



CENTRAL EUROPEAN OLYMPIAD IN INFORMATICS

Münster, Germany
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Day 1: **pearls**

Input File: pearls.in
Output File: -
Source File: pearls.pas/.c/.cpp

100 Points
Time limit: 1 s
Memory limit: 16 MB

Solution

The task is about finding winning moves. A winning move is a move which leads the player to a win, if he plays optimal in the following moves.

If we suppose that both players play optimally, each move is either a winning move or a losing move as there is no draw in this game.

A move is further a winning move, if a dwarf can give the necklace to a dwarf of his tribe who has a winning move, or if a dwarf can give the necklace to a dwarf of the opposite tribe who does not have a winning move.

A game situation is determined by the state of the necklace and the current dwarf:

$$\begin{aligned} \text{hasWinningMove}(c_1 c_2 \dots c_k, d) := \\ \exists d' \in d.\text{list}(c_1) : d.\text{color} = d'.\text{color} \wedge \text{hasWinningMove}(c_2 \dots c_k, d') \quad \vee \\ d.\text{color} \neq d'.\text{color} \wedge \neg \text{hasWinningMove}(c_2 \dots c_k, d') \end{aligned}$$

$\text{hasWinningMove}(c_1, d)$ is always true, as there is only the diamond left.

With this information you can easily compute who is going to win and which moves to take, by using bottom-up dynamic programming. The dimensions of the dynamic programming matrix are the indices in the necklace ($1 \dots L$) and the dwarfs ($1 \dots N$). The main loop is over the necklace indices in reverse order, i.e. $L, \dots, 1$.