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## Tools for parsimonious edge-colouring of graphs with maximum degree three

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#### Abstract

The notion of a  $\delta$ -minimum edge-colouring was introduced by J-L. Fouquet (in its french PhD Thesis [2]). Here we present some structural properties of  $\delta$ -minimum edge-colourings, some of them being first presented in the above thesis. The paper serves as an auxiliary tool for another paper submitted by the authors to Graphs and Combinatorics.

### 1 Introduction

Throughout this note, we shall be concerned with connected graphs with maximum degree 3. We know by Vizing's theorem [6] that these graphs can be edge-coloured with 4 colours. Let  $\phi : E(G) \to \{\alpha, \beta, \gamma, \delta\}$  be a proper edgecolouring of G. It is often of interest to try to use one colour (say  $\delta$ ) as few as possible. When it is optimal following this constraint, we shall say that such a parsimonious edge-colouring  $\phi$  is  $\delta$ -minimum . In [1] we gave without proof (in French) results on  $\delta$ -minimum edge-colourings of cubic graphs. Some of them have been obtained later and independently by Steffen [4] and [5]. The purpose of Section 2 is to give with their proofs those results as structural properties of  $\delta$ -minimum edge-colourings.

An edge colouring of G using colours  $\alpha, \beta, \gamma, \delta$  is said to be  $\delta$ -improper provided that adjacent edges having the same colours (if any) are coloured with  $\delta$ . It is clear that a proper edge colouring of G is a particular  $\delta$ -improper edge colouring, obviously that is also true for a  $\delta$ -minimum edge-colouring. For a proper or  $\delta$ -improper edge colouring  $\phi$  of G, it will be convenient to denote  $E_{\phi}(x)$   $(x \in \{\alpha, \beta, \gamma, \delta\})$  the set of edges coloured with x by  $\phi$ . For  $x, y \in \{\alpha, \beta, \gamma, \delta\}$ ,  $x \neq y, \phi(x, y)$  is the partial subgraph of G spanned by these two colours, that is  $E_{\phi}(x) \cup E_{\phi}(y)$  (this subgraph being a union of paths and even cycles where the colours x and y alternate). Since any two  $\delta$ -minimum edge-colourings of G have the same number of edges coloured  $\delta$  we shall denote by s(G) this number (the colour number as defined in[4]).

As usual, for any undirected graph G, we denote by V(G) the set of its vertices and by E(G) the set of its edges. A strong matching C in a graph G is a matching C such that there is no edge of E(G) connecting any two edges of C, or, equivalently, such that C is the edge-set of the subgraph of G induced on the vertex-set V(C).

## 2 Structural properties of $\delta$ -minimum edge-colourings

The graph G considered in the following series of Lemmas will have maximum degree 3.

**Lemma 1** [2] Any 2-factor of G contains at least s(G) disjoint odd cycles.

**Proof** Assume that we can find a 2-factor of G with k < s(G) odd cycles. Then let us colour the edges of this 2-factor with  $\alpha$  and  $\beta$ , except one edge (coloured  $\delta$ ) on each odd cycle of our 2-factor and let us colour the remaining edges by  $\gamma$ . We get hence a new edge colouring  $\phi$  with  $E_{\phi}(\delta) < s(G)$ , impossible.

**Lemma 2** [2] Let  $\phi$  be a  $\delta$ -minimum edge-colouring of G. Any edge in  $E_{\phi}(\delta)$  is incident to  $\alpha$ ,  $\beta$  and  $\gamma$ . Moreover each such edge has one end of degree 2 and the other of degree 3 or the two ends of degree 3.

**Proof** Any edge of  $E_{\phi}(\delta)$  is certainly adjacent to  $\alpha, \beta$  and  $\gamma$ . Otherwise this edge could be coloured with the missing colour and we should obtain an edge colouring  $\phi'$  with  $|E_{\phi'}(\delta)| < |E_{\phi}(\delta)|$ .

Lemma 3 below was proven in [3], we give its proof for the sake of completeness.

**Lemma 3** [3] Let  $\phi$  be a  $\delta$ -improper colouring of G then there exists a proper colouring of  $G \phi'$  such that  $E_{\phi'}(\delta) \subseteq E_{\phi}(\delta)$ 

**Proof** Let  $\phi$  be a  $\delta$ -improper edge colouring of G. If  $\phi$  is a proper colouring, we are done. Hence, assume that uv and uw are coloured  $\delta$ . If d(u) = 2 we can change the colour of uv to  $\alpha, \beta$  or  $\gamma$  since v is incident to at most two colours in this set.

If d(u) = 3 assume that the third edge uz incident to u is also coloured  $\delta$ , then we can change the colour of uv for the same reason as above.

If uz is coloured with  $\alpha, \beta$  or  $\gamma$ , then v and w are incident to the two remaining colours of the set  $\{\alpha, \beta, \gamma\}$  otherwise one of the edges uv, uw can be recoloured with the missing colour. W.l.o.g., consider that uz is coloured  $\alpha$  then v and w are incident to  $\beta$  and  $\gamma$ . Since u has degree 1 in  $\phi(\alpha, \beta)$  let P be the path of  $\phi(\alpha, \beta)$  which ends on u. We can assume that v or w (say v) is not the other end vertex of P. Exchanging  $\alpha$  and  $\beta$  along P does not change the colours incident to v. But now uz is coloured  $\beta$  and we can change the colour of uv to  $\alpha$ .

In each case, we get hence a new  $\delta$ -improper edge colouring  $\phi_1$  with  $E_{\phi_1}(\delta) \subsetneq E_{\phi}(\delta)$ . Repeating this process leads us to construct a proper edge colouring of G with  $E_{\phi'}(\delta) \subseteq E_{\phi}(\delta)$  as claimed.  $\Box$ 

**Lemma 4** [2] Let  $\phi$  be a  $\delta$ -minimum edge-colouring of G. For any edge  $e = uv \in E_{\phi}(\delta)$  with  $d(v) \leq d(u)$  there is a colour  $x \in \{\alpha, \beta, \gamma\}$  present at v and a colour  $y \in \{\alpha, \beta, \gamma\} - \{x\}$  present at u such that one of connected components

of  $\phi(x, y)$  is a path of even length joining the two ends of e. Moreover, if d(v) = 2, then both colours of  $\{\alpha, \beta, \gamma\} - \{x\}$  satisfy the above assertion.

**Proof** Without loss of generality suppose that  $x = \gamma$  is present at v and  $\alpha, \beta$  are present at u (see Lemma 2). Then u is an endvertex of paths in both  $\phi(\alpha, \gamma)$  and  $\phi(\beta, \gamma)$ , while there is  $y \in \{\alpha, \beta\}$  such that v is an endvertex of a path in  $\phi(x, y)$ . Without loss of generality assume that both u and v are endvertices of paths in  $\phi(\alpha, \gamma)$ . If these paths are disjoint, we exchange the colours  $\alpha$  and  $\gamma$  on the path with endvertex u and then recolour e with  $\alpha$ ; this yields a contradiction to the  $\delta$ -minimality of  $\phi$ . To conclude the proof note that if d(v) = 2, then v is an endvertex of paths in both  $\phi(\alpha, \gamma)$  and  $\phi(\beta, \gamma)$ .

An edge of  $E_{\phi}(\delta)$  is in  $A_{\phi}$  when its ends can be connected by a path of  $\phi(\alpha, \beta)$ ,  $B_{\phi}$  by a path of  $\phi(\beta, \gamma)$  and  $C_{\phi}$  by a path of  $\phi(\alpha, \gamma)$ . From Lemma 4 it is clear that if d(u) = 3 and d(v) = 2 for an edge  $e = uv \in E_{\phi}(\delta)$ , the  $A_{\phi}$ ,  $B_{\phi}$  and  $C_{\phi}$  are not pairwise disjoint; indeed, if the colour  $\gamma$  is present at the vertex v, then  $e \in A_{\phi} \cap B_{\phi}$ .

When  $e \in A_{\phi}$  we can associate to e the odd cycle  $C_{A_{\phi}}(e)$  obtained by considering the path of  $\phi(\alpha, \beta)$  together with e. We define in the same way  $C_{B_{\phi}}(e)$  and  $C_{C_{\phi}}(e)$  when e is in  $B_{\phi}$  or  $C_{\phi}$ .

**Lemma 5** [2] If G is a cubic graph then  $|A_{\phi}| \equiv |B_{\phi}| \equiv |C_{\phi}| \equiv s(G) \pmod{2}$ .

**Proof**  $\phi(\alpha, \beta)$  contains  $2|A_{\phi}| + |B_{\phi}| + |C_{\phi}|$  vertices of degree 1 and must be even. Hence we get  $|B_{\phi}| \equiv |C_{\phi}| \pmod{2}$ . In the same way we get  $|A_{\phi}| \equiv |B_{\phi}| \pmod{2}$  leading to  $|A_{\phi}| \equiv |B_{\phi}| \equiv |C_{\phi}| \equiv s(G) \pmod{2}$ In the following lemma we consider an edge in  $A_{\phi}$ , an analogous result holds true whenever we consider edges in  $B_{\phi}$  or  $C_{\phi}$  as well.

**Lemma 6** [2] Let  $\phi$  be a  $\delta$ -minimum edge-colouring of G and let e be an edge in  $A_{\phi}$  then for any edge  $e' \in C_{A_{\phi}}(e)$  there is a  $\delta$ -minimum edge-colouring  $\phi'$  such that  $E_{\phi'}(\delta) = E_{\phi}(\delta) - \{e\} \cup \{e'\}, e' \in A_{\phi'}$  and  $C_{A_{\phi}}(e) = C_{A_{\phi'}}(e')$ . Moreover, each edge outside  $C_{A_{\phi}}(e)$  but incident with this cycle is coloured  $\gamma$ ,  $\phi$  and  $\phi'$  only differ on the edges of  $C_{A_{\phi}}(e)$ .

**Proof** By exchanging colours  $\delta$  and  $\alpha$  and  $\delta$  and  $\beta$  successively along the cycle  $C_{A_{\phi}}(e)$ , we are sure to obtain an edge colouring preserving the number of edges coloured  $\delta$ . Since we have supposed that  $\phi$  is  $\delta$ -minimum,  $\phi$  is proper. At each step, the resulting edge colouring remains to be  $\delta$ -minimum and hence proper. Hence, there is no edge coloured  $\delta$  incident with  $C_{A_{\phi}}(e)$ , which means that every such edge is coloured with  $\gamma$ .

We can perform these exchanges until e' is coloured  $\delta$ . In the  $\delta$ -minimum edge-colouring  $\phi'$  hence obtained, the two ends of e' are joined by a path of  $\phi(\alpha, \beta)$ . Which means that e' is in  $A_{\phi}$  and  $C_{A_{\phi}}(e) = C_{A_{\phi'}}(e')$ .

For each edge  $e \in E_{\phi}(\delta)$  (where  $\phi$  is a  $\delta$ -minimum edge-colouring of G) we can associate one or two odd cycles following the fact that e is in one or two sets from among  $A_{\phi}$ ,  $B_{\phi}$  or  $C_{\phi}$ . Let  $\mathcal{C}$  be the set of odd cycles associated to edges in  $E_{\phi}(\delta)$ .

**Lemma 7** [2] For each cycle  $C \in C$ , there are no two consecutive vertices with degree two.

**Proof** Otherwise, we exchange colours along C in order to put the colour  $\delta$  on the corresponding edge and, by Lemma 2, this is impossible in a  $\delta$ -minimum edge-colouring.

**Lemma 8** [2] Let  $e_1, e_2 \in E_{\phi}(\delta)$ , such that  $e_1 \neq e_2$  and let  $C_1, C_2 \in \mathcal{C}$  be such that  $C_1 \neq C_2$ ,  $e_1 \in E(C_1)$  and  $e_2 \in E(C_2)$  then  $C_1$  and  $C_2$  are (vertex) disjoint.

**Proof** If  $e_1$  and  $e_2$  are contained in the same set  $A_{\phi}$ ,  $B_{\phi}$ , or  $C_{\phi}$ , we are done since their respective ends are joined by an alternating path of  $\phi(x, y)$  for some two colours x and y in  $\{\alpha, \beta, \gamma\}$ .

Without loss of generality assume that  $e_1 \in A_{\phi}$  and  $e_2 \in B_{\phi}$ . Assume moreover that there exists an edge e such that  $e \in C_1 \cap C_2$ . We have hence an edge  $f \in C_1$  with exactly one end on  $C_2$ . We can exchange colours on  $C_1$  in order to put the colour  $\delta$  on f, which is impossible by Lemma 6.

**Lemma 9** [2] Let  $e_1 = uv_1$  be an edge of  $E_{\phi}(\delta)$  such that  $v_1$  has degree 2 in G. Then  $v_1$  is the only vertex in N(u) of degree 2 and for any other edge  $e_2 \in E_{\phi}(\delta)$ ,  $\{e_1, e_2\}$  induces a  $2K_2$ .

**Proof** We have seen in Lemma 2 that  $uv_1$  has one end of degree 3 while the other has degree 2 or 3. Hence, we have d(u) = 3 and  $d(v_1) = 2$ . Let  $v_2$  and  $v_3$  the other neighbours of u. We can suppose without loss of generality that  $uv_2$  is coloured  $\alpha$ ,  $uv_3$  is coloured  $\beta$  and, finally  $v_1$  is incident to an edge coloured  $\gamma$ , say  $v_1v_4$ .

Assume first that  $d(v_2) = 2$ . An alternating path of  $\phi(\beta, \gamma)$  using the edge  $uv_3$  ends with the vertex  $v_1$ , moreover  $v_2$  is incident to an edge coloured  $\gamma$  since an alternating path of  $\phi(\alpha, \gamma)$  using the edge  $uv_2$  ends with  $v_1$  (see Lemma 4), then, exchanging the colours along the component of  $\phi(\beta, \gamma)$  containing  $v_2$  allows us to colour  $uv_2$  with  $\gamma$  and  $uv_1$  with  $\alpha$ . A new edge colouring  $\phi'$  so obtained is such that  $|E_{\phi'}(\delta)| \leq |E_{\phi}(\delta)| - 1$ , impossible.

Thus,  $d(v_2) = 3$ , and, by symmetry,  $d(v_3) = 3$ . We know that  $e_1 \in B_{\phi} \cap C_{\phi}$ (see Lemma 4). By Lemma 6, since  $e_1 \in C_{B_{\phi}}(e_1)$  the edges incident to  $v_3$  in  $C_{B_{\phi}}(e_1)$  are coloured with  $\beta$  and  $\gamma$  and the third edge incident to  $v_3$  is coloured  $\alpha$ . Similarly the vertex  $v_2$  being on  $C_{B_{\phi}}(e_1)$  is incident to colours  $\alpha$  and  $\gamma$  and the third edge incident to  $v_2$  is coloured  $\beta$ . Moreover the vertex  $v_4$  being on both cycles  $C_{B_{\phi}}(e_1)$  and  $C_{C_{\phi}}(e_1)$  is incident to colours  $\alpha$ ,  $\beta$ ,  $\gamma$ . Hence no edge coloured  $\delta$  can be incident to  $v_2$  nor  $v_3$  nor  $v_4$ . It follows that for any edge  $e_2$  in  $E_{\phi}(\delta) - \{e_1\}$ , the set  $\{e_1, e_2\}$  induces a  $2K_2$ .

**Lemma 10** [2] Let  $e_1$  and  $e_2$  be two edges of  $E_{\phi}(\delta)$   $e_1 \neq e_2$ . If  $e_1$  and  $e_2$  are contained in two distinct sets of  $A_{\phi}, B_{\phi}, C_{\phi}$  then  $\{e_1, e_2\}$  induces a  $2K_2$  otherwise  $e_1, e_2$  are joined by at most one edge.

**Proof** By Lemma 9 one can suppose that all vertices incident with  $e_1$ ,  $e_2$  are of degree 3 so that there is exactly one of the sets  $A_{\phi}$ ,  $B_{\phi}$  or  $C_{\phi}$  containing  $e_1$  (respectively  $e_2$ ).

Assume in the first stage that  $e_1 \in A_{\phi}$  and  $e_2 \in B_{\phi}$ . Since  $C_{A_{\phi}}(e_1)$  and  $C_{B_{\phi}}(e_2)$  are disjoint by Lemma 8, we know that  $e_1$  and  $e_2$  have no common vertex. The edges having exactly one end in  $C_{A_{\phi}}(e_1)$  are coloured  $\gamma$  while those having exactly one end in  $C_{B_{\phi}}(e)$  are coloured  $\alpha$ . Hence there is no edge between  $e_1$  and  $e_2$  as claimed.

Assume in the second stage that  $e_1 = u_1v_1, e_2 = u_2v_2 \in A_{\phi}$ . Since  $C_{A_{\phi}}(e_1)$ and  $C_{B_{\phi}}(e_2)$  are disjoint by Lemma 8, we can consider that  $e_1$  and  $e_2$  have no common vertex. The edges having exactly one end in  $C_{A_{\phi}}(e_1)$  or  $C_{B_{\phi}}(e_2)$ are coloured  $\gamma$ . Assume that  $u_1u_2$  and  $v_1v_2$  are edges of G. We may suppose without loss of generality that  $u_1$  and  $u_2$  are incident to  $\alpha$  while  $v_1$  and  $v_2$  are incident to  $\beta$  (if necessary, colours  $\alpha$  and  $\beta$  can be exchanded on  $C_{A_{\phi}}(e_1)$  and  $C_{B_{\phi}}(e_2)$ ). We know that  $u_1u_2$  and  $v_1v_2$  are coloured  $\gamma$ . Let us colour  $e_1$  and  $e_2$ with  $\gamma$  and  $u_1u_2$  with  $\beta$  and  $v_1v_2$  with  $\alpha$ . We get a new edge colouring  $\phi'$  where  $|E_{\phi'}(\delta)| \leq |E_{\phi}(\delta)| - 2$ , contradiction since  $\phi$  is a  $\delta$ -minimum edge-colouring.  $\Box$ 

**Lemma 11** [2] Let  $e_1, e_2$  and  $e_3$  be three distinct edges of  $E_{\phi}(\delta)$  contained in the same set  $A_{\phi}, B_{\phi}$  or  $C_{\phi}$ . Then  $\{e_1, e_2, e_3\}$  induces a subgraph with at most four edges.

**Proof** Without loss of generality assume that  $e_1 = u_1v_1, e_2 = u_2v_2$  and  $e_3 = u_3v_3 \in A_{\phi}$ . From Lemma 10 we have just to suppose that (up to the names of vertices)  $u_1u_3 \in E(G)$  and  $v_1v_2 \in E(G)$ . Possibly, by exchanging the colours  $\alpha$  and  $\beta$  along the 3 disjoint paths of  $\phi(\alpha, \beta)$  joining the ends of each edge  $e_1, e_2$  and  $e_3$ , we can suppose that  $u_1$  and  $u_3$  are incident to  $\beta$  while  $v_1$  and  $v_2$  are incident to  $\alpha$ . Let  $\phi'$  be obtained from  $\phi$  when  $u_1u_3$  is coloured with  $\alpha$ ,  $v_1v_2$  with  $\beta$  and  $u_1v_1$  with  $\gamma$ . It is easy to check that  $\phi'$  is a proper edge-colouring with  $|E_{\phi'}(\delta)| \leq |E_{\phi}(\delta)| - 1$ , contradiction since  $\phi$  is a  $\delta$ -minimum edge-colouring.

Let us summarize most of the above results in a single Theorem.

**Theorem 12** Let G be a graph of maximum degree 3 and  $\phi$  be a  $\delta$ -minimum colouring of G. Then the following hold.

- 1.  $E_{\phi}(\delta) = A_{\phi} \cup B_{\phi} \cup C_{\phi}$  where an edge e in  $A_{\phi}$  ( $B_{\phi}$ ,  $C_{\phi}$  respectively) belongs to a uniquely determined cycle  $C_{A_{\phi}}(e)$  ( $C_{B_{\phi}}(e)$ ,  $C_{C_{\phi}}(e)$  respectively) with precisely one edge coloured  $\delta$  and the other edges being alternately coloured  $\alpha$  and  $\beta$  ( $\beta$  and  $\gamma$ ,  $\alpha$  and  $\gamma$  respectively).
- 2. Each edge having exactly one vertex in common with some edge in  $A_{\phi}$ ( $B_{\phi}$ ,  $C_{\phi}$  respectively) is coloured  $\gamma$  ( $\alpha$ ,  $\beta$ , respectively).
- 3. The multiset of colours of edges of  $C_{A_{\phi}}(e)$  ( $C_{B_{\phi}}(e)$ ,  $C_{C_{\phi}}(e)$  respectively) can be permuted to obtain a (proper)  $\delta$ -minimum edge-colouring of G in which the colour  $\delta$  is moved from e to an arbitrarily prescribed edge.
- 4. No two consecutive vertices of  $C_{A_{\phi}}(e)$  ( $C_{B_{\phi}}(e)$ ,  $C_{C_{\phi}}(e)$  respectively) have degree 2.
- 5. The cycles from 1 that correspond to distinct edges of  $E_{\phi}(\delta)$  are vertexdisjoint.

6. If the edges  $e_1, e_2, e_3 \in E_{\phi}(\delta)$  all belong to  $A_{\phi}$  ( $B_{\phi}, C_{\phi}$  respectively), then the set  $\{e_1, e_2, e_3\}$  induces in G a subgraph with at most 4 edges.

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