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Computing Chebyshev knot diagrams

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ABSTRACT

A Chebyshev curve $\mathcal{C}(a, b, c, \varphi)$ has a parametrization of the form $x(t) = T_a(t)$; $y(t) = T_b(t)$; $z(t) = T_c(t + \varphi)$, where a, b, c are integers, $T_n(t)$ is the Chebyshev polynomial of degree n and $\varphi \in \mathbf{R}$. When $\mathcal{C}(a, b, c, \varphi)$ has no double points, it defines a polynomial knot. We determine all possible knots when a, b and c are given.

Keywords

Zero dimensional systems, Chebyshev curves, Lissajous knots, polynomial knots, factorization of Chebyshev polynomials, minimal polynomial

1. INTRODUCTION

It is known that every knot may be obtained from a polynomial embedding $\mathbf{R} \rightarrow \mathbf{R}^3$ ([19, 5]).

Chebyshev knots are polynomial analogue to Lissajous knots that have been studied by many authors (see [1, 2, 8, 9, 13]). All knots are not Lissajous (for example the trefoil and the figure-eight knot). In [10], it is proved that any knot $K \subset \mathbf{R}^3$ is a Chebyshev knot, that is to say there exist positive integers a, b, c and a real φ such that K is isotopic to the curve

$$\mathcal{C}(a, b, c, \varphi) : x = T_a(t), y = T_b(t), z = T_c(t + \varphi),$$

where T_n is the Chebyshev polynomial of degree n . This is our motivation for the study of curves $\mathcal{C}(a, b, c, \varphi)$, $\varphi \in \mathbf{R}$.

In [10], the proof uses theorems on braids by Hoste, Zirbel and Lamm ([8, 13]), and a density argument (Kronecker theorem).

In [12], we developed an effective method to enumerate all the knots $\mathcal{C}(a, b, c, \varphi)$, $\varphi \in \mathbf{R}$ where $a = 3$ or $a = 4$, a and b coprime. This method was based on continued fraction expansion theory in order to get the minimal b , on resultant computations in order to determine the critical values φ for which $\mathcal{C}(a, b, c, \varphi)$ is singular, and on multi-precision interval arithmetic to determine the knot type of $\mathcal{C}(a, b, c, \varphi)$. Our goal was to give an exhaustive list of the minimal parametrizations for the first 95 rational knots with less than 10 crossings. We obtained in [12] almost every minimal parametrizations. For 6 of these knots, we know the minimal b and we know that the corresponding c must be > 300 .

In this paper, we develop a more efficient algorithm. It will give the parametrization of the 6 missing knots in [12] and

allows to compute all diagrams corresponding to $\mathcal{C}(a, b, c, \varphi)$, $\varphi \in \mathbf{R}$. One motivation is first to achieve the exhaustive list of certified minimal Chebyshev parametrizations for the first 95 rational knots. Another is to provide a certified tool for the study of polynomial curves topology.

Let us first recall some basic facts about knots.

Knot diagrams

The projection of the Chebyshev space curve $\mathcal{C}(a, b, c, \varphi)$ on the (x, y) -plane is the Chebyshev curve $\mathcal{C}(a, b) : x = T_a(t)$, $y = T_b(t)$. If a and b are coprime integers, the curve $\mathcal{C}(a, b, c, \varphi)$ is singular if and only if it has some double points. It is convenient to consider the polynomials in $\mathbf{Q}[s, t, \varphi]$

$$P_n = \frac{T_n(t) - T_n(s)}{t - s}, \quad Q_n = \frac{T_n(t + \varphi) - T_n(s + \varphi)}{t - s}. \quad (1)$$

We see that $\mathcal{C}(a, b, c, \varphi)$ is a knot iff $\{(s, t), P_a(s, t) = P_b(s, t) = Q_c(s, t, \varphi) = 0\}$ is empty.

We shall study the diagram of the curve $\mathcal{C}(a, b, c, \varphi)$, that is to say the plane projection $\mathcal{C}(a, b)$ onto the (x, y) -plane and the nature (under/over) of the crossings over the double points of $\mathcal{C}(a, b)$. There are two cases of crossing: the right twist and the left twist (see [15] and Figure 1). In [12],

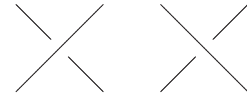


Figure 1: The right twist and the left twist

we showed that the nature of the crossing over the double point $A_{\alpha, \beta}$ corresponding to parameters $(t = \cos(\alpha + \beta), s = \cos(\alpha - \beta), \alpha = \frac{i\pi}{a}, \beta = \frac{j\pi}{b})$, is given by the sign of

$$D(s, t, \varphi) = Q_c(s, t, \varphi)P_{b-a}(s, t, \varphi). \quad (2)$$

$D(s, t, \varphi) > 0$ if and only if the crossing is a right twist.

Note that the crossing points of the Chebyshev curve $\mathcal{C}(a, b) : x = T_a(t)$, $y = T_b(t)$ lie on the $(b-1)$ vertical lines $T'_b(x) = 0$ and on the $(a-1)$ horizontal lines $T'_a(y) = 0$. We can represent the knot diagram of $\mathcal{C}(a, b, c, \varphi)$ by a billiard trajectory (see [10]), which is a pure combinatorial object. As an example, consider the knots $\bar{5}_2 = \mathcal{C}(4, 5, 7, 0)$, $\bar{5}_2 = \mathcal{C}(5, 6, 7, 0)$, $\bar{4}_1 = \mathcal{C}(3, 5, 7, 0)$. We can represent their diagrams by the billiard trajectories in Figure 2.

When $a = 3$ or $a = 4$, we obtain the diagrams in the Conway normal form. In this case, the knot is rational and can be identified very easily by its Schubert fraction (see [15,

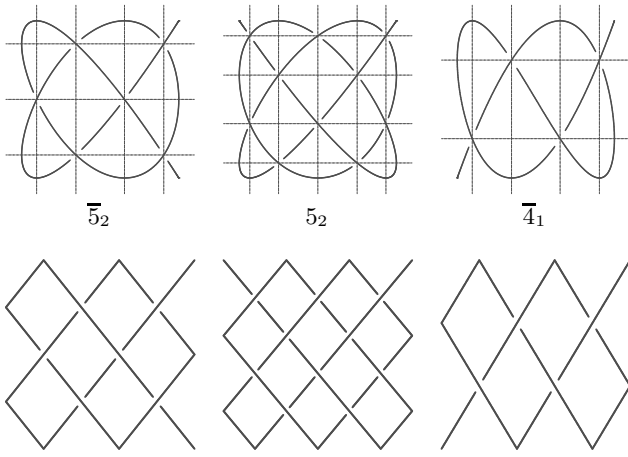


Figure 2: Billiard trajectories

11]). When $b > a \geq 5$, the problem of classification is much more difficult. Nevertheless, the knowledge of the diagrams allows the computation of all the classical invariants, like the Conway, Alexander and Jones polynomials (see [15]).

Summary

Our goal is to compute all diagrams of $\mathcal{C}(a, b, c, \varphi)$, where a, b, c are given integers and $\varphi \in \mathbf{R}$. From the algorithmic point of view, the description of the Chebyshev knots is strongly connected to the resolution of:

$$\mathcal{V}_{a,b,c} = \{P_a(s, t) = 0, P_b(s, t) = 0, Q_c(s, t, \varphi) = 0\}. \quad (3)$$

We first want to determine the set $\mathcal{Z}_{a,b,c}$ of critical values of φ for which the curve $\mathcal{C}(a, b, c, \varphi)$ is singular. Because $\deg_{\varphi} Q_n = n - 1$ and the leading term of Q_n is $2^{n-1}n\varphi^{n-1}$, we showed in [12], that $\mathcal{V}_{a,b,c}$ is 0-dimensional and has at most $(a-1)(b-1)(c-1)$ points. We deduced that $|\mathcal{Z}_{a,b,c}| \leq \frac{1}{2}(a-1)(b-1)(c-1)$.

Let $\mathcal{Z}_{a,b,c} = \{\varphi_1, \dots, \varphi_n\}$. The type of the knot $\mathcal{C}(a, b, c, \varphi)$ is given by its diagram which is constant when φ is in $(\varphi_i, \varphi_{i+1})$, because the crossings do not change in this interval. In order to get all possible knots $\mathcal{C}(a, b, c, \varphi)$, we only need sample points r_i in each $(\varphi_i, \varphi_{i+1})$ and to compute the diagram of $\mathcal{C}(a, b, c, r_i)$.

We can determine a polynomial $R_{a,b,c} \in \mathbf{Z}[\varphi]$ such that $\mathcal{Z}_{a,b,c} = Z(R)$. It can be defined by $(R) = \langle P_a, P_b, Q_c \rangle \cap \mathbf{Q}[\varphi]$ and may be obtained with Gröbner bases ([3]). In [12], we optimized the computation by an ad-hoc elimination based on the properties of the curves for $a = 3$ or $a = 4$. Gröbner bases could fully be substituted by some resultant computations in $\mathbf{Z}[X, \varphi]$, the systems being generic enough. However, this leads to solve systems of very high degree.

In the present paper we decompose the system by working on some (real cyclotomic) extension fields. We show that the system (3) is equivalent to the resolution of $\frac{1}{2}(a-1)(b-1)\lfloor \frac{c}{2} \rfloor$ second-degree polynomials with coefficients in $\mathbf{Q}(\cos \frac{\pi}{a}, \cos \frac{\pi}{b}, \cos \frac{\pi}{c})$. This result is deduced from geometric properties of the implicit Chebyshev curves.

We show some properties of these extensions that allow to simplify the computations. We can represent the coefficients of the polynomials by intervals and certify the resolution. We then easily and independently obtain the roots of the

second-order polynomials and the main difficulty becomes to compare them. A formal method would consist in computing their minimal polynomials over \mathbf{Q} , which is equivalent to the resolution of (3). We use multi-precision interval arithmetic for coding the algebraic numbers $\cos \frac{k\pi}{n}$ as well as the solutions φ we get. If the two intervals are disjoint, the roots are distinct. If not, we can certify whether the resultant of the two second-order polynomials equals 0 or not by Euclidean division.

In section 2, we first describe the Chebyshev polynomials and the link between their factorizations and the minimal polynomials of $\cos \frac{k\pi}{n}$. This allows us to represent efficiently the elements of $\mathbf{Q}(\cos \frac{\pi}{a}, \cos \frac{\pi}{b}, \cos \frac{\pi}{c})$. Along the way, we give an explicit factorization of the Chebyshev polynomials.

In section 3, we recall the definition of Lissajous curves and we give their implicit equations. We study the affine implicit Chebyshev curves $T_n(x) = T_m(y)$ and show that they have $\lfloor \frac{(n,m)}{2} \rfloor + 1$ irreducible components, $\lfloor \frac{(n,m)-1}{2} \rfloor$ being Lissajous curves.

This allows us to deduce an explicit factorization of $R_{a,b,c}$ as the product of second-degree polynomials $P_{\alpha,\beta,\gamma}$, in section 4.

We show in section 5, how to obtain $\mathcal{Z}_{a,b,c}$, the set of roots of $R_{a,b,c}$, with their multiplicities. The general algorithm is described in 6. This allows us to sample all Chebyshev knots $\mathcal{C}(a, b, c, \varphi)$, $\varphi \in \mathbf{R}$, by choosing a rational number r in each component of $\mathbf{R} - \mathcal{Z}_{a,b,c}$.

In section 7, we find an exhaustive and complete list of the minimal parametrization for the first 95 rational knots. The worst case appears with the knot $10_{33} = \mathcal{C}(4, 13, 856, 1/328)$, with $\deg R_{a,b,c} = 15390$. We discuss the efficiency of our algorithms and compare with those of [12].

2. CHEBYSHEV POLYNOMIALS

Chebyshev polynomials and their algebraic properties play a central role here. The curves we will study are defined by Chebyshev polynomials. The algebraic extensions we will consider are spanned by their roots and we need to know their factors. In this section we recall some classical properties of Chebyshev polynomials. We will also show the link between their effective factorization in $\mathbf{Q}[t]$ and the minimal polynomial of $\cos \frac{k\pi}{n}$.

The Chebyshev polynomials of the *first kind* are defined by the second-order linear recurrence

$$T_0 = 1, T_1 = t, T_{n+1} = 2tT_n - T_{n-1}. \quad (4)$$

$T_n \in \mathbf{Z}[t]$ and satisfies the identity $T_n(\cos \theta) = \cos n\theta$, and more generally $T_n \circ T_m = T_{nm}$. We have

$$T_n = 2^{n-1} \prod_{k=0}^{n-1} (t - \cos \frac{(2k+1)\pi}{2n}).$$

Let V_n be the Chebyshev polynomials of the *second kind* defined by the second-order linear recurrence (the same as in (4))

$$V_0 = 0, V_1 = 1, V_{n+1} = 2tV_n - V_{n-1}.$$

$V_n \in \mathbf{Z}[t]$ and satisfies $V_n(\cos \theta) = \frac{\sin n\theta}{\sin \theta}$. We have

$$V_n = 2^{n-1} \prod_{k=1}^{n-1} (t - \cos \frac{k\pi}{n}),$$

and therefore $V_d | V_n$ when $d | n$. Let us summarize some useful results in the following

LEMMA 1. *We have the following properties:*

- $T'_n(t) = 0 \Rightarrow T_n(t) = \pm 1$
- $T_n(t) = \pm 1 \Rightarrow T'_n(t) = 0$ or $t = \pm 1$.
- $T_n(t) = y$ has n real solutions iff $|y| < 1$.
- $T_n(t) = 1$ has $\lfloor \frac{n}{2} \rfloor$ real solutions.
- $T_n(t) = -1$ has $\lfloor \frac{n-1}{2} \rfloor$ real solutions.

PROOF. From $T'_n = nV_n$, we deduce that $t \mapsto T_n(t)$ is monotonic when $|t| \geq \cos \frac{\pi}{n}$, that T_n has $n-1$ local extrema for $t_k = \cos \frac{k\pi}{n}$ and $T_n(t_k) = (-1)^k$. \square

Minimal Polynomial of $\cos \frac{k\pi}{n}$

Let $\zeta_n = e^{\frac{2i\pi}{n}}$. It is well known ([20]) that the degree of $\mathbf{Q}(\zeta_n)$ is $\varphi(n)$ where φ is the Euler function. $\mathbf{Q}(\cos \frac{2\pi}{n}) = \mathbf{Q}(\zeta_n) \cap \mathbf{R}$ and the minimal polynomial over \mathbf{Q} of $\cos \frac{2\pi}{n}$ has degree $\frac{1}{2}\varphi(n)$ when $n > 1$. Its roots are $\cos \frac{2k\pi}{n}$ where k is coprime with n . Consequently, the minimal polynomial M_n of $\cos \frac{\pi}{n}$ has degree $\frac{1}{2}\varphi(2n)$, when $n > 1$. Its roots are $t_k = \cos \frac{k\pi}{n}$ where $(k, n) = 1$, and k is odd. $M_n(-t)$ is the minimal polynomial of $\cos \frac{2\pi}{n}$. The leading coefficient of M_n is $2^{\varphi(2n)/2}$.

Remark. $\cos \frac{k\pi}{n} \in \mathbf{Q}$ iff $\frac{1}{2}\varphi(2n) = 1$ or $n = 1$, that is $n = 1, 2, 3$. In this case we get $2 \cos \frac{k\pi}{n} \in \mathbf{Z}$.

We deduce the following

PROPOSITION 2. *Let P_n be defined by $P_0 = 1, P_1 = 2t-1, P_{n+1} = 2tP_n + P_{n-1}$. Then we have $(-1)^n P_n(-T_2) = V_{2n+1}$ and*

$$P_n = \prod_{d|2n+1} M_d \quad (5)$$

PROOF. We have $P_0(-T_2) = V_1, P_1(-T_2) = -2T_2 - 1 = -V_3$. The sequences V_{2n+1} and $(-1)^n P_n(-T_2)$ satisfy the same recurrence formula: $V_{2n+3} = 2T_2 V_{2n+1} - V_{2n-1}$. Let $d = 2m+1$ be a divisor of $2n+1$ and consider $t = \cos \frac{\pi}{d} = -\cos 2\frac{m\pi}{2m+1}$. We have $(-1)^n P_n(t) = V_{2n+1}(\cos \frac{m\pi}{2m+1}) = 0$. Thus $M_d | P_n$ and we conclude using the fact that $\sum_{d|2n+1} \deg M_d = 1 + \frac{1}{2} \sum_{d|2n+1, d>1} \varphi(2d) = n = \deg P_n$. \square

LEMMA 3. *We have $M_{2^k m} = M_m(T_{2^k})$ if m is odd.*

PROOF. We have $M_m \circ T_{2^k}(\cos \frac{\pi}{2^k m}) = 0$ and $(2^k, m) = 1$ so $M_{2^k m} | M_m(T_{2^k})$. We conclude since $M_{2^k m}$ and $M_m(T_{2^k})$ have same leading term. \square

The relations between the minimal polynomial of $\cos \frac{2\pi}{n}$ and the factorization of $T_{\lfloor \frac{n}{2} \rfloor + 1} - T_{\lfloor \frac{n}{2} \rfloor}$ is known ([20]). Formula (5) together with Lemma 3 give also an algorithm to compute M_n .

The number of factors of T_n is known ([7]). We give here the relation between the Chebyshev polynomials T_n and V_n and the polynomials M_n .

PROPOSITION 4. **Factorization of T_n and V_n .**

We have the following factorizations in irreducible factors

$$V_{2^k(2m+1)} = \prod_{d|2m+1} \left(\prod_{i=1}^k M_d(T_{2^i}) \right) \cdot M_d(t) M_d(-t)$$

$$T_{2^k(2m+1)} = \frac{1}{2} \prod_{d|2m+1} M_d(T_{2^{k+1}})$$

where M_n is the minimal polynomial of $\cos \frac{\pi}{n}$.

PROOF. The factorization of V_n is obtained by comparing its roots with those of $M_d(\pm t)$, when $d | n$. Let d be an odd divisor of n . We write $n = 2^k \cdot d_1 \cdot d$, where d_1 is odd. $\cos \frac{d_1 \pi}{2^n} = \cos \frac{\pi}{2^{k+1}d}$ is a root of T_n so $M_{2^{k+1}d} | T_n$. We deduce the factorization by comparing the leading terms. \square

3. CHEBYSHEV AND LISSAJOUS CURVES

The following proposition will explain the notions of Lissajous and Chebyshev curves.

PROPOSITION 5. *The parametric curve*

$$\mathcal{C} : x = \cos(at), y = \cos(bt + \varphi), t \in \mathbf{C},$$

where a, b are coprime integers (a odd) and $\varphi \in \mathbf{R}$ admits the equation

$$T_b(x)^2 + T_a(y)^2 - 2 \cos(a\varphi) T_b(x) T_a(y) - \sin^2(a\varphi) = 0. \quad (6)$$

1. *If $a\varphi \neq k\pi$, this equation is irreducible. \mathcal{C} is called a Lissajous curve. Its real part is 1-1 parametrized for $t \in [0, 2\pi]$.*

2. *If $a\varphi = k\pi$, this equation is equivalent to $T_b(x) = (-1)^k T_a(y)$. \mathcal{C} is called a Chebyshev curve. It can be parametrized by $x = T_a(t), y = (-1)^k T_b(t)$.*

PROOF. Let $(x, y) \in \mathcal{C}$. We have $T_b(x) = \cos(bat), T_a(y) = \cos(bat + a\varphi)$. Let $\lambda = a\varphi, \theta = abt$. We get $T_a(y) = \cos(\theta + \lambda)$ so $(1 - \cos^2 \theta) \sin^2 \lambda = (\cos \theta \cos \lambda - T_a(y))^2$, that is $(1 - T_a(x))^2 \sin^2 \lambda = (T_a(x) \cos \lambda - T_a(y))^2$, and we deduce our Equation (6).

Conversely, suppose that (x, y) satisfies (6). Let $x = \cos(at)$ where $t \in \mathbf{C}$. We also have $x = \cos a(t + \frac{2k\pi}{a})$. We have $T_b(x) = \cos \theta$. $A = T_a(y)$ is solution of the second-degree equation

$$A^2 - 2 \cos(a\varphi) \cos \theta A - \sin^2(a\varphi) = 0.$$

Consequently, we get $T_a(y) = \cos(\theta \pm a\varphi) = T_a(\cos(\pm bt + \varphi))$. We deduce that $y = \cos(\pm bt + \varphi + \frac{2h\pi}{a}), h \in \mathbf{Z}$. Changing t by $-t$, we can suppose that

$$x = \cos at, y = \cos(bt + \varphi + \frac{2h\pi}{a}).$$

By choosing k such that $kb + h \equiv 0 \pmod{a}$, we get $x = \cos at', y = \cos(bt' + \varphi)$, where $t' = t + \frac{2k\pi}{a}$.

If $a\varphi \not\equiv 0 \pmod{\pi}$. Suppose that Equation (6) factors in $P(x, y)Q(x, y)$. We can suppose, for analyticity reasons, that $P(\cos(at), \cos(bt + \varphi)) = 0$, for $t \in \mathbf{C}$. The curve \mathcal{C} intersects the line $y = 0$ in $2b$ distinct points so $\deg_x P \geq 2b$. Similarly, $\deg_y P \geq 2a$ so that Q is a constant which proves that the equation is irreducible.

If $\cos a\varphi = (-1)^k$, the equation becomes $T_b(x) - (-1)^k T_a(y) = 0$. In this case the curve admits the announced parametrization (see [6, 10] for more details). \square

Remark. If $a = b = 1$, we obtain the Lissajous ellipses. They are the first curves studied by Lissajous ([14]). Let $\mu \not\equiv 0 \pmod{\pi}$. The curve

$$\mathcal{E}_\mu : x^2 + y^2 - 2\cos(\mu)xy - \sin^2(\mu) = 0$$

is an ellipse inscribed in the square $[-1, 1]^2$. It admits the parametrization $x = \cos t$, $y = \cos(t + \mu)$.

The following notation will be useful. Let $E_\mu(x, y) = x^2 + y^2 - 2\cos(\mu)xy - \sin^2(\mu)$ when $\mu \not\equiv 0 \pmod{\pi}$ and $E_0 = x - y$, $E_\pi = x + y$. The Equation (6) is equivalent to $E_{a\varphi}(T_b(x), T_a(y)) = 0$. This shows that the real part of the curve \mathcal{C} (Equation (6)) is inscribed in the square $[-1, 1]^2$. Using Proposition 5 we recover the classical following result.

COROLLARY 6. *The Lissajous curve $x = \cos(at)$, $y = \cos(bt + \varphi)$, ($a\varphi \not\equiv 0 \pmod{\pi}$) has $2ab - a - b$ singular points which are real double points.*

PROOF. singular points of \mathcal{C} satisfy Equation (6) and the system

$$\begin{aligned} T'_b(x)(T_b(x) - T_a(y) \cos a\varphi) &= 0, \\ T'_a(y)(T_a(y) - T_b(x) \cos a\varphi) &= 0. \end{aligned}$$

Suppose that $T'_b(x) = T'_a(y) = 0$ then $T_a^2(y) = T_b^2(x) = 1$ and Equation 6 is not satisfied. Suppose that $T_b(x) - T_a(y) \cos a\varphi = T_a(y) - T_b(x) \cos a\varphi = 0$, then $T_b(x) = T_a(x) = 0$ and Equation 6 is not satisfied. We thus have either $T'_b(x) = 0$ and $T_a(y) - T_b(x) \cos a\varphi = 0$ that gives $(b-1) \times a$ real points because of the classical properties of Chebyshev polynomials, or $T'_a(y) = 0$ and $T_b(x) - T_a(y) \cos a\varphi = 0$ that gives $b \times (a-1)$ real double points. \square

Remark. The study of the double points of Lissajous curves is classical (see [2] for their parameters values). The study of the double points of Chebyshev curves is simpler (see [10]).

COROLLARY 7. *The affine implicit curve $T_n(x) = T_m(y)$ has $\lfloor \frac{n-1}{2} \rfloor \lfloor \frac{m-1}{2} \rfloor + \lfloor \frac{n}{2} \rfloor \lfloor \frac{m}{2} \rfloor$ singular points that are real double points.*

PROOF. The singular points satisfy either $T_n(x) = T_m(y) = 1$ or $T_n(x) = T_m(y) = -1$ and we conclude using Lemma 1. \square

THEOREM 8. Factorization of $T_n(x) - T_m(y)$.

Let $m = ad$, $n = bd$, $(a, b) = 1$ and a odd. We have the factorization

$$T_n(x) - T_m(y) = 2^{d-1} \prod_{k=0}^{\lfloor \frac{d}{2} \rfloor} C_k(x, y)$$

where

$$C_k(x, y) = E_{\frac{2ak\pi}{d}}(T_b(x), T_a(y))$$

is the irreducible equation of the curve $C_k : x = \cos(at)$, $y = \cos(bt + \frac{2k\pi}{d})$, given in Proposition 5.

PROOF. Let \mathcal{C} be the curve $T_n(x) = T_m(y)$. We easily get $C_k \subset \mathcal{C}$ and $C_k \neq C_{k'}$. When $k = 0$, C_0 admits the equation $T_b(x) - T_a(x) = 0$. If $2k = d$, C_k admits the equation $T_b(x) + T_a(y) = 0$. In the other cases, the dominant term in x of C_k is $2^{2b-2}x^{2b}$. If d is even, we deduce that the dominant

term of $\prod_{k=0}^{\lfloor \frac{d}{2} \rfloor} C_k(x, y)$ is $2^{(b-1)d}x^{2n}$ and we get our result in this case. If d is odd, we get the same result. \square

COROLLARY 9. *Let $d = \gcd(a, b)$. The curve $T_b(x) = T_a(y)$ has $\lfloor \frac{d}{2} \rfloor + 1$ components. $\lfloor \frac{d-1}{2} \rfloor$ of them are Lissajous curves.*

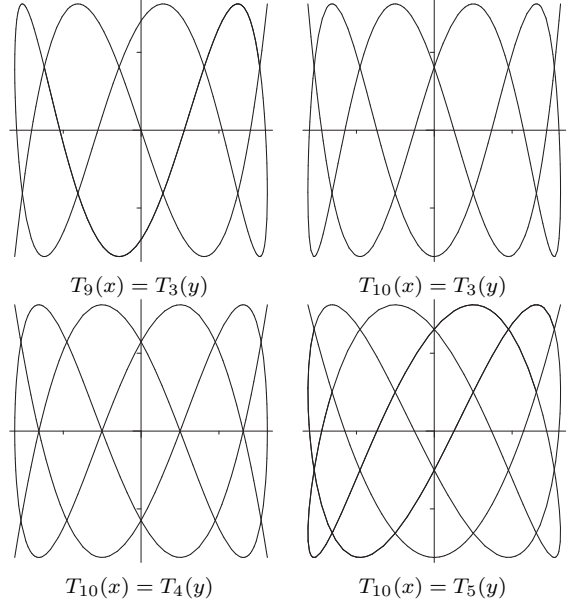


Figure 3: Implicit Chebyshev curves

Theorem 8 is particularly interesting when $m = n = d$ and $a = b = 1$. In this case the curve $T_n(x) = T_n(y)$ is a union of ellipses and some lines. It will be useful for the determination of the double points of Chebyshev space curves. We have

$$\frac{T_n(t) - T_n(s)}{t - s} = 2^{n-1} \prod_{k=1}^{\lfloor \frac{n}{2} \rfloor} E_{\frac{2k\pi}{n}}(s, t). \quad (7)$$

The curve $\frac{T_n(t) - T_n(s)}{t - s} = 0$ has $\lfloor \frac{n}{2} \rfloor$ irreducible components. Note that $\mathcal{E}_{\frac{2k\pi}{n}}$ and $\mathcal{E}_{\frac{2l\pi}{n}}$ intersect at the point $(t, s) = (\cos(\frac{k\pi}{n} + \frac{l\pi}{n}), \cos(\frac{k\pi}{n} - \frac{l\pi}{n}))$ and its symmetric with respect to the lines $s = -t$ and $s = t$. We recover the parametrization of the double points of $x = T_a(t)$, $y = T_b(t)$ that will be very useful for the description of Chebyshev space curves.

PROPOSITION 10 ([10, 12]). *Let a and b be nonnegative coprime integers, a being odd. Let the Chebyshev curve \mathcal{C} be defined by $x = T_a(t)$, $y = T_b(t)$. The pairs (t, s) giving a crossing point are*

$$t = \cos(\frac{j\pi}{b} + \frac{i\pi}{a}), \quad s = \cos(\frac{j\pi}{b} - \frac{i\pi}{a})$$

where $1 \leq i \leq \frac{1}{2}(a-1)$, $1 \leq j \leq b-1$.

4. CRITICAL VALUES

A polynomial $R_{a,b,c} \in \mathbf{Z}[\varphi]$ for which $\mathcal{Z}_{a,b,c} = Z(R)$ can be defined by $\langle R \rangle = \langle P_a, P_b, Q_c \rangle \cap \mathbf{Q}[\varphi]$ and may be obtained with Gröbner bases ([3]).

Example. When $a = 3$, $b = 4$, $c = 5$, we find that

$$R_{a,b,c} = (80\varphi^4 + 60\varphi^2 - 1) \cdot$$

$$(6400\varphi^8 - 3200\varphi^6 + 560\varphi^4 - 80\varphi^2 + 1).$$

There are exactly 6 critical values that are symmetrical

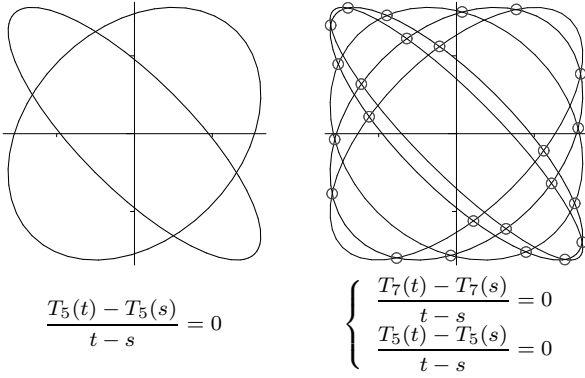


Figure 4: Double points in the parameters space

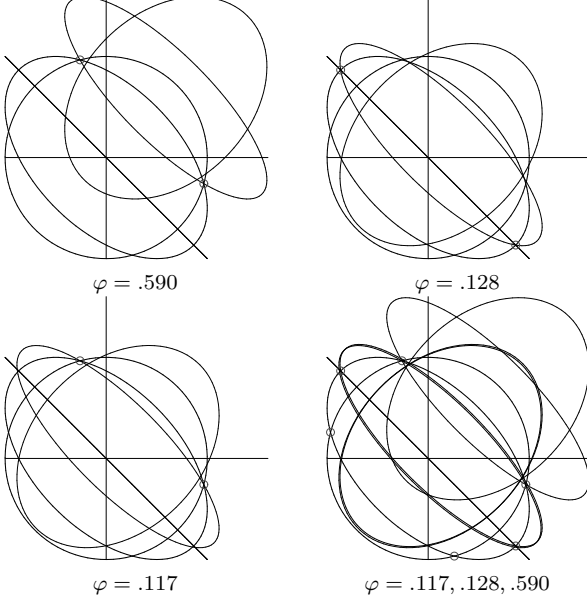


Figure 5: $P_3 = 0, P_4 = 0, Q_5 = 0$

about the origin. For these values of φ , the curve $Q_5(s, t, \varphi) = 0$, which is translated from the curve $P_5(s, t) = 0$ by the vector (φ, φ) , meets the points $\{P_3 = 0, P_4 = 0\}$.

In this part, we use the properties of Chebyshev curves obtained in section 3. We give an explicit formula for the polynomial $R_{a,b,c}$ as a product of univariate polynomials of degree 1 or 2 with coefficients in $\mathbf{Q}(\cos \frac{\pi}{a}, \cos \frac{\pi}{b}, \cos \frac{\pi}{c})$.

PROPOSITION 11. *Let a, b be nonnegative coprime integers and c be an integer. Suppose that a is odd. Let $R_{a,b,c}(\varphi)$ be the polynomial*

$$\prod_{i=1}^{\frac{a-1}{2}} \prod_{j=1}^{b-1} Q_c(\cos(\frac{j}{b} + \frac{i}{a})\pi, \cos(\frac{j}{b} - \frac{i}{a})\pi, \varphi). \quad (8)$$

$R_{a,b,c} \in \mathbf{Z}[\varphi]$ and $\mathcal{C}(a, b, c, \varphi)$ is singular iff $R_{a,b,c}(\varphi) = 0$.

PROOF. $\varphi \in \mathcal{Z}_{a,b,c}$ iff there exists (s, t) such that $P_a(s, t) = P_b(s, t) = 0$ and $Q_c(s, t, \varphi) = 0$. This conditions are equivalent to have $t = \cos(\frac{j\pi}{b} + \frac{i\pi}{a})$ and $s = \cos(\frac{j\pi}{b} - \frac{i\pi}{a})$ and $Q_c(s, t, \varphi) = 0$, for some $1 \leq i \leq \frac{a-1}{2}$ and $1 \leq j \leq b-1$, from Proposition 10.

$Q_c(s, t, \varphi)$ is a symmetric polynomial of $\mathbf{Z}[\varphi][t, s]$. Let $\alpha_i = \frac{i\pi}{a}$, $\beta_j = \frac{j\pi}{b}$ and $s = \cos(\alpha_i + \beta_j)$, $t = \cos(\alpha_i - \beta_j)$. From

$s + t = 2 \cos \alpha_i \cos \beta_j$ and $st = \cos^2 \alpha_i + \cos^2 \beta_j - 1$, we deduce that $Q_c(s, t, \varphi)$ belongs to $\mathbf{Z}[\varphi, \cos \alpha_i][\cos \beta_j]$.

$$R_i = \prod_{j=1}^{b-1} Q_c(\cos(\alpha_i + \beta_j), \cos(\alpha_i - \beta_j), \varphi)$$

belongs to $\mathbf{Z}[\varphi, \cos \alpha_i]$ because the $\cos \beta_j$ are the roots of $V_b \in \mathbf{Z}[t]$. From $Q_c(-s, -t, \varphi) = Q_c(s, t, -\varphi)$ we deduce

that $\prod_{i=1}^{\frac{a-1}{2}} R_i(-\varphi)R_i(\varphi) = \prod_{i=1}^{a-1} R_i(\varphi) \in \mathbf{Z}[\varphi]$. We thus have $R_{a,b,c}^2 \in \mathbf{Z}[\varphi]$ and so it is for $R_{a,b,c}$. \square

Let $s = \cos(\alpha + \beta)$, $t = \cos(\alpha - \beta)$. Using Theorem 8 and Formula (7), we get

$$Q_c(s, t, \varphi) = 2^{c-1} \prod_{k=1}^{\lfloor \frac{c}{2} \rfloor} E_{2k\pi/n}(s, t).$$

Let us consider $P_{\alpha,\beta,\gamma} = \frac{1}{\sin^2 \gamma} E_{2\gamma}(s + \varphi, t + \varphi)$. For $\gamma \neq \frac{\pi}{2}$, $P_{\alpha,\beta,\gamma}$ is

$$\varphi^2 + 2\varphi \cos \alpha \cos \beta + \frac{(\cos^2 \alpha - \cos^2 \gamma)(\cos^2 \beta - \cos^2 \gamma)}{\sin^2 \gamma}$$

and

$$P_{\alpha,\beta,\frac{\pi}{2}} = \varphi + \cos \alpha \cos \beta.$$

We therefore obtain

$$Q_c(\cos(\alpha + \beta), \cos(\alpha - \beta), \varphi) = K \prod_{k=1}^{\lfloor \frac{c}{2} \rfloor} P_{\alpha,\beta,\frac{k\pi}{c}}(\varphi)$$

with $K = 2^{c-1} \prod_{k=1}^c 2 \sin \frac{k\pi}{c} = c2^{c-1}$. We get therefore

$$R_{a,b,c}(\varphi) = K^{\frac{1}{2}(a-1)(b-1)} \prod_{k=1}^{\lfloor \frac{c}{2} \rfloor} \prod_{i=1}^{\frac{a-1}{2}} \prod_{j=1}^{b-1} P_{\frac{i\pi}{a}, \frac{j\pi}{b}, \frac{k\pi}{c}}(\varphi).$$

We have written $R_{a,b,c}$ as the product of second or first-degree polynomials $P_{\alpha,\beta,\gamma}$ in $\mathbf{Q}(\cos \frac{\pi}{a}, \cos \frac{\pi}{b}, \cos \frac{\pi}{c})$.

5. COMPUTING THE CRITICAL VALUES

Our strategy consists in first computing separately the real roots of each $P_{\alpha,\beta,\gamma}$ and then combining these roots to get those of $R_{a,b,c}$. A straightforward approach would be to use interval arithmetic to approximate the various trigonometric expressions, but this would fail when $R_{a,b,c}$ has multiple roots, unless we cannot ensure if some discriminant or some resultant are null.

5.1 Real roots of $P_{\alpha,\beta,\gamma}$

Let $\alpha = \frac{i\pi}{a}$, $\beta = \frac{j\pi}{b}$ and $\gamma = \frac{k\pi}{c}$ with $1 \leq i \leq \frac{a-1}{2}$, $1 \leq j \leq b-1$, $1 \leq k \leq \lfloor \frac{c-1}{2} \rfloor$.

If $\gamma = \frac{\pi}{2}$, the unique root of $P_{\alpha,\beta,\frac{\pi}{2}}$ is $-\cos \alpha \cos \beta$. If $\gamma \neq \frac{\pi}{2}$,

the discriminant of $P_{\alpha,\beta,\gamma}$ is $4 \cos^2 \gamma \left(1 - \frac{\sin^2 \alpha \sin^2 \beta}{\sin^2 \gamma}\right)$. It has the same sign as

$$\sin^2 \gamma - \sin^2 \alpha \sin^2 \beta \quad (9)$$

The knowledge of the sign of (9) then gives explicit formulas for the real roots of $P_{\alpha,\beta,\gamma}$.

5.2 Multiplicity of 0

PROPOSITION 12. The multiplicity of $\varphi = 0$ in $R_{a,b,c}$ is

$$\frac{a-1}{2}((b,c) - 1) + \lfloor \frac{b}{2} \rfloor ((a,c) - 1).$$

PROOF. We have to examine whenever $\varphi = 0$ is a root of $P_{\alpha,\beta,\gamma}$ where $\alpha = \frac{i\pi}{a}$, $\beta = \frac{j\pi}{b}$ and $\gamma = \frac{k\pi}{c}$. Here a is odd so $\cos \alpha \neq 0$. Thus, $\varphi = 0$ is a root of $P_{\alpha,\beta,\frac{\pi}{2}}$ if and only if

$$\cos \alpha \cos \beta \quad (10)$$

is null and when $\gamma \neq \frac{\pi}{2}$, $\varphi = 0$ is a root of $P_{\alpha,\beta,\gamma}$ if and only if the following expression is null

$$(\cos^2 \alpha - \cos^2 \gamma)(\cos^2 \beta - \cos^2 \gamma). \quad (11)$$

- If $\gamma = \beta = \frac{\pi}{2}$, $\varphi = 0$ is a root for $i = 1, \dots, \frac{a-1}{2}$.
- If $\gamma \neq \frac{\pi}{2}$, $\varphi = 0$ is a root of $P_{\alpha,\beta,\gamma}$ if and only if $\sin^2 \gamma = \sin^2 \alpha$ or $\sin^2 \gamma = \sin^2 \beta$, that is $ic = ka$ or $jc = kb$ or $(b-j)c = kb$. The root $\varphi = 0$ is obtained for $i = \lambda \frac{a}{(a,c)}$, $k = \lambda \frac{c}{(a,c)}$, $\lambda = 1, \dots, \frac{(a,c)-1}{2}$ and it is double when $\beta = \frac{\pi}{2}$. It is also obtained for $j = \mu \frac{b}{(b,c)}$, $k = \mu \frac{c}{(a,c)}$, $\mu = 1, \dots, (b,c) - 1$. We obtain $\lfloor \frac{b}{2} \rfloor ((a,c) - 1) + ((b,c) - 1)(a-1)/2$.

We thus obtain the result. \square

Remark. We find that 0 is not a critical value if and only if a , b and c are pairwise coprime integers. This result was first proved by Comstock ([4], 1897), who found the number of crossing points of the harmonic curve parametrized by $x = T_a(t)$, $y = T_b(t)$, $z = T_c(t)$.

5.3 Non null multiple roots of $R_{a,b,c}$

It may happen that $R_{a,b,c}$ has multiple root φ . Several cases may occur.

► $P_{\alpha,\beta,\gamma}$ has a double root if and only if $\text{Disc}(P_{\alpha,\beta,\gamma}) = 0$, that is to say $\sin^2 \gamma = \sin^2 \alpha \sin^2 \beta$. The double root is $\varphi = -\cos \alpha \cos \beta$.

► $P_{\alpha,\beta,\gamma_1}$ and $P_{\alpha,\beta,\gamma_2}$ have a common root. In this case $P_{\alpha,\beta,\gamma_1} = P_{\alpha,\beta,\gamma_2}$, that is to say

$$(\sin^2 \gamma_1 - \sin^2 \gamma_2)(\sin^2 \gamma_1 \sin^2 \gamma_2 - \sin^2 \alpha \sin^2 \beta) \quad (12)$$

is null.

► $P_{\alpha_1,\beta_1,\gamma_1}$ and $P_{\alpha_2,\beta_2,\gamma_2}$ have a common root.

The first two cases are related to the equation

$$\sin r_1 \pi \sin r_2 \pi = \sin r_3 \pi \sin r_4 \pi \quad (13)$$

where $r_i \in \mathbf{Q}$. All the solutions of Equation (13) are known (see [16]). There is a one-parameter infinite family of solutions corresponding to

$$\sin \frac{\pi}{6} \sin \theta = \sin \frac{\theta}{2} \sin \left(\frac{\pi}{2} - \frac{\theta}{2} \right),$$

and a finite number of solutions listed in [16]. We deduce from a careful study of the Equation (13):

PROPOSITION 13. Let $\alpha = \frac{i\pi}{a}$, $\beta = \frac{j\pi}{b}$ and $\gamma = \frac{k\pi}{c}$, where $(a,b) = 1$ and a is odd. $P_{\alpha,\beta,\gamma}$ has a double root iff $\beta = \frac{\pi}{2}$ and $\sin \gamma = \sin \alpha$. In this case, the double root is $\varphi = 0$.

and

PROPOSITION 14. Let $\alpha = \frac{i\pi}{a}$, $\beta = \frac{j\pi}{b}$ and $\gamma_1 = \frac{k_1\pi}{c}$, $\gamma_2 = \frac{k_2\pi}{c}$, where $(a,b) = 1$ and a is odd. Then $P_{\alpha,\beta,\gamma_1}$ and $P_{\alpha,\beta,\gamma_2}$ have a common root φ iff there are equal and

$$1. \sin \alpha = \sin \gamma_1, \sin \beta = \sin \gamma_2.$$

In this case, the roots are $\varphi = 0$ and $\varphi = -2 \cos \alpha \cos \beta$.

$$2. \sin \beta = \frac{1}{2}, \sin \gamma_1 = \sin \frac{1}{2} \alpha, \sin \gamma_2 = \cos \frac{1}{2} \alpha.$$

In this case the common roots are $\varphi = -\cos(\alpha \pm \frac{\pi}{6})$.

$$3. \sin \gamma_2 = \sin \gamma_1.$$

► In case when $\alpha_1 \neq \alpha_2$ or $\beta_1 \neq \beta_2$, $P_{\alpha_1,\beta_1,\gamma_1}$ and $P_{\alpha_2,\beta_2,\gamma_2}$ have a common root if $\text{Res}_\varphi(P_{\alpha_1,\beta_1,\gamma_1}, P_{\alpha_2,\beta_2,\gamma_2}) = 0$. This resultant can be expanded and its sign is the one of:

$$\begin{aligned} & ((\cos^2 \alpha_1 - \cos^2 \gamma_1)(\cos^2 \beta_1 - \cos^2 \gamma_1) \sin^2 \gamma_2 - \\ & (\cos^2 \alpha_2 - \cos^2 \gamma_2)(\cos^2 \beta_2 - \cos^2 \gamma_2) \sin^2 \gamma_1)^2 \\ & - 4(\cos \alpha_1 \cos \beta_1 - \cos \alpha_2 \cos \beta_2) \sin^2 \gamma_1 \sin^2 \gamma_2 \times \\ & ((\cos^2 \alpha_1 - \cos^2 \gamma_1)(\cos^2 \beta_1 - \cos^2 \gamma_1) \cos \alpha_2 \cos \beta_2 \sin^2 \gamma_2 - \\ & (\cos^2 \alpha_2 - \cos^2 \gamma_2)(\cos^2 \beta_2 - \cos^2 \gamma_2) \cos \alpha_1 \cos \beta_1 \sin^2 \gamma_1). \end{aligned} \quad (14)$$

It would be interesting to get an arithmetic condition asserting that this resultant is null.

5.4 Computing the diagrams

Let $\varphi \in \mathbf{R}$. φ may be a rational number $r \in \mathbf{Q} - \mathcal{Z}_{a,b,c}$ or an algebraic number given by a polynomial whom it is a root and an isolating interval. The main step is the computation of the crossing nature at the double point $A_{\alpha,\beta}$ corresponding to parameters $(t = \cos \alpha + \beta, s = \cos \alpha - \beta)$, where $\alpha = \frac{i\pi}{a}$, $\beta = \frac{j\pi}{b}$. There are two cases to consider.

1. We know the roots $\varphi_1 \leq \dots \leq \varphi_m$ of $Q_c(s, t, \varphi)$. If $\varphi < \varphi_1$ then $n = 0$ otherwise let $n = \max\{k, \varphi > \varphi_k\}$. We have $\text{sign}(Q_c(s, t, \varphi)) = (-1)^n$.
2. We do not know the roots of $Q_c(s, t, \varphi)$. We compute $Q_c(s, t, \varphi)$ using the recurrence formula:

$$\begin{aligned} Q_0 &= 0, Q_1 = 1, Q_2 = 2S + 4\varphi, \\ Q_3 &= -4T + 12\varphi S + 4S^2 + 12\varphi^2 - 3, \\ Q_{n+4} &= 2(S + 2\varphi)(Q_{n+3} + Q_{n+1}) \\ &\quad - 2(2\varphi^2 + 2T + 2\varphi S + 1)Q_{n+2} - Q_n. \end{aligned}$$

where $S = s + t = 2 \cos \alpha \cos \beta$ and $T = st = \cos^2 \alpha + \cos^2 \beta - 1$ (see [12]). We work formally in $\mathbf{Q}[u, v]/\langle M, N \rangle$ where M, N are the minimal polynomials of $u = \cos \alpha$, $v = \cos \beta$.

The sign of the crossing is

$$\begin{aligned} D(s, t, \varphi) &= Q_c(s, t, \varphi) P_{b-a}(s, t, \varphi) \\ &= (-1)^{i+j} \sin \frac{ib\pi}{a} \sin \frac{ja\pi}{b} Q_c(s, t, \varphi) \\ &= (-1)^{i+j+\lfloor \frac{ib}{a} \rfloor + \lfloor \frac{ja}{b} \rfloor} Q_c(s, t, \varphi). \end{aligned}$$

6. THE ALGORITHM

We want to compute all the real roots $\varphi_1 < \dots < \varphi_n$ of $R_{a,b,c}$ that factors in $\frac{1}{2}(a-1)(b-1)\lfloor \frac{c}{2} \rfloor$ polynomials $P_{\alpha_i,\beta_j,\gamma_k}$. We precisely want non overlapping intervals $[a_m, b_m]$ for these roots in order to chose sample rational points $r_0 < a_1$, $b_i < r_i < a_{i+1}$, $b_n < r_n$.

At some stages, one may need to compute the sign of $\text{Disc}(P_{\alpha,\beta,\gamma})$ (expression (9)) or $\text{Res}(P_{\alpha_1,\beta_1,\gamma_1}, P_{\alpha_2,\beta_2,\gamma_2})$ (expressions (12) and (14)) in order to decide whether two roots are distinct or not. This information is required for

two reasons. We first want to be sure that we get all the roots and secondly, we will need to know all the roots of $Q_c(\cos(\alpha_i + \beta_j), \cos(\alpha_i - \beta_j), \varphi)$ with their multiplicities in order to determine the nature of the crossing over the corresponding double point in the diagram (section 5.4).

The signs of (10) and (11), may be evaluated by simple arithmetic considerations on α, β, γ .

Isolate and Refine. A very first step is to get accurate isolating intervals with rational bounds for $\cos \alpha_i, \cos \beta_j$ and $\cos \gamma_k$ to perform interval arithmetic for the real roots of $P_{\alpha_i, \beta_j, \gamma_k}$.

Such intervals can be computed by performing algorithms based on Descartes rule of signs (see for example [18]) on the used Chebyshev polynomials V_n . Algorithms like in [18] can easily solve such polynomials for very high degrees (several thousands) with a large accuracy. The computation of the required isolating intervals can then be performed as a pre-processing for the global algorithm.

From now, we denote by **Isolate**(P, acc) the function that isolates the roots of a univariate polynomial P with rational coefficients by means of intervals with rational bounds for a given accuracy acc (maximal length of the intervals). This function provides non overlapping intervals that contains a unique real root of P (and such that each real root of P is contained in one of the intervals).

Note that if more accuracy is required for some intervals, it is easy to refine them from the isolating intervals provided by the function **Isolate**: just evaluating P at some points, without running again the **Isolate** function with a higher value for acc). We name by **Refine**(I, P, acc) the function that decreases the length of the interval I to get an accuracy $\leq acc$, knowing that the interval isolates a real root of P .

IsolateP. Thanks to Proposition 13, one can compute the roots of $P_{\alpha, \beta, \gamma}$ with an appropriate accuracy, using multi-precision interval arithmetic for the evaluations. We will use the function **IsolateP**($\alpha, \beta, \gamma, acc$) that returns a (possibly empty) list of $(\alpha, \beta, \gamma, [u, v])$ corresponding to isolating intervals $[u, v]$ for the roots.

SignTest. When two isolating intervals $[u_1, v_1]$ and $[u_2, v_2]$ corresponding to $(\alpha_1, \beta_1, \gamma_1)$ and $(\alpha_2, \beta_2, \gamma_2)$ are such that $[u_1, v_1] \cap [u_2, v_2] \neq \emptyset$, we first use a filter (named **SignTest** in the sequel) which consists in using multi-precision interval arithmetic for the evaluation of $\text{Res}_\varphi(P_{\alpha_1, \beta_1, \gamma_1}, P_{\alpha_2, \beta_2, \gamma_2})$ (expressions 12) and (14).

Thanks to Proposition 14, we know by arithmetic considerations when (12) is null. We know also the corresponding common roots and we change $[u_i, v_i]$ to $[u_1, v_1] \cap [u_2, v_2]$.

Expression (14) is

$$P = ((C_1^2 - C_5^2)(C_3^2 - C_5^2)(1 - C_6^2) - (C_2^2 - C_6^2)(C_4^2 - C_6^2)(1 - C_5^2))^2 - 4(C_1C_3 - C_2C_4)(1 - C_5^2)(1 - C_6^2) \times ((C_1^2 - C_5^2)(C_3^2 - C_5^2)C_2C_4(1 - C_6^2) - (C_2^2 - C_6^2)(C_4^2 - C_6^2)C_1C_3(1 - C_5^2)),$$

where $C_1 = \cos \alpha_1, C_2 = \cos \alpha_2, C_3 = \cos \beta_1, C_4 = \cos \beta_2, C_5 = \cos \gamma_1, C_6 = \cos \gamma_2$.

Given isolating intervals with rational bounds that contain the values of the required $C_i, i = 1, \dots, 6$, the function **SignTest**($\alpha_1, \alpha_2, \beta_1, \beta_2, \gamma_1, \gamma_2$) straightforwardly evaluates P . If the resulting interval is $[0, 0]$ or do not contains 0, one can decide the sign of the input, otherwise, the function returns **FAIL**.

FormalNullTest. In case of failure of **SignTest**, one has

to decide if the input is null or not, which is the goal of the function **FormalNullTest** we now describe.

Let us write $\alpha_1 = \frac{i_1\pi}{a_1}, \alpha_2 = \frac{i_2\pi}{a_2}, \beta_1 = \frac{j_1\pi}{b_1}, \beta_2 = \frac{j_2\pi}{b_2}, \gamma_1 = \frac{k_1\pi}{c_1}, \gamma_2 = \frac{k_2\pi}{c_2}$. Let m be the smallest common multiple of a_1, a_2, b_1, b_2, c_1 and c_2 . According to the definitions of T_n , we have $C_i = T_{n_i}(\cos \frac{\pi}{m})$. Since M_m is the minimal polynomial of $\cos \frac{\pi}{m}$, the expression $P(C_i, \dots, C_6)$ is null if and only if $P(T_{n_1}, \dots, T_{n_6}) = 0$ in $\mathbf{Q}[t]/\langle M_m(t) \rangle$.

DoubleTest. Our function first performs the **SignTest**. If it returns an interval with bounds of same sign, then the sign of the tested expression is the sign of the two bounds of the interval. Otherwise, we run the **FormalNullTest**. If this test returns 0 then the expression is null. Otherwise, we decrease the lengths of the intervals that represent the values of $\cos \frac{k\pi}{m}$ by calling the function **Refine** until the **SignTest** does not **FAIL** (the fact that the sign of the expression to be tested is known not to be 0 ensures that this process will end).

The global algorithm. We proceed in three steps :

(0) We isolate the roots of some Chebyshev polynomials using the **Isolate** black-box with an arbitrary accuracy.

(1) We compute separately the roots of the $P_{\alpha, \beta, \gamma}$ by using **IsolateP**.

(2) We then consider the list of these roots and observe carefully the overlapping intervals. For any pair of overlapping interval, we decide whether corresponding resultants are null or not using **DoubleTest**. If the corresponding roots are equal then we change their isolating intervals by taking their intersection.

From these disjoint intervals with rational bounds, we straightforwardly get the roots with their multiplicities. We thus deduce the sample points r_0, \dots, r_n we need. Furthermore, for each $\alpha_i = \frac{i\pi}{a}, \beta_j = \frac{j\pi}{b}$, we know the roots with their multiplicities of $Q_c(t, s, \varphi)$, where $t = \cos(\alpha_i + \beta_j)$ and $s = \cos(\alpha_i - \beta_j)$. This information is helpful for knowing the crossing nature at the point A_{α_i, β_j} (section 5.4).

7. EXPERIMENTS

In the appendix of [12], we gave parametrizations of every rational knot as $\mathcal{C}(3, b, c, \varphi)$ and $\mathcal{C}(4, b, c, \varphi)$ where (b, c) were minimal for the lexicographic order ($c \leq 300$). For 6 knots we knew the minimal b and that $c > 300$. With the method we developed here, we recover all the minimal parametrizations we gave in [12] but also for the 6 missing knots. The following knots admit the parametrizations:

$$\begin{aligned} 9_5 &= \mathcal{C}(3, 13, 326, 1/85), & 10_3 &= \mathcal{C}(4, 13, 348, 1/138), \\ 10_{30} &= \mathcal{C}(4, 13, 306, 1/738), & 10_{33} &= \mathcal{C}(4, 13, 856, 1/328), \\ 10_{36} &= \mathcal{C}(3, 14, 385, 1/146), & 10_{39} &= \mathcal{C}(3, 14, 373, 1/182). \end{aligned}$$

For example, one deduces that there is no parametrization of 9_5 as Chebyshev knots with $(a, b, c)_{<_{\text{lex}}} (3, 13, 326)$.

$R_{3,14,385}$ has degree 4992. It has 2883 real roots. All are simple roots except 0 that is of multiplicity 6.

$R_{4,13,856}$ has degree 15390 and 9246 real roots (0 has multiplicity 18). We get 2050 non trivial knots, 83 of them are distinct, and 63 have less than 10 crossings. The total running time — critical values with their multiplicities, sampling of 1442 values, computing knot invariant — was 450" (MAPLE 13, on Laptop, 3Gb of RAM, 3MHz).

Outside the intrinsic combinatorial aspects of the problem, the complexity of our algorithm essentially depends on the

FormalNullTest. In the worst case $d = abc$ and $\deg M_d = \frac{1}{2}(a-1)(b-1)(c-1)$, when a, b, c are prime integers, the most difficult computation consists in deciding if the expression 14 is null or not which is equivalent to testing if a univariate polynomial of degree at most $4d$ is null modulo M_d or not.

These computations can be speed up a lot since they can be performed modulo a prime integer: all the considered polynomials have a power of two as leading coefficient and we just need to test if one polynomial is null modulo another one.

In this challenging experiments, we never had to run the **FormalNullTest**, the **SignTest** being always sufficient, thanks to the filters given by propositions 13 and 14 and to a good (experimental) choice initial choice of accuracy when computing the prerequisites running the **Isolate** algorithm.

8. CONCLUSION

The method we developed in this paper allows us to compute Chebyshev knot diagrams for high values of a, b and c . Our experience with small a and b shows that the difficult cases (multiple roots of $R_{a,b,c}$ we found) were predictable. There are certainly some specific reasons connected with arithmetic properties and the structure of cyclic extensions.

The main difference with the algorithm described in [12] and the computation of $R_{a,b,c}$ as a polynomial of degree $\frac{1}{2}(a-1)(b-1)(c-1)$, is that it came as a resultant of a polynomial of degree $(c-1)$ in (X, φ) and a polynomial of degree $\frac{1}{2}(a-1)(b-1)$ in X with coefficients in a unique field extension. The example described in this section can be considered as the extremal case, in terms of degree, to be solved using methods from the state of the art when running [12] while it can be solved in few minutes with the method proposed in this article.

From the point of view of knot theory, it is proved in [11], that rational knots with N crossings can be parametrized by polynomials of degrees $(3, b, c)$ where $b + c \leq 3N$ which is far better than the results we obtain here. But we challenged to give, as it was done with Lissajous knots in [1], an exhaustive and certified list of minimal parametrizations. We consider that it might be one step in the computing of polynomial curves topology.

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