{e}_1, \hat{x}) admits a smallest solution which is less than k . Let us call p this solution. Then $p > 0$ and from Lemma 2.2 (ii), $k - p$ is the smallest solution of $(r, \hat{x}_{p,k})$.

Then let us consider the cyclical permutation of $\tilde{x}_{q,k}$ at p . It is a cyclical permutation of \tilde{x} that we shall denote by $\tilde{x}_{1,k}$. Note that \hat{x} and $\tilde{x}_{q,k}$ only differ from the last jump of $\hat{x}^{m,1}$ and $\tilde{x}_{q,k}^{m,1}$, more specifically, $\tilde{x}_{q,k}(k_m) - \tilde{x}_{q,k}(k_m - 1) = \hat{x}(k_m) - \hat{x}(k_m - 1) + 1$. Then from the above constructions, we can see that $\tilde{x}_{1,k}$ is obtained as follows:

$$\tilde{x}_{1,k}^{(i)}(k) = \begin{cases} \hat{x}_{p,k}^{(i)}(k), & k \leq k_i - p_i, i \neq m, \\ \hat{x}_{p,k}^{(m)}(k), & k < k_m - p_m, \\ \hat{x}_{p,k}^{(i)}(k_i - p_i) + \hat{x}^{(i)}(k - (k_i - p_i)), & k_i - p_i \leq k \leq k_i, i \neq m, \\ \hat{x}_{p,k}^{(m)}(k_m - p_m) + \hat{x}^{(m)}(k - (k_m - p_m)) + \mathbf{e}_1, & k_m - p_m \leq k \leq k_m. \end{cases}$$

Since $k - p$ is the smallest solution of $(r, \hat{x}_{p,k})$ and since $\tilde{x}_{1,k}$ is strictly greater than $\hat{x}_{p,k}$ at point $k - p$, there is no solution to $(r, \tilde{x}_{1,k})$, on the (multidimensional) interval $[0, k - p]$. Moreover, also from the construction of $(r, \tilde{x}_{1,k})$, since p is the smallest solution of (\mathbf{e}_1, \hat{x}) , the only solution of $(r, \tilde{x}_{1,k})$ on the interval $[k - p, k]$ is k . Therefore, the smallest solution of $(r, \tilde{x}_{1,k})$ is k .

On the other hand, note the following identity

$$(x_{1,k}^{i,j}(k), i, j = 1, \dots, d, k = 1, \dots, k_i - p_i) = (\hat{x}_{p,k}^{i,j}(k), i, j = 1, \dots, d, k = 1, \dots, k_i - p_i),$$

which can be seen directly from the definition of \hat{x} . It follows that $k - p$ is the smallest solution of the system $(r, x_{1,k})$. But since $p > 0$, the sequence $x_{1,k}$ is not a good permutation of x , and the inequality (4.20) is proved.

Let $q = (0, \dots, 0, k_m - 1, 0, \dots, 0)$, where $k_m - 1$ is the m -th coordinate of q and set $y := \tilde{x}_{q,k}$, then by applying the same arguments as above to the chain y , we obtain that

$$N_{r,\tilde{y}} \geq N_{r,y} = N_{r,\tilde{x}} \geq N_{r,x},$$

with obvious notations. But by reiterating k_m times this operation, we obtain again the chain x . This shows that equality holds in (4.20), that is $N_{r,\tilde{x}} = N_{r,x}$.

Finally, let $z \in S_d$ be a chain with length k , such that $z^{i,j}(k_i) = x^{i,j}(k_i)$, for all $i, j = 1, \dots, d$, and such that for all $i = 1, \dots, d$, the sequences $z^{i,j}$, $j = 1, \dots, d$ have at most one positive jump that occurs at the same point in $\{0, \dots, k_i\}$. Then it is easy to see that the chain x can be obtained after several cyclical permutations and iterations of the transformation (4.19) applied to z , at any coordinate. Therefore $N_{r,z} = N_{r,x}$. Assume that

there is a good cyclical permutation z' of z . Then (r, z') is a simple system. Therefore, from Lemma 4.4 and Lemma 4.5, $N_{r, z'} = N_{r, z} = N_{r, x} = \det(x^{i,j}(k_i))$. \square

Lemma 4.7. *Let $x \in S_d^r$ with finite length and let \bar{x} be the sequence of \bar{S}_d^r which is associated to x , as in (2.8). Then x and \bar{x} have the same number of good cyclical permutations.*

Proof. Let k and n be the respective lengths of \bar{x} and x . In particular, we have $n_i = \tau_{k_i}^{(i)}$, $i = 1, \dots, d$. Let $q \leq k - 1_d$ be such that $\bar{x}_{q,k}$ is a good cyclical permutation of \bar{x} . Then clearly, there is $p = (p_1, \dots, p_d) \leq n - 1_d$, such that $p_i = \tau_{q_i}^{(i)}$. Set $y = x_{p,n}$ and let us check that

$$(4.21) \quad \bar{y} = \bar{x}_{q,k}.$$

Define $\theta^{(i)}(m) = \min\{v : x_{p_i, n_i}^{i,i}(v) = -m\} = \min\{v : y^{i,i}(v) = -m\}$ and let $\tau_{q_i, k_i}^{(i)}$ be the cyclical permutation of the sequence $\tau^{(i)}$, as defined in (3.13). Then from the construction of y , we can check that

$$\theta^{(i)}(m) = \tau_{q_i, k_i}^{(i)}(m), \quad m \leq k_i,$$

from which we derive (4.21). Moreover, since \bar{y} is a good cyclical permutation of \bar{x} , we deduce from Lemma 2.6 that y is a good cyclical permutation of x .

Conversely, let $q < n - 1_d$ such that $x_{q,n}$ is a good cyclical permutation of x . Then from part (i) of Lemma 2.2, we must have $n_i = \min\{n : x_{q_i, n_i}^{i,i}(n) = -k_i\}$. Therefore, there exists $p < k$, such that $q_i = \tau_{p_i}^{(i)}$. Again, by setting $y = x_{q,n}$, we check that $\bar{y} = \bar{x}_{p,k}$ and we deduce from Lemma 2.6 that \bar{y} is a good cyclical permutation of \bar{x} . \square

Then we end the proof of Lemma 3.3. Let $x \in S_d$ be as in this lemma, that is x has finite length $n = (n_1, \dots, n_d) \in \mathbb{N}^d$ and n is a solution of the system (r, x) , where $r = (r_1, \dots, r_d) \in \mathbb{Z}_+^d$ is such that $r > 0$. Let $k'_i = -\min_{0 \leq n \leq n_i} x_n^{i,i}$ and set $n'_i = \tau_{k'_i}^{(i)}$. Then $n' \leq n$, so that we can define the cyclical permutation $y = x_{n', n}$ of x . By construction, the sequence \bar{y} has length k , where $k_i = -x^{i,i}(n_i)$ and $\bar{y}^{i,j}(k_i) = x^{i,j}(n_i)$, $i, j = 1, \dots, d$. Note that $k_i \geq 1$, from the assumption $x^{i,i}(n_i) \neq 0$. Moreover, k is a solution of the system (r, \bar{y}) . So, thanks to Lemma 4.6, the number of good cyclical permutations of \bar{y} is $\det(x^{i,j}(n_i))$ and from Lemmas 2.6 and 4.7 this is also the number of good cyclical permutations of y , the latter being clearly the number of good cyclical permutations of x . \square

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