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Chapter Number

Defects in Product Line Models and How to Identify Them

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1. Introduction

The history of software and system development shows that abstraction plays a major role in making complexity manageable (Bosch 2000). Thus, abstracting the common and variable artefacts of an undefined collection of products and organising them into a model may be a good option to manage the complexity of a product line. Product line models improve decision-making processes. In addition, the representation of PLMs in different views improves communication of the actors participating in the product line management (Finkelstein et al. 1992). Nuseibeh et al. (1994) describe views as partial representations of a system and its domain.

Several approaches have been found in literature to represent commonality and variability of a product line. Most of the approaches use features (Kang et al. 1990) as the central concept of product line models. However, other modelling approaches exist like Orthogonal Variability Models (OVM, cf. Pohl et al. 2005), Dopler variability models (Dhungana et al. 2010), Textual Variability Language (TVL, cf. Boucher et al. 2010 and Classen et al. 2010), and constraint-based product line language (Djebbi et al. 2007, Mazo et al. 2011e; Salinesi et al. 2010b; 2011).

Quality assurance of PLMs has recently been a prominent topic for researchers and practitioners in the context of product lines. Identification and correction of PLMs defects, is vital for efficient management and exploitation of the product line. Defects that are not identified or not corrected will inevitably spread to the products created from the product line, which can drastically diminish the benefits of the product line approach (Von der Maßen and Lichter 2004, Benavides 2007). Besides, product line modeling is an error-prone activity. Indeed, a product line specification represents not one, but an undefined collection of products that may even fulfil contradictory requirements (Lauenroth et al. 2010). The aforementioned problems enforce the urgent need of early identification and correction of defects in the context of product lines.

Product line models quality has been an intensive research topic over the last ten years (Von der Maßen & Lichter 2004; Zhang et al. 2004; Batory 2005; Czarnecki & Pietroszek 2006; Benavides 2007; Janota & Kiniry 2007; Lauenroth & Pohl 2007; Trinidad et al. 2008; Van den Broek & Galvão 2009; Elfaki et al. 2009; Kim et al. 2011; Liu et al. 2011). Usually, to guaranty a certain level of quality of a model, this one must be verified against a collection of criteria

1 and then, these defects must be corrected. Verifying PLMs entails finding undesirable
2 properties, such as redundancies, anomalies or inconsistencies (Von der Maßen *et al.* 2004).
3 It is widely accepted that manual verification is already tedious and error-prone (Benavides
4 *et al.* 2005). This is even worst when several (often millions) of products are represented
5 altogether in a single specification. Several approaches to automate verification of PLMs
6 have been proposed in order to overcome this limitation. However, despite the relative
7 success of these approaches, there is still a number of pending issues that have motivated
8 the proposal developed in this chapter:

- 9 1. Quality assurance techniques from the development of single systems cannot be
10 directly applied to product line specifications because these specifications contain
11 variability. As shows the example presented by Lauenroth *et al.* (2010), a product line
12 may contain requirements R and $\neg R$ at the same time. When a traditional technique is
13 used for verifying this specification, even though those requirements are not included
14 for the same product, a contradiction would be identified since the requirements R and
15 $\neg R$ cannot be fulfilled together. Therefore, it is necessary to take into account the
16 variability of the product line to check whether contradictory requirements can really
17 be part of the same product.
- 18 2. The current state of the art on verification is mainly focused on feature models (Kang *et al.*
19 1990). Only properties that can be evaluated over feature models represented as
20 boolean expressions are considered in these works. This brushes aside the non-boolean
21 elements of the more sophisticated product line specification formalisms (e.g., integer
22 cardinalities, attributes and complex constraints; cf. Mazo *et al.* 2011d, Salinesi *et al.*
23 2010b, 2011). Current approaches restrict the verification operations to those that can be
24 solved by boolean solvers. The verification is guided by the pre-selected technology and
25 not by the verification requirements themselves. As a result, verification techniques are
26 designed for a limited number of formalisms. These verification techniques are
27 inadequate for many of the existing formalisms, included some used in an industrial
28 context (Djebbi *et al.* 2007; Dhungana *et al.* 2010).
- 29 3. Inadequate support for multi-model specification. The size and complexity of industrial
30 product line models motivates the development of this one by heterogeneous teams
31 (Dhungana *et al.* 2006; Segura 2008). Nevertheless, existing tools provide only little
32 support for integrating the models developed by different teams and the subsequent
33 verification of the global model and configurations of products from that model. For
34 instance, a global model that integrates two models must itself have no defects resulting
35 from the integration.

36 Also in the context of PLs specified with several models, we have identified in our literature
37 review a weak support for verifying the global view of the product line. A product line
38 model has to change over time and in multi-model PLs a change on one of the models can
39 make the global view inconsistent. To the best of our knowledge, existing tools do not
40 provide automated mechanisms for detecting errors on the global PLM as a result of the
41 changes in the different models of the PLM.

42 This chapter addresses the fourth problem situations aforementioned. To tackle these
43 situations, we present in Section 2 the most relevant concepts used in this chapter, a
44 literature review of related works and the running example to be used in the rest of the
45 chapter. Section 3 presents our typology of verification criteria, which is developed in

1 Section 4 for the case of single-view product line models, and Section 5 for the case of multi-
2 view product line models. Section 6 presents the evaluation of the approach presented in
3 this chapter.

4 **2. Background and running example**

5 This section presents a literature review on verification of product line models and the
6 corresponding analysis regarding the gaps and challenges identified in each approach. This
7 section also presents a UNIX product line and the corresponding model of the whole or a
8 part of the PL in three different PL modeling languages. The UNIX PL will be used in the
9 rest of this chapter as our running example.

10 **2.1 Verification of product line models**

11 Verifying PLMs entails several aspects. On the one hand, a product line model,
12 independently of the language used to express it, must respect certain properties associated
13 with the domain of product lines. On the other hand, certain properties are associated with
14 the concepts used in the language in which it is expressed. Therefore, some properties of
15 PLMs are independent of the language while other ones are particular to each language.
16 Thus, product line models can be verified from two different points of view. This chapter
17 proposes an approach for PLM verification (Von der Maßen & Lichter 2004; Lauenroth &
18 Pohl 2007; Mendonça *et al.* 2009) in which the engineer selects the verification operations that
19 he/she want to use according to the language in which the model(s) to be verified are
20 specified. In this approach, verification consists in “finding undesirable properties, such as
21 redundant or contradictory information” (Trinidad *et al.* 2008). For instance, PLMs should
22 not be void (i.e., they should allow to configure more than one product) and for the
23 languages with the concept of *optionality*, elements modeled as optional must be really
24 optional (i.e., they should not appear in all the products configured from the PLM).

25 **2.2 Related work**

26 Von der Maßen & Lichter (2004) present an approach to identify redundancies, anomalies
27 and inconsistencies. According to the authors, a feature model contains redundancy, “if at
28 least one semantic information is modeled in a multiple way”; anomalies, “if potential
29 configurations are being lost, though these configurations should be possible”; and
30 inconsistencies, “if the model includes contradictory information”. Several cases of
31 redundancies, anomalies and inconsistencies on FMs are identified. In order to validate the
32 approach, the authors use RequiLine, a tool that allows detecting inconsistencies on the
33 domain and on the product configuration level (Von der Maßen & Lichter 2003). The
34 approach was evaluated in “a small local software company” and “in a global player of the
35 automotive industry”. However no information about the automating detection of
36 redundancies and anomalies, no details about the sizes of the models or about the
37 technology used to automate the approach or about the results obtained were provided.

38 Whereas Batory (2005) used grammar and propositional formulas to represent basic FMs
39 and enable truth maintenance systems and SAT solvers to identify contradictory (or
40 inconsistency) predicates to verify that a given combination of features effectively defines a
41 product. In the same line as Batory, Hemakumar (2008) proposed a dynamic solution to find

1 contradictions, “where errors can be detected during usage and silently reported to model
2 designers”. The author proposes an incremental consistency algorithm that incrementally
3 verifies some contradiction properties. The approach consists in verify that a model is
4 contradiction-free if it is k -contradiction free for all k where $0 < k \leq n$ (A feature model is k -
5 contradiction free if every selection of k features does not expose a contradiction, for
6 example: “unconditionally” dead features are exposed when $k=1$). When $k=n$, where n is the
7 number of user selectable features, the model has been proven to be contradiction free.
8 However, the incremental consistency algorithm has important practical limits because it is
9 limited to “verify contradiction freedom of models with about 20 or fewer features”.

10 In (Benavides *et al.* 2005a; 2005b; 2006; 2007; Trinidad *et al.* 2008), authors transform FODA
11 models with and without attributes into Boolean expressions. These expressions are
12 executed on Constraint Satisfaction Problem (CSP), Satisfiability (SAT) and Binary Decision
13 Diagrams (BDD) solvers in order to execute analysis and verification operations over feature
14 models. In (Benavides *et al.* 2006) the relationships of the FM are represented as *ifThenElse*
15 constrains on CPS. Despite the originality of this proposal, the constraint representing a
16 feature cardinality (m,n) between the father feature A and its child B (according to their
17 notation: *ifThenElse*($A=0;B=0;B$ in $\{n,m\}$)) does not consider that the feature A can itself have
18 a feature cardinality, and in this case the semantic of feature cardinalities is not well
19 represented in the constraint. Authors performed a comparative test between two off-the-
20 shelf CSP Java solvers (JaCoP and Choco). The experiment was executed on five FMs with
21 up to 52 features and in both solvers. The time to get one solution seemed to be linear and
22 the time to get all solutions seemed to be exponential.

23 Janota & Kiniry (2007) have formalized in higher-order logic (HOL) a “feature model meta-
24 model” that integrates properties of several feature modeling approaches such as attributes
25 and cardinalities. Once the model represented in HOL, author have formulated HOL
26 expressions for root selectivity, existence of a path of selected features from the root to a
27 feature that has been selected, and cardinality satisfaction of a selected feature that each
28 feature model must respect. The approach has been implemented in *Mobius program*
29 *verification environment*, an Eclipse-based platform for designing, testing, performing various
30 kinds of verification of Java programs and bytecode. Nevertheless, the paper does not
31 provide evidence about the evaluation of the approach, its scalability and its applicability to
32 real cases.

33 Trinidad *et al.* (2008) mapped FMs into CSP in order to find and diagnose three types of
34 errors: (i) “dead features” are non-selectable features (features that do not appear in any
35 product); (ii) “false optional features”, which are features that in spite of being modeled as
36 optional, are always chosen whenever their parents are chosen; and (iii) “void models”, a
37 feature model is said to be void if no product can be defined from it. The goal of Trinidad *et*
38 *al.* is to detect the above three errors and provide explanations for the cause of these errors.
39 In order to achieve the first goal, authors transform the FM into a CSP expression and then,
40 to query the Choco solver (by means of the FaMa tool) to find the errors. The approach has
41 been evaluated on five FMs up to 86 features. Unfortunately, no details about the scalability
42 and the efficiency of the approach and tool are provided.

43 Van der Storm (2007) transformed feature diagrams into BDDs in order to check
44 configurations, obtain valid configurations and check consistency of the feature diagram.

1 Checking the consistency of the feature diagram consists in checking the satisfiability for the
2 BDD logical formulas. Unfortunately, neither details about implementation nor performance
3 nor scalability of the approach are provided in the paper.

4 Yan et al. (2009) proposed an approach that consists in eliminating verification-irrelevant
5 features and constraints from FMs in order to reduce the problem size of verification, and
6 alleviate the state-space explosion problem. The authors carried out an experiment in which
7 they generated FMs with up to 1900 features. The authors verified the consistency of models
8 and showed that verification is faster when the redundant features had been eliminated. The
9 problem with this approach is that it only considers as redundant, the constraints that
10 contain redundant features, whereas it does not consider typical redundancies such as
11 domain overlapping or cyclic relationships (Salinesi *et al.* 2010; Mazo *et al.* 2011). Besides, (i)
12 the validation of the approach was done with in-house and random build features models,
13 which does not guaranty that the approach works with real world feature models; and (ii)
14 the details about the formalisation and implementation of the approach are not revealed.

15 Van den Broek & Galvão (2009) analyze FODA product line models using generalized
16 feature trees. In their approach they translate FMs into feature trees plus additional
17 constraints. Once FMs represented in the functional programming language Miranda, they
18 detect the existence of products (void models), dead features and minimal set of conflicting
19 constraints. In FMs with cross-tree constraints, the function to find the number of products
20 belongs to $O(N*2^M)$, where N is the number of features and M is the number of cross-tree
21 constraints. Unfortunately, no evaluation of the theoretical calculations of efficiency is
22 reported in the paper. The approach was validated with a feature tree of 13 features and two
23 cross-tree constraints, which is not enough to evaluate the scalability and the usability of the
24 approach on industrial models.

25 Elfaki *et al.* (2009) propose to use FOL to detect dead features and inconsistencies due to
26 contradictions between include-type and exclude-type relationships in FMs. The innovative
27 point of their work is the suggestion of expressions dealing with both individuals and sets of
28 features.

29 SPLOT (Mendonca *et al.* 2009b) is a Web-based reasoning and configuration system for
30 feature models supporting group-cardinalities instead of alternative and or-relations. The
31 system maps feature models into propositional logic formulas and uses boolean-based
32 techniques such as BDD and SAT solvers to verify the validity of models (not void) and find
33 dead features.

34 2.3 Running example

35 The example taken in this chapter is that of the UNIX operating system, initially presented
36 in (Mazo *et al.* 2011d). UNIX was first developed in the 1960s, and has been under constant
37 development ever since. As other operating systems, it is a suite of programs that makes
38 computers work. In particular, UNIX is a stable, multi-user and multi-tasking system for
39 many different types of computing devices such as servers, desktops, laptops, down to
40 embedded calculators, routers, or even mobile phones. There are many different versions of
41 UNIX, although they share common similarities. The most popular varieties of UNIX are
42 Sun Solaris, Berkeley (BSD), GNU/Linux, and MacOS X.

1 The UNIX operating system is made up of three parts: the kernel, the shell and the
2 programs; and two constituent elements: files and processes. Thus, these three parts consist
3 in a collection of files and processes allowing interaction among the parts. The kernel of
4 UNIX is the hub of the operating system: it allocates time and memory to programs and
5 handles the file-store and communications in response to system calls. The shell acts as an
6 interface between the user and the kernel, interprets the commands (programs) typed in by
7 users and arranges for them to be carried out. As an illustration of the way the shell, the
8 programs and the kernel work together, suppose a user types *rm myfile* (which has the effect
9 of removing the file *myfile*). The shell searches the file-store for the file containing the
10 program *rm*, and then requests the kernel, through system calls, to execute the program *rm*
11 on *myfile*. The process *rm* removes *myfile* using a specific system-call. When the process *rm*
12 *myfile* has finished running, the shell gives the user the possibility to execute further
13 commands.

14 As for any product line, our example emphasizes the common and variable elements of the
15 UNIX family and the constraints among these elements. This example is built from our
16 experience with UNIX operating systems and it does not pretend to be exhaustive, neither
17 on the constituent elements nor on the constraints among these elements. The idea with this
18 PL is, for instance, to look at what utility programs or what kinds of interfaces are available
19 for a particular user. This PL is composed of the following six constraints:

20 Constraint 1. UNIX can be installed or not and the installation can be from a CDROM, a
21 USB device or from the NET.

22 Constraint 2. UNIX provides several hundred UTILITY PROGRAMS for each user. The
23 collection of UTILITY PROGRAMS varies even when the UNIX product is full-
24 configured.

25 Constraint 3. The SHELL is a kind of UTILITY PROGRAM. Different USERS may use
26 different SHELLS. Initially, each USER has a default shell, which can be overridden or
27 changed by users. Some common SHELLS are:

- 28 • Bourne shell (SH)
- 29 • TC Shell (TCSH)
- 30 • Bourne Again Shell (BASH)

31 For the sake of simplicity will consider only two users in this running example:
32 ROOT_USER and GUEST_USER.

33 Constraint 4. Some functions accomplished by the UTILITY PROGRAMS are:

- 34 • EDITING (mandatory and requires USER INTERFACE)
- 35 • FILE MAINTENANCE (mandatory and requires USER INTERFACE)
- 36 • PROGRAMMING SUPPORT (optional and requires USER INTERFACE)
- 37 • ONLINE INFO (optional and requires USER INTERFACE)

38 Constraint 5. The USER INTERFACE can be GRAPHICAL and/or TEXTUAL.

39 Constraint 6. The GRAPHICAL interface is characterized by a WIDTH RESOLUTION
40 and a HEIGHT RESOLUTION that can have the following couples of values [800,600],
41 [1024,768] and [1366,768].

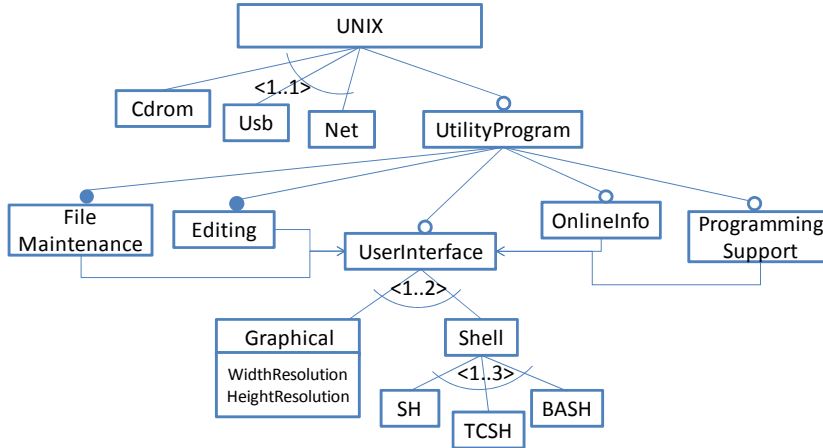
2.3.1 Representation of the UNIX product line as a feature model

Feature Models (FMs) were first introduced in 1990 as a part of the *Feature-Oriented Domain Analysis (FODA)* method (Kang *et al.* 2002) as a means to represent the commonalities and variabilities of PLs. Since then, feature modeling has become a de facto standard adopted by the software product line community and several extensions have been proposed to improve and enrich their expressiveness. A FM defines the valid combinations of features in a PL, and is depicted as a graph-like structure in which nodes represent features, and edges the relationships between them (Kang *et al.* 2002). Two of these extensions are cardinalities (Riebisch *et al.* 2002; Czarnecki *et al.* 2005) and attributes (Streitferdt *et al.* 2003; White *et al.* 2009). Although there is no consensus on a notation to define attributes, most proposals agree that an attribute is a variable with a name, a domain and a value. Attributes are integers, enumerations, and boolean values representing important properties of a feature; as for instance the price, the cost, the width, the height or the time spent to build the corresponding feature. In this chapter we use the group cardinalities grouping bundles of features (cf. *Cdrom*, *Usb* and *Net* in Figure 1). We use the semantic of feature models proposed by (Schobbens *et al.* 2007).

The elements of the FM notation used in this chapter are presented and exemplified as follows:

- *Feature*: A feature is a prominent or distinctive user-visible aspect, quality, or characteristic of a software system (Kang *et al.* 1990). For the sake of simplicity FMs usually comport only the name of the feature; for instance *Editing* in Figure 1. Every FM must have one root, which is called *root feature* and identifies the product line; for example *UNIX* in Figure 1.
- *Attribute*: Although there is no consensus on a notation to define attributes, most proposals agree that an attribute is a variable with a name (*Name*), a domain (*Domain*), and a value (consistent with the domain) at a given configuration time. From a technical point of view an attribute must to be differentiated from the other ones by an identifier (*IdAttribute*). For instance in Figure 1, *WidthResolution* and *HeightResolution* are two attributes with a domain determined by the constraint at the bottom of the model.
- *Mandatory*: Given two features *F1* and *F2*, *F1* father of *F2*, a mandatory relationship between *F1* and *F2* means that if the *F1* is selected, then *F2* must be selected too and vice versa. For instance in Figure 1, features *UtilityProgram* and *Editing* are related by a mandatory relationship.
- *Optional*: Given two features *F1* and *F2*, *F1* father of *F2*, an optional relationship between *F1* and *F2* means that if *F1* is selected then *F2* can be selected or not. However, if *F2* is selected, then *F1* must also be selected. For instance in Figure 1, features *UNIX* and *UtilityProgram* are related by an optional relationship.
- *Requires*: Given two features *F1* and *F2*, *F1* requires *F2* means that if *F1* is selected in product, then *F2* has to be selected too. Additionally, it means that *F2* can be selected even when *F1* is not. For instance, *Editing* requires *UserInterface* (cf. Figure 1).
- *Exclusion*: Given two features *F1* and *F2*, *F1* excludes *F2* means that if *F1* is selected then *F2* cannot to be selected in the same product. This relationship is bi-directional: if *F2* is selected, then *F1* cannot to be selected in the same product.
- *Group cardinality*: A group cardinality is an interval denoted $\langle n..m \rangle$, with *n* as lower bound and *m* as upper bound limiting the number of child features that can be part of a

1 product when its parent feature is selected. If one of the child features is selected, then
 2 the father feature must be selected too. For instance in Figure 1, *Cdrom*, *Usb* and *Net* are
 3 related in a $\langle 1..1 \rangle$ group cardinality.



4 Graphical \rightarrow relation([(WidthResolution, HeightResolution), {[800, 600], [1024, 768], [1366, 768]}])

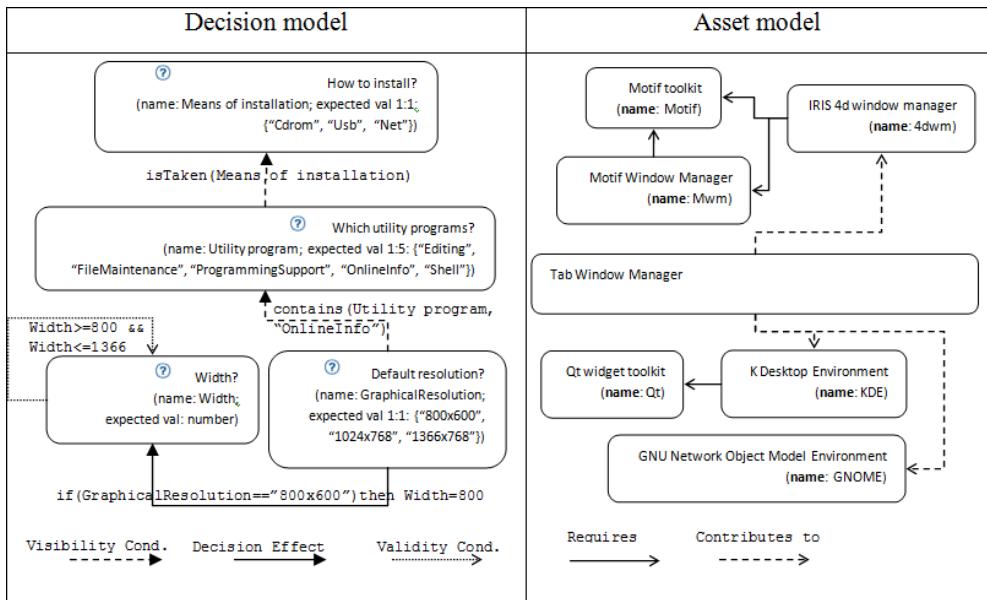
5 Fig. 1. User model of the UNIX operating system family of our running example

6 Figure 1 corresponds to the feature representation of the user model of our running
 7 example. In this model, a user has the possibility to install a UNIX system using one of the
 8 following options: a CD ROM, an USB devise or a network. In addition, users have the
 9 possibility to install or not utility programs for file maintenance, edition, online access, and
 10 user interface. The user interface may be graphical or command-line (Shell) based; there are
 11 three options of command-line interface: SH, TCSH and BASH. The utility programs for
 12 user interface, online information and programming support are optional features.

13 2.3.2 Representation of the UNIX product line as a dopler variability model

14 The Decision-oriented (Dopler) variability modeling language focuses on product derivation
 15 and aims at supporting users configuring products. In Dopler variability models (Dhungana
 16 *et al.* 2010a; 2010b), the product line's problem space is defined using *decision models* whereas
 17 the solution space is specified using *asset models*. An example of Dopler model is presented
 18 in Figure 2. This figure depicts the installation of a UNIX operating system (decision model)
 19 and the associated packages (asset model) that can be selected if the UNIX system is
 20 installed with a graphical interface. The decision model is composed of four decisions. The
 21 first one proposes one of three ways to install a UNIX operating system (with a CD ROM,
 22 with a USB or with the Net). The solution of this decision implies the solution of a second
 23 decision in which the user must select the utility programs to be installed in the particular
 24 UNIX system; in that regard, five utility programs are proposed: one tool for editing, one for
 25 file maintenance, one for programming, one for online information access and one shell. If
 26 the choice contains the utility program for online information, the user must decide what
 27 kind of graphical resolution will be configured and several choices are proposed: *800x600*,

1 1024x768, 1366x768. Depending of each selection, the values of the variables corresponding
 2 to the width and height resolution will be assigned automatically by means of several
 3 decision effects; for instance in Figure 2: *if(GraphicalResolution==800x600) then Width=800*. To
 4 finish, the assignation of the width and height resolution must respect a certain number of
 5 validity conditions like for instance: *Width ≥ 800* and *Width ≤ 1366*. The asset model is
 6 composed of seven graphical user interfaces and libraries that can be used in a UNIX
 7 graphical interface. The *Tab Window Manager* asset is available for all UNIX implementations
 8 with a graphical interface and requires the asset *Motif*; the others assets are optional. The
 9 IRIS 4d window manager is based on *Mwm* and *Motif* and therefore requires all of them in
 10 order to work in the same way as the *KDE* asset requires the *Qt widget toolkit* to work.



11
 12 Fig. 2. Example of Dopler Model: Installation of a UNIX System

13 A **decision model** consists of a set of decisions (e.g., *Which utility programs?* with two
 14 attributes: **name** and **expected values**) and dependencies among them (i.e., the **Visibility**
 15 **condition** *isTaken(Means of installation)* forcing the answer of the decision *Utility program* if
 16 the decision *Means of installation* is taken). **Assets** allow defining an abstract view of the
 17 solution space to the degree of details needed for subsequent product derivation. In a
 18 domain-specific metamodel attributes and dependencies can be defined for the different
 19 types of assets. Decisions and assets are linked with **inclusion conditions** defining
 20 traceability from the solution space to the problem space (e.g., the asset *Tab Window Manager*
 21 must be included in the solution space if the option *OnlineInfo* of the decision *Utility program*
 22 is selected in a particular configuration). In our integration approach, these inclusion
 23 conditions are constraints that will be added to the collection of constraints representing the
 24 decision and asset model. Once these constraints are added, both viewpoints of the PL are
 25 integrated, and the model is ready to be verified against the typology of verification criteria
 26 presented in this chapter.

2.3.3 Representation of the UNIX product line as a constraint program

Constraint Programming (CP) emerged in the 1990's as a successful paradigm to tackle complex combinatorial problems in a declarative manner (Van Hentenryck 1989). CP extends programming languages with the ability to deal with logical variables of different domains (e.g. integers, reals or booleans) and specific declarative relations between these variables called *constraints*. These constraints are solved by specialized algorithms, adapted to their specific domains and therefore much more efficient than generic logic-based engines. A constraint is a logical relationship among several variables, each one taking a value in a given domain of possible values. A constraint thus restricts the possible values that variables can take.

In modern Constraint Programming languages (Diaz & Codognet 2001; Schulte & Stuckey 2008), many different types of constraints exist and are used to represent real-life problems: arithmetic constraints such as $X + Y < Z$, symbolic constraints like $atmost(N, [X1, X2, X3], V)$ which means that at most N variables among $[X1, X2, X3]$ can take the value V , global constraints like $alldifferent(X1, X2, \dots, Xn)$ meaning that all variables should have different values, and reified constraints that allow the user to reason about the truth-value of a constraint. Solving constraints consists in first reducing the variable domains by propagation techniques that will eliminate inconsistent value within domains and then finding values for each constrained variable in a labeling phase, that is, iteratively grounding variables (fixing a value for a variable) and propagating its effect onto other variable domains (by applying again the same propagation-based techniques). The labeling phase can be improved by using heuristics concerning the order in which variables are considered as well as the order in which values are tried in the variable domains. Consult (Schulte & Stuckey 2008) for more details. Mazo et al. (2011e) present a constraint system to represent product line models by means of abstract constraints where the domain is an argument of the system.

Our running example can also be represented as a constraint program according to the method proposed by Salinesi et al. (2010; 2011) and Mazo et al. (2011d). The resulting model is presented in the following table, where the first column corresponds to each constraint of our example and the second column its representation as a constraint program.

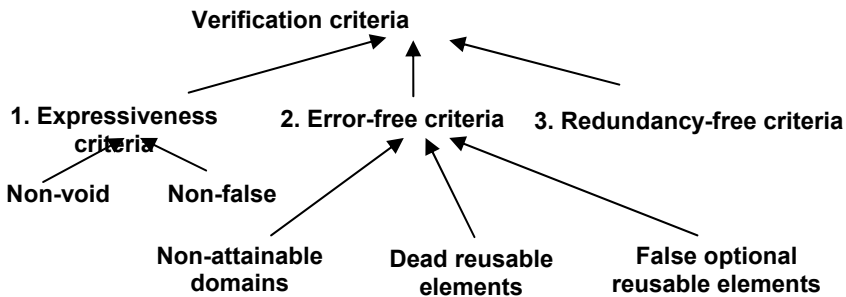
Constraint	CP Representation
C. 1	$UNIX \leq Cdrom + Usb + Net \leq UNIX$
C. 2	$UtilityProgram \leq UNIX$
C. 3	$Shell = UtilityProgram \wedge$ $Shell \Rightarrow ((1 * ROOT_USER \leq ROOT_USERSH + ROOT_USERTCSH +$ $ROOT_USERBASH \leq 3 * ROOT_USER) \wedge (1 * GUEST_USER \leq$ $GUEST_USERSH + GUEST_USERTCSH + GUEST_USERBASH \leq 3 *$ $GUEST_USER))$
C. 4	$Editing = UtilityProgram \wedge$ $Editing \Rightarrow UserInterface \wedge$ $FileMaintenance = UtilityProgram \wedge$ $FileMaintenance \Rightarrow UserInterface \wedge$ $ProgrammingSupport \leq UtilityProgram \wedge$ $ProgrammingSupport \Rightarrow UserInterface \wedge$

Constraint	CP Representation
	$\text{OnlineInfo} \leq \text{UtilityProgram} \wedge$ $\text{OnlineInfo} \Rightarrow \text{UserInterface} \wedge$ $\text{UserInterface} \leq \text{UtilityProgram}$
C. 5	$1 * \text{UserInterface} \leq \text{Graphical} + \text{Textual} \leq 2 * \text{UserInterface}$
C. 6	$\text{Graphical} = 1 \Leftrightarrow (\text{WidthResolution} = W1 \wedge \text{HeightResolution} = H1) \wedge$ $\text{Graphical} = 0 \Leftrightarrow (\text{WidthResolution} = 0 \wedge \text{HeightResolution} = 0) \wedge$ $\text{relation}([W1, H1], [[800, 600], [1024, 768], [1366, 768]])$

1 Table 1. UNIX PL represented as a constraint program

2 **3. Typology of verification criteria**

3 Verifying PLMs entails several aspects. On the one hand, a product line model,
 4 independently of the language used to express it, must respect certain properties associated
 5 with the domain of product lines. On the other hand, certain properties are associated with
 6 the fact that each PLM respects the syntactic rules of the language in which it is expressed.
 7 Therefore, some properties of PLMs are independent of the language while other ones are
 8 particular to each language. In light of this observation, this chapter proposes a typology of
 9 PLM verification criteria adapted from the initial version presented in (Salinesi et al. 2010a).
 10 The typology presented in Figure 1 is structure in two levels; the top level represents the
 11 three categories of verification criteria and the bottom level represents the corresponding
 12 operations of the two criteria with more than one operation. This figure indicates that not all
 13 PLM verification criteria are equivalent: some are a result of the specification of the PL with
 14 a metamodel, whereas others can be used to verify PL specifications independent of the
 15 formalism used when they were specified. Besides, some criteria help verifying the ability of
 16 PLM to generate all the desired products and only them, whereas others are interested in the
 17 quality of PLMs, independently of their semantics (i.e., the collection of all possible products
 18 that can be generated from it). This is for example the case with the respect of certain rules
 19 providing formality (i.e., absence of ambiguity) at the PLM.



20
 21 Fig. 3. Typology of verification criteria on PLMs

22 The outcomes of the typology are multiple:

- 1 a. the typology classify the criteria semantic, allowing the identification of similarities and
2 differences among the criteria;
- 3 b. the typology makes easier to identify some defects for which no verification criterion is
4 available in the literature. Redundancy of relationships among reusable elements is an
5 example of defect for which no verification criterion has been defined in the literature
6 (at least to our knowledge).
- 7 c. the classification behind the typology makes it easier the proposition of a standard and
8 reusable approach to verify the domain-specific criteria of PLMs; and
- 9 d. the typology can be used to select the criteria that one wants to use to verify a PLM
10 according to the impact that these criteria have or the expected level of quality of a
11 particular PLM.

12 The following sections use the typology of verification criteria presented in Figure 3 to
13 develop the verification approach proposed in this chapter.

14 4. Single-model verification

15 In order to verify models against the verification criteria identified and calcified in the
16 former section, it is necessary to represent PLMs in a way that is (a) expressiveness-enough
17 to represent the semantics (i.e. the collection of products that can be configured from the
18 PLM) of PLMs, (b) consistent with the formalization of the criteria, and (c) easy to parse
19 with analysis tools. Experience shows that the semantic of every PLM can be represented as
20 a collection of variables over different domains and constrains among these variables. While
21 the variables specify what can vary from a configuration to another one, constraints express
22 under the form of restrictions what combinations of values are allowed in the products.

23 This section will show how to represent the semantic of PLMs with a constraint based
24 approach, and to verify each and every criterion shown in the typology of the former section
25 on a PLM. The approach will be applied to our feature model example to show how to
26 navigate between the generic specifications of the criteria. The genericity of the approach
27 will be shown by providing examples with other formalisms (cf. Section 5).

28 Verifying PLMs is about looking for undesirable properties such as redundant or
29 contradictory information. This chapter proposes three domain-specific verification criteria:
30 expressiveness, error-free and redundancy-free. Each domain-specific verification criterion
31 is defined, formalized and exemplified with our running example (cf. Figure 1 and Table 1)
32 as follows.

33 **2.1. Expressiveness:** every PLM must allow configuring more than one product, i.e., the
34 model must be not void and the model must be expressive enough to allow configure
35 more than one product (Benavides *et al.* 2005). In case the PLM allows configuring only
36 one product, the PLM, even if it is not considered as a void model, is not expressive
37 enough to be a PLM. Indeed, the purpose of PLMs is to represent at least two products
38 -or there is not reuse. Two verification operations can be used to implement this
39 criterion:

- 40 a. **Non-void PLMs.** This operation takes a PLM as input and returns “Void PLM” if the
41 PLM does not define any products. Two alternative techniques have been proposed so
42 far to implement this operation: calculate the number of products (Van den Broek &

Galvão 2009) or ask for a product configuration that meets the constraints of a FM (Benavides *et al.* 2005; Trinidad *et al.* 2008). Our proposal follows along the lines of the latter alternative and is formalized in the following algorithm. It consists in determining if there is at least one product that can be generated by means of a query to an off-the-shelf solver. If the PLM is not void, the solver will return one valid product or *false* otherwise.

```

7 Non-void_PLM(PLM M, Solver S) {
8     S.charge(M);
9     Answer = S.getOneSolution();
10    If (Answer ≠ "false") {
11        Write (Answer);
12    }
13    Else {
14        Write ("Void PLM");
15    }
16 }

```

The execution of this algorithm over the running example gives as result that our UNIX PL is non-void.

- b. **Non-*false* PLMs.** This operation takes a PLM as input and returns "*False PLM*" if at most one valid product can be configured with it. Although this operation could also help detect when PLMs are void (our precedent operation), the converse is not true. The two operations have then a separate implementation. Our approach consists in asking the solver to generate two products in order to decide if the PLM is false. The algorithm proposed to automate this operation is as follows:

```

25 Non-false_PLM(PLM M, Solver S) {
26     S.charge(M);
27     Answer1 = S.getOneSolution();
28     If (Answer1 ≠ "false") {
29         Answer2 = S.getNextSolution();
30         If (Answer2 ≠ "false") {
31             Write (Answer1, Answer2);
32         }
33         Else {
34             Write ("False PLM");
35         }
36     }
37     Else {
38         Write ("False PLM");
39     }
40 }

```

The execution of this algorithm over the running example gives as result that our UNIX PL is a non-*false* PLM.

- 2.2. **Error-free.** The Dictionary of Computing defines an error as "A discrepancy between a computed, observed, or measured value or condition, and the true, specified, or theoretically correct value or condition" (Howe 2010). In PLMs, an error represents a

discrepancy between what the engineer want to represent and the result obtained from the model. For instance, this is the case when the engineer includes a new reusable element (in a given domain) in a PLM, but this element never appears in a product. The error-free criterion can be verified by means of three operations: the first one allows identifying the non-attainable domain values of PLM's reusable elements; the second one allows identifying the dead elements, i.e. elements of the PL that are never used in a product; the third one allows identifying the reusable elements modeled as optional but that appear in all the products of the PL. These operations are presented as follows:

- c. **Non-attainable domains:** This operation takes a PLM and a collection of reusable elements as input (all of them by default) and returns the reusable elements that cannot attain one of the values of their domain. Reusable elements can have domains represented as particular values (e.g., 800), intervals of values (e.g., [0..5]), or collections of values (e.g., {0, 800, 1024, 1280}). A non-attainable value of a domain is the value of an element that never appears in any product of the product line. For example, if a reusable element R has the domain [0..1], value 1 is non-attainable if R can never be integrated in a product line it never take the value of 1. Non-attainable values are clearly undesired since they give the user a wrong idea about domain of reusable elements. The approach presented in this chapter can assess the attainability of any reusable elements for all (or parts of) their domain values. This operation was also implemented by Trinidad *et al.* (2008), but only for boolean domains on FMs. Our proposal goes a step further by offering an algorithm for any domain as e.g. needed when using attributes or features whit individual cardinality. Our algorithm to automate this operation evaluates the domain of each variables of the PLM. For each vale of the domain, the algorithm requests the solver at hand for a solution. If the solver gives a solution for all the values of the variable's domain, the variable is erased from the list of reusable elements with non-attainable domains. Otherwise, the variable, representing a reusable element, is affected with the non-attainable value(s) and kept in the list of reusable elements with non-attainable domains. In each product obtained from the solver, all the variables of the PLM are affected with a particular value of the corresponding domain. Thus, this algorithm takes advantage of that fact and records the answers given by the solver in order to avoid achieving useless requests testing the attainability of domain values that have already been obtained in precedent tests. The corresponding algorithm is as follows:

```

NonAttainableDomains(PLM M, Solver S) {
    S.charge(M);
    For (each variable V ∈ M) {
        For(each Di ∈ domain of V AND not in {PrecedentProducts}){
            Product = S.getOneSolution("V = Di");
            If (Product = "false") {
                Write ("The domain " + Di + " of " + V + " is non-
attainable");
            }
            Else {
                PrecedentProducts += Product;
            }
        }
    }
}

```

```

1      }
2    }
3  }

```

For instance in our running example, if when asking for a product with *WidthResolution=800* we get a product

```

6 P1 = [UNIX=1, Cdrom=1, Usb=0, Net=0, UtilityProgram=1, FileMaintenance=1, Editing=1,
7 UserInterface=1, Graphical=1, WidthResolution=800, HeightResolution=600, Shell=1, SH=1,
8 TCSH=1, BASH=0, OnlineInfor=0, ProgrammingSupport=0].

```

This means both that *WidthResolution* can attain the value of 800, and that the rest of variables can attain the values assigned by the solver. Thus, for instance, it is not necessary to ask if the variable *UNIX* can attain the value of 1 or if *HeightResolution* can attain the value of 600.

- d. **Dead-free reusable elements:** A reusable element is dead if it cannot appear in any product of the product line. This operation takes as input a PLM and a collection of reusable elements, and it returns the set of dead reusable elements, or false if there is none in the input list. Reusable elements can be dead because: (i) they are excluded by an element that appears in all products (also known as full-mandatory or core reusable elements, c.f. Von der Maßen & Lichter 2004; Benavides *et al.* 2005; Trinidad *et al.* 2008; Van den Broek & Galvão 2009); and (ii) they are wrongly constrained (e.g., an attribute of the feature is > 5 and < 3 at the same time, or a group cardinality is wrong defined). Elfaki *et al.* (2009) detect dead features by searching only for predefined cases, i.e. defined dead features in the domain-engineering process. Trinidad *et al.* (2006, 2008) detect dead features by finding all products and then searching for unused features. Van den Broek and Galvão (2009) detect dead features by transforming the FM into a generalized feature tree, and then searching the feature occurrences that cannot be true. To the better of our knowledge there is not details in literature about the way in which the other references have implemented this operation. Our approach evaluates each non-zero value of each reusable element's domain, and reuses each solution obtained from the solver in order to avoid useless computations. If a reusable element cannot attain any of its non-zero values, then the reusable element is dead. The reuse of the solutions previously obtained makes our dead artefacts detection technique scalable as showed below, by contrasts to the state of the art. The corresponding algorithm is presented as follows:

```

34 DeadReusableElements(PLM M, Solver S) {
35     S.charge(M);
36     DeadElementsList = all variables of M;
37     For (each variable V ∈ DeadElementsList) {
38         Product = S.getOneSolution("V > 0");
39         If (Product = "false") {
40             Write ("The variable " + V + " is dead");
41         }
42         Else {
43             Erase V and all the other non-zero variables obtained in Product from
44             DeadElementsList;
45         }
46     }
47 }

```


Our algorithm first creates a list of the reusable elements whose dead or non-dead condition is yet to be assessed. For example:

```
deadElements=[UNIX, Cdrom, Usb, Net, UtilityProgram, FileMaintenance, Editing,
UserInterface, Graphical, WidthResolution, HeightResolution, Shell, SH, TCSH, BASH,
OnlineInfor, ProgrammingSupport].
```

Then, our algorithm queries for a configuration based on reusable elements for which we still ignore if they are dead or not, and sieves the selected (and thus alive) elements from this list. For example, to know if *UtilityProgram* is dead or not, it is sufficient to query the solver for a product with *UtilityProgram=1*, which provides a product

```
P1 = [UNIX=1, Cdrom=1, Usb=0, Net=0, UtilityProgram=1, FileMaintenance=1, Editing=1,
UserInterface=1, Graphical=1, WidthResolution=800, HeightResolution=600, Shell=1, SH=1,
TCSH=1, BASH=0, OnlineInfor=0, ProgrammingSupport=0].
```

This means not only that the reusable element *UtilityProgram* is not dead, but also that the other elements with values different from 0 are not dead. Therefore these elements can be sieved from the list of dead elements. The test can be repeated until all elements are sieved. For example querying for products with *Usb =1*, the solver provides another product which means that this reusable element is not dead either. According to our algorithm, the variable *Usb*, and all the other non-zero variables, must be erased from the list of dead elements. At this point the list of dead elements is empty, which means that there are no dead elements in the product line model.

The purpose of the aforementioned list is to reduce the number of queries. For instance in this example, only two queries were necessary to evaluate all reusable elements. In contrast, 17 queries would have been required in the current state of the art algorithm. However, it is not possible to calculate in advance how many queries would be needed, or even, to guaranty that the minimal number of queries will be executed, as this depends on the configuration generated by the solver.

- e. **False optional reusable elements:** a reusable element is false optional if it is included in all the products of the product line despite being declared optional (Von der Maßen & Lichter 2004; Benavides *et al.* 2005; Trinidad *et al.* 2008). This operation takes a PLM and a collection of reusable elements modeled as optional as input, and returns the set of false optional reusable elements, or *false* if no one exists. Trinidad *et al.* (2006, 2009) detect false optional features based on finding all products and then searching for common features among those which are not assigned as common. To verify if an optional reusable element is false optional, we query for a product that does not contain the reusable element at hand (setting the feature's value to 0). If there is no such product, then the reusable element we are evaluating is indeed false optional.

```
FalseOptionalReusableElements(PLM M, Solver S) {
```

```
    S.charge(M);
```

```
    FalseOptionalElementsList = all optional elements of M;
```

```
    For (each variable V ∈ FalseOptionalElementsList) {
```

```
        Product = S.getOneSolution("V = 0");
```

```
        If (Product = "false") {
```

```
            Write (V + " is false optional");
```

```
        }
```

```
    } Else {
```

```
        Erase V and all the other variables with a Zero affectation into
        Product, from DeadElementsList;
```

```
1         }  
2     }  
3 }
```

4 For example if we want to know whether the optional reusable component *Usb* is false
5 optional or not, it is sufficient to request for a product without this component
6 (*Support_usb=0*). The solver, in this case, returns the product $P1 = [UNIX=1, Cdrom=1,$
7 $Usb=0, Net=0, UtilityProgram=1, FileMaintenance=1, Editing=1, UserInterface=1,$
8 $Graphical=1, WidthResolution=800, HeightResolution=600, Shell=1, SH=1, TCSH=1,$
9 $BASH=0, Onlinelnfor=0, ProgrammingSupport=0]$, which means that this optional
10 reusable element can take the value of 0, it is, be effectively optional.

11 **2.3. Redundancy-free:** according to the Oxford dictionary something redundant is
12 something “able to be omitted without loss of meaning or function” (Oxford University
13 2008). Therefore, redundancy in a PLM is about the presence of reusable elements and
14 variability constraints among them that can be omitted from the PLM without loss of
15 semantic on the PLM. Redundant constraints in FMs are undesired because, although
16 they do not alter the space of solutions, they may consume extra computational effort in
17 derivation and analysis operations (Yan *et al.* 2009), and they are likely to generate
18 inconsistencies when the PL evolves. For the sake of evolution, it is certainly better
19 detect and correct these redundancies. In order to detect them in a PLM this chapter
20 proposes an operation that takes a PLM and a constraint as input and returns *true* if
21 removing the constraint does not change the space of solutions.

22 Three alternatives can be implemented to check if a relationship is redundant or not.
23 The naïve algorithm consists in calculating all the products of the PLM with the
24 constraint to check; then, remove the constraint; and calculate all the solutions of the
25 new model. If both results are equal (i.e. exact the same products can be configured
26 with and without the constraint), then the constraint is redundant. This approach is
27 computationally very expensive as it requires (a) to compute all configurations twice
28 and (b) to perform an intersection operation between two potentially very large sets
29 (e.g. 10^{21} configurations for the Renault PLM according to Dauron & Astesana (2010)).
30 Not only this algorithm is not scalable, it is typically unfeasible. The second algorithm,
31 proposed by Yan *et al.* (2009) defines a redundant constraint of a PLM as a constraint in
32 which a redundant reusable element takes part. This approach consists in calculating
33 the redundant reusable elements on feature models – features disconnected from the
34 FM – and then the redundant constraint in this approach are those in which the
35 redundant features take part. Though it yields a solution, this algorithm is not
36 sufficiently general: indeed, only these trivial cases of redundancy are considered. The
37 approach proposed in this chapter is based on the fact that if a system is consistent, then
38 the system plus a redundant constraint is consistent too. Therefore, negating the
39 allegedly redundant relation implies contradicting the consistency of the system and
40 thus rendering it inconsistent (Mazo *et al.* 2011a). This approach is more efficient, and
41 thus more scalable, when applied on large models. Our algorithm is in two steps: first,
42 it tries to obtain a solution with the set of constraints. Then, if a solution exists, we
43 negate the constraint we want to check. In the case where no solution is found, the
44 inspected constraint turns out to be redundant. This alternative to find redundant
45 constraints can be formalized as follows:

```

1   If (at least 1 product can be configured from PLM M under a collection of constraints C
2   = {C1,...,Ci}) {
3       Write (C | = M);
4       Let take Cr ∈ C a constraint to be evaluated;
5       If (C without Cr | = M AND C ∪ ¬Cr | ≠ M) {
6           Write (Cr is redundant);
7       }
8       Else{
9           Write (Cr is not redundant);
10      }
11  }

```

12 For example, to check if the constraint $UNIX \geq UtilityProgram$ (cf. Table 1) is redundant
13 or not, it is sufficient to query the solver for a product. Then, if a product is found, the
14 algorithm proceeds to replace the constraint by its negation ($UNIX < UtilityProgram$)
15 and ask again for a product. If the solver does not give a solution (as is the case for our
16 running example), one can infer that the constraint ($UNIX \geq UtilityProgram$) is not
17 redundant.

18 5. Multi-model verification

19 Multi-model modeling allows tackling various models and aspects of a system, in particular
20 in the presence of stakeholders with multiple viewpoints (executives, developers,
21 distributors, marketing, architects, testers, etc.; cf. Nuseibeh *et al.* 1994). For example, a
22 UNIX product line can be composed of several models, each one developed by a different
23 team or developing a particular view of the PL. Thus, while the team responsible of the
24 kernel develops a model, the team responsible of the user interface develops another model.
25 Motivated by the fact that (a) this practice is current in industry (Dhungana *et al.* 2010); (b)
26 even if each individual model is consistent, once the models are integrated, they can easily
27 be inconsistent; and (c) the lacks in current state of the art in multi-model PL verification,
28 this chapter proposes a method to verify multi-model PLs. This method is composed of
29 fourth steps: (i) the base models' semantic should be transformed into constraint programs;
30 (ii) once these base models transformed into CP, they may be integrate using the integration
31 strategies and rules appropriates for each language (cf. Mazo *et al.* 2011a for further details
32 about integration of Dopler models, and Mazo *et al.* 2011d for further details about
33 integration of constraint-based PLMs; and (iii) once the base models integrated, the
34 collection of verification criteria, proposed in Section 4 for single models, can be applied on
35 the integrated model in the same manner as for single models.

36 The application of these verification criteria over the Dopler model depicted in Figure 2 and
37 the explanation regarding the minor variants are presented as follows:

- 38 1. **Non-void model.** This model is not a void because it allows configure at least one
39 product; for instance $C1 = \{USB, Editing, ProgrammingSupport, Shell\}$
- 40 2. **Non-false model.** This model is not a false because it allows configure more than two
41 products; for instance: $C2 = \{Cdrom, Editing, OnlineInfo, Shell, Twm, KDE, Qt,$
42 $GraphicalResolution = "800x600", Width = 800\}$ and $C3 = \{USB, Editing\}$.
- 43 3. **Non-attainable validity conditions' and domains' values.** This operation either (i)
44 takes a collection of decisions as input and returns the decisions that cannot attain

1 one or more values of its validity condition; or (ii) takes a collection of assets as input
2 and returns the assets that cannot attain one of the values of its domain. A non-
3 attainable value of a validity condition or a domain is a value that can never be taken
4 by a decision or an asset in a valid product. Non-attainable values are undesired
5 because they give the user a wrong idea of the values that decisions and assets
6 modeled in the product line model can take. In our example of Figure 2, the validity
7 condition $Width \geq 800 \ \&\& \ Width \leq 1366$ determines a very large range of values that
8 can take the variable *Width*, however this variable can really take three values: 800,
9 1024 and 1366 which means that values like 801, 802,..., 1023, 1025, ..., 1365 are not
10 attainable values.

- 11 4. **Dead reusable elements.** In Dopler language, the reusable elements are Decisions and
12 Assets. This operation takes a collection of decisions and assets as input and returns the
13 set of dead decisions and assets (if some exist) or *false* otherwise. A decision is dead if it
14 never becomes available for answering it. An asset is dead if it cannot appear in any of
15 the products of the product line. The presence of dead decisions and assets in product
16 line models indicates modeling errors and intended but unreachable options. A
17 decision can become dead (i) if its visibility condition can never evaluate to true (e.g., if
18 contradicting decisions are referenced in a condition); (ii) a decision value violates its
19 own visibility condition (e.g., when setting the decision to true will in turn make the
20 decision invisible); or (iii) its visibility condition is constrained in a wrong way (e.g., a
21 decision value is $> 5 \ \&\& \ < 3$ at the same time). An asset can become dead (i) if its
22 inclusion depends on dead decisions, or (ii) if its inclusion condition is false and it is not
23 included by other assets (due to *requires* dependencies to it). Dead variables in CP are
24 variables than can never take a valid value (defined by the domain of the variable) in
25 the solution space. Thus, our approach consists in evaluating each non-zero value of
26 each variable's domain. If a variable cannot attain any of its non-zero values, the
27 variable is considered dead. For instance, in the Dopler model of Figure 2, there are not
28 dead decisions or assets.
- 29 5. **Redundancy-free.** In the asset model (cf. the right side of Figure 2) the asset *4dwn*
30 requires *MwM*, which at the same time requires the asset *Motif*, therefore the
31 dependency *4dwm* requires *Motif* is redundant according to the redundancy-free
32 algorithm presented in Section 4.

33 It is worth noting that the domain-specific operation "false optional-free reusable elements"
34 is not applicable in Dopler models due to the fact that this language does not have explicitly
35 the concept of optional. Decisions and assets are optional in Dopler models according to the
36 evaluation of the visibility conditions (in the case of decisions) and inter-assets
37 dependencies in the case of assets

38 6. Validation

39 We performed a series of experiments to evaluate the verification approach proposed in this
40 chapter. The goal was to measure the effectiveness or precision of the defect's detection, the
41 computational scalability and the usability of the approach to verify different kinds of
42 product line models. These measurements are presented in the next sections, grouped by the
43 kind of product line models used to evaluate our approach.

6.1 Single-view models

We assessed the feasibility, precision and scalability of our approach with 46 models, out of which 44 were taken from the SPLOT repository (Mendonca *et al.* 2009b) and the other two models are the Vehicle movement control system (Salinesi *et al.* 2010b) and the Stago model (Salinesi *et al.* 2011). The sizes of the models are distributed as follows: 32 models of sizes from 9 to 49 features, 4 from 50 to 99, 5 from 100 to 999 and 6 from 1000 to 2000 features. The six largest feature models that we have were not considered in this experiment due to the fact that the solver used does not accept more that 5000 variables. Note that SPLOT models do not have attributes, on the contrary to our two industrial models. Therefore artificial attributes were introduced in a random way, in order to have models with 30%, 60% or 100% of their features with attributes. In order to do that, we created a simple tool¹ that translates models from SPLOT format to constraint programs, and we integrate next the artificial attributes. In order to test that the transformation respects the semantic of each feature model, we compared the results of our models without attributes with the results obtained with the tools SPLOT (Mendonca *et al.* 2009b) and FaMa (Trinidad *et al.* 2008b). In both comparisons we obtained the same results in all the shared functions: detection of void models, dead features, and false optional features. These results show that our transformation algorithm respects the semantic of initial models.

6.1.2 Precision of the detection

Not only must the transformation of FMs into CPs be correct but also the detection of defects. As aforementioned, we compared the results obtained with our tool VariaMos against these obtained with two other tools: SPLOT and FaMa. These comparisons were made over models without attributes due to the fact that original models taken from SPLOT, and also available for FaMa, do not have attributes. In these comparisons we find the same results, for the common verification functions on the three tools, but due to the fact that our own models contain attributes and group cardinalities $\langle m..n \rangle$, for any m and n bellowing to non negative integer numbers, a manual inspection were necessary. A manual inspection on two samples of 28 and 56 features showed that our approach identify the 100% of the anomalies with 0% false positive.

6.1.3 Computational scalability

The execution time of the verification operations in our tool shows that the performance obtained with our approach is acceptable in realistic situations; because in the worst case, users can execute any verification operation less than 19 seconds for models up to 2000 features. Figure 4 shows the execution time of each one of the six verification operations in the 50 models. In Figure 4 each plot corresponds to a verification operation: Figure 4(1) corresponds to operation 1, Figure 4(2) corresponds to operation 2 and so on. Times in the Y axis are expressed in milliseconds (ms) and X axis corresponds to the number of features. It is worth noting that most of the results overlap the other ones; we avoid the use of a logarithmic scale in the X axis, to keep the real behaviour of the results.

¹ parserSPLOTmodelsToCP.rar available at: <https://sites.google.com/site/raulmazo/>

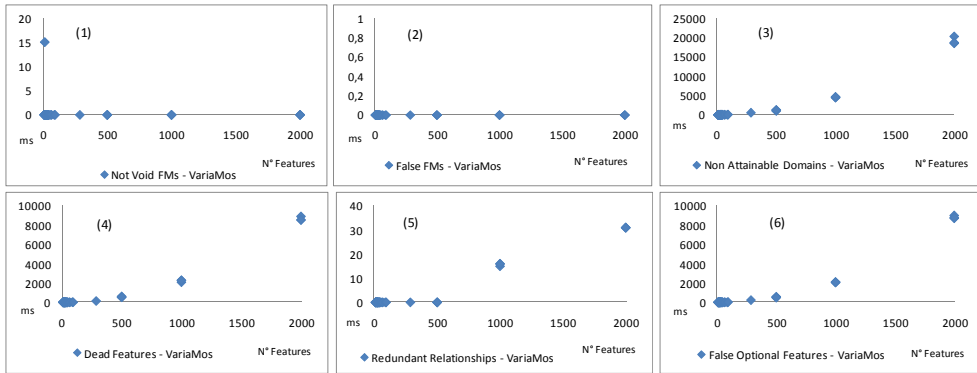


Fig. 4. Execution time of the six verification operations, per number of features

Let us now present the results in more detail. For the models with sizes between 9 and 100 features our approach verified all operations in less than 1 second on average. For the models with sizes between 101 and 500 features verified dead features and false optional features in 0,4 seconds, 1 second to calculate the non attainable domains and 0 milliseconds in the rest of verification operations. It is worth noting that our solver does not provide time measures of microseconds (10^{-6} seconds); thus, 0 milliseconds (10^{-3} seconds) must be interpreted as less than 1 millisecond. In general, over the 46 FMs, the execution time to detect dead features, false optional features and non attainable domains is inferior than 8,68, 8,82 and 19,09 seconds respectively. For the rest of verification operations, the execution time is inferior to 0,02 seconds even for the largest models. Following the projection of our results, our approach is able to be used in larger FMs with a quadratic increase, in the worst of cases, of the time to execute any verification operation proposed in this paper. To finish, the verification operations like redundant relationships, false feature models and void feature models are executed in less than 0,03 seconds. According to the results of our experiment, we can conclude that our verification approach presented in this chapter is scalable to large FMs.

3.6 The case multi-view models

We also tested our verification approach with two Doppler variability models (Mazo *et al.* 2011a). In both models, we seeded 33 defects in the DOPLER model and 22 defects in the camera model. The defects cover different types of problems to show the feasibility of the verification approach. For instance, the decision *Wizard_height* cannot take the values 1200, 1050, 1024 and 768 and the asset *VAI_Configuration_DOPLER* cannot take the value 1 (is never included for any product), even if these values take part in the corresponding variables' domain. Furthermore, we measured the execution time of applying the approach for both models for the different verification operations as presented below.

Applying our verification approach to the DOPLER model has shown that the model is not void and can generate 23016416 products. However, we discovered 18 defects related with non-attainable domain values and 15 dead decisions and assets (these together are the 33 defects we have seeded before). By applying our verification approach on the digital camera

1 model we obtained that the model is not void and can generate 442368 products. In this
 2 model, we discovered 11 defects related with non-attainable domain values as well as 11
 3 dead decisions and assets (these together are the 22 defects we have seeded before). It is
 4 noteworthy that the same number of defects was identified in a manual verification of both
 5 models. The automated verification found all of the seeded defects in the DOPLER model
 6 and all of the seeded defects in the camera model.

7 Table 2 shows the number of defects found and the execution time (in milliseconds)
 8 corresponding to the verification operations on the models. No defects were found
 9 regarding the “Void model”, “False model” and “Redundant relationships” operations and
 10 the execution time was less than 1 millisecond for each one of these operations in each
 11 model. The model transformations from Dopler models to constraint programs took about 1
 12 second for each model.
 13

		Void model	False model	Non-attainable domains	Dead Decisions and Assets	Redundant relationships
DOPLER 81 Variables	Defects	No	No	18	15	No
	Time	0	0	125	47	0
Camera 39 Variables	Defects	No	No	11	11	No
	Time	0	0	16	15	0

14 Table 2. Results of model verifications: Execution time (in milliseconds) and number of
 15 defects found with each verification operation.

16 In the same way as for the single-view models, the results obtained on multi-view models
 17 allow concluding that the verification approach presented in this chapter is scalable to
 18 medium Dopler models and give promising expectations on large Dopler models.

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