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An improved lower bound for $(1, \leq 2)$ -identifying codes in the king grid*

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Abstract

We call a subset C of vertices of a graph G a $(1, \leq \ell)$ -identifying code if for all subsets X of vertices with size at most ℓ , the sets $\{c \in C \mid \exists u \in X, d(u, c) \leq 1\}$ are distinct. The concept of identifying codes was introduced in 1998 by Karpovsky, Chakrabarty and Levitin. Identifying codes have been studied in various grids. In particular, it has been shown that there exists a $(1, \leq 2)$ -identifying code in the king grid with density $\frac{3}{7}$ and that there are no such identifying codes with density smaller than $\frac{5}{12}$. Using a suitable frame and a discharging procedure, we improve the lower bound by showing that any $(1, \leq 2)$ -identifying code of the king grid has density at least $\frac{47}{111}$.

Keywords: Identifying codes, king grid, density, discharging method

1 Introduction

Let $G = (V, E)$ be a simple undirected graph with vertex set V and edge set E . The closed neighbourhood of a vertex $v \in V$, which consists of the vertex itself and all the adjacent vertices, is denoted by $N[v]$. In addition, we write $N[X] = \bigcup_{v \in X} N[v]$ for $X \subseteq V$. In [8], a subset $C \subseteq V$ is called a $(1, \leq \ell)$ -*identifying code* if

$$N[X] \cap C \neq N[Y] \cap C$$

for any two distinct subsets $X \subseteq V$ and $Y \subseteq V$ of size at most ℓ (one of them can be the empty set). The elements of a code C are called *codewords*.

Identifying codes were introduced by Karpovsky, Chakrabarty and Levitin in [7] and can be applied to locate objects in sensor networks [14]. A network is modelled by a graph and a sensor can check its closed neighbourhood. It gives an alarm if it detects at least one of the sought objects there. Suppose we have a $(1, \leq \ell)$ -identifying code C in the graph, and we place the sensors to the vertices corresponding to the codewords of C . Then, knowing the set of alarming sensors $A \subseteq C$, we can determine where the objects are (assuming that there are at most ℓ of them). Indeed, we have $A = N[X] \cap C$ for some subset X of size at most ℓ , and because these are unique, we can determine X — the set of vertices where the objects are.

Of course, we would like to use as few sensors as possible, so our aim is to find identifying codes with smallest possible cardinality. For infinite grids, we define below the measure *density* for this purpose.

Identifying codes have been considered, for example, in the following infinite grids: the square grid, the triangular grid, the king grid and the hexagonal mesh [1, 2, 3, 4, 5, 6, 9, 10, 11]. For more papers concerning the topic of identification, see [12]. In this paper, we focus on $(1, \leq 2)$ -identifying codes in the king grid. The king grid has vertex set \mathbb{Z}^2 and two vertices are adjacent if and only if the Euclidean distance between them is at most $\sqrt{2}$. Hence, the closed neighbourhood of a vertex consists of nine vertices (see Figure 1).

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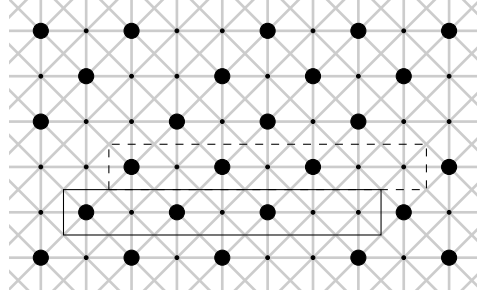


Figure 1: The infinite king grid with a $(1, \leq 2)$ -identifying code of density $\frac{3}{7}$.

In [8], a construction of a $(1, \leq 2)$ -identifying code in the infinite king grid with density $\frac{3}{7}$ is given (see Figure 1). Recently, in [13], a lower bound of $\frac{5}{12}$ was obtained. In this paper, we show that our approach which utilizes certain *frames*, improves this lower bound to $\frac{47}{111}$.

2 Preliminaries

We call *code* any subset of \mathbb{Z}^2 . As usual, the *density* $D(C)$ of a code $C \subseteq \mathbb{Z}^2$ is defined by:

$$D(C) = \limsup_{n \rightarrow \infty} \frac{|C \cap Q_n|}{|Q_n|}$$

where $Q_n = \{(x, y) \in \mathbb{Z}^2 \mid |x| \leq n, |y| \leq n\}$.

The next theorem gives the previously known bounds for the optimal density of a $(1, \leq 2)$ -identifying code in the king grid.

Theorem 1 ([5, 13]). *Let d be the optimal density of a $(1, \leq 2)$ -identifying code in the king grid. Then $0.41\overline{6} = \frac{5}{12} \leq d \leq \frac{3}{7} = 0.4286\dots$*

To prove bounds on the density of a code C , one method is to use a finite subset X of \mathbb{Z}^2 . For any vertex $v \in \mathbb{Z}^2$, we denote by $v + X$ or $X + v$ the set $\{v + u \mid u \in X\}$, and by $n(X, C)$ the average number of codewords of C in the translations of X : $n(X, C) = \limsup_{n \rightarrow \infty} \frac{\sum_{v \in Q_n} |(v + X) \cap C|}{|Q_n|}$. If we have some knowledge about $n(X, C)$, we can get results on the density of C with the following proposition.

Proposition 2. *Let C be a code of \mathbb{Z}^2 and X a nonempty finite subset of \mathbb{Z}^2 , then:*

$$D(C) = \frac{n(X, C)}{|X|}.$$

Proof. Let h be an integer such that $X \subseteq Q_h$. Let n be a positive integer. A vertex of $(v + X) \cap C$, for $v \in Q_{n+h}$ is either in $C \cap Q_n$, or in $Q_{n+2h} \setminus Q_n$. Among all the sets $(v + X) \cap C$, with $v \in Q_{n+h}$, each vertex of $C \cap Q_n$ is counted exactly $|X|$ times and each vertex of $Q_{n+2h} \setminus Q_n$ is counted at most $|X|$ times. Hence we have:

$$|X| |C \cap Q_n| \leq \sum_{v \in Q_{n+h}} |(v + X) \cap C| \leq |X| (|C \cap Q_n| + |Q_{n+2h} \setminus Q_n|).$$

By dividing every term by $|Q_n|$, we obtain:

$$|X| \cdot \frac{|C \cap Q_n|}{|Q_n|} \leq \frac{\sum_{v \in Q_{n+h}} |(v + X) \cap C|}{|Q_n|} \leq |X| \cdot \frac{|C \cap Q_n|}{|Q_n|} + |X| \cdot \frac{|Q_{n+2h} \setminus Q_n|}{|Q_n|}.$$

Since h is a fixed integer, we have: $\limsup_{n \rightarrow \infty} \frac{|Q_{n+2h} \setminus Q_n|}{|Q_n|} = 0$. As $n \rightarrow \infty$ we obtain:

$$\limsup_{n \rightarrow \infty} \frac{\sum_{v \in Q_{n+h}} |(v+X) \cap C|}{|Q_n|} = |X| \cdot D(C).$$

To end the proof, one can easily verify that $\limsup_{n \rightarrow \infty} \frac{\sum_{v \in Q_{n+h}} |(v+X) \cap C|}{|Q_n|} = n(X, C)$.

□

The identifying code problem can be seen as a covering problem. Indeed, a code C is a $(1, \leq \ell)$ -identifying code if and only if, for any two distinct subsets of vertices, X and Y , of size at most ℓ , the symmetric difference $N[X] \Delta N[Y]$ is covered by at least one element of C . The following theorem, from [5], reduces (in the case of the king grid) the family of sets that a code has to cover to be a $(1, \leq 2)$ -identifying code.

Theorem 3 ([5], Theorem 2.1). *A code C is a $(1, \leq 2)$ -identifying code in the king grid if and only if, for all $(x, y) \in \mathbb{Z}^2$, the sets*

1. $\{(x, y), (x+3, y), (x, y+3), (x+3, y+3)\},$
2. $\{(x, y), (x+1, y), (x+2, y)\}$
3. $\{(x, y), (x, y+1), (x, y+2)\}$

each contain at least one codeword of C .

Note that Theorem 3 together with Proposition 2 directly implies that the density of a $(1, \leq 2)$ -identifying code of the king grid is at least $\frac{1}{3}$. To improve this lower bound (and the one of Theorem 1), we will use, as the set X of Proposition 2, the set of Figure 2 (without loss of generality, one can set the left-lower corner of this set to be the origin). We call *frame* any set isomorphic to the set of Figure 2. The aim of this paper is to give a lower bound on the average number of codewords in a frame. More precisely, we will use a discharging procedure to show that the average number of codewords in a frame is at least $5 + \frac{3}{37}$. Hence, we will need to consider the *neighbourhood* of a frame. There is a natural bijection between all the frames and the set \mathbb{Z}^2 . We will also consider the *frames lattice* over the set of all frames, where the distance between two frames will be the distance in the king grid between the two corresponding vertices. As an example, the *2-ball* of a frame F is the set of frames

$$\{F + (x, y) \mid (x, y) \in \mathbb{Z}^2, \max\{|x|, |y|\} \leq 2\}.$$

We will often study a frame together with all the vertices at distance at most two of some vertex of the frame. Therefore, we will use the notation \mathbf{Xy} with $\mathbf{X} \in \mathbf{A-H}$ and $\mathbf{y} \in \mathbf{a-h}$, for a vertex on line \mathbf{X} and row \mathbf{y} with respect to the coordinates of Figure 2. The four vertices in positions **Cc**, **Cf**, **Fc** and **Ff** are called the *corners* of the frame. A *side* of a frame is a set of four vertices of the frame lying in the same line or column.

The conditions of Theorem 3 can be reformulated using our terminology. A code C is a $(1, \leq 2)$ -identifying code in the king grid, if and only if for each frame F :

Condition 1: At least one corner of F is a codeword of C .

Condition 2: Each set of three consecutive vertices on a side of F contains at least one codeword of C .

In the following, we assume that C is a $(1, \leq 2)$ -identifying code of the king grid. Given an integer k , a k -*frame* (resp. k^+ -*frame*) is a frame containing exactly k (resp. at least k) vertices of C .

Lemma 4. *Let F be a frame, then F is a 4^+ -frame. If F is a 4-frame, then all its codewords are in the corners of F .*

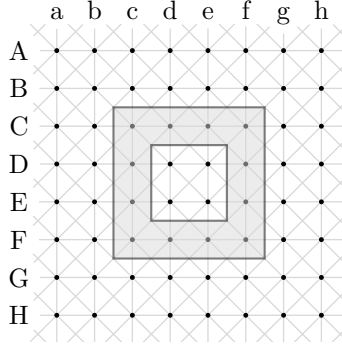


Figure 2: Notation for vertices in the neighbourhood of a frame.

Proof. It is easily observed that any frame F can be partitioned into four sets of three consecutive vertices on a line or column. Hence, by Condition 2, F contains at least four codewords.

Assume that F contains exactly four codewords. By Condition 1, one of the codewords must be in a corner, say in position **Cc**. Then, by Condition 2, there must be another codeword on line **C**. If it was on column **d** or **e**, one could partition the remaining vertices into three sets of three consecutive vertices on a line or column and F would be a 5^+ -frame, a contradiction. So there is a codeword in position **Cf**. Using similar arguments, the two other codewords must be in positions **Fc** and **Ff**. \square

Observation 5. *If a frame has three non-corner codewords on two of its sides, then it is a 6^+ -frame.*

Proof. Assume first that the three non-corner codewords lie on opposite sides. Without loss of generality, we can assume there are in positions **Cd**, **Ce** and **Fd**. By Condition 1, there must be one codeword in a corner. Assume it is in column **c**. Then, by Condition 2, there must be another codeword in column **c** and one in column **f**, leading to a 6^+ -frame. If the corner codeword is in column **f**, then there would be two codewords in column **f** and one in column **c**, leading again to a 6^+ -frame.

We assume now that the three non-corner codewords are on adjacent sides. Without loss of generality, we can assume they are not in the set S composed by line **C** and column **f**. By Condition 2, there are at least two codewords in S . By Condition 1, there is a codeword in a corner. If there is a codeword in position **Fc** then F is a 6^+ -frame. Otherwise, there is a corner codeword in S and then there must be three codewords in S , leading to a 6^+ -frame. \square

Observation 6. *If a frame F has a corner codeword c and two non-corner codewords on one of the sides of F that c does not belong to, then F is a 6^+ -frame.*

Proof. Without loss of generality, we assume that there are three codewords in positions **Cc**, **Fd** and **Fe**. By Condition 2, there must be another codeword in column **c**, another codeword in line **C** and one among positions **Df**, **Ef** and **Ff**. Hence, F is a 6^+ -frame. \square

Lemma 7. *Let F be a 4-frame. Then the four frames $F + \{(0, 2), (0, -2), (2, 0), (-2, 0)\}$ are 6^+ -frames.*

Proof. By Lemma 4, the four codewords of F are its four corners. By Condition 1 applied on $F + (1, 0)$, there is a codeword on either position **Cg** or position **Fg**. Hence by Observation 5, $F + (2, 0)$ is a 6^+ -frame. The claim is obtained by symmetry. \square

Lemma 7 gives a short proof of a result first given in [13]. The proof is a warm-up for our main result.

Theorem 8 ([13]). *The density of any $(1, \leq 2)$ -identifying code of the king grid is at least $\frac{5}{12}$.*

Proof. Let C be a $(1, \leq 2)$ -identifying code of the king grid. We use the discharging method to show that the average number of codewords in a frame is at least 5. Since the size of each frame is 12, the result will follow by Proposition 2.

In the beginning, each k -frame has charge k . Then, each 6^+ -frame F gives charge $\frac{1}{4}$ to each 4-frame among frames $F + \{(0, 2), (0, -2), (2, 0), (-2, 0)\}$. By Lemma 4, there exist only 4^+ -frames. By Lemma 7, each 4-frame receives charge 1 and each 6^+ -frame gives at most charge 1 away. So after the discharging process, each frame has at least charge 5. Each frame only gives charge to vertices at distance at most 2, hence the average number of codewords of C in a frame is at least 5 and we are done. \square

In the following, we will prove our main result:

Theorem 9. *The density of any $(1, \leq 2)$ -identifying code of the king grid is at least $\frac{47}{111} = 0.\overline{423}$.*

To prove this theorem, we will use a similar technique than the one of the proof of Theorem 8, but giving a final charge of more than 5 to each frame. Hence, a charge of 5 can be seen as the reference value for the charge of a frame. We say that the *charge excess* of a k -frame is $k - 5$. The *charge excess* within a subset S of frames is the sum of the charge excesses of all the frames of S .

3 Structural properties of code C

We now prove some results on the structure of C in the viewpoint of the frames. We call 4-benefactor a 6^+ -frame F having a 4-frame among frames $F + \{(0, 2), (0, -2), (2, 0), (-2, 0)\}$.

Lemma 10. *Let F be a 4-benefactor. Then F has 6^+ -frames in each corner position of its 2-ball being at distance 2 from the 4-frames among $F + \{(0, 2), (0, -2), (2, 0), (-2, 0)\}$. Moreover, if F is a 6-frame, there is a unique 4-frame among frames $F + \{(0, 2), (0, -2), (2, 0), (-2, 0)\}$. If F is a 7^+ -frame, there are at most two 4-frames among frames $F + \{(0, 2), (0, -2), (2, 0), (-2, 0)\}$.*

Proof. The first part of the claim is a direct consequence of Lemma 7.

There is at least one 4-frame among $F + \{(0, 2), (0, -2), (2, 0), (-2, 0)\}$, without loss of generality, we can assume that $F + (-2, 0)$ is a 4-frame. Then $F + (2, 0)$ cannot be a 4-frame since then, no corner of F is a codeword, contradicting Condition 1 on F . Suppose now that there is another 4-frame among $F + \{(0, 2), (0, -2)\}$, say $F + (0, 2)$. As before, $F + (0, -2)$ cannot be a 4-frame. By Condition 1 applied on F , there is a codeword in position \mathbf{Ff} . By Condition 1 applied on $F + (-1, 0)$ and $F + (0, 1)$, there are two codewords among positions \mathbf{Ce} , \mathbf{Fe} , \mathbf{Ec} and \mathbf{Ef} . Hence F is a 7^+ -frame and we are done. \square

6^+	X	X	X	X
X	Y	Z	Z	X
4	X	F	Z	X
X	Y	Z	Z	X
6^+	X	X	X	X

Figure 3: 2-ball around a 4-benefactor 6-frame.

The 6^+ -frames in the corner positions of a 4-benefactor 6^+ -frame described in the previous lemma are called the *co-benefactors* of F . The next lemma is valid for other cases by symmetry.

Lemma 11. *Let F be a 4-benefactor 6-frame oriented as in Figure 3. Then:*

- *Frame F has at least two 6^+ -frames F_1, F_2 in its 2-ball (in addition to its co-benefactors).*

- If F has no 6^+ -frame in a Y - or Z -position, then F has at least three 6^+ -frames in its 2-ball (in addition to its co-benefactors), or one of its co-benefactors is a 7^+ -frame and $F + (-1, 0)$ is not a 6^+ -frame.

Proof. In this proof, we will often use Conditions 1 and 2 without explicitly referring to them.

By Condition 1 applied on $F + (-1, 0)$, there is a codeword in either **Ce** or **Fe**. Without loss of generality, we may assume that there is a codeword in position **Ce** (the other case follows by symmetry).

Since F is not a 7^+ -frame, there is no codeword in position **Fe** and only one codeword among positions **Dc** and **Ec**.

Assume there is a codeword in position **Ec** and no codeword in position **Dc**. Due to Condition 2, there is a codeword in **De**. Hence, by Observation 5, $F + (-1, 1)$ is a 6^+ -frame. If $F + (-1, 0)$ is a 6^+ -frame, we are done. Otherwise, $F + (2, 1)$ is a 6^+ -frame and we are also done.

Assume now that there is no codeword in position **Ec**. Then by Condition 2, there is a codeword in position **Dc**.

If there is a codeword in position **De**, $F + (-1, 0)$ is a 6^+ -frame and, by Observation 6, $F + (-1, 1)$ is a 6^+ -frame and we are done. Hence we may suppose that there is no codeword in position **De**. This implies that there is a codeword in position **Df**, and no codeword in position **Ef**.

If there is a codeword in position **Cf**, by Observation 5, $F + (0, 1)$ is a 6^+ -frame. Moreover, by Observation 6, $F + (1, 0)$ is a 6^+ -frame and we are done. So, we may assume there is no codeword in position **Cf**, which implies that there is a codeword in position **Ff**.

If there is a codeword in position **Gd**, by Observation 5, $F + (-1, -1)$ and $F + (1, -2)$ are 6^+ -frames and we are done. Hence we may assume that there is no codeword in position **Gd**.

If there is a codeword in position **Fg**, $F + (1, 0)$ is a 6^+ -frame and by Observation 5, $F + (2, 0)$ is a 6^+ -frame and we are done. So we may assume that there is no codeword in position **Fg**.

If there is a codeword in position **Db**, $F + (-1, 0)$ is a 6-frame and by Observation 5, $F + (-2, -1)$ is a 6^+ -frame. If $F + (-1, -1)$ is a 6^+ -frame we are done, otherwise, by Observation 5, $F + (0, -2)$ is a 6^+ -frame. Hence we may assume that there is no codeword in position **Db**.

If there is no codeword in position **Bf**, then there is a codeword in position **Bc**, and by Observation 5, $F + (-2, 1)$ is a 6^+ -frame. Then, there must be a codeword among positions **Bd** and **Be**, implying that $F + (-1, 1)$ is a 6^+ -frame. Hence we may assume that there is a codeword in position **Bf**.

If there is a codeword in position **Be**, by Observation 5, $F + (1, 1)$ and $F + (2, 2)$ are 6^+ -frames and we are done. Hence we assume that there is no codeword in position **Be**.

If there is no codeword in position **Bg**, by Observation 5, $F + (2, -1)$ and $F + (2, 2)$ are 6^+ -frames. Now, if there is a codeword in **Bd**, then by Observation 5, $F + (1, 2)$ is a 6^+ -frame and we are done. Otherwise, there is a codeword in **Bc** and **Ba**, and $F + (-2, 1)$ is a 6-frame, and we are done. So, assume that there is a codeword in position **Bg**.

If there is a codeword in position **Dg**, then $F + (1, 0)$ is a 6^+ -frame and by Observation 5, $F + (2, -1)$ is a 6^+ -frame. So we assume there is no codeword in position **Dg**, which implies that there is a codeword in position **Gg** by Condition 1.

If there is a codeword in position **Gf**, by Observation 5, $F + (0, -2)$ and $F + (2, -1)$ are 6^+ -frames. Now, if there is a codeword in position **Ge**, $F + (1, -1)$ is a 6^+ -frame and we are done. Otherwise, there is a codeword in position **Gb** and $F + (-2, -1)$ is a 6^+ -frame. So we assume that there is no codeword in position **Gf**, which implies that there is a codeword in position **Ge**.

If there is a codeword in position **Gb**, $F + (-1, -1)$ and $F + (-2, -1)$ are 6-frames and we are done. So we assume that there is no codeword in position **Gb**.

If there is a codeword in position **Bd**, by Observation 5, $F + (1, 2)$ is a 6^+ -frame. Moreover $F + (1, 1)$ is a 6^+ -frame so we are done. So we assume there is no codeword in position **Bd**. Then there is a codeword in both positions **Ba** and **Bc**, implying that $F + (-2, 1)$ is a 6-frame.

If there is a codeword in position **Cg**, $F + (1, 0)$ is a 6-frame and we are done. So, we assume there is no codeword in **Cg**. Then, by Observation 5, $F + (2, 1)$ is a 6-frame.

Now, if $F + (1, -2)$ is a 6^+ -frame we are done. Otherwise, there is no codeword in position **Hd** and exactly one codeword among positions **He** and **Hf**. Since $F + (0, -2)$ is a 5^+ -frame, there is a codeword in position **Hc**. By Condition 1, there is a codeword in position **Ha**. This implies that $F + (-2, -2)$ is a 7-frame and finishes the case analysis. \square

A frame is called *1-poor* (resp. *2-poor*) if it is a 5-frame having at most one 6-frame in its 2-ball (resp. no 6^+ -frame at distance 1 and at most two 6-frames at distance 2).

Lemma 12. *Let F be a 5-frame, then F has at least one 6^+ -frame in its 2-ball. Moreover, one of the following properties holds:*

- F has a 6-frame at distance 1 and another 6^+ -frame in its 2-ball,
- F has a total charge excess of at least 3 within its 2-ball, with at least two 6^+ -frames,
- F is 1-poor and, up to symmetry, the configuration around F is depicted on Figure 5,
- F is 2-poor and, up to symmetry, the configuration around F is depicted on Figure 6.

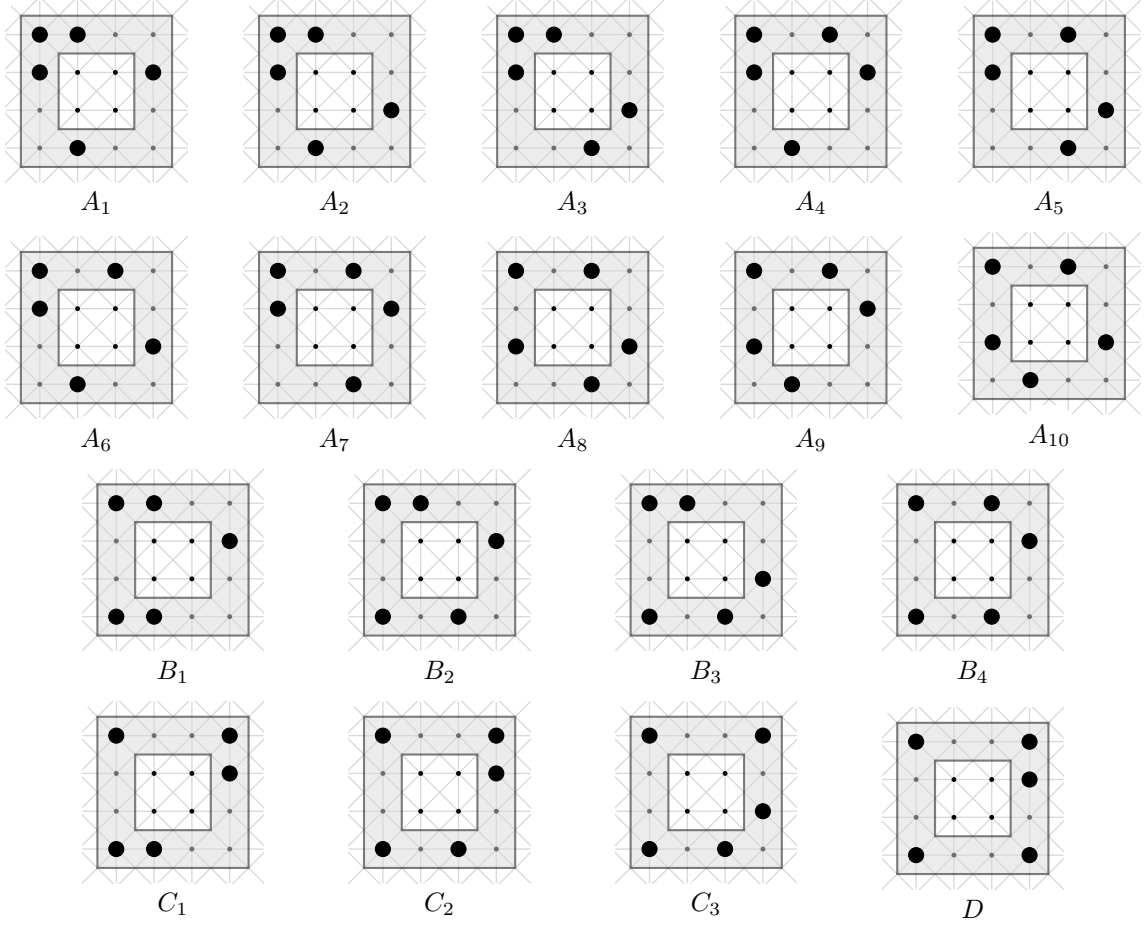


Figure 4: All the possibilities for a 5-frame.

Proof. Let F be 5-frame. We make a case analysis according to the configurations of F depicted on Figure 4. Note that cases A_1 to A_{10} represent cases where F has only one corner codeword, cases B_1 to B_4 are those cases where F has exactly two corner codewords on the same side, cases C_1 to C_3 are the cases where F has three corner codewords, and case D is the case where F has four corner codewords. The cases which are symmetric to those of Figure 4 will follow from the same arguments. Moreover, note that if F had exactly two corner codewords on opposite corners, by Condition 2, each side of F would contain an additional codeword and F would be a 6^+ -frame.

Case A_1 . By Observation 5, $F + (-1, 0)$ and $F + (0, 1)$ both are 6^+ -frames and we are done.

Case A_2 . By Observation 5, $F + (-1, 0)$ is a 6^+ -frame. By Observation 6, $F + (0, 1)$ is a 6^+ -frame and we are done.

Case A_3 . By Observation 6, $F + (-1, 0)$ and $F + (0, 1)$ both are 6^+ -frames and we are done.

Case A_4 . By Observation 5, $F + (0, 1)$ is a 6^+ -frame. Now, by Condition 2, there is a codeword on either position **Dd** or **Ed**. In either case, by Observation 5, $F + (-1, 2)$ or $F + (1, -1)$ respectively, is a 6^+ -frame and we are done.

Case A_5 . By Observation 6, $F + (0, 1)$ is a 6^+ -frame. Now, by Condition 2, there is a codeword on either position **Dd** or **De**. In either case, by Observation 5, $F + (-1, -1)$ or $F + (2, 1)$ respectively, is a 6^+ -frame and we are done.

Case A_6 . By Observation 6, $F + (0, 1)$ is a 6^+ -frame. Now, by Condition 2, there is a codeword in either position **Dd** or **Ed**. In both cases, by Observation 5, $F + (-1, -1)$ or $F + (1, -1)$ respectively, is a 6^+ -frame and we are done.

Case A_7 . By Observation 5, $F + (0, 1)$ is a 6^+ -frame. If $F + (-1, -1)$ is a 6^+ -frame, we are done. Otherwise, from the assumptions there are already four non-corner codewords in $F + (-1, -1)$, hence $F + (-1, -1)$ is a 5-frame and, by Condition 1, all its further non-corner positions contain no codeword. In particular, there is no codeword in both positions **Dd** and **Gd**. By Condition 1, there is a codeword on at least one of the positions **Dg** and **Gg**. Hence, by Observation 5, $F + (2, -1)$ is a 6^+ -frame and we are done.

Case A_8 . Assume there is no codeword in position **Dd**. Then, $F + (-1, 0)$ and $F + (2, 1)$ are 6^+ -frames and we are done. So, assume there is a codeword in position **Dd**.

If there is a codeword in position **Ed**, by Observation 5, $F + (-2, 0)$ and $F + (1, 0)$ are 6^+ -frames. Hence we may assume there is no codeword in position **Ed**, and by symmetry, no codeword in position **De** either.

If there is a codeword in position **Ee**, by Observation 5, $F + (1, 1)$ and $F + (-1, -1)$ are 6^+ -frames. Hence we may assume there is no codeword in position **Ee**.

If there is a codeword in position **Be**, by Observation 5, $F + (1, 1)$ and $F + (2, 2)$ are 6^+ -frames. Hence we may assume there is no codeword in position **Be**, and by symmetry, no codeword in position **Eb** either. By Condition 1, this implies that there is a codeword in position **Bb**.

If there is a codeword in position **eg**, by Observation 5, $F + (1, 0)$ and $F + (2, 1)$ are 6^+ -frames. Hence we may assume there is no codeword in position **eg**, and by symmetry, no codeword in position **Ge** either.

If there is no codeword in position **Bd**, by Observation 5, $F + (0, 2)$ and $F + (-2, 1)$ are 6^+ -frames. If $F + (1, 1)$ is a 6^+ -frame, we are done. Otherwise, by Observation 5, $F + (2, 1)$ is a 6^+ -frame and we are done too. Hence we may assume there is a codeword in position **Bd**, and by symmetry, another codeword in position **Db**.

If there is a codeword in position **Bc**, $F + (0, 1)$ is a 6-frame and by Observation 5, $F + (-1, 1)$ is a 6^+ -frame, so we are done. Hence we may assume there is no codeword in position **Bc**, and by symmetry, no codeword in position **Cb** either.

If there is a codeword in position **Gb**, $F + (-1, -1)$ and $F + (-1, -2)$ are 6^+ -frames. Hence we may assume there is no codeword in position **Gb**, and by symmetry, no codeword in position **Bg** either.

If there is no codeword in position **Gf**, there is, by Condition 1, a codeword in position **Gc** and then $F + (-1, -1)$ is a 6-frame. Now, if there is no codeword in position **Fg**, $F + (1, -1)$ is a 4-frame and by Lemma 7, $F + (1, 1)$ is a 6^+ -frame and we are done. Otherwise, by Observation 5, $F + (1, -2)$ is a 6^+ -frame. Hence we may assume that there is a codeword in position **Gf**, and by symmetry, another codeword in position **Fg**.

If there is a codeword in position **Gg**, $F + (1, -1)$ is a 6-frame and by Observation 5, $F + (1, -2)$ is a 6^+ -frame, so we are done. Hence we may assume there is no codeword in position **Gg**.

If there is a codeword in position **Gc**, $F + (0, -1)$ and $F + (-1, -1)$ is a 6^+ -frame, so we are done. Hence we may assume there is no codeword in positions **Gc** and **Cg**.

If $F + (-1, 2)$ is a 5-frame, by Condition 2, there is no codeword in position **Ae** and a codeword on exactly one of the positions **Ac** and **Ad**. Then there must be a codeword in position **Af** because

$F + (0, 2)$ has a fifth codeword. But then $F + (2, 2)$ is a 6^+ -frame. Hence, there is a 6^+ -frame among $F + (-1, 2)$ and $F + (2, 2)$. By symmetry, there is also a 6^+ -frame among $F + (-2, 1)$ and $F + (-2, -2)$.

If $F + (-1, -2)$ is a 6^+ -frame, we are done. Otherwise, there is no codeword in position **Hd**, which implies that there is a codeword in position **Hg**, as well as one codeword among positions **He** and **Hf**. By symmetry, if $F + (2, 1)$ is a 5 -frame, there is a codeword in position **Gh**, as well as a codeword among positions **eh** and **Fh**. Then, $F + (2, -2)$ is a 6^+ -frame and we are done.

Case A_9 . By Condition 2, there are at least two codewords among positions **Dd**, **De**, **Ed** and **Ee**. If there are two codewords on the same column or line, (resp. **d**, **D**, **e**, **E**), then by Observation 5, resp. $F + (1, 0)$, $F + (0, -1)$, $F + (-1, 0)$, $F + (0, 1)$, is a 6^+ -frame. Hence, if there are three codewords among positions **Dd**, **De**, **Ed** and **Ee**, there are two 6^+ -frames at distance 1 of F and we are done.

Therefore, we assume that there are exactly two codewords among positions **Dd**, **De**, **Ed** and **Ee**. Then by Condition 2, they cannot lie on a same line or column. If there are codewords in positions **De** and **Ed**, by Observation 5, $F + (1, -1)$ and $F + (-1, 1)$ are 6^+ -frames and we are done. Hence, we assume that there are codewords in positions **Dd** and **Ee**.

If there is a codeword in position **Eb**, by Observation 5, $F + (-2, 1)$ is a 6^+ -frame, and by Condition 2, $F + (-1, 1)$ is a 6^+ -frame as well. Hence we may assume there is no codeword in position **Eb**.

If there is a codeword in position **Cb**, by Observation 5, $F + (-2, 0)$ is a 6^+ -frame, and by Condition 2, $F + (-1, 0)$ is a 6^+ -frame as well. Hence we may assume there is no codeword in position **Cb**.

If there is a codeword in position **Gd**, $F + (1, -1)$ is necessarily a 6^+ -frame, and by Observation 5, $F + (1, -2)$ is a 6^+ -frame. Hence we may assume there is no codeword in position **Gd**.

By symmetry, we can assume that there is no codeword in any of the positions **Be**, **Bc** and **Dg**. Using Condition 2, this implies that there are codewords in positions **Db** and **Bd**.

If there is a codeword in position **Gg**, $F + (1, -1)$ is necessarily a 6^+ -frame, and by Observation 5, $F + (2, -1)$ is a 6^+ -frame. Hence we may assume there is no codeword in position **Gg**.

If there is a codeword in position **Ge**, $F + (1, -1)$ is necessarily a 6^+ -frame. Now, if there is a codeword in position **Fb**, $F + (-1, -1)$ is necessarily a 6 -frame too. Otherwise, either $F + (-2, -1)$, $F + (-1, 1)$ or $F + (-2, 2)$ is a 6^+ -frame, and we are done. Hence we may assume there is no codeword in position **Ge** and by symmetry, there is no codeword in position **Eg** either.

If there is a codeword in position **Fb** (resp. in position **Bf**), $F + (-1, 0)$ (resp. $F + (0, 1)$) is necessarily a 6 -frame. If there is a codeword in both positions, we are done. Hence by symmetry, without loss of generality, we can assume that there is no codeword in position **Fb**. This implies that $F + (-2, -1)$ is a 6^+ -frame. If there is a 6^+ -frame at distance 1 of F , we are done. Otherwise, $F + (1, 2)$ is a 6^+ -frame. If $F + (-2, 2)$ is a 6^+ -frame, we are done. Otherwise, $F + (-2, -1)$ and $F + (1, 2)$ are both necessarily 7^+ -frames and we are done too.

Case A_{10} . Assume that there are no codewords among positions **Ed** and **Ee**. Then, by Condition 2, there are codewords on both positions **Dd** and **De**, and by Observation 5, $F + (0, -1)$ and $F + (2, 1)$ are 6^+ -frames.

Now, assume that there is a codeword in position **Ed** or **Ee**. Then, by Condition 2, $F + (0, 1)$ is a 6^+ -frame. If there is a codeword in position **Ed**, by Observation 5, $F + (1, -1)$ is a 6^+ -frame. If there is a codeword in position **Ee**, by Observation 5, $F + (1, 1)$ is a 6^+ -frame. In both cases we are done.

Case B_1 . By Observation 5, $F + (-1, 0)$ is a 6^+ -frame. By Condition 1, there is a codeword among positions **Cb** and **Fb**. Then, by Observation 5, $F + (-2, 0)$ is a 6^+ -frame and we are done.

Case B_2 . By Observation 5, $F + (-1, 0)$ is a 6^+ -frame. By Condition 2, there is a codeword among positions **Dd** and **Ed**. If there is a codeword in position **Dd** (resp. **Ed** and not in **Dd**), by Observation 5, $F + (1, 1)$ (resp. $F + (1, -1)$) is a 6^+ -frame and we are done.

Case B_3 . By Observation 5, $F + (-1, 0)$ is a 6^+ -frame. If $F + (1, 1)$ is a 6^+ -frame, we are done. Otherwise, $F + (1, 1)$ is a 5-frame and it has a unique corner codeword in position **Ed**. By Condition 2, this implies that there is a codeword in position **Bc**, and then by Observation 5, $F + (0, 2)$ is a 6^+ -frame and we are done.

Case B_4 . Assume first that there are no codewords among positions **De** and **Ee**. Then by Condition 2, there must be a codeword on both positions **Dd** and **Ed**. Then by Observation 5, $F + (1, 0)$ and $F + (-2, 0)$ are 6^+ -frames.

If there is a codeword in position **De** or **Ee**, then $F + (-1, 0)$ is a 6^+ -frame. If there is a codeword in position **De** (resp. **Ee**), by Observation 5, $F + (1, -1)$ (resp. $F + (2, -1)$) is a 6^+ -frame and we are done.

Case C_1 . By Observation 5, both $F + (-1, 0)$ and $F + (0, 1)$ are 6^+ -frames.

Case C_2 . By Observation 5, $F + (0, 1)$ is a 6^+ -frame. By Condition 1, there is a codeword among positions **Bc** and **Bf**. Then, by Observation 5, $F + (0, 2)$ is a 6^+ -frame.

Case C_3 . Assume first that there is no codeword in position **De**. Then, there must be codewords in both positions **Dd** and **Ee**. Then, by Observation 5, $F + (1, 1)$ and $F + (-1, -1)$ are both 6^+ -frames. Hence we can assume that there is a codeword in position **De**, and, by symmetry in position **Ed**.

If there is a codeword in position **Dd**, by Observation 5, $F + (0, 2)$ and $F + (0, -1)$ are 6^+ -frames and we are done. Hence we may assume there is no codeword in position **Dd**.

If there is a codeword in position **Ee**, by Observation 5, $F + (0, 1)$ and $F + (2, 0)$ are 6^+ -frames and we are done. Hence we may assume there is no codeword in position **Ee**.

If there is a codeword in position **Bc** or **Bf**, by Observation 5, $F + (0, 2)$ is a 6^+ -frame and $F + (0, 1)$ must be a 6^+ -frame too, so we are done. Hence we may assume there is no codeword in positions **Bc** and **Bf** and, by symmetry, on positions **Cb** and **Fb**.

If there is a codeword in position **Be**, then $F + (0, 1)$ is a 6^+ -frame. Now, if either $F + (-1, 1)$ or $F + (-2, 2)$ is a 6^+ -frame, we are done. Otherwise, there is no codeword in position **Ca** and hence there is a codeword in position **Fa** by Condition 1. This implies that $F + (-2, -1)$ is a 6^+ -frame, so we are done. Hence we may assume there is no codeword in position **Be** and, by symmetry, in position **Eb**.

If there is a codeword in position **Ed**, by Observation 5, $F + (2, 1)$ and $F + (2, -2)$ are 6^+ -frames. If $F + (1, 0)$ is a 6^+ -frame, we are done. Otherwise, there is no codeword in position **Dg**, and a codeword among positions **Cg** and **Fg**. If there is a codeword in position **Fg**, by Observation 5, $F + (1, -1)$ is a 6^+ -frame. If there is a codeword in position **Cg**, by Observation 5, $F + (2, 0)$ is a 6^+ -frame. Hence we may assume there is no codeword in position **Ed** and, by symmetry, in position **Ge**.

If there is a codeword in position **Gc**, by Observation 5, $F + (0, -2)$ and $F + (-2, -1)$ are 6^+ -frames. If there are only two 6-frames and no additional 6^+ -frame in the 2-ball of F , F is 2-poor, then, using the same techniques than previously, one can check that the codewords around F are fixed as in Figure 6 and we are done. Otherwise, there is a total charge excess of at least 3 within the 2-ball of F and we are done too. Hence we may assume there is no codeword in position **Gc** and, by symmetry, in position **Cg**.

We note that $F + (1, -1)$ is a 6^+ -frame. If $F + (1, -1)$ is a 6-frame and if it is the only 6-frame within the 2-ball of F , then, one can check that F is 1-poor and that the codewords around F are fixed as in Figure 5, and so, we are done. Otherwise, either $F + (1, -1)$ is a 7^+ -frame and then $F + (1, -2)$ is a 6^+ -frame, or $F + (1, -1)$ is a 6-frame and there exists another 6^+ -frame within the 2-ball of F . In both cases we are done.

Case D . By Observation 5, $F + (0, 1)$ is a 6^+ -frame. By Condition 1, there is a codeword among positions **Bc** and **Bf**. Then, by Observation 5, $F + (0, 2)$ is a 6^+ -frame. \square

Lemma 13. *If F is a 1-poor frame, then F has a 6-frame F_6 at distance 1 such that:*

- F_6 has at most two 1-poor frames in its 2-ball and each of them is at distance 1 of F_6 ,

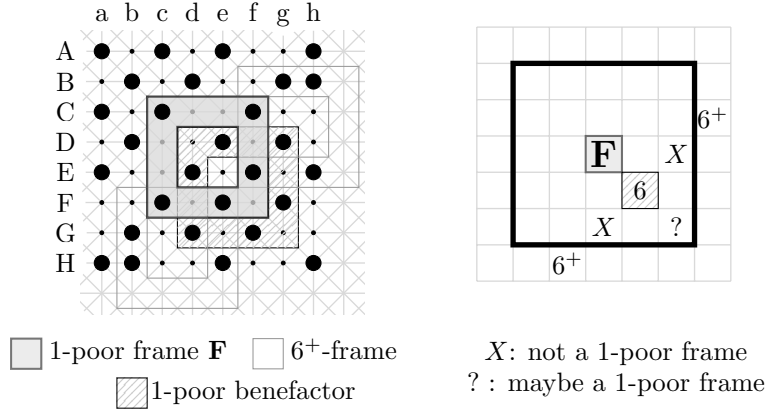


Figure 5: Neighbourhood of a 1-poor frame in the vertex lattice (on the left) and in the frame lattice (on the right).

- F_6 has at least two 6⁺-frames at distance 2,
- F_6 is not a 4-benefactor.

Proof. By Lemma 12, if F is a 1-poor frame, the neighbourhood of F is fixed and depicted in Figure 5 (up to symmetry). The only 6⁺-frame in the neighbourhood of F is $F_6 = F + (1, -1)$ and contains exactly six codewords. The first point of the lemma follows from the fact that a 1-poor frame must have a 6-frame in a corner of its 1-ball, however by using the left part of Figure 5 the positions marked by "X" on the right part of the figure cannot host a 1-poor frame. The second point follows using Observation 5 on the left part of Figure 5. The last point follows by noting that a 4-benefactor must have two non-codeword corners in the same column or line, which is not the case here. \square

A frame playing the role of F_6 in the previous lemma is called a 1-poor-benefactor.

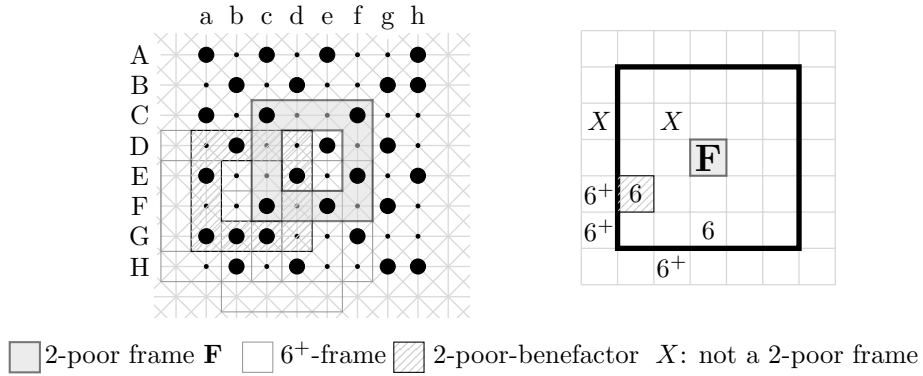


Figure 6: Neighbourhood of a 2-poor frame in the vertex lattice (on the left) and in the frame lattice (on the right).

Lemma 14. *If F is a 2-poor frame, then F has two 6-frames at distance 2. One of them, say F_6 , has the following properties:*

- F_6 is neither a 4-benefactor nor a 1-poor-benefactor,
- F_6 has at least four 6⁺ frames in its 2-ball,

- F_6 has only one 2-poor frame in its 2-ball.

Proof. By Lemma 12, if F is a 2-poor frame, the neighbourhood of F is fixed and depicted in Figure 6 (up to symmetry). Let $F_6 = F + (-2, -1)$. The first point of the lemma follows from the fact that, by Lemma 12, a 1-poor benefactor must have a 6-frame in its 2-ball (and only this 6-frame in its 2-ball) and that a 4-benefactor must have two non-codeword corners in the same line or column (hence, the only possibility would be that $F + (-2, 1)$ is a 4-frame which is not true). The second point follows by using Observation 5 on the left part of Figure 6. For the third point, observe that the only possible locations for another 2-poor frame in the 2-ball of F_6 are those marked by an "X" on the right part of Figure 6: $F + (-3, 1)$ and $F + (-1, 1)$. However, both have two 6^+ -frames in their 2-ball that do not match the codeword configuration of a 2-poor frame as given on the left part of Figure 6. \square

A frame playing the role of F_6 in the previous lemma is called a *2-poor-benefactor*.

4 The discharging procedure

We are now ready to describe the discharging procedure which leads to our result. In the beginning, each k -frame has charge k . Let $\alpha = \frac{1}{37}$. We apply the following rules:

1. Let F be a 4-benefactor 6-frame without any 7^+ -frame among its co-benefactors. Assume F is oriented as in Figure 3 (other cases follow by symmetry). We consider two subrules:
 - (a) If there is a 6^+ -frame in a Y -position of Figure 3, say in position $F + (-1, 1)$ (the other case will be covered by symmetry), and no other 6^+ -frames in Y - and Z -positions of Figure 3, then F gives charge $\frac{3\alpha+1}{4}$ to the 4-frame $F + (-2, 0)$, charge α to frame $F + (0, 1)$, charge 2α to the other 5-frames in Z -positions of Figure 3, charge $\frac{3\alpha}{2}$ to the other 5-frame in Y -position of Figure 3 and charge α to all the other 5-frames in the 2-ball of F .
 - (b) Otherwise, F gives charge $\frac{3\alpha+1}{4}$ to the 4-frame $F + (-2, 0)$, charge 2α to 5-frames in Z -positions of Figure 3, charge $\frac{3\alpha}{2}$ to 5-frames in Y -positions of Figure 3 and charge α to all the other 5-frames in the 2-ball of F .
2. Let F be a 4-benefactor 6-frame with a 7^+ -frame among its co-benefactors. Assume F is oriented as in Figure 3, and that $F + (-2, 2)$ is a 7^+ -frame (other cases follow by symmetry). Then F gives the same charges than in Rule 1b, except for the frame $F + (-1, 1)$ which receives no charge from F .
3. A 4-benefactor 7^+ -frame F gives charge $\frac{3\alpha+1}{4}$ to the 4-frames among $F + \{(0, 2), (0, -2), (2, 0), (-2, 0)\}$, charge 3α to frames at distance 1 and charge 2α to frames at distance 2.
4. A 1-poor-benefactor, recall that it is a 6-frame, gives charge 3α to the 1-poor frames in its 2-ball, 2α to the other 5-frames at distance 1 and charge α to the other 5-frames at distance 2.
5. A 2-poor-benefactor, recall that it is a 6-frame, gives charge 2α to the unique 2-poor frame in its 2-ball, 2α to the other 5-frames at distance 1 and charge α to the other 5-frames at distance 2.
6. Other 6-frames give charge 2α to 5-frames at distance 1 and charge α to 5-frames at distance 2.
7. Other 7^+ -frames give charge 3α to 5-frames at distance 1 and charge 2α to 5-frames at distance 2.

We note that the rules are not ambiguous. Indeed, by Lemmas 13 and 14, a 6-frame can be either a 4-benefactor, a 1-poor benefactor, a 2-poor benefactor or not a benefactor at all, but never two at the same time.

Lemma 15. *After the application of the discharging rules, each 7^+ -frame has charge at least $5 + 3\alpha$.*

Proof. Let F be a 7^+ -frame. It is sufficient to show that F gives at most charge $2 - 3\alpha = \frac{71}{37}$.

If F is a 4-benefactor, then by Lemma 10, it has at most two 4-frames among the frames $F + \{(0, 2), (0, -2), (2, 0), (-2, 0)\}$.

If F has two 4-frames among $F + \{(0, 2), (0, -2), (2, 0), (-2, 0)\}$ then it has at least three co-benefactors. Hence, F has at most eleven 5-frames at distance 2. By Rule 3, F gives at most charge $2 \cdot \frac{3\alpha+1}{4} + 8 \cdot 3\alpha + 11 \cdot 2\alpha = \frac{66}{37}$.

If F has only one 4-frame among $F + \{(0, 2), (0, -2), (2, 0), (-2, 0)\}$, then it has two co-benefactors and hence, at most thirteen 5-frames at distance 2. By Rule 3, F gives at most charge $\frac{3\alpha+1}{4} + 8 \cdot 3\alpha + 13 \cdot 2\alpha = \frac{60}{37}$.

Finally, if F is not a 4-benefactor, then by Rule 7, it gives at most charge $8 \cdot 3\alpha + 16 \cdot 2\alpha = \frac{56}{37}$. \square

Lemma 16. *After the application of the discharging rules, each 6-frame has charge at least $5 + 3\alpha$.*

Proof. Let F be a 6-frame. It is sufficient to show that F gives at most charge $1 - 3\alpha = \frac{34}{37}$.

Assume first that F is a 4-benefactor without any 7^+ -frame among its co-benefactors. If there were no other 6^+ -frames than the co-benefactors of F within its 2-ball, F would give, according to Rule 1, charge $\frac{3\alpha+1}{4} + 5 \cdot 2\alpha + 2 \cdot \frac{3\alpha}{2} + 14\alpha = \frac{37}{37}$. However, by Lemma 11, F has at least two non-co-benefactor 6^+ -frames in its 2-ball. Moreover by Lemma 11 as well, either one of those is on a Y - or Z -position, or there is a third non-co-benefactor 6^+ -frame in the 2-ball of F . We distinguish two cases. If there is an extra 6^+ -frame on a Y -position and one on an X -position (Rule 1a applies), then we save charge at least $\frac{3\alpha}{2} + \alpha$ from the extra frames and charge α from frame $F + (0, 1)$. Otherwise (Rule 1b applies), the total charge saved on the extra 6^+ -frames is at least 3α . Indeed, we save at least $2 \cdot \frac{3\alpha}{2}$ if the two extra 6^+ -frames are both on Y -positions, $2 \cdot 2\alpha$ if they are both on Z -positions, $2\alpha + \frac{3\alpha}{2}$ for a Z -position and a Y -position, $2\alpha + \alpha$ for a Z -position and an X -position, and $3 \cdot \alpha$ if there are three extra 6^+ -frames. In all cases, F gives at most $\frac{34}{37}$ of charge.

If F is a 4-benefactor with a 7^+ -frame among its co-benefactors, by Lemma 11, there are two 6^+ -frames in the 2-ball of F in addition to its two co-benefactors and they are at distance 2 from F . Hence, F has at most eleven 5-frames at distance 2. By Rule 2, F gives at most $\frac{3\alpha+1}{4} + 5 \cdot 2\alpha + \frac{3\alpha}{2} + 12\alpha = \frac{33.5}{37}$.

If F is a 1-poor-benefactor, by Lemma 13, F has at most two 1-poor frames in its 2-ball and each of them is at distance 1 of F . Moreover, there are at least two 6^+ -frames at distance 2 of F . We may assume that F has two 1-poor frames in its 2-ball since by Rule 4, F would give away more charge in this case. Now, by Rule 4, F gives at most $2 \cdot 3\alpha + 6 \cdot 2\alpha + 14 \cdot \alpha = \frac{32}{37}$.

If F is a 2-poor-benefactor, by Lemma 14, F has only one 2-poor frame and at least four 6^+ -frames in its 2-ball. We may assume that these frames are at distance 2 of F since by Rule 5, F would give away more charge in this case. By Rule 5, F gives at most $2\alpha + 8 \cdot 2\alpha + 11 \cdot \alpha = \frac{29}{37}$.

Finally, if F is not a benefactor, then by Rule 6, it gives at most charge $8 \cdot 2\alpha + 16 \cdot \alpha = \frac{32}{37}$. \square

Lemma 17. *After the application of the discharging rules, each 5-frame has charge at least $5 + 3\alpha$.*

Proof. It is enough to prove that each 5-frame receives charge 3α .

We first note that by our discharging rules, each 6-frame F_6 gives at least charge α to each 5-frame of its 2-ball, except in Rule 2 where F_6 does not give anything to one 5-frame F' at distance 1 of F_6 . However, F' has a 4-benefactor 7^+ -frame at distance 1, which, by Rule 3, gives charge 3α to F' .

Let F be a 5-frame. By the previous paragraph, we can consider that F receives at least charge α from each 6-frame in its 2-ball. By Rules 3 and 7, F receives at least charge 2α from each 7^+ -frame in its 2-ball. Hence, if F has a total charge excess of at least 3 within its 2-ball, with two 6^+ -frames, F receives at least charge 3α and we are done. Otherwise, by Lemma 12, F is either 1-poor, 2-poor or has a 6-frame at distance 1 and another one in its 2-ball.

If F is 1-poor, by Lemma 13, it has a 1-poor-benefactor in its 2-ball which, by Rule 4, gives charge 3α to F .

If F is 2-poor, by Lemma 14, it has a 2-poor-benefactor in its 2-ball which, by Rule 5, gives charge 2α to F . Moreover, by Lemma 14, F has another 6^+ -frame in its 2-ball which gives charge α to F .

Finally, suppose F has a 6-frame F_6 at distance 1 and another one, F'_6 , in its 2-ball. If F receives charge 2α from F_6 , we are done. Otherwise, by our discharging rules F_6 is necessarily a

4-benefactor. Without loss of generality, we can assume that F_6 is oriented as in Figure 3. Then F_6 gives always charge 2α to the 5-frames in Z -positions, except in Rule 1a where $F_6 + (0, 1)$ receives only charge α , but in this case, $F_6 + (0, 1)$ has three 6^+ -frames in its 2-ball, so $F \neq F_6 + (0, 1)$. Frame $F_6 + (-1, 0)$ has three 6^+ -frames in its 2-ball, so $F \neq F_6 + (-1, 0)$. Thus, we can assume that F is on a Y -position. Hence, F'_6 is also a 4-benefactor. In this case, by Rules 1 and 2, both F_6 and F'_6 give charge $\frac{3\alpha}{2}$ to F , and we are done. \square

Lemma 18. *After the application of the discharging rules, each 4-frame has charge at least $5 + 3\alpha$.*

Proof. By Lemma 7, a 4-frame F has four 4-benefactors. By Rules 1, 2 and 3, each 4-benefactor gives charge $\frac{3\alpha+1}{4}$ to F . Hence, F receives charge $1 + 3\alpha$ and ends with charge $5 + 3\alpha$. \square

After the application of our discharging rules, by Lemmas 15, 16, 17 and 18, each frame has charge at least $5 + 3\alpha = \frac{188}{37}$, therefore, the average number of codewords in each frame is at least $\frac{188}{37}$ (this is due to the fact that each frame gives charge to vertices at distance at most 2). There are twelve vertices in each frame, hence, by Proposition 2, we obtain Theorem 9.

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