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# ASYMPTOTIC PROPERTIES OF OPTIMAL TRAJECTORIES IN DYNAMIC PROGRAMMING

SYLVAIN SORIN, XAVIER VENEL, GUILLAUME VIGERAL

ABSTRACT. We show in a dynamic programming framework that uniform convergence of the finite horizon values implies that asymptotically the average accumulated payoff is constant on optimal trajectories. We analyze and discuss several possible extensions to two-person games.

## 1. PRESENTATION

Consider a dynamic programming problem as described in Lehrer and Sorin [1]. Given a set of states  $S$ , a correspondence  $\Phi$  from  $S$  to itself with non empty values and a payoff function  $f$  from  $S$  to  $[0, 1]$ , a feasible play at  $s \in S$  is a sequence  $\{s_m\}$  of states with  $s_1 = s$  and  $s_{m+1} \in \Phi(s_m)$ . It induces a sequence of payoffs  $\{f_m = f(s_m)\}$ ,  $m = 1, \dots, n, \dots$ . Recall that starting from a standard problem with random transitions and/or signals on the state, this presentation amounts to work on the set of probabilities on  $S$  and to consider expected payoffs.

Let  $v_n(s)$  (resp.  $v_\lambda(s)$ ) be the value of the  $n$  stage program  $G_n(s)$  (resp.  $\lambda$  discounted program  $G_\lambda(s)$ ) starting from state  $s$ . The **asymptotic approach** deals with asymptotic properties of the values  $v_n$  and  $v_\lambda$  as  $n$  goes to  $\infty$  or  $\lambda$  goes to 0.

The **uniform approach** focuses on properties of the strategies that hold uniformly in long horizons.  $v_\infty$  is the uniform value if for each  $\varepsilon > 0$  there exists  $N$  such that for each  $s \in S$ :

1) there is a feasible play  $\{s_m\}$  at  $s$  with

$$\frac{1}{n} \sum_{m=1}^n f(s_m) \geq v_\infty(s) - \varepsilon, \quad \forall n \geq N$$

2) for any feasible play  $\{s'_m\}$  at  $s$  and any  $n \geq N$

$$\frac{1}{n} \sum_{m=1}^n f(s'_m) \leq v_\infty(s) + \varepsilon.$$

Obviously the second approach is more powerful than the second (existence of a uniform value implies existence of an asymptotic value : the limit of  $v_n$  exists) but it is also more demanding: there are problems without uniform value where the asymptotic value exists (see Section 2). Note that the condition for the existence of a uniform value implies that the average accumulated payoff on optimal trajectories remains close to the value.

We will prove that a similar phenomenon holds true under conditions that are stronger than the existence of an asymptotic value but weaker than the existence of a uniform value.

Say that the dynamic programming problem is **regular** if :

- i)  $\lim v_n(s) = v(s)$  exists for each  $s \in S$ .
- ii) the convergence is uniform.

This condition was already introduced and studied in Lehrer and Sorin [1] (see Section 2).

We consider the following property **P**:

For any  $\varepsilon > 0$ , there exists  $n_0$ , such that for all  $n \geq n_0$ , for any state  $s$  and any feasible play  $\{s_m\}$   $\varepsilon$ -optimal for  $G_n(s)$  and for any  $t \in [0, 1]$ :

$$(1) \quad 3\varepsilon \geq \frac{1}{n} \left( \sum_{m=1}^{\lfloor tn \rfloor} f_m \right) - tv(s) \geq -3\varepsilon.$$

where  $\lfloor tn \rfloor$  stands for the integer part of  $tn$ .

This condition says that the average payoff remains close to the value on every almost-optimal trajectory with long duration (but the trajectory may depend on this duration). It also implies a similar property on every time interval.

## 2. EXAMPLES AND COMMENTS

1) The existence of the asymptotic value is not enough to control the payoff as required in property **P**.

An example is given in Lehrer and Sorin [1] (Section 2), where both  $\lim v_n$  and  $\lim v_\lambda$  exist on  $S$  but where the asymptotic average payoff is not constant on the unique optimal trajectory, nor on  $\varepsilon$ -optimal trajectories: in  $G_{2n}$ , an optimal play will induce  $n$  times 0 then  $n$  times 1 while  $v = 1/2$ .

Note that this example is not regular: the convergence of  $v_n$  to  $v$  is not uniform.

2) Recall that in the framework of dynamic programming, regularity is also equivalent to uniform convergence of  $v_\lambda$  (and with the same limit), see Lehrer and Sorin [1] (Section 3).

Note also that this regularity condition is not sufficient to obtain the existence of a uniform value, see Monderer and Sorin [2] (Section 2).

3) General conditions for regularity can be found in Renault [5].

## 3. MAIN RESULT

**Theorem 3.1.** *Assume that the program is regular, then **P** holds.*

### Proof

Let us start with the upper bound inequality in (1).

The result is clear for  $t \leq \varepsilon$  (recall that the payoff is in  $[0, 1]$ ). Otherwise let  $n_1$  large enough so that  $n \geq n_1$  implies  $\|v_n - v\| \leq \varepsilon$  by uniform convergence. Then the required inequality holds for  $n \geq n_2$  with  $\lfloor \varepsilon n_2 \rfloor \geq n_1$ .

Consider now the lower bound inequality in (1).

The result holds for  $t \geq 1 - \varepsilon$  by the  $\varepsilon$ -optimal property of the play, for  $n \geq n_1$ . Otherwise we use the following lemma from Lehrer and Sorin [1] (Proposition 1).

**Lemma 3.1.** *Both  $\limsup v_n$  and  $\limsup v_\lambda$  decrease on feasible histories.*

In particular, starting from  $s_{\lfloor tn \rfloor}$  the value of the program for the last  $n - \lfloor tn \rfloor$  stages is at most  $v(s_{\lfloor tn \rfloor}) + \varepsilon$  for  $n \geq n_2$ , by uniform convergence, hence less than the initial  $v(s) + \varepsilon$ , using the previous Lemma. Since the play is  $\varepsilon$ -optimal in  $G_n(s)$ , this implies that

$$(2) \quad \sum_{m=1}^{\lfloor tn \rfloor} f_m + (n - \lfloor tn \rfloor)(v(s) + \varepsilon) \geq n(v_n(s) - \varepsilon) \geq n(v(s) - 2\varepsilon)$$

hence the required inequality. ■

## 4. EXTENSIONS

### 4.1. Discounted case.

A similar result holds for the program  $G_\lambda$  corresponding to the evaluation  $\sum_{m=1}^{\infty} \lambda(1 - \lambda)^{m-1} f_m$ . Explicitly, one introduces the property **P'**:

For any  $\varepsilon > 0$ , there exists  $\lambda_0$ , such that for all  $\lambda \leq \lambda_0$ , for any state  $s$  and any feasible play  $\{s_m\}$   $\varepsilon$ -optimal for  $G_\lambda(s)$  and for any  $t \in [0, 1]$ :

$$(3) \quad 3\varepsilon \geq \sum_{m=1}^{n(t;\lambda)} \lambda(1-\lambda)^{m-1} f_m - tv(s) \geq -3\varepsilon.$$

where  $n(t; \lambda) = \inf\{p \in \mathbb{N}; \sum_{m=1}^p \lambda(1-\lambda)^{m-1} \geq t\}$ . Stage  $n(t; \lambda)$  corresponds to the fraction  $t$  of the total duration of the program.

**Theorem 4.1.** *Assume that the program is regular, then  $\mathbf{P}'$  holds.*

**Proof**

The proof follows the same lines than the proof of Theorem 3.1. Recall that by regularity both  $v_n$  and  $v_\lambda$  converge uniformly to  $v$ . Moreover the discounted sums  $(1-\lambda)^{-N} \sum_{m=1}^N \lambda(1-\lambda)^{m-1} f_m$  belong to the convex hull of the averages  $\frac{1}{n} \sum_{m=1}^n f_m; 1 \leq n \leq N$ . The counterpart of equation (2) is now

$$(4) \quad \sum_{m=1}^{n(t;\lambda)} \lambda(1-\lambda)^{m-1} f_m + (1-t)(v(s) + \varepsilon) \geq (v_\lambda(s) - \varepsilon) \geq v(s) - 2\varepsilon$$

■

#### 4.2. Continuous time.

Similar results holds in the following set-up:  $v_T(x)$  is the value of the control problem  $\Gamma_T$  with control set  $U$  where the state variable in  $X$  is governed by a differential equation (or more generally a differential inclusion)

$$\dot{x}_t = f(x_t, u_t)$$

starting from  $x$  at time 0. The real payoff function is  $g(x, u)$  and the evaluation is given by:

$$\frac{1}{T} \int_0^T g(x_t, u_t) dt.$$

Regularity in this framework amounts to uniform convergence (on  $X$ ) of  $V_T$  to some  $V$ . (Sufficient conditions for regularity can be found in Quincampoix and Renault [4]). The corresponding property is now  $\mathbf{P}''$ :

For any  $\varepsilon > 0$ , there exists  $T_0$ , such that for all  $T \geq T_0$ , for any state  $x$  and any feasible trajectory  $\varepsilon$ -optimal for  $\Gamma_T(x)$  and for any  $\theta \in [0, 1]$ :

$$(5) \quad 3\varepsilon \geq \frac{1}{T} \int_0^{\theta T} g(x_t, u_t) dt - \theta V(x) \geq -3\varepsilon.$$

**Theorem 4.2.** *Assume that the optimal control problem is regular, then  $\mathbf{P}''$  holds.*

**Proof**

Follow exactly the same lines than the proof of Theorem (2). ■

Finally the same tools can be used for an evaluation of the form  $\lambda \int_0^{+\infty} e^{-\lambda t} g(x_t, u_t) dt$ .

## 5. TWO-PLAYER ZERO-SUM GAMES

In trying to extend this result to a two-person zero-sum framework, several problems occurs.

### 5.1. Optimal strategies on both sides.

First it is necessary, to obtain good properties on the trajectory, to ask for optimality on both sides.

For example in the Big Match with no signals,

	$\alpha$	$\beta$
$a$	$1^*$	$0^*$
$b$	$0$	$1$

where a  $*$  denotes an absorbing payoff, the optimal strategy of player 1 in the “asymptotic game” on  $[0, 1]$  is to play “ $a$  before time  $t$ ” with probability  $t$ , see Sorin [6] Section 5.3.2. Obviously, if there is no restrictions on player 2’s moves the average payoff will not be constant. However, the optimal strategy of player 2 is “always  $(1/2, 1/2)$ ” hence time independent on  $[0, 1]$ . It thus induces a constant payoff and it is easy to see that the property is robust to small perturbations in the evaluation of the payoff.

### 5.2. Player 1 controls the transition.

Consider a repeated game with finite characteristics (states, moves, signals, ...) and use the recursive formula corresponding to the canonical representation with entrance laws being consistent probabilities on the universal belief space, see Mertens, Sorin and Zamir [3], Chapters III.1, IV.3. This representation preserves the values but in the auxiliary game, if player 1 controls the transition an optimal strategy of player 2 is to play a stage by stage best reply. Hence the model reduces to the dynamic programming framework and the results of the previous sections apply. A simple example corresponds to a game with incomplete information on one side where asymptotically an optimal strategy of the uniform player 1 is a splitting at time 0, while player 2 can obtain  $u(p_t)$  at time  $t$  where  $u$  is the value of the non-revealing game and  $p_t$  the martingale of posteriors at time  $t$ , see Sorin [6], 3.7.2.

### 5.3. Example.

Back to the general framework of two person zero-sum repeated games, the following example shows that in addition one has to strengthen the conditions on the pair of  $\varepsilon$ -optimal strategies. We exhibit a game having a uniform value  $v$  but for some state  $s$  with  $v(s) = 0$  one can construct, for each  $n$ , optimal strategies in  $\Gamma_n(s)$  inducing essentially a constant payoff 1 during the first half of the game.

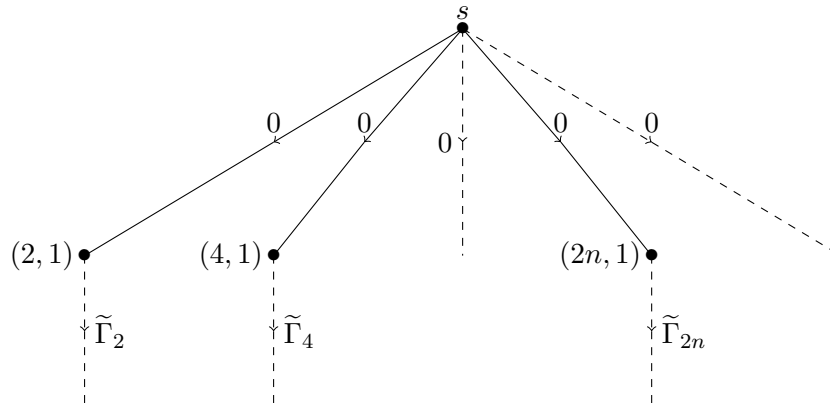
Starting from the initial state  $s$ , the tree representing the game  $\Gamma$  has countably many subgames  $\tilde{\Gamma}_{2n}$ , the transition being controlled by player 1 (with payoff 0). In  $\tilde{\Gamma}_{2n}$  there are at most  $n$  stages before reaching an absorbing state. At each of these stages of the form  $(2n, m), m = 1, \dots, n$ , the players plays a “jointly controlled” process leading either to a payoff 1 and the next stage  $(2n, m + 1)$  (if they agree) or an absorbing payoff  $x_{2n,m}$  with  $(m - 1) + (2n - (m - 1))x_{2n,m} = 0$ , otherwise. Hence every feasible path of length  $2n$  in  $\tilde{\Gamma}_{2n}$  gives a total payoff 0. Obviously the uniform value exists since each player can stop the game at each node, inducing the same absorbing payoff. The representation is as follows:

Notice that in the  $2n + 1$  stage game, after a move of player 1 to  $\tilde{\Gamma}_{2n}$ , any play is compatible with optimal strategies, in particular those leading to the sequence of payoffs  $2n$  times 0 or  $n$  times 1 then  $n$  times  $-1$ .

### 5.4. Conjectures.

A natural conjecture is that in any regular game (i.e. where  $v_n$  converges uniformly to  $v$ ): for any  $\varepsilon > 0$ , there exists  $n_0$ , such that for all  $n \geq n_0$ , for any initial state  $s$ , there exists a couple  $(\sigma_n, \tau_n)$  of  $\varepsilon$ -optimal strategies in  $G_n(s)$  such that for any  $t \in [0, 1]$ :

$$(6) \quad 3\varepsilon \geq \frac{1}{n} \mathbf{E}_{\sigma_n, \tau_n}^s \left( \sum_{m=1}^{[tn]} f_m \right) - tv(s) \geq -3\varepsilon.$$

FIGURE 1. The game  $\Gamma$  starting from state  $s$ 

	C	A		C	A		C	A		C	A		C	A		C	A
C	$\frac{1}{\gamma}$	$0^*$	C	$\frac{1}{\gamma}$	$x_{2n,2}^*$	C	$\frac{1}{\gamma}$	$x_{2n,m}^*$	C	$\frac{1}{\gamma}$	$x_{2n,n}^*$	C	$-1^*$	$-1^*$	A	$-1^*$	$-1^*$
A	$0^*$	$0^*$	A	$x_{2n,2}^*$	$x_{2n,2}^*$	A	$x_{2n,m}^*$	$x_{2n,m}^*$	A	$x_{2n,n}^*$	$x_{2n,n}^*$	A	$-1^*$	$-1^*$	A	$-1^*$	$-1^*$
	$(2n, 1)$		$(2n, 2)$	$\dots$	$(2n, m)$	$\dots$	$(2n, n)$	$-1^*$									

FIGURE 2. The subgame  $\tilde{\Gamma}_{2n}$  starting from state  $(2n, 1)$ 

where  $[tn]$  stands for the integer part of  $tn$  and  $f_m$  is the payoff at stage  $m$ .

A more elaborate conjecture would rely on the existence of an asymptotic game  $\Gamma^*$  played in continuous time on  $[0, 1]$  with value  $v$  (as in Section 5.1). We use the representation of the repeated game as a stochastic game through the recursive structure as above, see Mertens, Sorin, Zamir [3], Chapter IV. The condition is now the existence of a couple of strategies  $(\sigma, \tau)$  in the asymptotic game that would depend only on the time  $t \in [0, 1]$  and on the current state  $s$  such that for any  $\varepsilon > 0$ , there exists  $\eta$  with the following property: in any repeated game where the (relative) weight of stage  $m$  is  $\alpha_m$ , with  $\{\alpha_m\}$  decreasing and less than  $\eta$ , thus defining a partition  $\Pi$  of  $[0, 1]$ , the strategies  $(\sigma_\Pi, \tau_\Pi)$  induced in the repeated game by  $(\sigma, \tau)$  satisfies (6).

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