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Tom Hirschowitz

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# CARTESIAN CLOSED 2-CATEGORIES AND PERMUTATION EQUIVALENCE IN HIGHER-ORDER REWRITING

TOM HIRSCHOWITZ

ABSTRACT. We propose a semantics for permutation equivalence in higher-order rewriting. This semantics takes place in cartesian closed 2-categories, and is proved sound and complete.

## 1. INTRODUCTION

It is known since the end of the 80's that 2-categories with finite products provide a semantics for term rewriting [3]. *Higher-order rewriting* [10, 17, 14, 15] is a framework for specifying rewrite systems on terms with variable binding. Many results from standard term rewriting have been generalised to higher-order rewriting, notably normalisation or confluence results. An important tool for confluence results is the notion of *permutation equivalence*, which was generalised to the higher-order case by Bruggink [1]. He defines a calculus of *proof terms* for specifying reductions in a higher-order rewrite system.

We here propose a categorical semantics for a variant of this calculus, in terms of *cartesian closed 2-categories*. We first define *cartesian closed 2-signatures*, which generalise higher-order rewrite systems, and organise them into a category  $\text{Sig}$ . We then construct an adjunction

$$(1.1) \quad \text{Sig} \begin{array}{c} \xrightarrow{\mathcal{H}} \\ \perp \\ \xleftarrow{\mathcal{W}} \end{array} 2\text{CCCat},$$

where  $2\text{CCCat}$  is the category of small cartesian closed 2-categories. From a given higher-order rewrite system  $S$ , the functor  $\mathcal{H}$  constructs a cartesian closed 2-category, whose 2-cells are Bruggink's proof terms modulo permutation equivalence, which we prove is the free cartesian closed 2-category generated by  $S$ .

We review a number of examples and non-examples, and sketch an extension to deal with the latter.

**Related work.** Our cartesian closed 2-signatures are a 2-dimensional refinement of *cartesian closed sketches* [16, 4, 9]. Bruggink's calculus of permutation equivalence is close in spirit to Hilken's 2-categorical semantics of the simply-typed  $\lambda$ -calculus [7], but technically different and generalised to arbitrary higher-order rewrite systems. Capriotti [2] proposes a semantics of so-called *flat* permutation equivalence in sesquicategories. More related work is discussed in Section 4.2.

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## 2. CARTESIAN CLOSED SIGNATURES AND CATEGORIES

We start by recalling the well-known, or at least folklore, adjunction between what we here call *(cartesian closed) 1-signatures* and cartesian closed categories.

For any set  $X$ , define *types* over  $X$  by the grammar:

$$A, B, \dots \in \mathcal{L}_0(X) ::= x \mid 1 \mid A \times B \mid B^A,$$

with  $x \in X$ .

**Proposition 1.**  $\mathcal{L}_0$  defines a monad on  $\mathbf{Set}$ .

Let the set of *sequents* over a set  $X$  be  $\mathcal{S}_0(X) = \mathcal{L}_0(X)^* \times \mathcal{L}_0(X)$ , i.e., sequents are pairs of a list of types and a type. The assignment  $X \mapsto \mathcal{S}_0(X)$  extends to an endofunctor on  $\mathbf{Set}$ .

**Definition 1.** A 1-signature consists of a set  $X_0$  of sorts, and an  $\mathcal{S}_0(X_0)$ -indexed set  $X_1$  of operations, or equivalently a map  $X_1 \rightarrow \mathcal{S}_0(X_0)$ .

A morphism of 1-signatures  $(X_0, X_1) \rightarrow (Y_0, Y_1)$  is a pair  $(f_0, f_1)$  where  $f_i: X_i \rightarrow Y_i$  such that

$$\begin{array}{ccc} X_1 & \xrightarrow{f_1} & Y_1 \\ \downarrow & & \downarrow \\ \mathcal{S}_0(X_0) & \xrightarrow{\mathcal{S}_0(f_0)} & \mathcal{S}_0(Y_0) \end{array}$$

commutes. Morphisms compose in the obvious way, and we have:

**Proposition 2.** Composition of morphisms is associative and unital, and hence 1-signatures and their morphisms form a category  $\mathbf{Sig}_1$ .

There is a well-known adjunction

$$\mathbf{Sig}_1 \begin{array}{c} \xrightarrow{\mathcal{H}_1} \\ \perp \\ \xleftarrow{\mathcal{W}_1} \end{array} \mathbf{CCCat}$$

between 1-signatures and the category  $\mathbf{CCCat}$  of small cartesian closed categories (with chosen structure) and (strict) cartesian closed functors.

The functor  $\mathcal{W}_1$  sends a cartesian closed category  $\mathcal{C}$  to the signature with sorts  $\mathcal{C}_0$ , its set of objects, and with operations  $A_1, \dots, A_n \rightarrow A$  the set  $\mathcal{C}(\llbracket A_1 \times \dots \times A_n \rrbracket, \llbracket A \rrbracket)$ , where  $\llbracket - \rrbracket$  denotes the function  $\mathcal{L}_0(\mathcal{C}_0) \rightarrow \mathcal{C}_0$  defined by induction:

$$(2.1) \quad \begin{aligned} \llbracket c \rrbracket &= c & c \in \mathcal{C}_0 \\ \llbracket 1 \rrbracket &= 1 \\ \llbracket A \times B \rrbracket &= \llbracket A \rrbracket \times \llbracket B \rrbracket \\ \llbracket B^A \rrbracket &= \llbracket B \rrbracket^{\llbracket A \rrbracket}. \end{aligned}$$

Conversely, given a 1-signature  $X$ , consider the simply-typed  $\lambda$ -calculus with base types in  $X_0$  and constants in  $X_1$ . Terms modulo  $\beta\eta$  form a category  $\mathcal{H}_1(X)$  with objects all types over  $X_0$  and morphisms  $A \rightarrow B$  all terms of type  $B$  with one free variable of type  $A$ .

A less often formulated observation, which is useful to us, is that the adjunction  $\mathcal{H}_1 \dashv \mathcal{W}_1$  decomposes into two adjunctions

$$\begin{array}{ccccc}
& & \mathcal{L}_1 & & \mathcal{F}_1 \\
& \swarrow & \xrightarrow{\quad} & \searrow & \xrightarrow{\quad} \\
\text{Sig}_1 & \xrightarrow{\quad} & \mathcal{L}_1\text{-Alg} & \xrightarrow{\quad} & \text{CCCat}, \\
& \nwarrow & \xleftarrow{\quad} & \swarrow & \xleftarrow{\quad} \\
& & \mathcal{U}_1 & & \mathcal{V}_1
\end{array}$$

where  $\mathcal{L}_1\text{-Alg}$  is the category of algebras for the monad  $\mathcal{L}_1$  defined as follows (and  $\mathcal{L}_1$  is shorthand for the functor  $X \mapsto (\mathcal{L}_1(X), \mu)$ ).

For any 1-signature  $X$ , let  $\mathcal{L}_1(X)$  denote the 1-signature with

- as sorts the set  $X_0$ , and
- as operations  $\Gamma \vdash A$  the  $\lambda$ -terms  $\Gamma \vdash M : A$ , modulo  $\beta\eta$ .

$\mathcal{L}_1$  extends to an endofunctor on  $\text{Sig}_1$ , whose action on morphisms of 1-signatures  $X \xrightarrow{f} Y$  substitutes constants  $c \in X_1$  with  $f_1(c)$ . We obtain

**Proposition 3.**  $\mathcal{L}_1$  is a monad on  $\text{Sig}_1$ .

The functor  $\mathcal{V}_1$  sends any cartesian closed category  $\mathcal{C}$  to the  $\mathcal{L}_1$ -algebras  $(\mathcal{C}_0, \mathcal{C}_1)$  defined as follows. First,  $\mathcal{C}_0$  is the set of objects of  $\mathcal{C}$ . It has a canonical  $\mathcal{L}_0$ -algebra structure, say  $h_0 : \mathcal{L}_0(\mathcal{C}_0) \rightarrow \mathcal{C}_0$ , obtained by interpreting type constructors in  $\mathcal{C}$  as in (2.1). Extending this to contexts  $G$  by  $h_0(G) = \prod_i h_0(G_i)$ , let the operations in  $\mathcal{C}_1(G, A)$  be the 1-cells in  $\mathcal{C}(h_0(G), h_0(A))$ . Beware: the domain and codomain of such an operation are really  $G$  and  $A$ , not  $h_0(G)$  and  $h_0(A)$ . Similarly, interpreting the  $\lambda$ -calculus in  $\mathcal{C}$ , the 1-signature  $(\mathcal{C}_0, \mathcal{C}_1)$  has a canonical  $\mathcal{L}_1$ -algebra structure, say  $h_1 : \mathcal{L}_1(\mathcal{C}_0, \mathcal{C}_1) \rightarrow (\mathcal{C}_0, \mathcal{C}_1)$ :

$$\begin{aligned}
h_1(G \vdash x_i : G_i) &= \pi_i \\
h_1(G \vdash () : 1) &= ! \\
h_1(G \vdash c(M_1, \dots, M_n)) &= c \circ (h_1(M_1), \dots, h_1(M_n)) \\
h_1(G \vdash \lambda x : A. M : B^A) &= \varphi(h_1(G, x : A \vdash M : B)) \\
h_1(G \vdash MN : B) &= ev \circ (h_1(M), h_1(N)) \\
h_1(G \vdash (M, N) : A \times B) &= (h_1(M), h_1(N)) \\
h_1(G \vdash \pi M : A) &= \pi \circ M \\
h_1(G \vdash \pi' M : A) &= \pi' \circ M,
\end{aligned}$$

where  $!$  is the unique morphism  $h_0(G) \rightarrow 1$ ,  $\varphi$  is the bijection  $\mathcal{C}(h_0(G, A), h_0(B)) \cong \mathcal{C}(h_0(G), h_0(B^A))$ , and  $ev$  is the structure morphism  $h_0(B^A \times A) \rightarrow h_0(B)$ .

$\mathcal{L}_1$ -algebras are much like cartesian closed categories whose objects are freely generated by their set of sorts. A perhaps useful analogy here is with multicategories  $\mathcal{M}$ , seen as being close to monoidal categories whose objects are freely generated by those of  $\mathcal{M}$  by tensor and unit. Here, the functor  $\mathcal{F}_1$  sends any  $\mathcal{L}_1$ -algebra  $(X, h)$  to the cartesian closed category with

- objects the types over  $X_0$ , i.e.,  $\mathcal{L}_0(X_0)$ ,
- morphisms  $A \rightarrow B$  the set of operations in  $X_1(A, B)$ .

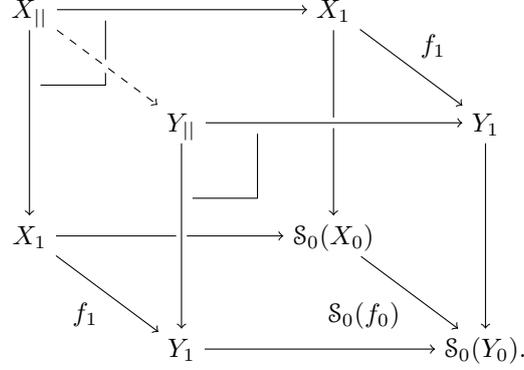
This canonically forms a cartesian closed category, with structure induced by the  $\mathcal{L}_1$ -algebra structure. We define it in more detail in dimension 2 in Section 7.2.

### 3. CARTESIAN CLOSED 2-SIGNATURES

Given a 1-signature  $X$ , let  $X_{||}$  denote the set of pairs of parallel operations, i.e., pairs of operations  $M, N$  above the same sequent. Otherwise said,  $X_{||}$  is the pullback

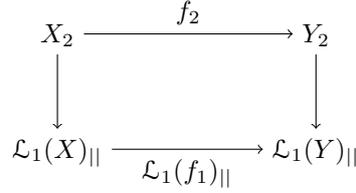
$$\begin{array}{ccc}
X_{||} & \xrightarrow{\quad} & X_1 \\
\downarrow & \lrcorner & \downarrow \\
X_1 & \xrightarrow{\quad} & \mathfrak{S}_0(X_0).
\end{array}$$

Any morphism  $f: X \rightarrow Y$  of 1-signatures yields a function  $f_{||}: X_{||} \rightarrow Y_{||}$ , via the dashed arrow (obtained by universal property of pullback) in



**Definition 2.** A 2-signature consists of a 1-signature  $X$ , plus a set  $X_2$  of reduction rules with a function  $X_2 \rightarrow \mathcal{L}_1(X)_{||}$ .

A morphism of 2-signatures  $(X, X_2) \rightarrow (Y, Y_2)$  is a pair  $(f, f_2)$  where  $f: X \rightarrow Y$  is a morphism of 1-signatures and  $f_2: X_2 \rightarrow Y_2$  makes the diagram



commute. We obtain:

**Proposition 4.** Composition of morphisms is associative and unital, and hence 2-signatures and their morphisms form a category  $\mathbf{Sig}$ .

#### 4. EXAMPLES

**4.1. Higher-order rewrite systems.** The prime example of a 2-signature is that for the pure  $\lambda$ -calculus: it has a sort  $t$  and operations

$$a: t \times t \rightarrow t \qquad \ell: t^t \rightarrow t,$$

with a reduction rule  $\beta$  above the pair

$$x: t^t, y: t \vdash a(\ell(x), y), x(y): t$$

in  $\mathcal{L}_1(\{t\}, \{\ell, a\})_{||}$ . Categorically, this will yield a 2-cell

$$\begin{array}{ccc} & \ell \times t & t \times t \\ & \nearrow & \searrow a \\ t^t \times t & & t. \\ & \searrow ev & \\ & & \end{array} \quad \Downarrow \beta$$

This is an example of a *higher-order rewrite system* in the sense of Nipkow [14]. Nipkow's definition is formally different, but his higher-order rewrite systems are in bijection with 2-signatures  $h: X_2 \rightarrow \mathcal{L}_1(X)_{||}$  such that for all rules  $r \in X_2$ , letting  $(\Gamma \vdash M, N: A) = h(r)$ :

- $M$  is not a variable,
- $A$  is a sort,
- each variable occurring in  $\Gamma$  occurs free in  $M$ .

These restrictions help formulating and proving decidability problems on higher-order rewrite systems, whose extension to our setting we leave open.

Let us now anticipate over our main results below and state our soundness and completeness theorem. Given a higher-order rewrite system  $X$ , i.e., a 2-signature satisfying the above conditions, let  $\mathcal{R}(X)$  be the following locally-preordered 2-category. It has:

- objects are types in  $\mathcal{L}_0(X_0)$ ;
- morphisms  $A \rightarrow B$  are  $\lambda$ -terms in  $\mathcal{L}_1(X)(A \vdash B)$ , modulo  $\beta\eta$ ;
- given two parallel morphisms  $M$  and  $N$ , there is one 2-cell  $M \rightarrow N$  exactly when there is a sequence of reductions  $M \rightarrow^* N$  in the usual sense [14].

**Proposition 5.**  $\mathcal{R}(X)$  is 2-cartesian closed.

$\mathcal{R}(X)$  and  $\mathcal{H}(X)$  have the same objects and morphisms. But because our inference rules for forming reductions are the same as deduction rules for proving the existence of a reduction in the usual sense, we may send any reduction  $P: M \rightarrow N$  to the unique reduction  $M \rightarrow N$  in  $\mathcal{R}(X)$ .

**Theorem 1** (Soundness and completeness). *This defines an identity-on-objects, identity-on-morphisms, locally full cartesian closed 2-functor  $\mathcal{R}(X) \xrightarrow{1} \mathcal{H}(X)$ .*

**4.2. Theories with binding.** Understanding reduction rules as equations, it is easy to define the free cartesian closed category generated by a 2-signature. This yields an adjunction

$$(4.1) \quad \text{Sig} \begin{array}{c} \xrightarrow{\mathcal{H}'} \\ \perp \\ \xleftarrow{\mathcal{W}'} \end{array} \text{CCCat.}$$

This adjunction provides a categorical semantics for theories with binding, which is more general than other approaches by Fiore and Hur [6], Hirschowitz and Maggesi [8], and Zsidó [18], and which is in line with Lambek's seminal paper [11].

If I understand correctly, the motivation for Fiore and Hur's subtle approach is the will to explain the  $\lambda$ -calculus by strictly less than itself. The present framework does not obey this specification, and instead tends to view the  $\lambda$ -calculus as a universal (parameterised) theory with binding.

We end this section by giving a formal construction of the adjunction (4.1). Cartesian closed categories form a full, reflective subcategory of  $2\text{CCCat}$ , via the functor  $\mathcal{J}: 2\text{CCCat} \rightarrow \text{CCCat}$  sending a cartesian closed 2-category  $\mathcal{C}$  to the cartesian closed category with:

- objects those of  $\mathcal{C}$ ,
- morphisms those of  $\mathcal{C}$ , modulo the congruence generated by  $f \sim g$  iff there exists a 2-cell  $f \rightarrow g$ .

Here,  $\mathcal{J}(\mathcal{C})$  is thought of as the free cartesian closed 2-category with trivial 2-cells (i.e., 0 or 1). The desired adjunction is obtained by composing the adjunctions

$$\text{Sig} \begin{array}{c} \xrightarrow{\mathcal{H}} \\ \perp \\ \xleftarrow{\mathcal{W}} \end{array} 2\text{CCCat} \begin{array}{c} \xrightarrow{\mathcal{J}} \\ \perp \\ \xleftarrow{\quad} \end{array} \text{CCCat.}$$

**4.3. Non-examples.** Non-examples are given by calculi whose reduction semantics is defined on terms modulo a so-called *structural congruence*, e.g., CCS [12], or the  $\pi$ -calculus [5, 13].

For example, consider the CCS term  $(a \mid 0) \mid \bar{a}$ . In CCS, it is *structurally* equivalent to  $(a \mid \bar{a}) \mid 0$ , which then reduces to  $0 \mid 0$ .

In order to account for this, we would have to consider a 2-signature with reduction rules for structural congruence, here  $(M_1 \mid M_2) \mid M_3 \rightarrow M_1 \mid (M_2 \mid M_3)$  for associativity, and  $M \mid N \rightarrow N \mid M$  for commutativity. But then, these reductions count as proper reductions, which departs from the desired computational behaviour. For example, the term  $a \mid a$  has an infinite reduction sequence, using commutativity.

Anticipating the development in the next sections, a potential solution is to extend 2-signatures to *2-theories*. For any 2-signature  $X$ , let  $X_{\parallel}$  denote the set of pairs of reduction rules  $r, s$  with a common type  $G \vdash M \rightarrow N : A$ . A 2-theory is a 2-signature  $X$ , together with a set of equations between parallel reductions, i.e., a subset  $X_3$  of  $\mathcal{L}(X)_{\parallel}$ .

The main adjunction announced above (1.1) extends to an adjunction between 2-theories and cartesian closed 2-categories. Using equations, we may specify that any reduction  $M \rightarrow M$  using only structural rules be the identity on  $M$ , and consider the computational behaviour of a 2-category to consist of its non-invertible morphisms, as proposed by Hilken [7]. A question is whether for a given calculus this can be done with finitely many equations.

## 5. A 2-LAMBDA-CALCULUS

We now begin the construction of Adjunction (1.1). We start in this section by defining a monad  $\mathcal{L}$  on  $\text{Sig}$ , which we will use to factor Adjunction (1.1) as

$$\text{Sig} \begin{array}{c} \xrightarrow{\mathcal{L}} \\ \perp \\ \xleftarrow{\mathcal{U}} \end{array} \mathcal{L}\text{-Alg} \begin{array}{c} \xrightarrow{\mathcal{F}} \\ \perp \\ \xleftarrow{\mathcal{V}} \end{array} 2\text{CCCat},$$

where:

- $\mathcal{L}\text{-Alg}$  is the category of  $\mathcal{L}$ -algebras,
- $\mathcal{L} : \text{Sig} \rightarrow \mathcal{L}\text{-Alg}$  is a shortcut for  $X \mapsto (\mathcal{L}^2 X \xrightarrow{\mu} \mathcal{L} X)$ ,
- $\mathcal{U}(\mathcal{L} X \xrightarrow{h} X) = X$ ,
- $2\text{CCCat}$  is the category of cartesian closed 2-categories, which we define in Section 6.

The left-hand adjunction holds by  $\mathcal{L}$  being a monad, thus we concentrate in Section 7 on establishing the right-hand one.

But for now, let us define the monad  $\mathcal{L}$ .

**5.1. Syntax.** Given a 2-signature  $X = ((X_0, X_1), h : X_2 \rightarrow \mathcal{L}_1(X)_{\parallel})$  (actually  $\mathcal{L}_1(X)$  is  $\mathcal{L}_1(X_0, X_1)$ ), we construct a new 2-signature  $\mathcal{L}(X)$ , whose reduction rules represent reduction sequences in the “higher-order rewrite system” defined by  $X$ , modulo permutation equivalence. The 2-signature  $\mathcal{L}(X)$  has the same base 1-signature  $(X_0, X_1)$ , and as reduction rules the terms of a  $2\lambda$ -calculus (in the sense of Hilken [7]) modulo permutation equivalence, which we now define.

First, terms, called *reductions*, are defined by induction in Figure 1. The typing judgement has the shape  $\Gamma \vdash P : M \rightarrow N : A$ , where  $A$  is a type in  $\mathcal{L}_0(X_0)$ ,  $\Gamma$  is a list of pairs of a variable and a type, with no variable appearing more than once,  $M$  and  $N$  are terms of type  $\Gamma \vdash A$  modulo  $\beta\eta$ , and  $P$  is a reduction. In the sequel, we often forget the variables in such pairs  $(\Gamma \vdash A)$ , and identify them with sequents in  $\mathcal{S}_0(X_0)$ .

When clear from context, we abbreviate substitutions  $[M_1/x_1, \dots, M_n/x_n]$  by  $[M_1, \dots, M_n]$ . For a context  $G$ ,  $G_i$  denotes its  $i$ th type. Also, for  $(M, N) \in \mathcal{L}_1(X)_{\parallel}$ , we let  $X(M, N)$  be the set of all reduction rules  $r \in X_2$  such that  $h(r) = (M, N)$ . We write  $X(\Gamma \vdash M, N : A)$  to indicate the common type of  $M$  and  $N$ . Similarly,  $X(G \vdash A)$  denotes the set of operations in  $X_1$  above  $G \vdash A$ .

$$\begin{array}{c}
\dfrac{\dots \quad \Gamma \vdash P_i : M_i \rightarrow N_i : G_i \quad \dots}{\Gamma \vdash r(P_1, \dots, P_n) : M[M_1, \dots, M_n] \rightarrow N[N_1, \dots, N_n] : A} \quad (r \in X(G \vdash M, N : A)) \\
\\
\dfrac{\Gamma \vdash P : M_1 \rightarrow M_2 : A \quad \Gamma \vdash Q : M_2 \rightarrow M_3 : A}{\Gamma \vdash P ;_{M_2} Q : M_1 \rightarrow M_3 : A} \quad \Gamma, x : A, \Delta \vdash x : x \rightarrow x : A \\
\\
\Gamma \vdash () : () \rightarrow () : 1 \\
\\
\dfrac{\Gamma \vdash P_1 : M_1 \rightarrow N_1 : G_1 \quad \dots \quad \Gamma \vdash P_n : M_n \rightarrow N_n : G_n}{\Gamma \vdash c(P_1, \dots, P_n) : c(M_1, \dots, M_n) \rightarrow c(N_1, \dots, N_n) : A} \quad (c \in X_1(G \vdash A)) \\
\\
\dfrac{\Gamma, x : A \vdash P : M \rightarrow N : B}{\Gamma \vdash \lambda x : A. P : \lambda x : A. M \rightarrow \lambda x : A. N : B^A} \\
\\
\dfrac{\Gamma \vdash P : M \rightarrow M' : B^A \quad \Gamma \vdash Q : N \rightarrow N' : A}{\Gamma \vdash PQ : MN \rightarrow M'N' : B} \\
\\
\dfrac{\Gamma \vdash P : M \rightarrow M' : A \quad \Gamma \vdash Q : N \rightarrow N' : B}{\Gamma \vdash (P, Q) : (M, N) \rightarrow (M', N') : A \times B} \\
\\
\dfrac{\Gamma \vdash P : M \rightarrow N : A \times B}{\Gamma \vdash \pi_{A,B} P : \pi_{A,B} M \rightarrow \pi_{A,B} N : A} \quad \dfrac{\Gamma \vdash P : M \rightarrow N : A \times B}{\Gamma \vdash \pi'_{A,B} P : \pi'_{A,B} M \rightarrow \pi'_{A,B} N : B}
\end{array}$$

FIGURE 1. Reductions

5.2. **Substitution.** Next, we define substitution, which has “type”

$$(5.1) \quad \dfrac{\Gamma \vdash Q : N \rightarrow N' : \Delta \quad \Delta \vdash P : M \rightarrow M' : A}{\Gamma \vdash P[Q] : M[N] \rightarrow M'[N'] : A},$$

i.e., given a reduction  $P$  and a tuple of reductions  $Q$ , it produces a reduction of the indicated type, which we denote  $P[Q]$ . Here, we denote by  $\Gamma \vdash Q : N \rightarrow N' : \Delta$  a tuple of reductions  $\Gamma \vdash Q_i : N_i \rightarrow N'_i : \Delta_i$ , for  $1 \leq i \leq |\Delta|$ .

The definition is a bit tricky:

- first we define *left whiskering*, which has “type”

$$\dfrac{\Gamma \vdash Q : N \rightarrow N' : \Delta \quad \Delta \vdash M : A}{\Gamma \vdash M[Q] : M[N] \rightarrow M'[N'] : A},$$

- then we define *right whiskering*, which has “type”

$$\dfrac{\Gamma \vdash N : \Delta \quad \Delta \vdash P : M \rightarrow M' : A}{\Gamma \vdash P[N] : M[N] \rightarrow M'[N] : A},$$

(where  $N$  denotes a tuple),

- then we define substitution by

$$P[Q] = (P[N] ;_{M'[N]} M'[Q]).$$

There is of course another legitimate definition, namely

$$P[Q] = (M[Q] ;_{M[N']} P[N']).$$

The two will be equated by permutation equivalence in the next section.

Left whiskering is defined by induction, with  $\Delta = (x_1 : A_1, \dots, x_n : A_n)$  and  $Q = (Q_1, \dots, Q_n)$ , by:

$$\begin{aligned}
() [Q] &= () \\
x_i [Q] &= Q_i \\
c(M_1, \dots, M_p) [Q] &= c(M_1 [Q], \dots, M_p [Q]) \\
(\lambda x : B. M) [Q] &= \lambda x : B. (M [Q, x]) \quad (\text{for } x \notin \text{dom}(\Delta)) \\
(MN) [Q] &= (M [Q] N [Q]) \\
(M, N) [Q] &= (M [Q], N [Q]) \\
(\pi_{A,B} M) [Q] &= \pi_{A,B} (M [Q]) \\
(\pi'_{A,B} M) [Q] &= \pi'_{A,B} (M [Q]).
\end{aligned}$$

Right whiskering is defined by induction, with  $\Delta = (x_1 : A_1, \dots, x_n : A_n)$  and  $N = (N_1, \dots, N_n)$ , by:

$$\begin{aligned}
(r(P_1, \dots, P_p)) [N] &= r(P_1 [N], \dots, P_p [N]) \\
(P_1 ;_{M''} P_2) [N] &= (P_1 [N] ;_{M'' [N]} P_2 [N]) \\
() [N] &= () \\
x_i [N] &= N_i \\
c(P_1, \dots, P_p) [N] &= c(P_1 [N], \dots, P_p [N]) \\
(\lambda x : B. P') [N] &= \lambda x : B. (P' [N, x]) \quad (\text{for } x \notin \text{dom}(\Delta)) \\
(P_1 P_2) [N] &= (P_1 [N] P_2 [N]) \\
(P_1, P_2) [N] &= (P_1 [N], P_2 [N]) \\
(\pi_{A,B} P') [N] &= \pi_{A,B} (P' [N]) \\
(\pi'_{A,B} P') [N] &= \pi'_{A,B} (P' [N]).
\end{aligned}$$

**Definition 3.** Let  $P[Q] = (P[N] ;_{M'[N]} M'[Q])$ .

**Proposition 6.** Given reductions  $P$  and  $Q$  as above, the capture-avoiding substitution  $P[Q]$  is a well-typed reduction  $\Gamma \vdash P[Q] : M[N] \rightarrow M'[N'] : A$ .

Similarly, there is a weakening operation with “type”

$$\frac{\Gamma \vdash P : M \rightarrow N : A}{\Gamma, x : B \vdash P : M \rightarrow N : A.} \quad (x \notin \Gamma)$$

**5.3. Permutation equivalence.** We now define *permutation equivalence* on reductions, by the equations in Figures 3 and 4, in Appendix A. The *congruence* rules in Figure 3 are bureaucratic: they just say that permutation equivalence is a congruence. The *category* rules make reductions of a given type  $\Gamma \vdash A$  into a category. In Figure 4, the *beta* and *eta* rules mirror the term-level beta and eta rules. Finally, the *lifting* rules lift composition of reductions towards toplevel.

So,  $\mathcal{L}(X)$  has sorts  $X_0$ , operations  $X_1$ , and as reduction rules in  $\mathcal{L}(X)(G \vdash M, N : A)$  all reductions  $G \vdash P : M \rightarrow N : A$ , modulo the equations.

This easily extends to:

**Proposition 7.**  $\mathcal{L}$  is a functor  $\text{Sig} \rightarrow \text{Sig}$ .

Now, consider  $\mathcal{L}\mathcal{L}(X)$ . We define a mapping  $\mu_X : \mathcal{L}\mathcal{L}(X) \rightarrow \mathcal{L}(X)$ , by induction on reductions. The typing rule for reduction rules specialises to:

$$\frac{(R \in \mathcal{L}(X)(G \vdash M, N : A)) \quad \Gamma \vdash P_1 : M_1 \rightarrow N_1 : G_1 \quad \dots \quad \Gamma \vdash P_n : M_n \rightarrow N_n : G_n}{\Gamma \vdash R(P_1, \dots, P_n) : M[M_1, \dots, M_n] \rightarrow N[N_1, \dots, N_n] : A}$$

We set  $\mu(R(P_1, \dots, P_n)) = R[\mu(P_1), \dots, \mu(P_n)]$ . The other cases just propagate the substitution:

$$\begin{aligned}
P ; Q &\mapsto \mu(P) ; \mu(Q) \\
x &\mapsto x \\
() &\mapsto () \\
c(P_1, \dots, P_n) &\mapsto c(\mu(P_1), \dots, \mu(P_n)) \\
\lambda x : A. P &\mapsto \lambda x : A. \mu(P) \\
PQ &\mapsto \mu(P)\mu(Q) \\
(P, Q) &\mapsto (\mu(P), \mu(Q)) \\
\pi P &\mapsto \pi(\mu(P)) \\
\pi' P &\mapsto \pi'(\mu(P)).
\end{aligned}$$

**Lemma 1.** *This defines a natural transformation  $\mu : \mathcal{L}^2 \rightarrow \mathcal{L}$ , which makes the diagram*

$$\begin{array}{ccc}
\mathcal{L}^3 & \xrightarrow{\mathcal{L}\mu} & \mathcal{L}^2 \\
\mu\mathcal{L} \downarrow & & \downarrow \mu \\
\mathcal{L}^2 & \xrightarrow{\mu} & \mathcal{L}
\end{array}$$

*commute.*

Similarly, there is a natural transformation  $\eta : id \rightarrow \mathcal{L}$ , sending each  $r \in X(G \vdash M, N : A)$  to the reduction  $G \vdash r(x_1, \dots, x_n) : M \rightarrow N : A$ , and we have:

**Lemma 2.** *The diagram*

$$\begin{array}{ccccc}
& & \eta\mathcal{L} & & \\
& & \rightarrow & & \\
\mathcal{L} & & & \mathcal{L}^2 & \xleftarrow{\mathcal{L}\eta} & \mathcal{L} \\
& \searrow & & \downarrow \mu & \swarrow & \\
& & & \mathcal{L} & & 
\end{array}$$

*commutes.*

**Corollary 1.**  *$(\mathcal{L}, \mu, \eta)$  is a monad on  $\text{Sig}$ .*

A crucial result is:

**Proposition 8.** *For all  $\Gamma \vdash Q : N \rightarrow N' : \Delta$  and  $\Delta \vdash P : M \rightarrow M' : A$ , we have:*

$$\Gamma \vdash P[Q] \equiv (M[Q] ;_{M[N']} P[N']) : M[N] \rightarrow M'[N'] : A.$$

*Proof.* We proceed by induction on  $P$ . Most cases are bureaucratic. Consider for instance  $P = c(P_1, \dots, P_p)$ . Then, by definition:

$$P[Q] = (c(P_1[N], \dots, P_p[N]) ;_{c(M'_1[N], \dots, M'_p[N])} c(M'_1[Q], \dots, M'_p[Q])).$$

By the third lifting rule, this is  $\equiv$ -related to

$$c(P_1[N] ;_{M'_1[N]} M'_1[Q], \dots, P_p[N] ;_{M'_p[N]} M'_p[Q]).$$

By  $p$  applications of the induction hypothesis, we obtain

$$c(M_1[Q] ;_{M_1[N']} P_1[N'], \dots, M_p[Q] ;_{M_p[N']} P_p[N']),$$

which by lifting again yields the desired result:

$$c(M_1[Q], \dots, M_p[Q]) ;_{c(M_1[N'], \dots, M_p[N'])} c(P_1[N'], \dots, P_p[N']).$$

The case where something actually happens is  $P = r(P_1, \dots, P_p)$ , with  $r \in X(G \vdash M_0, M'_0 : A)$  and each  $\Delta \vdash P_i : M_i \rightarrow M'_i : G_i$ . Then, the left-hand side is

$$r(P_1[N], \dots, P_p[N]) ;_{M_0[M_1, \dots, M_p][N]} M'_0[M'_1, \dots, M'_p][Q].$$

By lifting, omitting indices of vertical compositions, we have

$$r(P_1[N], \dots, P_p[N]) \equiv r(M_1[N], \dots, M_p[N]); M'_0[P_1[N], \dots, P_p[N]].$$

Observing that  $M'_0[M'_1, \dots, M'_p][Q] = M'_0[M'_1[Q], \dots, M'_p[Q]]$ , the whole is  $\equiv$ -related to

$$\begin{aligned} & r(M_1[N], \dots, M_p[N]); \\ & M'_0[P_1[N], \dots, P_p[N]]; \\ & M'_0[M'_1[Q], \dots, M'_p[Q]], \end{aligned}$$

i.e., by lifting (inductively):

$$\begin{aligned} & r(M_1[N], \dots, M_p[N]); \\ & M'_0[(P_1[N]; M'_1[Q]), \dots, (P_p[N]; M'_1[Q])]. \end{aligned}$$

This is by induction hypothesis  $\equiv$ -related to

$$\begin{aligned} & r(M_1[N], \dots, M_p[N]); \\ & M'_0[(M_1[Q]; P_1[N']), \dots, (M_1[Q]; P_p[N'])], \end{aligned}$$

i.e., by lifting again to

$$\begin{aligned} & r(M_1[N], \dots, M_p[N]); \\ & M'_0[M_1[Q], \dots, M_1[Q]]; \\ & M'_0[P_1[N'], \dots, P_p[N']]. \end{aligned}$$

The second lifting rule then yields

$$\begin{aligned} & r(M_1[Q], \dots, M_p[Q]); \\ & M'_0[P_1[N'], \dots, P_p[N']]. \end{aligned}$$

And then the first lifting rule yields

$$\begin{aligned} & M_0[M_1[Q], \dots, M_1[Q]]; \\ & r(M_1[N'], \dots, M_p[N']); \\ & M'_0[P_1[N'], \dots, P_p[N']], \end{aligned}$$

so, by the second lifting rule again:

$$\begin{aligned} & M_0[M_1[Q], \dots, M_1[Q]]; \\ & r(P_1[N'], \dots, P_p[N']), \end{aligned}$$

i.e., the right-hand side. □

## 6. CARTESIAN CLOSED 2-CATEGORIES

**6.1. Definition.** In a 2-category  $\mathcal{C}$ , a diagram  $A \xleftarrow{p} C \xrightarrow{q} B$  is a *product diagram* iff for all object  $D$ , the induced functor

$$\mathcal{C}(D, C) \xrightarrow{\Delta} \mathcal{C}(D, C) \times \mathcal{C}(D, C) \xrightarrow{\mathcal{C}(D, p) \times \mathcal{C}(D, q)} \mathcal{C}(D, A) \times \mathcal{C}(D, B)$$

is an isomorphism. Because this family of functors is 2-natural in  $D$ , the inverse functors will also be 2-natural.

Similarly, an object  $1$  of  $\mathcal{C}$  is *terminal* iff for all  $D$  the unique functor

$$\mathcal{C}(D, 1) \xrightarrow{!} 1$$

is an isomorphism (where the right-hand  $1$  is the terminal category).

**Definition 4.** A 2-category with finite products, or *fp 2-category*, is a 2-category  $\mathcal{C}$ , equipped with a terminal object and a 2-functor

$$\mathcal{C} \times \mathcal{C} \xrightarrow{\times} \mathcal{C},$$

plus, for all  $A$  and  $B$ , a product diagram

$$A \xleftarrow{p} A \times B \xrightarrow{q} B.$$

In such an fp 2-category  $\mathcal{C}$ , given objects  $A$  and  $B$ , an *exponential* for them is a pair of an object  $B^A$  and a morphism  $ev: A \times B^A \rightarrow B$ , such that for all  $D$ , the functor

$$\begin{array}{ccc}
 & \mathcal{C}(A, A) \times \mathcal{C}(D, B^A) & \xrightarrow{\times} & \mathcal{C}(A \times D, A \times B^A) \\
 (id_{A^!}, id) \nearrow & & & \searrow \mathcal{C}(A \times D, ev) \\
 \mathcal{C}(D, B^A) & & & \mathcal{C}(A \times D, B)
 \end{array}$$

is an isomorphism. As above, because this family of functors is 2-natural in  $D$ , the inverse functors will also be 2-natural.

**Definition 5.** A cartesian closed 2-category, or *cartesian closed 2-category*, is an fp 2-category, equipped with a choice of exponentials for all pairs of objects. The category  $2CCat$  has cartesian closed 2-categories as objects, and strictly structure-preserving functors between them as morphisms.

## 7. MAIN ADJUNCTION

**7.1. Right adjoint.** Given a cartesian closed 2-category  $\mathcal{C}$ , define  $\mathcal{V}(\mathcal{C}) = (\mathcal{C}_0, \mathcal{C}_1, \mathcal{C}_2)$  as follows. First, let as in Section 2  $(\mathcal{C}_0, \mathcal{C}_1) = \mathcal{V}_1(\mathcal{C})$ , and recall the canonical  $\mathcal{L}_0$  and  $\mathcal{L}_1$ -algebra structures  $h_0$  and  $h_1$ . Let then the reduction rules in  $\mathcal{C}_2(G \vdash M, N : A)$  be the 2-cells in  $\mathcal{C}(h_0(G), h_0(A))(h_1(M), h_1(N))$ , abbreviated to  $\mathcal{C}(G, A)(M, N)$  in the sequel.

This signature  $\mathcal{V}\mathcal{C}$  has a canonical  $\mathcal{L}$ -algebra structure  $h_2: \mathcal{L}(\mathcal{V}\mathcal{C}) \rightarrow \mathcal{V}\mathcal{C}$ , which we define by induction over terms in Figure 2. In the case for  $\lambda$ ,  $\varphi$  denotes the structure isomorphism  $\mathcal{C}((\prod \Gamma) \times A, B) \cong \mathcal{C}(\prod \Gamma, B^A)$ .

In order for the definition to make sense as a morphism  $\mathcal{L}(\mathcal{V}\mathcal{C}) \rightarrow \mathcal{V}\mathcal{C}$ , we have to check its compatibility with the equations. We have first:

**Lemma 3.** For all  $\Delta \vdash Q : N \rightarrow N' : \Gamma$  and  $\Gamma \vdash P : M \rightarrow M' : A$  in  $\mathcal{L}(\mathcal{V}\mathcal{C})$ ,

$$\begin{array}{ccc}
 \Delta \begin{array}{c} \xrightarrow{M[N]} \\ \Downarrow h_2(P[Q]) \\ \xrightarrow{M'[N']} \end{array} A & = & \Delta \begin{array}{c} \xrightarrow{N} \\ \Downarrow h_2(Q) \\ \xrightarrow{N'} \end{array} \Gamma \begin{array}{c} \xrightarrow{M} \\ \Downarrow h_2(P) \\ \xrightarrow{M'} \end{array} A.
 \end{array}$$

*Proof.* By induction on  $P$  and the axioms for cartesian closed 2-categories.  $\square$

**Lemma 4.** Any two equated reductions are mapped to the same 2-cell in  $\mathcal{C}$ .

*Proof.* We proceed by induction on the proof of the considered equation. The rules of Figure 3 hold because, in  $\mathcal{C}$ , vertical composition is associative and unital, and equality is a congruence. The beta rule is less easy, so we spell it out.

The left-hand reduction is interpreted in  $\mathcal{C}$  as

$$\begin{array}{ccc}
 & (\varphi M, N) & \\
 & \curvearrowright & \\
 \prod \Gamma & \Downarrow (\varphi P, Q) & B^A \times A \xrightarrow{ev} B \\
 & \curvearrowleft & \\
 & (\varphi M', N') & 
 \end{array}$$

which is equal to

$$\begin{aligned}
& (G \vdash x_i : x_i \rightarrow x_i : G_i) \mapsto (id_{\pi_i} : \pi_i \rightarrow \pi_i : \prod G \rightarrow G_i) \\
& (G \vdash () : () \rightarrow () : 1) \mapsto (id_! : ! \rightarrow ! : \prod G \rightarrow 1) \\
& (\Gamma \vdash c(P_1, \dots, P_n) : c(M_1, \dots, M_n) \rightarrow c(N_1, \dots, N_n) : A) \mapsto \\
& \quad (M_1, \dots, M_n) \\
& \quad \prod \Gamma \begin{array}{c} \curvearrowright \\ \Downarrow P \\ \curvearrowleft \end{array} \prod G \xrightarrow{c} A \quad (c \in \mathcal{C}_1(G, A), P = (P_1, \dots, P_n)) \\
& \quad (N_1, \dots, N_n) \\
& (\Gamma \vdash r(P_1, \dots, P_n) : M[M_1, \dots, M_n] \rightarrow N[N_1, \dots, N_n] : A) \mapsto \\
& \quad (M_1, \dots, M_n) \\
& \quad \prod \Gamma \begin{array}{c} \curvearrowright \\ \Downarrow P \\ \curvearrowleft \end{array} \prod G \begin{array}{c} \xrightarrow{M} \\ \Downarrow r \\ \xrightarrow{N} \end{array} A \quad (P = (P_1, \dots, P_n)) \\
& \quad (N_1, \dots, N_n) \\
& (G \vdash P ;_{M_2} Q : M_1 \rightarrow M_3 : A) \mapsto \prod G \begin{array}{c} \xrightarrow{M_1} \\ \Downarrow P \\ \Downarrow Q \\ \xrightarrow{M_3} \end{array} A \\
& (\Gamma \vdash \lambda x : A. P : \lambda x : A. M \rightarrow \lambda x : A. N : B^A) \mapsto \varphi(P : M \rightarrow N : (\prod \Gamma) \times A \rightarrow B) \\
& \quad (M, N) \\
& (\Gamma \vdash PQ : MN \rightarrow M'N' : B) \mapsto \prod \Gamma \begin{array}{c} \curvearrowright \\ \Downarrow (P, Q) \\ \curvearrowleft \end{array} B^A \times A \xrightarrow{ev} B \\
& \quad (M', N') \\
& (\Gamma \vdash (P, Q) : (M, N) \rightarrow (M', N') : A \times B) \mapsto \prod \Gamma \begin{array}{c} \curvearrowright \\ \Downarrow (P, Q) \\ \curvearrowleft \end{array} A \times B \\
& \quad (M', N') \\
& (\Gamma \vdash \pi_{A,B} P : \pi_{A,B} M \rightarrow \pi_{A,B} N : A) \mapsto \prod \Gamma \begin{array}{c} \xrightarrow{M} \\ \Downarrow P \\ \xrightarrow{N} \end{array} A \times B \xrightarrow{\pi} A \\
& (\Gamma \vdash \pi'_{A,B} P : \pi'_{A,B} M \rightarrow \pi'_{A,B} N : B) \mapsto \prod \Gamma \begin{array}{c} \xrightarrow{M} \\ \Downarrow P \\ \xrightarrow{N} \end{array} A \times B \xrightarrow{\pi'} B
\end{aligned}$$

FIGURE 2. The  $\mathcal{L}$ -algebra structure on  $\mathcal{V}(\mathcal{C})$

$$\begin{array}{ccccc}
& & (id, N) & & \varphi M \times A \\
& \curvearrowright & \rightarrow & \curvearrowright & \\
\Pi \Gamma & & \Downarrow (id, Q) & & \Downarrow \varphi P \times A \\
& \curvearrowleft & \rightarrow & \curvearrowleft & \\
& & (id, N') & & \varphi M' \times A
\end{array}
\quad
\begin{array}{c}
B^A \times A \xrightarrow{ev} B
\end{array}$$

which is in turn equal (by cartesian closedness of  $\mathcal{C}$ ) to:

$$\begin{array}{ccccc}
& & (id, N) & & M \\
& \curvearrowright & \rightarrow & \curvearrowright & \\
\Pi \Gamma & & \Downarrow (id, Q) & & \Downarrow P \\
& \curvearrowleft & \rightarrow & \curvearrowleft & \\
& & (id, N') & & M'
\end{array}
\quad
\begin{array}{c}
B
\end{array}$$

and hence to the right-hand side of the equation by Lemma 3. The other beta and eta rules similarly hold by the properties of products, internal homs, and terminal object in  $\mathcal{C}$ .

The lifting rules hold by (particular cases of) the interchange law in  $\mathcal{C}$  and functoriality of the structural isomorphisms

$$\mathcal{C}(A \times B, C) \cong \mathcal{C}(B, C^A) \quad \text{and} \quad \mathcal{C}(C, A \times B) \cong \mathcal{C}(C, A) \times \mathcal{C}(C, B),$$

which concludes the proof.  $\square$

This assignment extends to cartesian closed functors and we have:

**Proposition 9.**  $\mathcal{V}$  is a functor  $2\text{CCCat} \rightarrow \text{Sig}$ .

**7.2. Left adjoint.** Given an  $\mathcal{L}$ -algebra  $h: \mathcal{L}(X) \rightarrow X$ , we now construct a cartesian closed 2-category  $\mathcal{F}(X, h)$ . It has:

- objects the types in  $\mathcal{L}_0(X_0)$ ;
- 1-cells  $A \rightarrow B$  the terms in  $\mathcal{L}_1(X_0, X_1)(A, B)$ ;
- 2-cells  $M \rightarrow N: A \rightarrow B$  the reduction rules in  $X_2(M, N)$ .

We then must define the cartesian closed 2-category structure, and we start with the 2-category structure. Composition of 1-cells  $A \xrightarrow{M} B \xrightarrow{N} C$  is defined to be  $A \xrightarrow{N[M]} C$ . Vertical composition of 2-cells

$$\begin{array}{ccc}
& M_1 & \\
& \curvearrowright & \\
A & \xrightarrow{M_2} & B \\
& \curvearrowleft & \\
& M_3 & \\
& \Downarrow \beta & \\
& \Downarrow \alpha & \\
& \curvearrowright & \\
& M_4 & 
\end{array}$$

is given by  $h(\eta(\alpha);_{M_2} \eta(\beta))$ .

Horizontal composition of 2-cells

$$(7.1) \quad \begin{array}{ccccc}
& M & & N & \\
& \curvearrowright & & \curvearrowright & \\
A & \xrightarrow{\quad} & B & \xrightarrow{\quad} & C \\
& \curvearrowleft & & \curvearrowleft & \\
& M' & & N' & \\
& \Downarrow \alpha & & \Downarrow \beta & 
\end{array}$$

is obtained as  $h(\beta(\eta(\alpha)))$ .

To show that this yields a 2-category structure, the only non obvious point is the interchange law. We deal with it using the following series of results. First, consider the *left whiskering*

$$\begin{array}{ccccc}
& & M & & \\
& \curvearrowright & \downarrow \alpha & \curvearrowleft & \\
A & & & & B \xrightarrow{N} C \\
& \curvearrowleft & & \curvearrowright & \\
& & M' & & 
\end{array}$$

of a 2-cell  $\alpha$  by a 1-cell  $N$ , i.e., the composition  $id_N \circ \alpha = h((h(N))(\eta(\alpha)))$ .

**Lemma 5.** *We have:  $h((h(N))(\eta(\alpha))) = h(N[\eta(\alpha)])$ .*

*Proof.* Indeed, consider the term  $N(\eta(\eta(\alpha)))$  in  $\mathcal{L}(\mathcal{L}(X))$ . Its images by  $h \circ \mathcal{L}(h)$  and  $h \circ \mu$  coincide, and are respectively  $h((h(N))(\eta(\alpha)))$ , i.e.,  $id_N \circ \alpha$ , and  $h(N[\eta(\alpha)])$ .  $\square$

Similarly, consider the *right whiskering*

$$\begin{array}{ccccc}
& & & & N \\
& & & \curvearrowright & \\
A \xrightarrow{M} & B & & \downarrow \gamma & C \\
& & & \curvearrowleft & \\
& & & & N'
\end{array}$$

of a 2-cell  $\gamma$  by a 1-cell  $M$ , i.e., the composition  $\gamma \circ id_N = h(\gamma(\eta(h(M))))$ .

**Lemma 6.** *We have:  $h(\gamma(\eta(h(M)))) = h(\gamma(M))$ .*

*Proof.* Consider  $(\eta\gamma)(\eta M)$  in  $\mathcal{L}(\mathcal{L}(X))$ . Its images by  $h \circ \mathcal{L}(h)$  and  $h \circ \mu$  coincide, and are respectively  $h(\gamma(\eta(h(M))))$  and  $h(\gamma(M))$ .  $\square$

Now, we prove that the two sensible ways of mimicking horizontal composition using whiskering coincide with actual horizontal composition:

**Lemma 7.** *For any cells as in (7.1),*

$$(\beta \circ id_M) ; (id_{N'} \circ \alpha) = \beta \circ \alpha = (id_N \circ \alpha) ; (\beta \circ id_{M'}).$$

*Proof.* Consider first the reduction  $\eta(\beta(M)) ; \eta(N'[\eta(\alpha)])$  in  $\mathcal{L}(\mathcal{L}(X))$ . Taking  $h \circ \mathcal{L}(h)$  and  $h \circ \mu$  as above respectively yields

- $h(\eta(h(\beta(M)))) ; \eta(h(N'[\alpha]))$ , and
- $h(\beta(M) ; N'[\eta(\alpha)]) = h(\beta(\eta(\alpha)))$ ,

hence the left-hand equality. Then consider  $\eta(N[\eta(\alpha)]) ; \eta(\gamma(M'))$ . Evaluating as before yields the right-hand equality.  $\square$

Finally, consider any configuration like:

$$\begin{array}{ccccc}
& & M & & \\
& \curvearrowright & \downarrow \alpha & \curvearrowleft & \\
A & \xrightarrow{M'} & & \downarrow \beta & B \xrightarrow{N} C \\
& \curvearrowleft & & & \\
& & M'' & & 
\end{array}$$

**Lemma 8.** *We have  $(id_N \circ \alpha) ; (id_N \circ \beta) = id_N \circ (\alpha ; \beta)$ .*

*Proof.* Consider  $\eta(N[\eta(\alpha)]) ; \eta(N[\eta(\beta)])$ . Evaluating yields equality of

- $h(\eta(h(N[\eta(\alpha)]))) ; \eta(h(N[\eta(\beta)]))$ , i.e., the left-hand side, and
- $h(N[\eta(\alpha)] ; N[\eta(\beta)])$ , i.e.,  $h(N[\eta(\alpha) ; \eta(\beta)])$  by lifting.

But now consider  $N[\eta(\alpha) ; \eta(\beta)]$ . Evaluating yields equality of

- $h(N[\eta(\alpha) ; \eta(\beta)])$ , as above, and
- $h(N[\eta(h(\eta(\alpha) ; \eta(\beta)))])$ , i.e.,  $h(N[\eta(\alpha) ; \beta])$  (where  $\alpha ; \beta$  denotes vertical composition in our candidate 2-category), i.e., the right-hand side.  $\square$

**Lemma 9.** *The interchange law holds, i.e., for all reduction rules as in*

$$\begin{array}{ccc}
& M_1 & N_1 \\
A \begin{array}{c} \curvearrowright \\ \xrightarrow{M_2} \\ \curvearrowleft \end{array} & \Downarrow \alpha & \begin{array}{c} \curvearrowright \\ \xrightarrow{N_2} \\ \curvearrowleft \end{array} B \\
& M_3 & N_3 \\
& \Downarrow \beta & \Downarrow \theta \\
& & C,
\end{array}$$

we have

$$(\gamma; \theta) \circ (\alpha; \beta) = (\gamma \circ \alpha); (\theta \circ \beta).$$

*Proof.* By the previous results, we have

$$\begin{aligned}
& (\gamma; \theta) \circ (\alpha; \beta) \\
&= ((\gamma; \theta) \circ M_1); (N_3 \circ (\alpha; \beta)) \\
&= (\gamma \circ M_1); (\theta \circ M_1); (N_3 \circ \alpha); (N_3 \circ \beta) \\
&= (\gamma \circ M_1); (N_2 \circ \alpha); (\theta \circ M_2); (N_3 \circ \beta) \\
&= (\gamma \circ \alpha); (\theta \circ \beta).
\end{aligned}$$

□

Now, let us show cartesian closedness. We have a bijection of hom-sets  $\mathcal{L}_1(X)(C \vdash A \times B) \cong \mathcal{L}_1(X)(C \vdash A) \times \mathcal{L}_1(X)(C \vdash B)$ , given by

$$\begin{aligned}
\mathcal{L}_1(X)(C \vdash A \times B) &\rightarrow \mathcal{L}_1(X)(C \vdash A) \times \mathcal{L}_1(X)(C \vdash B) \\
M &\mapsto \pi M, \pi' M
\end{aligned}$$

and

$$\begin{aligned}
\mathcal{L}_1(X)(C \vdash A) \times \mathcal{L}_1(X)(C \vdash B) &\rightarrow \mathcal{L}_1(X)(C \vdash A \times B) \\
M, N &\mapsto (M, N).
\end{aligned}$$

These are mutually inverse thanks to the beta and eta rules for products in the simply-typed  $\lambda$ -calculus.

On 2-hom-sets, we have

$$\begin{aligned}
\mathcal{L}(X)(C \vdash M, N: A \times B) &\rightarrow \mathcal{L}(X)(C \vdash \pi M, \pi N: A) \times \mathcal{L}(X)(C \vdash \pi' M, \pi' N: B) \\
P &\mapsto \pi P, \pi' P
\end{aligned}$$

and (omitting  $C$ )

$$\begin{aligned}
\mathcal{L}(X)(M_1, N_1: A) \times \mathcal{L}(X)(M_2, N_2: B) &\rightarrow \mathcal{L}(X)((M_1, M_2), (N_1, N_2): A \times B) \\
P_1, P_2 &\mapsto (P_1, P_2),
\end{aligned}$$

which are mutually inverse thanks to the beta and eta rules for products in Figure 4. We use these to define the desired isomorphism  $(u, v)$

$$X_2(C \vdash M, N: A \times B) \cong X_2(C \vdash \pi M, \pi N: A) \times X_2(C \vdash \pi' M, \pi' N: B),$$

as in the diagrams

$$\begin{array}{ccc}
X_2(M, N) & \xrightarrow{u} & X_2(\pi M, \pi N) \times X_2(\pi' M, \pi' N) \\
\eta \downarrow & & \uparrow h \times h \\
\mathcal{L}(X)(M, N) & \xrightarrow{\cong} & \mathcal{L}(X)(\pi M, \pi N) \times \mathcal{L}(X)(\pi' M, \pi' N)
\end{array}$$

and

$$\begin{array}{ccc}
X_2(\pi M, \pi N) \times X_2(\pi' M, \pi' N) & \xrightarrow{v} & X_2(M, N) \\
\eta \times \eta \downarrow & & \uparrow h \\
\mathcal{L}(X)(\pi M, \pi N) \times \mathcal{L}(X)(\pi' M, \pi' N) & \xrightarrow{\cong} & \mathcal{L}(X)(M, N).
\end{array}$$

Starting from  $r \in X_2(M, N)$ , we obtain

$$v(u(r)) = h(\eta(h(\pi(\eta(r))))), \eta(h(\pi'(\eta(r))))).$$

But consider  $(\eta(\pi\eta(r)), \eta(\pi'\eta(r)))$  in  $\mathcal{L}(\mathcal{L}X)$ ; its images by  $h \circ \mathcal{L}h$  and  $h \circ \mu$  are respectively:

- $h(\eta(h(\pi(\eta r))), \eta(h(\pi'(\eta r))))$ , and
- $h(\pi\eta(r), \pi'\eta(r))$ , i.e.,  $h(\eta(r))$ , i.e.,  $r$ ,

which must be equal because  $h$  is an  $\mathcal{L}$ -algebra, hence  $v \circ u = id$ .

Conversely, starting from  $(r, s) \in X_2(M_1, M_2) \times X_2(N_1, N_2)$ , we obtain the pair with components

$$h(\pi(\eta(h(\eta(r), \eta(s)))))) \quad \text{and} \quad h(\pi'(\eta(h(\eta(r), \eta(s))))).$$

Considering  $\pi(\eta(\eta(r), \eta(s))) \in \mathcal{L}(\mathcal{L}(X))$ , its images by  $h \circ \mathcal{L}(h)$  and  $h \circ \mu$  are respectively:

- $h(\pi(\eta(h(\eta(r), \eta(s))))))$ , and
- $h(\pi(\eta(r), \eta(s))) = h(\eta(r)) = r$ .

As above, they must be equal, and by symmetry the second component is  $s$ , and we have proved  $u \circ v = id$ . Similar reasoning for the terminal object and internal homs leads to:

**Proposition 10.** *This yields a cartesian closed 2-category structure on  $\mathcal{C}$ .*

This extends to morphisms of  $\mathcal{L}$ -algebras, so we have constructed a functor  $\mathcal{F}: \mathcal{L}\text{-Alg} \rightarrow 2\text{CCCat}$ .

**7.3. Adjunction.** Consider any  $\mathcal{L}$ -algebra  $(X, h)$ . What does  $(Y, k) = \mathcal{V}(\mathcal{F}(X, h))$  look like? Sorts in  $Y_0$  are types in  $\mathcal{L}_0(X_0)$ . Operations  $Y_1(G \vdash A)$  are terms in  $\mathcal{L}_1(X_0, X_1)(\prod G \vdash A)$ . Reduction rules in  $Y_2(G \vdash M, N: B)$  are reductions in  $\mathcal{L}(X)(\prod G \vdash M', N': B)$ , where  $M' = M[\pi_1 x/x_1, \dots, \pi_n x/x_n]$  (and similarly for  $N'$ ).

Let  $\eta_X$  send:

- each sort  $\iota \in X_0$  to the type  $\iota \in \mathcal{L}_0(X_0)$ ,
- each operation  $c \in X(G \vdash A)$  to the term  $c(\pi_1 x, \dots, \pi_n x)$ , and
- each reduction rule  $r \in X_2(G \vdash M, N: A)$  to the reduction  $x: \prod G \vdash r(\pi_1 x, \dots, \pi_n x): M' \rightarrow N': A$ .

**Theorem 2.** *This  $\eta$  is a natural transformation which is the unit of an adjunction*

$$\begin{array}{ccc} & \mathcal{F} & \\ & \curvearrowright & \\ \mathcal{L}\text{-Alg} & \perp & 2\text{CCCat} \\ & \curvearrowleft & \\ & \mathcal{V} & \end{array}$$

*Proof.* Consider any morphism  $f: (X, h) \rightarrow \mathcal{V}(\mathcal{C})$ , and let  $(Y, k) = \mathcal{V}(\mathcal{F}(X, h))$  and  $\mathcal{V}(\mathcal{C}) = (\mathcal{C}_0, \mathcal{C}_1, h_2: \mathcal{C}_2 \rightarrow \mathcal{C}_1)$ . We now define a uniquely determined cartesian closed functor  $f': \mathcal{F}(X, h) \rightarrow \mathcal{C}$  making the triangle

$$\begin{array}{ccc} X & \xrightarrow{\eta_X} & \mathcal{V}(\mathcal{F}(X)) \\ & \searrow f & \downarrow \mathcal{V}(f') \\ & & \mathcal{V}(\mathcal{C}) \end{array}$$

commute.

On objects, it is determined by induction: on sorts by  $f_0$ , and on type constructors by the requirement that  $f'$  be cartesian closed. On morphisms, it is similarly determined by  $f_1$  and  $f'$  being cartesian closed. On 2-cells, define  $f'$  to be  $f_2: X_2(A \vdash M, N: B) \rightarrow \mathcal{C}(f'(A), f'(B))(f'(M), f'(N))$ , which is also the only possible choice from  $f$ .

We thus only have to show that  $f'$  is cartesian closed, which follows by  $f$  being a morphism of  $\mathcal{L}$ -algebras. For example, to show that binary products of reductions are preserved, consider  $r \in X_2(C \vdash M_1, M_2: A)$  and  $s \in X_2(C \vdash N_1, N_2: B)$ . Their product in  $\mathcal{F}(X)$  is obtained by considering the atomic reductions  $x: C \vdash r(x): M_1 \rightarrow M_2: A$  and  $x: C \vdash s(x): N_1 \rightarrow N_2: B$  and taking  $h(r(x), s(x))$ , which is sent by  $f_2$  to  $f_2(h(r(x), s(x)))$ . But, because  $f$  is a morphism of  $\mathcal{L}$ -algebras, this is the same as  $h_2((f_2(r))(x), (f_2(s))(y))$ , which is by definition (i.e., Figure 2) the product  $(f_2(r), f_2(s))$  in  $\mathcal{C}$ .  $\square$

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## APPENDIX A. EQUATIONS ON REDUCTIONS

*Current address:* CNRS, Université de Savoie

Congruence	
$\frac{\Gamma \vdash P : M \rightarrow N : A}{\Gamma \vdash P \equiv P : M \rightarrow N : A}$	$\frac{\Gamma \vdash P \equiv Q : M \rightarrow N : A}{\Gamma \vdash Q \equiv P : M \rightarrow N : A}$
$\frac{\Gamma \vdash P_1 \equiv P_2 : M \rightarrow N : A \quad \Gamma \vdash P_2 \equiv P_3 : M \rightarrow N : A}{\Gamma \vdash P_1 \equiv P_3 : M \rightarrow N : A}$	
$\frac{\Gamma \vdash P \equiv P' : M_1 \rightarrow M_2 : A \quad \Gamma \vdash Q \equiv Q' : M_2 \rightarrow M_3 : A}{\Gamma \vdash (P ;_{M_2} Q) \equiv (P' ;_{M_2} Q') : M_1 \rightarrow M_3 : A}$	
$(r \in X(G \vdash M, N : A))$	
$\frac{\Gamma \vdash P_1 \equiv Q_1 : M_1 \rightarrow N_1 : G_1 \quad \dots \quad \Gamma \vdash P_n \equiv Q_n : M_n \rightarrow N_n : G_n}{\Gamma \vdash r(P_1, \dots, P_n) \equiv r(Q_1, \dots, Q_n) : M[M_1, \dots, M_n] \rightarrow N[N_1, \dots, N_n] : A}$	
$(c \in X(G \vdash A))$	
$\frac{\Gamma \vdash P_1 \equiv Q_1 : M_1 \rightarrow N_1 : G_1 \quad \dots \quad \Gamma \vdash P_n \equiv Q_n : M_n \rightarrow N_n : G_n}{\Gamma \vdash c(P_1, \dots, P_n) \equiv c(Q_1, \dots, Q_n) : c(M_1, \dots, M_n) \rightarrow c(N_1, \dots, N_n) : A}$	
$\frac{\Gamma, x : A \vdash P \equiv Q : M \rightarrow N : B}{\Gamma \vdash (\lambda x : A. P) \equiv (\lambda x : A. Q) : \lambda x : A. M \rightarrow \lambda x : A. N : B^A}$	
$\frac{\Gamma \vdash P \equiv P' : M \rightarrow M' : B^A \quad \Gamma \vdash Q \equiv Q' : N \rightarrow N' : A}{\Gamma \vdash (PQ) \equiv (P'Q') : MN \rightarrow M'N' : B}$	
$\frac{\Gamma \vdash P \equiv P' : M \rightarrow M' : A \quad \Gamma \vdash Q \equiv Q' : N \rightarrow N' : B}{\Gamma \vdash (P, Q) \equiv (P', Q') : (M, N) \rightarrow (M', N') : A \times B}$	
$\frac{\Gamma \vdash P \equiv Q : M \rightarrow N : A \times B}{\Gamma \vdash (\pi_{A,B} P) \equiv (\pi_{A,B} Q) : \pi_{A,B} M \rightarrow \pi_{A,B} N : A}$	
$\frac{\Gamma \vdash P \equiv Q : M \rightarrow N : A \times B}{\Gamma \vdash (\pi'_{A,B} P) \equiv (\pi'_{A,B} Q) : \pi'_{A,B} M \rightarrow \pi'_{A,B} N : A}$	
Category	
$\frac{\Gamma \vdash P_1 : M_1 \rightarrow M_2 : A \quad \Gamma \vdash P_2 : M_2 \rightarrow M_3 : A \quad \Gamma \vdash P_3 : M_3 \rightarrow M_4 : A}{\Gamma \vdash (P_1 ;_{M_2} (P_2 ;_{M_3} P_3)) \equiv ((P_1 ;_{M_2} P_2) ;_{M_3} P_3) : M_1 \rightarrow M_4 : A}$	
$\frac{\Gamma \vdash P : M \rightarrow N : A}{\Gamma \vdash (P ;_N N) \equiv P : M \rightarrow N : A}$	$\frac{\Gamma \vdash P : M \rightarrow N : A}{\Gamma \vdash (M ;_M P) \equiv P : M \rightarrow N : A}$

FIGURE 3. Equations on reductions (Congruence and category)

Beta and eta	
$\frac{\Gamma, x : A \vdash P : M \rightarrow M' : B \quad \Gamma \vdash Q : N \rightarrow N' : A}{\Gamma \vdash ((\lambda x : A. P)Q) \equiv P[Q/x] : (\lambda x : A. M)N \rightarrow M'[N'/x] : B}$	
$\frac{\Gamma \vdash P \equiv M \rightarrow N : B^A}{\Gamma \vdash P \equiv \lambda x : A. (Px) : M \rightarrow N : B^A} \quad (x \notin \Gamma)$	
$\frac{\Gamma \vdash P : M_1 \rightarrow M_2 : A \quad \Gamma \vdash Q : N_1 \rightarrow N_2 : B}{\Gamma \vdash \pi(P, Q) \equiv P : \pi(M_1, N_1) \rightarrow M_2 : A}$	
$\frac{\Gamma \vdash P : M_1 \rightarrow M_2 : A \quad \Gamma \vdash Q : N_1 \rightarrow N_2 : B}{\Gamma \vdash \pi'(P, Q) \equiv P : \pi'(M_1, N_1) \rightarrow N_2 : A}$	
$\frac{\Gamma \vdash P : (M_1, N_1) \rightarrow (M_2, N_2) : A \times B}{\Gamma \vdash P \equiv (\pi P, \pi' P) : (M_1, N_1) \rightarrow (M_2, N_2) : A \times B}$	$\frac{\Gamma \vdash P : M \rightarrow N : 1}{\Gamma \vdash P \equiv () : M \rightarrow N : 1}$
Lifting	
$\frac{(r \in X(\Gamma \vdash \langle M_1, M_2 \rangle : A)) \quad \Delta \vdash P : N_1 \rightarrow N_2 : \Gamma \quad \Delta \vdash Q : N_2 \rightarrow N_3 : \Gamma}{\Gamma \vdash r(P ;_{N_2} Q) \equiv M_1[P] ;_{M_1[N_2]} r(Q) : M_1[N_1] \rightarrow M_2[N_3] : A}$	
$\frac{(r \in X(\Gamma \vdash \langle M_1, M_2 \rangle : A)) \quad \Delta \vdash P : N_1 \rightarrow N_2 : \Gamma \quad \Delta \vdash Q : N_2 \rightarrow N_3 : \Gamma}{\Gamma \vdash r(P ;_{N_2} Q) \equiv r(P) ;_{M_2[N_2]} M_2[Q] : M_1[N_1] \rightarrow M_2[N_3] : A}$	
$\frac{\Gamma \vdash P : M_1 \rightarrow M_2 : G \quad \Gamma \vdash Q : M_2 \rightarrow M_3 : G}{\Gamma \vdash (c(P ;_{M_2} Q)) \equiv (c(P) ;_{c(M_2)} c(Q)) : M_1 \rightarrow M_3 : A} \quad (c \in X(G \vdash A))$	
$\frac{\Gamma, x : A \vdash P : M_1 \rightarrow M_2 : B \quad \Gamma, x : A \vdash Q : M_2 \rightarrow M_3 : B}{\Gamma \vdash (\lambda x : A. (P ;_{M_2} Q)) \equiv ((\lambda x : A. P) ;_{\lambda x : A. M_2} (\lambda x : A. Q)) : \lambda x : A. M_1 \rightarrow \lambda x : A. M_3 : B^A}$	
$\frac{\Gamma \vdash P : M_1 \rightarrow M_2 : B^A \quad \Gamma \vdash P' : M_2 \rightarrow M_3 : B^A \quad \Gamma \vdash Q : N_1 \rightarrow N_2 : A \quad \Gamma \vdash Q' : N_2 \rightarrow N_3 : A}{\Gamma \vdash ((P ;_{M_2} P')(Q ;_{N_2} Q')) \equiv ((PQ) ;_{M_2 N_2} (P'Q')) : M_1 N_1 \rightarrow M_3 N_3 : B}$	
$\frac{\Gamma \vdash P : M_1 \rightarrow M_2 : A \quad \Gamma \vdash P' : M_2 \rightarrow M_3 : A \quad \Gamma \vdash Q : N_1 \rightarrow N_2 : B \quad \Gamma \vdash Q' : N_2 \rightarrow N_3 : B}{\Gamma \vdash ((P ;_{M_2} P'), (Q ;_{N_2} Q')) \equiv ((P, Q) ;_{(M_2, N_2)} (P', Q')) : (M_1, N_1) \rightarrow (M_3, N_3) : A \times B}$	
$\frac{\Gamma \vdash P : M_1 \rightarrow M_2 : A \times B \quad \Gamma \vdash Q : M_2 \rightarrow M_3 : A \times B}{\Gamma \vdash (\pi_{A,B}(P ;_{M_2} Q)) \equiv (\pi_{A,B} P ;_{\pi_{A,B} M_2} \pi_{A,B} Q) : M_1 \rightarrow M_3 : A}$	
$\frac{\Gamma \vdash P : M_1 \rightarrow M_2 : A \times B \quad \Gamma \vdash Q : M_2 \rightarrow M_3 : A \times B}{\Gamma \vdash (\pi'_{A,B}(P ;_{M_2} Q)) \equiv (\pi'_{A,B} P ;_{\pi'_{A,B} M_2} \pi'_{A,B} Q) : M_1 \rightarrow M_3 : B}$	

FIGURE 4. Equations on reductions (beta-eta and lifting)