Pairing-based algorithms for jacobians of genus 2 curves with maximal endomorphism ring

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Abstract. Using Galois cohomology, Schmoyer characterizes cryptographic non-trivial self-pairings of the ℓ-Tate pairing in terms of the action of the Frobenius on the ℓ-torsion of the Jacobian of a genus 2 curve. We apply similar techniques to study the non-degeneracy of the ℓ-Tate pairing restrained to subgroups of the ℓ-torsion which are maximal isotropic with respect to the Weil pairing. First, we deduce a criterion to verify whether the jacobian of a genus 2 curve has maximal endomorphism ring. Secondly, we derive a method to construct horizontal (ℓ, ℓ)-isogenies starting from a jacobian with maximal endomorphism ring.

1 Introduction

A central problem in elliptic and hyperelliptic curve cryptography is that of constructing an elliptic curve or an abelian surface having a given number of points on their Jacobian. The solution to this problem relies on the computation of the Hilbert class polynomial for a quadratic imaginary field in the genus one case. The analogous genus 2 case needs the Igusa class polynomials for quartic CM fields. There are three different methods to compute these polynomials: the analytic algorithm [16], the p-adic algorithm [7] and a Chinese Remainder Theorem-based algorithm [5]. The last one relies heavily on an algorithm for determining endomorphism rings of the jacobians of genus 2 curves over prime fields.

Eisenträger and Lauter [5] gave the first algorithm for computing endomorphism rings of Jacobians of genus 2 curves over finite fields. They assume that the endomorphism ring is an order in a primitive quartic CM field, i.e. a purely imaginary quadratic extension field of a real quadratic field with no proper imaginary quadratic fields. Given a CM quartic field K such that the real quadratic field K_0 has class number 1, the algorithm computes a set of generators of an order $\mathcal{O}$ in the CM field and tests whether these generators are endomorphisms of J, in order to decide whether the order $\mathcal{O}$ is the endomorphism ring $\mathrm{End}(J)$ or not. In view of application to the CRT method for computing Igusa class polynomials, Freeman and Lauter bring a series of improvements to this algorithm,

in the particular case where we need to decide whether $\mathrm{End}(J)$ is the maximal order or not.

Note that the Eisenträger-Lauter CRT method for computing class polynomials searches for curves defined over some prime field $\mathbb{F}_p$ and belonging to a certain isogeny class. Once such a curve is found, the algorithm keeps the curve only if it has maximal endomorphism ring. This search is rather expensive and ends only when all curves having maximal endomorphism ring were found. Recent research in the area [1, 15, 3] has shown that we can significantly reduce the time of this search by using *horizontal isogenies*, i.e. isogenies between jacobians having the same endomorphism ring. Indeed, once a Jacobian with maximal endomorphism ring is found, many others can be generated from it by computing horizontal isogenies.

In this paper, we propose a new method for checking if the endomorphism ring is locally maximal at ℓ, for $\ell > 2$ prime. Our method relies on the computation of the Tate pairing. We study subgroups of the ℓ-power torsion which are maximal isotropic with respect to the Weil pairing and such that the Tate pairing restricted to these subgroups is $k_{\ell,J}$-degenerate (in the sense of Definition 1). We show that the computation of $k_{\ell,J}$ suffices to check whether the endomorphism ring is locally maximal at ℓ, in many cases. Moreover, we give a method to distinguish kernels of horizontal (ℓ, ℓ)-isogenies from other isogenies of principally polarized abelian varities. Our main result is the following theorem.

Theorem 1. *Let H be a hyperelliptic curve defined over a finite field $\mathbb{F}_q$ and let J be its jacobian, whose endomorphism ring is a locally maximal order at ℓ of a CM-field K. Suppose that the Frobenius endomorphism is exactly divisible by ℓ^n, $n \in \mathbb{Z}$ and that the conditions in Lemma 3 are satisfied. Then a subgroup $G \subset J[\ell]$, which is maximal isotropic with respect to the Weil pairing, is the kernel of a descending isogeny if the Tate pairing is $k_{\ell,J}$-non-degenerate over $\bar{G} \times \bar{G}$, for $\bar{G} \subset J[\ell^n]$ such that $\ell^{n-1}\bar{G} = G$ and that $\bar{G}$ is maximal isotropic with respect to the ℓ^n-Weil pairing. The isogeny is horizontal if the Tate pairing is $k_{\ell,J}$-degenerate over $\bar{G} \times \bar{G}$.*

In view of application to the CRT method for Igusa polynomial computation, we deduce an algorithm to compute kernels of horizontal isogenies efficiently. This generalizes a result on horizontal ℓ-isogenies for genus 1 curves [8].

This paper is organised as follows. In Section 2 we recall briefly the Eisenträger-Lauter algorithm for computing endomorphism rings. In Section 3 we give the definition and properties of the Tate pairing. Section4 describes our algorithm for checking whether a Jacobian has locally maximal order at ℓ. Finally, in Section 5 we show that we can compute kernels of horizontal (ℓ, ℓ)-isogenies by some Tate pairing calculations.

Notation and assumptions. In this paper, we assume that principally polarized abelian surfaces are *simple*, i.e. not isogenous to a product of elliptic curves. A quartic CM field K is a totally imaginary quadratic extension of a totally real field. We denote by K_0 the real quadratic subfield of K and we assume that K_0 has class number 1. $K = \mathbb{Q}(\eta)$, with $\eta = i\sqrt{a + b\sqrt{d}}$ if $d \equiv 2, 3 \mod 4$ or

$\eta = i\sqrt{a + b\left(\frac{-1+\sqrt{d}}{2}\right)}$ if $d \equiv 1 \mod 4$. A CM-type Φ is a couple of pairwise non-complex conjugate embeddings of K in $\mathbb{C}$

$$\Phi(z) = (\phi_1(z), \phi_2(z)).$$

An abelian surface over $\mathbb{C}$ with complex multiplication by $\mathcal{O}_K$ is given by $A(\mathbb{C}) = \mathbb{C}^2/\Phi(\mathfrak{a}^{-1})$, where $\mathfrak{a}$ is an ideal of $\mathcal{O}_K$ and Φ is a CM type. This variety is said to be of CM-type (K, Φ). A CM-type (K, Φ) is primitive is Φ cannot be obtained as a lift of a CM-type of a CM-subfield of K. The principally polarized abelian variety $\mathbb{C}^2/\Phi(\mathfrak{a}^{-1})$ is simple if and only if its CM-type is *primitive* [13].

2 Computing the endomorphism ring of a jacobian

The endomorphism ring of an ordinary jacobian J over a finite field $\mathbb{F}_q$ is an order in a quartic CM field K such that

$$\mathbb{Z}[\pi, \bar{\pi}] \subset \mathrm{End}(J) \subset \mathcal{O}_K,$$

where $\mathbb{Z}[\pi, \bar{\pi}]$ denotes the order generated by π, the Frobenius endomorphism and by $\bar{\pi}$, the Verschiebung. We give a brief description of the Eisenträger-Lauter algorithm [5] which computes the endomorphism ring of J. For a fixed order $\mathcal{O}$ in the lattice of orders of K, the algorithm tests whether this order is contained in $\mathrm{End}(J)$. This is done by computing a $\mathbb{Z}$-basis for the order and checking whether the elements of this basis are endomorphisms of J or not. In order to test if $\alpha \in \mathcal{O}$ is an endomorphism, we write

$$\alpha = \frac{a + b\pi + c\pi^2 + d\pi^3}{n}, \tag{1}$$

with a, b, c, d, n some integers such that a, b, c, d have no commun factor with n (n is the smallest integer such that $n\alpha \in \mathbb{Z}[\pi]$). The LLL algorithm computes a sequence a, b, c, d, n such that α can be written as in Equation 1. In order to check whether α is an endomorphism or not, Eisenträger and Lauter [5] use the following result.

Lemma 1. *Let A be an abelian variety defined over a field k and n an integer coprime to the characteristic of k. Let $\alpha : A \to A$ be an endomorphism of A. Then $A[n] \subset Ker\ \alpha$ if and only if there is another endomorphism β of A such that $\alpha = n \cdot \beta$.*

Using Lemma 1, we get $\alpha \in \mathrm{End}(J)$ if and only if $a + b\pi + c\pi^2 + d\pi^3$ acts as zero on the n-torsion. Freeman and Lauter show that n divides the index $[\mathcal{O}_K : \mathbb{Z}[\pi]]$ (see [6, Lemma 3.3]). Since $\mathbb{Z}[\pi, \bar{\pi}]$ is 1 or p, we have that n divides $[\mathcal{O}_K : \mathbb{Z}[\pi, \bar{\pi}]]$ if $(n, p) = 1$. Moreover, Freeman and Lauter show that if n factors as $\ell_1^{d_1}\ell_2^{d_2}\ldots\ell_r^{d_r}$, it suffices to check if

$$\frac{a + b\pi + c\pi^2 + d\pi^3}{\ell_i^{d_i}},$$

for every prime factor ℓ_i in the factorization of n. The advantage of using this family of elements instead of α is that instead of working over the extension field generated by the coordinates of the n-torsion points, we may work over the field of definition of the $\ell_i^{d_i}$-torsion, for every prime factor ℓ_i. Freeman and Lauter prove the following result, which allows computing a bound for the degree of the smallest extension field over which the ℓ-torsion points are defined.

Proposition 1. *[6, Prop. 6.2] Let J be the Jacobian of a genus 2 curve over $\mathbb{F}_q$ and suppose that End(J) is isomorphic to the ring of integers $\mathcal{O}_K$ of the primitive quartic CM field K. Let $\ell \neq q$ be a prime number, and suppose $\mathbb{F}_{p^r}$ is the smallest field over which the points of $J[\ell]$ are defined. If ℓ is unramified in K, then r divides one of the following:*

(a) $\ell - 1$, if ℓ splits completely in K;
(b) $\ell^2 - 1$, if ℓ splits into two or three ideals in K;
(c) $\ell^3 - \ell^2 + \ell - 1$, if ℓ is inert in K.

If ℓ ramifies in K, then r divides one of the following:

(a) $\ell^3 - \ell^2$, if there is a prime over ℓ of ramification degree 3, or if ℓ is totally ramified in K and $\ell \geq 3$.
(b) $\ell^2 - \ell$, in all other cases where ℓ factors into four prime ideals in K (counting multiplicities).
(c) $\ell^3 - \ell$, if ℓ factors into two or three prime ideals in K (counting multiplicities).

Once we computed the extension field over which the ℓ-torsion is defined, the ℓ^d-torsion will be computed using the following result [6].

Proposition 2. *[6, Prop. 6.3] Let A be an ordinary abelian variety defined over a finite field $\mathbb{F}_q$ and let ℓ be a prime number not equal to the characteristic of $\mathbb{F}_q$. Let d be a positive integer. If the ℓ-torsion points of A are defined over $\mathbb{F}_q$, then the ℓ^d-torsion points are defined over $\mathbb{F}_{q^{\ell^{d-1}}}$.*

3 Background on the Tate pairing

Consider now H a hyperelliptic genus 2 curve defined over a finite field $\mathbb{F}_q$, with $q = p^n$, whose equation is

$$y^2 + h(x)y = f(x), \tag{2}$$

with $h, f \in \mathbb{F}_q[x]$, $\deg h \leq 2$, f monic and $\deg f \leq 4$. Let J be the jacobian of H and let $\bar{\mathbb{F}}_q$ the algebraic closure of $\mathbb{F}_q$ and by $G = \mathrm{Gal}(\bar{\mathbb{F}}_q/\mathbb{F}_q)$ the Galois group. Let $m \in \mathbb{N}$. The Weil pairing

$$W_m : J[m] \times \hat{J}[m] \to \mu_m$$

is a bilinear, non-degenerate map and it commutes with the action of G. If $\lambda : A \to \hat{A}$ is a polarization, then we define the Weil pairing as

$$\begin{aligned} W_m : J[m] \times J[m] &\to \mu_m \\ (P, Q) &\to W_m(P, \lambda(Q)). \end{aligned}$$

We denote by $H^i(G, J)$ the i-th Galois cohomology group, for $i \geq 0$. We consider the exact sequence $0 \to J[m] \to J(\bar{\mathbb{F}}_q) \to J(\bar{\mathbb{F}}_q) \to 0$. Then by taking Galois cohomology we get the connecting morphism

$$\begin{aligned} \delta : J(\mathbb{F}_q)/mJ(\mathbb{F}_q) = H^0(G, J)/mH^0(G, J) &\to H^1(G, J[m]) \\ P &\to F_P, \end{aligned}$$

where the map F_P is defined as follows

$$\begin{aligned} F_P : G &\to J(\bar{\mathbb{F}}_q)[m] \\ \sigma &\to \sigma(\bar{P}) - \bar{P}, \end{aligned}$$

where $\bar{P}$ is any point such that $m\bar{P} = P$. Suppose that the group of m-torsion points $J[m]$ (i.e. points of order m) is J-rational. Using the connecting morphism and the Weil pairing, the Tate pairing is defined as follows

$$\begin{aligned} J(\mathbb{F}_q)/mJ(\mathbb{F}_q) \times \hat{J}[m](\mathbb{F}_q) &\to H^1(G, \mu_m) \\ (P, Q) &\to [\sigma \to W_m(F_P(\sigma), Q)]. \end{aligned}$$

By Hilbert's Theorem 90 we have that $H^1(G, \mu_m) \simeq \mathbb{F}_q^*/\mathbb{F}_q^{*m}$. It follows that we may define a bilinear map

$$t_m(\cdot, \cdot) : J(\mathbb{F}_q)/mJ(\mathbb{F}_q) \times \hat{J}[m](\mathbb{F}_q) \to \mathbb{F}_q^*/\mathbb{F}_q^{*m}$$

For a fixed principal polarization $\lambda : J \to \hat{J}$ we define a pairing on J itself

$$\begin{aligned} t_m^{\lambda}(\cdot, \cdot) : J(\mathbb{F}_q)/mJ(\mathbb{F}_q) \times A[m](\mathbb{F}_q) &\to \mathbb{F}_q^*/\mathbb{F}_q^{*m} \\ (P, Q) &\to t_m(P, \lambda(Q)). \end{aligned}$$

Most often, if J has a distinguished principal polarization and there is no risk of confusion, we write simply $t_m(\cdot, \cdot)$ instead of $t_m^{\lambda}(\cdot, \cdot)$. Lichtenbaum [10] describes a version of the Tate pairing on Jacobian varieties. More precisely, suppose we have $m | \#J_H(\mathbb{F}_q)$ and denote by k the *embedding degree with respect to* m, i.e. the smallest integer $k \geq 0$ such that $m | q^k - 1$. Let $[D_1] \in J_H(\mathbb{F}_{q^k})$ and $[D_2] \in J_H[m](\mathbb{F}_{q^k})$ two divisor classes, and let $[D_1]$ be represented by D_1 and $[D_2]$ by D_2 such that $\mathrm{supp}(D_1) \cap \mathrm{supp}(D_2) = \emptyset$. Since $[D_2]$ has order m, there is a function f_{m,D_2} is such that $\mathrm{div}(f_{m,D_2}) = mD_2$. The Tate pairing of the divisor classes $[D_1]$ and $[D_2]$ is computed as

$$t_m(D_1, D_2) = f_{D_2}(D_1).$$

Moreover, in computational applications, it is convenient to work with a unique value of the pairing. We denote by $\mu_m \subset \mathbb{F}_{q^k}^*$ the subgroup of m-th roots of unity. Given that $\mathbb{F}_{q^k}^*/(\mathbb{F}_{q^k})^m \simeq \mu_m$, we use the *reduced Tate pairing*, given by

$$T_m(\cdot,\cdot) : J_H(\mathbb{F}_{q^k})/mJ_H(\mathbb{F}_{q^k}) \times J[m](\mathbb{F}_{q^k}) \to \mu_m$$
$$(P,Q) \to t_m(P,Q)^{(q^k-1)/m}.$$

The function $f_{m,D_1}(D_2)$ is computed using Miller's algorithm [11] in $O(\log m)$ operations in $\mathbb{F}_{q^k}$.

4 Pairings and endomorphism ring computation

In this section we relate some properties of the Tate pairing to the isomorphism class of the endomorphism ring of the Jacobian. Let ℓ be a prime odd number. We give a method to check whether the endomorphism ring is locally maximal at ℓ (i.e. the index $[\mathcal{O}_K : \mathcal{O}]$ is not divisible by ℓ) by computing a certain number of pairings.

Let H be a genus 2 curve defined over a finite field $\mathbb{F}_q$, J its jacobian and suppose that $J[\ell^n] \subseteq J(\mathbb{F}_q)$ and that $J[\ell^{n+1}] \nsubseteq J(\mathbb{F}_q)$, with ℓ different from p and $n \geq 1$. We denote by $\mathcal{W}$ the set of maximal isotropic subgroups in $J[\ell^n]$ with respect to the ℓ-Weil pairing and we define $k_{\ell,J}$ to be

$$k_{\ell,J} = \max_{G \in \mathcal{W}} \{k | \exists P, Q \in G \text{ and } T_{\ell^n}(P,Q) \in \mu_{\ell^k} \backslash \mu_{\ell^{k-1}}\}$$

Definition 1. *Let G be a rank 2 subgroup of $J[\ell^n]$ in $\mathcal{W}$. We say that the Tate pairing is $k_{\ell,J}$-non-degenerate (or simply non-degenerate) on $G \times G$ if its restriction*

$$T_{\ell^n} : G \times G \to \mu_{\ell^{k_{\ell,J}}}$$

is surjective. Otherwise, we say that the Tate pairing is $k_{\ell,J}$-degenerate (or simply degenerate) on $G \times G$. Moreover, for two divisor classes $D_1, D_2 \in G$, we say that they have non-degenerate pairing if $T_{\ell^n}(D_1, D_2)$ is a $\ell^{k_{\ell,J}}$th root of unity and degenerate otherwise.

Lemma 2. *The reduced Tate pairing defined as*

$$T_{\ell^n} : J[\ell^n] \times J[\ell^n] \to \mu_{\ell^n}$$

is $k_{\ell,J}$-antisymmetric, i.e. $T_{\ell^n}(D_1, D_2)T_{\ell^n}(D_2, D_1) \in \mu_{\ell^{k_{\ell,J}}}$, for all $D_1, D_2 \in J[\ell^n]$.

Proof. Indeed, assume that there are $D_1, D_2 \in J[\ell^n]$ such that $T_{\ell^n}(D_1, D_2)T_{\ell^n}(D_2, D_1) \in \mu_{\ell^n} \backslash \mu_{\ell^{k_{\ell,J}}}$. We denote by $G = \langle D_1, D_2 \rangle$ and by $r > k_{\ell,J}$ the largest integer such that $T_{\ell^n}(D_1, D_2)T_{\ell^n}(D_2, D_1)$ is an ℓ^kth primitive root of unity. Then the polynomial

$$\mathcal{P}(a,b) = \log T_{\ell^n}(P,P)a^2 + \log(T_{\ell^n}(P,Q)T_{\ell^n}(Q,P))ab + \log T_{\ell^n}(Q,Q)b^2,$$

where the log function is computed with respect to some fixed ℓ^n-th root of unity, is zero mod ℓ^{n-r-1} and non-zero mod ℓ^{n-r}. Dividing by ℓ^{n-r-1}, we may view $\mathcal{P}$ as a polynomial in $\mathbb{F}_\ell[a,b]$. Since $\mathcal{P}$ is a quadratic non-zero polynomial, it has at most two roots. These correspond to two divisor classes in G, with $k_{\ell,J}$-degenerate self-pairing. Hence, there is at least one divisor $D \in G$ such that $T_{\ell^n}(D,D)$ is a ℓ^r-th root of unity. Since there is at least one maximal isotropic subgroup $W \in \mathcal{W}$ with respect to the Weil pairing such that $\ell^{n-1}D \in W$, this contradicts the definition of $k_{\ell,J}$.

Let $\mathcal{O}$ be an order of K and let $\theta \in \mathcal{O}$. We define

$$v_{\ell,\mathcal{O}}(\theta) := \max_{m \geq 0}\{m : \theta \in \mathbb{Z} + \ell^m \mathcal{O}\}.$$

We denote by $1, \delta, \gamma, \eta$ a $\mathbb{Z}$-basis of $\mathcal{O}$ and and we write $\theta = a_1 + a_2\delta + a_3\gamma + a_4\eta$. Then we compute $v_{\ell,\mathcal{O}}$ as

$$v_{\ell,\mathcal{O}} = v_\ell(\gcd(a_2, a_3, a_4)). \tag{3}$$

Note that the equality in (3) is independent of the choice of the basis. We say that θ is divisible by $t \in \mathbb{Z}$ if we have $\theta \in t\mathcal{O}$. We say that θ is exactly divisible by ℓ^n if it is divisible by ℓ^n and it is not divisible by ℓ^{n+1}. The following lemma gives a criterion to check whether an order is locally maximal at ℓ or not.

Lemma 3. *Let $K := \mathbb{Q}(i\sqrt{a+b\sqrt{d}})$ be a quartic CM field, where $\eta = i\frac{\sqrt{a+b\sqrt{d}}}{2}$, if $d \equiv 1 \bmod 4$ and $\eta = i\sqrt{a+b\sqrt{d}}$, if $d \equiv 2,3 \bmod 4$. We assume that $a, b, d \in \mathbb{Z}$ and that d and $a^2 - b^2 d$ are square free. Assume that $K_0 = \mathbb{Q}(\sqrt{d})$ has class number 1. Let $\ell > 2$ a prime number that does not divide $lcm(a,b,d)$. Let $\mathcal{O}_K$ be the maximal order of K and $\mathcal{O}$ an order such that $[\mathcal{O}_K : \mathcal{O}]$ is divisible by ℓ. Let $\pi \in \mathcal{O}$ such that $N_{K/K_0}(\pi) \in \mathbb{Z}$ is not divisible by ℓ and that $v_{\ell,\mathcal{O}_K}(\pi) > 0$. We suppose that $\pi = a_1 + a_2\frac{-1+\sqrt{d}}{2} + (a_3 + a_4\frac{-1+\sqrt{d}}{2})$, if $d \equiv 1 \bmod 4$ and $\pi = a_1 + a_2\frac{-1+\sqrt{d}}{2} + (a_3 + a_4\sqrt{d})i\sqrt{a + b\frac{1+\sqrt{d}}{2}}$, if $d \equiv 2,3 \bmod 4$. If $v_\ell(a_3 - a_4) < \min(v_\ell(a_3), v_\ell(a_4))$, then $v_{\ell,\mathcal{O}}(\pi) < v_{\ell,\mathcal{O}_K}(\pi)$.*

Proof. We denote by $\mathcal{O}_1 = \mathcal{O}_{K_0} + \mathcal{O}_{K_0}\eta$. Since $\ell > 2$, it suffices to show that $v_{\mathcal{O}\cap\mathcal{O}_1}(\pi) < v_{\mathcal{O}_1}(\pi)$. We will therefore assume, without restricting the generality, that $\mathcal{O} \subset \mathcal{O}_1$. Consider the basis $1, \delta := \sqrt{d}, \gamma := i\sqrt{d}\sqrt{a+b\sqrt{d}}, \eta := i\sqrt{a+b\sqrt{d}}$ for $\mathcal{O}_1$. We write $\pi = a_1 + a_2\delta + a_3\gamma + a_4\eta$. By writing down the norm condition for $d \equiv 2,3 \bmod 4$

$$\left(a_1 + a_2\sqrt{d} + (a_3 + a_4\sqrt{d})i\sqrt{a+b\sqrt{d}}\right)\left(a_1 + a_2\sqrt{d} - (a_3 + a_4\sqrt{d})i\sqrt{a+b\sqrt{d}}\right) \in \mathbb{Z} \tag{4}$$

we get that

$$2a_1a_2 + a_3^2 b + a_4^2 bd + 2aa_3a_4 = 0. \tag{5}$$

Similarly, for $d \equiv 1 \bmod 4$, we have

$$-\frac{a_2^2}{2} + a_1 a_2 - \frac{a a_4^2}{2} + a a_3 a_4 + \frac{a_3^2 b}{2} + \frac{a_4^2(1+d)b}{2} = 0.$$

Since $\ell \nmid a_1$, equations (4) and (5) imply that $v_\ell(a_2) > \max(v_\ell(a_3), v_\ell(a_4))$. Since there is always an order $\mathcal{O}'$ such that $\mathcal{O} \subset \mathcal{O}' \subset \mathcal{O}_1$ such that $[\mathcal{O}_K : \mathcal{O}']$ is a power of ℓ, it suffices to prove the lemma in the case $[\mathcal{O}_1 : \mathcal{O}]$ is a power of ℓ. For the order $\mathcal{O}$, we choose $\{1, \delta', \gamma', \eta'\}$ a HNF basis with respect to $\{1, \delta, \gamma, \eta\}$. We denote by $(a_{i,j})_{1 \leq i,j \leq 4}$ the corresponding transformation matrix. Then $[\mathcal{O}_K : \mathcal{O}] = \prod_{1 \leq i \leq 4} a_{i,i}$. Note that neither η nor γ are in $\mathcal{O}$. Otherwise, $\mathcal{O}$ is the maximal order. Indeed, assume $\eta \in \mathcal{O}$. Since ℓ divides neither a nor b, it follows that $\sqrt{d} \in \mathcal{O}$. This implies that $\mathcal{O}$ is the maximal order. We consider the decomposition of π over the basis $\{1, \delta', \gamma', \eta'\}$

$$\pi = a_1' + a_2'\delta' + a_3'\gamma' + a_4'\eta', a_i' \in \mathbb{Z}.$$

Since $\eta \notin \mathcal{O}$, we know that a_{44} is ℓ. If a_{33} is divisible by ℓ, then $v_\ell(a_3') < v_\ell(a_3)$. If $a_{34} = 1$, then $a_4' = -(a_3 - a_4)/\ell$. If $a_{34} = 0$, then $a_4' = a_4/\ell$. If $a_{33} = 1$, it follows that $a_{34} = 1$ (otherwise we would have $\gamma \in \mathcal{O}$). Then $a_3' = a_3$ and $a_4' = -(a_3 - a_4)/\ell$. We conclude that $v_{\ell,\mathcal{O}}(\pi) < v_{\ell,\mathcal{O}_K}(\pi)$.

Since we know that $J[\ell^n]$ is $\mathbb{F}_q$-rational, while $J[\ell^{n+1}]$ is not, Lemma 1 implies that $\pi - 1$ is exactly divisible by ℓ^n. Moreover, the Frobenius matrix on the Tate module is the identity matrix $I_4 \bmod \ell^n$. The following lemma computes the precision up to which the Frobenius matrix on the Tate module is of the form λI_4, with $\lambda \in \mathbb{Z}$.

Lemma 4. *Let J be a abelian surface defined over a finite field $\mathbb{F}_q$ and π the Frobenius endomorphism. Then the largest integer m such that the matrix of the Frobenius endomorphism on the ℓ-Tate module is of the form*

$$\begin{pmatrix} \lambda & 0 & 0 & 0 \\ 0 & \lambda & 0 & 0 \\ 0 & 0 & \lambda & 0 \\ 0 & 0 & 0 & \lambda \end{pmatrix} \bmod \ell^m \tag{6}$$

is $v_{\ell,\mathcal{O}}(\pi)$, where $\mathcal{O}$ is the endomorphism ring of J.

Proof. Let m be the largest integer such that the matrix of the Frobenius on $J[\ell^m]$ has the form given in Equation (6). Let $\mathcal{O}$ be the endomorphism ring of J. We denote by $\{1, \delta, \gamma, \eta\}$ the $\mathbb{Z}$-basis of $\mathcal{O}$ and by $\pi = a_1 + a_2\delta + a_3\gamma + a_4\eta$ the decomposition of π over this basis. It is obvious that $m \geq v_\ell(\gcd(a_2, a_3, a_4))$. For the converse, we note that $\pi - \lambda$ kills the ℓ^m-torsion, hence we may write $\pi - \lambda = \ell^m \alpha$, with $\alpha \in \mathrm{End}(J)$. We write down the decomposition of α over the basis $\{1, \delta, \gamma, \eta\}$ and conclude that $\ell^m | \gcd(a_2, a_3, a_4)$. Hence $m \leq v_\ell(\gcd(a_2, a_3, a_4))$. We conclude that $m = v_\ell(\gcd(a_2, a_3, a_4))$, hence $m = v_{\ell,\mathcal{O}}$ by equation (3).

Using Galois cohomology, Schmoyer [12] computes the matrix of the Frobenius on the Tate module, up to a certain precision, if the self-pairings of the Tate pairing are degenerate. We use a similar approach and show that the precision up to which the Frobenius acts on the Tate module as a multiple of the identity is $2n - k_{\ell,J}$. Consequently, we recover information on the conductor of the endomorphism ring of J by computing $k_{\ell,J}$. For $m \in \mathbb{Z}$, we will use a *symplectic basis* of $J[\ell^m]$, i.e. a basis such that the matrix associated to the ℓ^m-Weil pairing is

$$\begin{pmatrix} 0 & I \\ -I & 0 \end{pmatrix} \bmod \ell^m. \tag{7}$$

Proposition 3. *Let H be a hyperelliptic curve defined over a finite field $\mathbb{F}_q$, J its jacobian. Suppose that the Frobenius endomorphism π is such that $\pi - 1$ is exactly divisible by ℓ^n, for $\ell \geq 3$ prime. Then*

$$v_{\ell,End(J)}(\pi) = 2n - k_{\ell,J}. \tag{8}$$

Proof. Let $\{Q_1, Q_{-1}, Q_2, Q_{-2}\}$ a symplectic basis for the ℓ^{2n}-torsion (whose matrix is given by Equation (7)) and let $\pi(Q_g) = \sum_{h=-2}^{2} a_{h,g} Q_h$, with $a_{h,g}$, $h, g \in \{-2, -1, 1, 2\}$ in $\mathbb{Z}$. By bilinearity, we have that

$$T_{\ell^n}(\ell^n Q_i, \ell^n Q_j) = W_{\ell^{2n}}(Q_i, \pi(Q_j) - Q_j) = W_{\ell^{2n}}(Q_i, \sum_{\substack{h=-2 \\ h\neq 0}}^{2} a_{h,j} Q_h - Q_j)$$

$$= W_{\ell^{2n}}(Q_i, Q_j)^{a_{j,j}-1} \prod_{h=-2\setminus\{0,j\}}^{2} W_{\ell^{2n}}(Q_i, Q_h)^{a_{h,j}}.$$

If $j \neq -i$, we have that $T_{\ell^n}(\ell^n Q_i, \ell^n Q_j) \in \mu_{\ell^{k_{\ell,J}}}$. It follows that

$$a_{-i,j} \equiv 0 \pmod{\ell^{2n-k_{\ell,J}}},$$

for $i \neq -j$. If $j = -i$, then $T_{\ell^n}(\ell^n Q_i, \ell^n Q_j) = W_{\ell^{2n}}(Q_i, Q_j)^{a_{j,j}-1}$. Since the Tate pairing is $k_{\ell,J}$-antisymmetric we get

$$a_{i,i} \equiv a_{-i,-i} \pmod{\ell^{2n-k}}.$$

It remains to prove that $a_{i,i} \equiv a_{j,j}$, for $i, j \in \{-2, -1, 1, 2\}$. Note that by Galois invariance, we have $W_{\ell^{2n}}(\pi(Q_i), \pi(Q_j)) = \pi(W_{\ell^{2n}}(Q_i, Q_j)) = W_{\ell^{2n}}(Q_i, Q_j)^q$. For $i = -j$ we have

$$W_{\ell^{2n}}(\pi(Q_i), \pi(Q_{-j})) = W_{\ell^{2n}}(\sum_{\substack{h=-2 \\ h\neq 0}}^{2} {}_{h,i} Q_h, \sum_{\substack{g=-2 \\ g\neq 0}}^{2} a_{g,-i} Q_g)$$

$$= \prod_{\substack{h=-2 \\ h\neq 0}}^{2} \prod_{\substack{g=-2 \\ g\neq 0}}^{2} W_{\ell^{2n}}(a_{h,i} Q_h, a_{g,-i} Q_g) = W_{\ell^{2n}}(Q_i, Q_{-i}) \prod_{\substack{h=-2 \\ h\neq 0,-i}}^{2} W_{\ell^{2n}}(a_{h,i} Q_i, a_{-i,-i} Q_{-i})$$

$$\cdot \prod_{\substack{g=-2 \\ g\neq 0,-i}}^{2} W_{\ell^{2n}}(a_{i,i} Q_i, a_{g,-i} Q_g) \prod_{\substack{s=-2 \\ s\neq 0,-i}}^{2} \prod_{\substack{t=-2 \\ t\neq 0,-i}}^{2} W_{\ell^{2n}}(Q_s, Q_t)^{a_{s,i} a_{t,-i}}$$

Since $\{Q_{-2}, Q_{-1}, Q_1, Q_2\}$ is a symplectic basis and that $a_{i,j} \equiv 0 \pmod{\ell^n}$, for $i \neq -j$, then

$$W_{\ell^{2n}}(\pi(Q_i), \pi(Q_j)) = W_{\ell^{2n}}^{a_{i,i}a_{-i,-i}}(Q_i, Q_j).$$

Since $a_{i,i} \equiv a_{-i,-i} \pmod{\ell^{2n-k_{\ell,J}}}$, it follows that

$$a_{i,i}^2 \equiv q \text{ for all } i \in \{-2, -1, 1, 2\}.$$

Since $a_{i,i} \equiv 1 \pmod{\ell^n}$, it follows that $a_{i,i} \equiv b \pmod{\ell^{2n-k_{\ell,J}}}$, for some $b \in \mathbb{Z}$. By Lemma 6, we have $2n - k_{\ell,J} \leq v_\ell(\pi)$. For the converse, we may assume that the matrix of the Frobenius endomorphism has the form given in Equation 6 mod ℓ^{2n-k}, where $k = 2n - v_{\ell,\mathcal{O}}(\pi)$. Let R, S be two points in $J[\ell^n]$ such that $W_\ell(\ell^{n-1}R, \ell^{n-1}S) = 1$. It suffices to show that $T_{\ell^n}(R, S)$ is k-degenerate. We write $\pi - 1 = a_1 + a_2\alpha + a_3\beta + a_4\theta$, where $1, \alpha, \beta, \theta$ are a $\mathbb{Z}$-basis of $\mathrm{End}(J)$. We take $\bar{S}$ such that $S = \ell^n\bar{S}$ and we get

$$\begin{aligned} &T_{\ell^n}(R, S) = W_{\ell^n}(R, (\pi - 1)(\bar{S}) = \\ &= W_{\ell^n}(R, S)^{\frac{a_1}{\ell^n}} W_{\ell^n}(R, (\frac{a_2}{\ell^{n-k}}\delta + \frac{a_3}{\ell^{n-k}}\gamma + \frac{a_4}{\ell^{n-k}}\eta)(S))^{\ell^{n-k}}. \end{aligned}$$

Since $W_\ell(\ell^{n-1}R, \ell^{n-1}S) = 1$ and $v_\ell(\gcd(a_2, a_3, a_4)) = \ell^{2n-k}$, we have $T_{\ell^n}(R, S) \in \mu_{\ell^k}$. Hence $k \geq k_{\ell,J}$. This concludes the proof.

Theorem 2. *Let H be a hyperelliptic curve defined over a finite field $\mathbb{F}_q$ and J its jacobian. Let π the Frobenius endomorphism, ℓ an odd prime, and let n be the greatest positive integer such that $\pi - 1$ is exactly divisible by ℓ^n. Then End(J) is a locally maximal order at ℓ if and only if $k_{\ell,J}$ equals $2n - v_{\ell,\mathcal{O}_K}(\pi)$.*

Proof. By Proposition 3, $k_{\ell,J}$ equals $2n - v_{\ell^n,\mathcal{O}}(\pi)$, where $\mathcal{O} \simeq \mathrm{End}(J)$. By Lemma 3, the value of $v_{\ell^n,\mathcal{O}_K}(\pi)$ uniquely characterizes orders which are locally maximal at ℓ.

Corollary 1. *Let H be a hyperelliptic curve defined over a finite field $\mathbb{F}_q$ and J its jacobian. Let $\pi = 1 + a_1 + a_2\delta + a_3\gamma + a_4\eta$ be the decomposition of the Frobenius over a $\mathbb{Z}$-basis of $\mathcal{O}_K$. Then $k_{\ell,J} > 0$ if and only if $v_\ell(\gcd(a_2, a_3, a_4)) < 2v_\ell \gcd(a_1, a_2, a_3, a_4)$.*

We conclude this section by giving in Algorithm 1 a computational method which verifies whether the jacobian J of a genus 2 curve has locally maximal endomorphism ring.

5 Application to horizontal isogeny computation

In this section, we are interested in computing *horizontal* isogenies, i.e. isogenies between Jacobians having the same endomorphism ring. Note that if $I : J_1 \to J_2$ is an isogeny such that J_1 has maximal endomorphism ring at ℓ, we distinguish

Algorithm 1 Checking whether the endomorphism ring is locally maximal

INPUT: A jacobian J of a genus 2 curve defined over $\mathbb{F}_q$ such that $J[\ell^n] \subset J(\mathbb{F}_q)$, the Frobenius π, a symplectic basis $(Q_1, Q_{-1}, Q_2, Q_{-2})$ for $J[\ell^n]$

OUTPUT: The algorithm outputs true if $\mathrm{End}(J)$ is maximal at ℓ if $v_{\ell,\mathcal{O}_K}(\pi) < 2n$.

```
Compute representatives for 𝒲
for all G ∈ 𝒲 do
    Let P, Q be two generators of G.
    Let a = T_{ℓ^n}(P, P), b = T_{ℓ^n}(Q, P) · T_{ℓ^n}(P, Q) and c = T_{ℓ^n}(P, P)
    Let Count = 0.
    repeat
        Let a = a^ℓ, b = b^ℓ and c = c^ℓ
        Let Count = Count + 1
    until a = 1 and b = 1 and c = 1
    if k_{ℓ,J} < Count then
        k_{ℓ,J} ← Count
    end if
    if k_{ℓ,J} = 0 then
        abort
    end if
    if k_{ℓ,J} = 2n − v_{ℓ,O_K} then
        return true
    else
        return false
    end if
end for
```

two cases: either $\mathrm{End}(J_2)$ is locally maximal at ℓ, or $\mathrm{End}(J_2) \subset \mathrm{End}(J_1)$. In the last case we say that the isogeny is *descending*.

Over the complex numbers, horizontal isogenies are given in terms of the action of the Shimura class group [13]. Let Φ be a CM-type and let A be an abelian surface over $\mathbb{C}$ with complex multiplication by $\mathcal{O}_K$, given by $A = \mathbb{C}^2/\Phi(I^{-1})$, where I is an ideal of $\mathcal{O}_K$. The surface is principally polarized if there is a purely imaginary $\xi \in \mathcal{O}_K$ with $\mathrm{Im}(\Phi_i(\xi)) > 0$, for $i \in \{1, 2\}$, and such that $\xi\mathfrak{D}_K = I\bar{I}$ (where $\mathfrak{D}_K$ is the different $\{\alpha \in \mathcal{O}_K : \mathrm{Tr}_{K/\mathbb{Q}}(\alpha\mathcal{O}_K) \subset \mathbb{Z}\}$. Computing horizontal isogenies is usually done by using the action of the Shimura class group [13]. Shimura defines $\mathfrak{C}(K)$ as

$$\{(\mathfrak{a}, \alpha) | \mathfrak{a} \text{ is a fractional } \mathcal{O}_K\text{-ideal with } \mathfrak{a}\bar{\mathfrak{a}} = (\alpha) \text{ with } \alpha \in K_0 \text{ totally positive}\},$$

where $(\mathfrak{a}, \alpha) \sim (\mathfrak{b}, \beta)$ if and only if there exists $u \in K^*$ with $\mathfrak{b} = u\mathfrak{a}$ and $\beta = u\bar{u}\alpha$. The action of $(\mathfrak{a}, \alpha) \in \mathfrak{C}(K)$ on an principally polarized abelian surface given by (I, ξ) is given by the ideal $(\mathfrak{a}I, \alpha\xi)$. This action is transitive and free [13, §14.6].

The kernel of the horizontal isogeny corresponding to $\mathfrak{a}$ is a subgroup of the ℓ-torsion invariate under the Frobenius endomorphism. Hence in order to compute the kernel, we need to compute the matrix of the Frobenius for some basis of the ℓ-torsion and then determine subspaces invariated by this matrix (see [2, Algorithm VI.3.4]). We show that, when a Jacobian with locally maximal

order at ℓ is given, we distinguish horizontal isogenies from descending ones, by doing computations similar to those in Section 4. The resulting algorithm, whose complexity is analysed in Section 6, computes kernels of horizontal isogenies with only a few pairing computations. We state the following lemma for jacobians of genus 2 curves over finite fields, which are the framework for this paper. We note that the result holds for abelian varieties.

Lemma 5. *(a) Let J_1, J_2 be jacobians of genus 2 curves defined over a finite field $\mathbb{F}_q$ and $I : J_1 \to J_2$ an isogeny defined over $\mathbb{F}_q$ which splits multiplication by d. Let $\lambda : J_1 \to \hat{J}_1$ a principal polarization. Then for $P \in J_1(K)$, $Q \in J_1[m](K)$ we have*

$$T_m^{\lambda_I}(I(P), I(Q)) = T_m^{\lambda}(P, Q)^d,$$

where $\lambda_I : J_2 \to \hat{J}_2$ is such that $I \circ \lambda_I \circ \check{I} = d \circ \lambda$.

(b) Let J_1, J_2 be jacobians of genus 2 curves defined over $\mathbb{F}_q$ and $I : J_1 \to J_2$ an isogeny defined over $\mathbb{F}_q$ which splits multiplication by m. Let $P \in J_1(K)$, $Q \in J_1[mm'](K)$ such that $I(Q)$ is a m'-torsion point.

$$T_{m'}^{\lambda_I}(I(P), I(Q)) = T_{mm'}^{\lambda}(P, Q)^m,$$

where λ_I is a principal polarization of J_2 such that $I \circ \lambda_I \circ \check{I} = m \circ \lambda$.

Proof. (a) It is easy to check that $\delta(I(P)) = I(\delta(P))$. Hence for $\sigma \in G_K$ we have

$$e_m(F_{I(P)}(\sigma), I(Q)) = e_m(I(F_P(\sigma)), I(Q)).$$

By using Milne [9, 13.2.b]

$$e_m^{\lambda_I}(I(F_P(\sigma)), I(Q)) = e_m^{\check{I} \circ \lambda_I \circ I}(F_P(\sigma), Q).$$

(b) The proof is immediate by using (a) and the fact that $T_{mm'}(I(P), I(Q)) = T_{m'}(I(P), I(Q))$.

We may now prove Theorem 2.

Proof of Theorem 1. We denote by $I : J \to J'$ the isogeny of kernel G. Suppose that G is such that the Tate pairing is non-degenerate over $G \times G$. Then by applying Lemma 5 we have

$$T_{\ell^{n-1}}(I(P_1), I(P_2)) \in \mu_{\ell^{k_{\ell,J}-1}} \backslash \mu_{\ell^{k_{\ell,J}-2}},$$

for $P_1, P_2 \in G$. If $J'[\ell^n]$ is not defined over $\mathbb{F}_q$, then its endomorphism ring cannot be maximal at ℓ, hence the isogeny is descending. Assume then that $J'[\ell^n]$ is defined over $\mathbb{F}_q$. Let $\bar{P}_1, \bar{P}_2 \in J'[\ell^n]$ be such that $I(P_1) = \ell\bar{P}_1$, $I(P_2) = \ell\bar{P}_2$. Then $T_{\ell^n}(\bar{P}_1, \bar{P}_2) \in \mu_{\ell^{k_{\ell,J}+1}} \backslash \mu_{\ell^{k_{\ell,J}}}$. We denote by $G' =< \bar{P}_1, \bar{P}_2 >$. The subgroup $\ell^{n-1}G'$ is maximal isotropic with respect to the Weil pairing. It follows that $k_{J'} \geq k_J + 1$. By Theorem 2, we deduce that the endomorphism ring of J' is not locally maximal at ℓ, hence the isogeny is descending.
Suppose now that the Tate pairing is degenerate over $G \times G$. We distinguish two

cases.

Case 1. Suppose that $J'[\ell^n]$ is defined over $\mathbb{F}_q$. With the same notations as above, we get that $T_{\ell^n}(\bar{P}_1, \bar{P}_2) \in \mu_{\ell^{k_{\ell,J}}}$. Let $L \subset J'[\ell^n]$ be a subgroup of rank 2 such that $\ell^{n-1}L$ is maximal isotropic with respect to the Weil pairing and consider $Q_1, Q_2 \in L \backslash G'$. Then $\ell^{n-1}Q_1, \ell^{n-1}Q_2 \in Ker\ I^\dagger$. Since $T_{\ell^{n-1}}(I^\dagger(Q_1), I^\dagger(Q_2)) \in \mu_{\ell^{k_{\ell,J}-2}}$, it follows that $T_{\ell^n}(P_1, P_2) \in \mu_{\ell^{k_{\ell,J}-1}}$. Hence $k_{J'} \leq k_{\ell,J}$. By Theorem 2, we conclude that the endomorphism ring of J' is locally maximal at ℓ.

Case 2. Suppose that $J'[\ell^n]$ is not defined over $\mathbb{F}_q$. Hence I is descending. We have

$$T_{\ell^{n-1}}(I(P_1), I(P_2)) \in \mu_{\ell^{k_{\ell,J}-2}}.$$

Let $L \subset J'[\ell^{n-1}]$ be a subgroup of rank 2 such that $\ell^{n-2}L$ is maximal isotropic with respect to the Weil pairing and consider $Q_1, Q_2 \in L \backslash G'$. Then $\ell^{n-2}Q_1, \ell^{n-2}Q_2 \in Ker\ I^\dagger$. Since $T_{\ell^{n-1}}(I^\dagger(Q_1), I^\dagger(Q_2)) \in \mu_{\ell^{k_{\ell,J}-1}}$, it follows that $T_{\ell^{n-1}}(Q_1, Q_2) \in \mu_{\ell^{k_{\ell,J}}}$. Hence $k_{J'} = k_{\ell,J}$, which contradicts the hypothesis that I is descending.

Algorithm 2 computes kernels of horizontal isogenies by checking, for each subgroup in $G \in \mathcal{W}$, whether the Tate pairing restricted to $\mathbb{G} \times G$ maps to $\mu_{\ell^{k_{\ell,J}}}$ surjectively or not.

Algorithm 2 Computing horizontal isogenies

INPUT: A jacobian J of a genus 2 curve defined over $\mathbb{F}_q$ such that $J[\ell^n] \subset J(\mathbb{F}_q)$ and End(J) is maximal at ℓ, a symplectic basis (P_1, P_2, P_3, P_4) for $J[\ell^n]$

OUTPUT: The algorithm outputs $\mathcal{A}$ the set of kernels of horizontal (ℓ, ℓ)-isogenies.

$\mathcal{A} \leftarrow \emptyset$
Compute representatives for $\mathcal{W}$
for all $G \in \mathcal{W}$ **do**
  Let P, Q be two generators of G.
  Let $a = T_{\ell^n}(P, P)$, $b = T_{\ell^n}(Q, P) \cdot T_{\ell^n}(P, Q)$ and $c = T_{\ell^n}(P, P)$
  Let Count $= 0$.
  repeat
    Let $a = a^\ell$, $b = b^\ell$ and $c = c^\ell$
    Let Count $=$ Count $+ 1$
  until $a = 1$ and $b = 1$ and $c = 1$
  if $k_{\ell,J} >$ Count **then**
    Add G to $\mathcal{A}$.
  end if
end for

6 Complexity analysis

In this section, we evaluate the complexity of Algorithm 1 and compare its performance to that of the Freeman-Lauter algorithm. Note that for a fixed $\ell > 2$, both algorithms perform computations in extension fields over which the ℓ^d-torsion, for a certain ℓ^d dividing $[\mathcal{O}_K : \mathbb{Z}[\pi, \bar{\pi}]]$, is rational.

Checking locally maximal endomorphism rings. In Freeman and Lauter's algorithm, in order to check if $\mathrm{End}(J)$ is locally maximal at ℓ, for $\ell > 2$, it suffices to check that $\sqrt{d}$ and η are endomorphisms of J (see [**?**, Lemma 6]). If $\pi = c_1 + c_2\sqrt{d} + (c_3 + c_4\sqrt{d})\sqrt{\eta}$[1] then we have

$$2c_2\sqrt{d} = \pi + \bar{\pi} - 2c_1 \tag{9}$$

$$(4c_2(c_3^2 - c_4^2))\eta = (2c_2c_3 - c_4(\pi + \bar{\pi} - 2c_1))(\pi - \bar{\pi}). \tag{10}$$

Hence, for a fixed $\ell > 2$ dividing the index $[\mathcal{O}_K : \mathbb{Z}[\pi, \bar{\pi}]]$, we need to consider an extension field over which $J[\ell^u]$ is defined, where u is the ℓ-adic valuation of the index. Meanwhile, Algorithm 1 performs computations over the smallest extension field containing the ℓ-torsion points, whose degree r is smaller than ℓ^3, by Proposition 1.

Notation. We denote by r the degree of the smallest extension field $\mathbb{F}_{q^r}$ such that the ℓ-torsion is $\mathbb{F}_{q^r}$-rational.

We suppose that $\pi^r - 1$ is exactly divisible by ℓ^n. First, we need to compute a basis for the ℓ^n-torsion. We assume that the zeta function of $J/\mathbb{F}_q$ and the factorization $\#J(\mathbb{F}_{q^r}) = \ell^s m$ are known in advance. In order to compute the generators of $J[\ell^n]$, we use Freeman and Lauter's probabilistic algorithm [6], which needs $O(rM(r)\log q)$ operations in $\mathbb{F}_q$. We then compute a symplectic basis of $J[\ell^n]$, by using an algorithm similar to Gram–Schmidt orthogonalization. In order to compute $k_{\ell,J}$ we need to compute all subgroups in $\mathcal{W}$. If $\{Q_1, Q_{-1}, Q_2, Q_{-2}\}$ is a symplectic basis for $J[\ell^n]$, then a subgroup of rank 2 generated by $\lambda_1 Q_1 + \lambda_{-1}Q_{-1} + \lambda_2 Q_2 + \lambda_{-2}Q_{-2}$ and $\lambda'_1 Q_1 + \lambda'_{-1}Q_{-1} + \lambda'_2 Q_2 + \lambda'_{-2}Q_{-2}$, with $\lambda_i, \lambda'_j \in \mathbb{F}_\ell$, $i, j \in \{-2, -1, 1, 2\}$, is maximal isotropic with respect to the Weil pairing if the following equation is satisfied

$$(\lambda_1\lambda'_{-1} - \lambda_{-1}\lambda'_1) + (\lambda_2\lambda'_{-2} - \lambda_{-2}\lambda_2) = 0. \tag{11}$$

For each subgroup in $\mathcal{W}$ we store the values of the λ_i, λ'_js. There are $\ell^3 + \ell^2 + \ell + 1$ subgroups in $\mathcal{W}$ (see [4, Lemma 6.1]). In order to compute the values of the Tate pairing on the 2 generators of G, we use the values of the Tate pairing $T_{\ell^n}(Q_i, Q_j)$ for $i, j \in \{1, -1, 2, -2\}$, and the bilinearity of the Tate pairing. The values $T_{\ell^n}(Q_i, Q_j)$ are computed once and for all in the beginning and this costs $O(rM(r)\log q)$ operations in $\mathbb{F}_q$, assuming the final exponentiation is the dominating part of the pairing computation. Finally, computing pairings for all subgroups in $\mathcal{W}$ and the value of $k_{\ell,J}$ costs $O(\ell^3 M(r))$ operations in $\mathbb{F}_q$. We conclude that the cost of Algorithm 1 is $O(M(r)(r\log q + \ell^3))$. The complexity of Freeman and Lauter's algorithm for endomorphism ring computation is dominated by the cost of computing the ℓ-Sylow group of the Jacobian defined over the extension field containing the ℓ^u-torsion, whose degree is $r + \ell^{u-r}$ (by Proposition 2). The costs of the two algorithms are given in Table 1.

[1] Note that we cannot always write π in this form, but if this is not case, we can always replace π by $2^s\pi$, for some $s \in \mathbb{Z}$.

Table 1. Cost for checking locally maximal endomorphism rings at ℓ

Freeman and Lauter	This work (Algorithm 1)
$O((r+\ell^{u-r})M(r+\ell^{u-r})\log q)$	$O(M(r)(r\log q+\ell^3))$

Computing horizontal isogenies. Both classical algorithms and our algorithm need to compute first a basis for the ℓ-torsion. As stated before, this costs $O(rM(r)\log q\ell^{s-4n}(-\log\epsilon))$. The classical algorithm (see [2, Algorithm VI.3.4]) computes subspaces which are invariant under the action of Frobenius. More precisely, this algorithm needs to compute the matrix of the Frobenius endomorphism (in $O(\ell^2)$ operations in $\mathbb{F}_{q^r}$ using a baby-step giant-step approach). We conclude that the overall complexity of this algorithm is $O(M(r)(r\log q\ell^{s-4d}+\ell^2))$. Algorithm 2 computes a symplectic basis of the ℓ^n-torsion, all subgroups in $\mathcal{G}/\sim$ (in and searches among these subgroups those which have degenerate Tate pairing. It has the same cost as Algorithm 1.

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