



Inter-Operator Spectrum Sharing from a Game Theoretical Perspective

Samson Lasaulce, M. Bennis, Merouane Debbah

► To cite this version:

Samson Lasaulce, M. Bennis, Merouane Debbah. Inter-Operator Spectrum Sharing from a Game Theoretical Perspective. *Eurasip Journal on Applied Signal Processing*, 2009, 2009 (8), pp.1-12. 10.1155/2009/295739 . hal-00446284

HAL Id: hal-00446284

<https://hal.science/hal-00446284>

Submitted on 12 Jan 2010

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

Research Article

Inter-Operator Spectrum Sharing from a Game Theoretical Perspective

Mehdi Bennis,¹ Samson Lasaulce,² and Merouane Debbah³

¹ Centre for Wireless Communications, University of Oulu, 90014 Oulu, Finland

² SUPELEC, CNRS, 91192 Gif-sur-Yvettes Cedex, France

³ Alcatel-Lucent Chair in Flexible Radio, SUPELEC, 91192 Gif-sur-Yvettes Cedex, France

Correspondence should be addressed to Mehdi Bennis, bennis@ee.oulu.fi

Received 15 February 2009; Revised 1 June 2009; Accepted 8 July 2009

Recommended by K. Subbalakshmi

We address the problem of spectrum sharing where competitive operators coexist in the same frequency band. First, we model this problem as a *strategic* non-cooperative game where operators *simultaneously* share the spectrum according to the Nash Equilibrium (NE). Given a set of channel realizations, several Nash equilibria exist which renders the outcome of the game unpredictable. Then, in a cognitive context with the presence of primary and secondary operators, the inter-operator spectrum sharing problem is reformulated as a *Stackelberg* game using hierarchy where the primary operator is the leader. The Stackelberg Equilibrium (SE) is reached where the best response of the secondary operator is taken into account upon maximizing the primary operator's utility function. Moreover, an extension to the multiple operators spectrum sharing problem is given. It is shown that the Stackelberg approach yields better payoffs for operators compared to the classical *water-filling* approach. Finally, we assess the goodness of the proposed distributed approach by comparing its performance to the centralized approach.

Copyright © 2009 Mehdi Bennis et al. This is an open access article distributed under the Creative Commons Attribution License, which permits unrestricted use, distribution, and reproduction in any medium, provided the original work is properly cited.

1. Introduction

Spectrum sharing between wireless networks improves the efficiency of spectrum usage where a migration toward flexible spectrum management is paramount to alleviate spectrum scarcity and its underutilization. In this respect and motivated by the ever-increasing demands for wireless services, several works have appeared in literature ([1–9] among many others) wherein interestingly both theoretical and practical (system-level) contributions, stemming from pricing [1, 7], opportunistic power control [10] to resource sharing [8], and others have been made.

In this paper, we study spectrum sharing between two *competing* operators operating in the same frequency band in which base stations communicate with their mobile terminals. In this case, a transmitter T_1 wants to send information to its mobile R_1 , while at the same time another base station T_2 (from a competitive operator) wants to transmit information to its mobile R_2 . These systems, therefore, share the same medium where the communication pairs (T_1, R_1) and (T_2, R_2) take place simultaneously and

on the same frequency band. This setup is known as the interference channel (IFC) ([11–15] to mention a few).

There is a great deal of work on the IFC channel using game theory. In [13], the problem of power allocation in a frequency-selective multiuser interference channel is studied. An iterative Water-Filling (WF) algorithm is proposed to efficiently reach the Nash equilibrium. Moreover, it is found that under suitable conditions, the iterative WF algorithm for the two-user Gaussian interference game converges to the unique Nash equilibrium from any starting point. In their scenario, the Nash equilibria lead to nonefficient and non pareto-optimal solutions. Similarly, in [11], the authors consider the problem of spectrum sharing on the IFC for *flat-fading* channels. The interference channel is viewed as a noncooperative game and the Nash equilibrium is characterized under a set of sufficient conditions. In [16], the authors investigate the problem of simultaneous water-filling solution for the gaussian IFC under *weak* interference. Motivated by the pareto-inefficiency of the water-filling approach, the authors propose a distributed algorithm to transform a

symmetric system from simultaneously waterfilled to a fair orthogonal signal space partitions.

In [17], the problem of two wireless networks operating on the same frequency band was considered. Pairs within a given network cooperate to schedule transmissions according to a random-access protocol where each network chooses an access probability for its users. In [18], the authors consider the problem of coordinating two competing multiple-antenna wireless systems in the Multiple Input Single Output (MISO) IFC. It turns out that if the systems do not cooperate, then the corresponding equilibrium rates are bounded regardless of how much transmit power the base stations have available. Also, Nash bargaining solutions were found to be close to the sum-rate bound. On the other hand, in [19–21], the authors study the problem of maximizing mutual information subject to mask constraints and transmit power, for both simultaneous and asynchronous cases. (Under this setup, some users are allowed to update their strategy more frequently than the others. And, they might even perform these updates using outdated information on the interference caused by others.) The existence of the Nash equilibrium is proven and sufficient conditions are given for the uniqueness. Finally, in [20], distributed iterative algorithms are proposed to reach the Nash equilibrium.

In most of these works, the existence of the Nash equilibrium is easily demonstrated, whereas the uniqueness is generally more complicated for which only sufficient conditions are given. Because of the very hard problem of the uniqueness of the Nash equilibrium points in the WF game, Nash bargaining (NBS) solutions were considered in [14]. However, NBS requires the knowledge of all channel state information which is not always possible in practice.

Within the same framework of spectrum sharing but under a different scenario, Stackelberg games [22] have been applied in the context of cognitive radios where the desirability of outcomes depends not only on their own actions but also on other cognitive radios. It is worth pointing out that the Stackelberg formulation naturally arises in some contexts of practical interest: (a) when primary and secondary systems share the spectrum, (b) when user have access to the medium in an asynchronous manner, (c) when operators deploy their networks at different times, and (d) when some nodes have more power than others such as the base station. Stackelberg is based on a *leader-follower* approach in which the leader plays his strategy before the follower, and then enforces it. In [23], a game theoretic framework has been proposed in the context of fading multiple-access channel where a Stackelberg formulation is proposed in which the base station is the designated game leader with the purpose to have a distributed allocation strategy approaching all corners of the capacity region. In [24], a two-level Stackelberg game is proposed for distributed relay selection and power control for multiuser cooperative networks. The objective is to jointly consider the benefits of source and relay nodes in which the source node is modeled as a buyer and the relay nodes as the sellers. Also, the energy-efficient power control problem is investigated in [25]. Moreover, in [26], the authors investigate a similar power allocation problem but *solely* focus on channel realizations in

which the Nash equilibrium of the game is unique. (The case with multiple Nash equilibria was not treated.) However, this work differs in that the Stackelberg approach is mainly motivated by the nonuniqueness of the Nash equilibrium and unpredictability of the game.

In essence, the fundamental questions we address in this paper are the following.

- (i) In the first operators' deployment scenario, if both operators *simultaneously* operate in a noncooperative (i.e., *selfish*) manner: what are their power allocation strategies across their carriers? Clearly, there is a conflict situation where a good strategy for the link (T_1, R_1) will generate interference for R_2 and viceversa. Hence, an equilibrium has to be found.
- (ii) Given any set of channel realizations, is it possible to predict the outcome of the game?, and if so, how to characterize the Nash equilibria regions? Is the Nash equilibrium *unique*?
- (iii) In the second scenario of operators' deployment in which primary and secondary operators coexist in the same spectral band, what is the outcome of the game when a hierarchy exists between operators? How does this approach compare with the selfish approach (classical water-filling)?
- (iv) How close is the distributed approach from the centralized (*sum-rate*) power allocation?

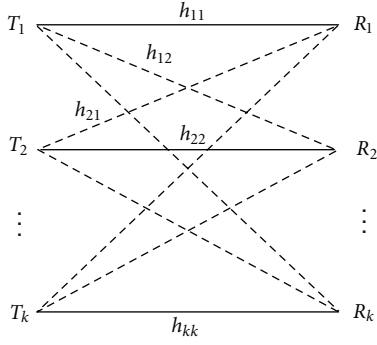
The paper is organized as follows: the system model is introduced in Section 2. In Section 3, the spectrum sharing game between operators is formulated using noncooperative game theory. In Section 4, a special case with two operators transmitting on two carriers is investigated to gain insights into the Nash equilibria regions. In Section 5, we formulate the interoperator spectrum sharing problem as a Stackelberg game where a hierarchy exists between operators as well as an extension to the multiple operators case. Section 6 provides a comparison between the distributed (selfish) and centralized approach. Finally, we conclude this work in Section 7.

2. System Model

We suppose that K transmitters share a frequency band composed of N carriers where each transmitter transmits in any combination of channels and at any time. (The terms transmitter and operator are interchangeably used throughout the paper.) On each carrier $n = 1 \dots N$, transmitter $k = 1 \dots K$ sends the information $x_k^n = \sqrt{p_k^n} s_k^n$, where s_k^n represents the transmitted data and p_k^n denotes the corresponding transmitted power of user k on carrier n . The received signal at the receiver i in carrier n can be expressed as

$$r_i^n = \sum_{j=1}^K h_{ji}^n x_j^n + w_i^n, \quad i = 1, \dots, K, n = 1, \dots, N, \quad (1)$$

where h_{ji}^n is the fading channel gain on the n th carrier between the pair (T_i, R_j) . In addition, the noise process w_i^n

FIGURE 1: K -user N -carriers interference channel under study.

is characterized by its received noise power on each carrier n , by σ_n^2 .

For transmitter i , the transmit power p_i^n is subject to its power constraint:

$$\sum_{n=1}^N p_i^n \leq \bar{P}_i, \quad i = 1, \dots, K. \quad (2)$$

At the receiver i , the signal to interference plus noise ratio (SINR) on carrier n is given by

$$\text{SINR}_i^n = \frac{p_i^n |h_{ii}^n|^2}{\sigma_n^2 + \sum_{j=1, j \neq i}^K p_j^n |h_{ji}^n|^2}. \quad (3)$$

Furthermore, assuming Gaussian codebook, the maximum achievable rate at receiver i is given by

$$R_i = \sum_{n=1}^N \log_2(1 + \text{SINR}_i^n). \quad (4)$$

3. Noncooperative Spectrum Sharing Game

In this section, we model the interoperator spectrum sharing problem from a noncooperative standpoint [27]. Figure 1 illustrates the spectrum sharing scenario under study for K operators and N carriers.

3.1. Game Formulation. The noncooperative spectrum sharing game is defined as $\Gamma^{\text{NCG}} \triangleq [\mathcal{K}, \{\mathcal{P}_i\}_{i \in \mathcal{K}}, \{\mathcal{U}_i\}_{i \in \mathcal{K}}]$. The players (from the set $\mathcal{K} \triangleq \{1, 2, \dots, K\}$) are defined as the different links with a strategy $p_i^n \in \mathcal{P}_i$ and the payoffs are the achievable rates on each link $u_i(p_i^n, p_{-i}^n) = R_i(p_i^n, p_{-i}^n) \in \mathcal{U}_i$, for $i = 1, \dots, K$ and $n = 1, \dots, N$. Each player competes against the others by choosing its transmit power (i.e., strategy) to maximize its own utility subject to some power constraints \bar{P}_i . In this work, we assume full channel state information in which operators know their fading channel gains as well as other's fading cross-channels.

Since the operators do not cooperate, the only reasonable outcome of the spectrum conflict is an operating point which constitutes a Nash equilibrium (NE) [28]. This is a point where none of the players can improve their utilities by unilaterally changing their strategies. One should note that a Nash equilibrium is not an optimal or even desirable outcome. However, it is an insightful point where one is likely to end up operating at, if players are not willing to cooperate.

In a noncooperative approach, operator i selfishly maximizes his utility function subject to the power constraint \bar{P}_i :

$$\begin{aligned} \max_{p_i^n} R_i &= \max_{p_i^n} \sum_{n=1}^N \log_2 \left(1 + \frac{|h_{ii}^n|^2 p_i^n}{\sigma_n^2 + \sum_{j \neq i}^K |h_{ji}^n|^2 p_j^n} \right) \\ \text{s.t.} \quad &\sum_{n=1}^N p_i^n \leq \bar{P}_i, \\ &p_i^n \geq 0. \end{aligned} \quad (5)$$

Furthermore, the channel realization set \underline{h} is defined as

$$\underline{h} = \{h_{ij}^n : i, j = 1, \dots, K, n = 1, \dots, N\}. \quad (6)$$

The solutions to (5) are given by the water-filling power allocation solutions:

$$p_i^n = \left(\frac{1}{\mu_i} - \frac{\sigma_n^2 + \sum_i |h_{-i,i}^n|^2 p_{-i}^n}{|h_{ii}^n|^2} \right)^+ \quad (7)$$

$$i = 1, \dots, K, \quad n = 1, \dots, N,$$

where $(x)^+ = \max\{x, 0\}$ and $\mu_i > 0$ is the Lagrangian multiplier chosen to satisfy the power constraint: $\sum_{n=1}^N p_i^n = \bar{P}_i$. Note that the equality follows from the concavity of the objective function in p_i .

4. A Special Case of Two Operators and Two Carriers

In order to gain insight into the properties of the Nash equilibria for our interoperator spectrum game, let us focus on the case where two operators transmit over two carriers (i.e., $K = N = 2$).

4.1. Notations. (i) For the ease of notation and readability that will prove helpful in the sequel, we introduce the following notations: $g_{ij}^n = \bar{P}_i |h_{ij}^n|^2 / \sigma_n^2$, $c_1 = g_{11}^1 / g_{11}^2$, and $c_2 = g_{22}^1 / g_{22}^2$.

(ii) The pair (α_1, α_2) means that user 1 transmits with power $(p_1^1, p_1^2) = (\alpha_1 \bar{P}_1, (1 - \alpha_1) \bar{P}_1)$ on carrier 1 and 2 while user 2 transmits with power $(p_2^1, p_2^2) = (\alpha_2 \bar{P}_2, (1 - \alpha_2) \bar{P}_2)$ on carriers 1 and 2, respectively.

Figure 2 depicts the space of the 9 Nash equilibria of the game obtained upon solving (7), the details of which are given in Appendix A. Given a set of channel realizations \underline{h} ,

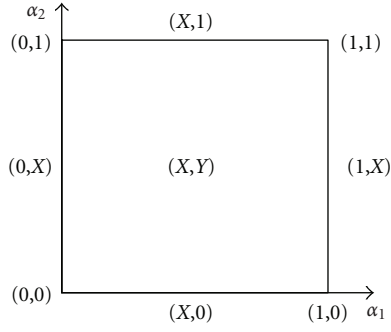


FIGURE 2: Illustration of the Nash equilibria space where (α_1, α_2) denotes the power allocation strategy for both operators 1 and 2, in the first carrier.

the game converges to different equilibrium points. Figure 3 illustrates one possible representation of the Nash equilibria space. Depending on the quantities $(1 + g_{21}^1)/(1 + g_{11}^2)$, $(1 + g_{11}^1)/(1 + g_{21}^2)$, and $(1 + g_{22}^1)/(1 + g_{12}^2)$, $(1 + g_{12}^1)/(1 + g_{22}^2)$ (see Appendix A), four different representation of the regions are possible. These regions are depicted in Figure 4 whose purpose is to reflect the 8-dimensional problem related to the channel realization set \underline{h} .

It turns out that given certain channel realizations, the Nash equilibrium is *unique* (white rectangle areas) while some of the grayish rectangle regions exhibit at least one Nash equilibria.

4.2. Existence of the Nash Equilibria. The existence of the Nash equilibria is proven using the theorem in [29] within the context of noncooperative concave games. Hence, the game defined in (5) admits at least one Nash equilibrium.

4.3. Uniqueness of the Nash Equilibria. In [13], the authors give sufficient conditions for the uniqueness of the NE but do not precisely state which NE are obtained for any given channel realization set \underline{h} . Therefore, building on these results, a full characterization of the Nash equilibria region for the 2-operators 2-carriers case is herein given. Beside, the proof of the uniqueness when both operators transmit in both carriers (full-spread) is given in Appendix B.

Remark 4.1. The operators were assumed to be noncooperative hence operating at the Nash Equilibrium was their best response in a selfish context (one-shot game). It was also shown that under certain channel realizations, the spectrum sharing game is predictable with a *unique* Nash equilibrium. However, in other regions and given a set of channel realizations, *nonunique* Nash equilibria exist. In this case, the spectrum sharing game is *no longer* predictable.

Remark 4.2. We note that the sufficient conditions given for the *flat-fading* case studied in [11] are depicted in Figure 3 for the low-interference regime (X, Y) where $g_{2,1}^1/g_{1,1}^1 < 1$ and $g_{1,2}^1/g_{2,2}^1 < 1$.

Remark 4.3. In the multioperators case, the results for the 2 operator and 2 carriers case carry over where the sufficient conditions for the uniqueness are given by

$$\forall n \in \{1, \dots, N\}, \quad \sum_{i=1, j \neq i}^K \frac{|h_{ji}^n|^2}{|h_{ii}^n|^2} < 1. \quad (8)$$

This result comes from the Karush-Kuhn-Tucker conditions of the optimization problem of the set of data rates. Additionally, the physical meaning of (8) is that the uniqueness of the Nash equilibrium is ensured if the links are sufficiently far from each other.

Remark 4.4. We note that when one of the cross-gain $|h_{-i,i}|^2 = 0$, the IFC becomes a Z-channel [30] where the NE exists and is *unique* (the characterization of the Nash equilibria region for the Z-channel follows the same lines as the IFC).

5. Introducing Hierarchy (Stackelberg Game Approach)

In this section, we look at the interoperator spectrum sharing problem where the concept of hierarchy is accounted for. This situation is inherent in situations where primary and secondary operators share the same spectrum. In what follows, we formulate and solve the interoperator spectrum sharing problem (with hierarchy) for the 2-operators 2-carriers case then provide insights for the case with more than 2 operators.

A *Stackelberg* game $\Gamma^{\text{SG}} \triangleq [\mathcal{K}, \{\mathcal{P}_i\}_{i \in \mathcal{K}}, \{\mathcal{U}_i\}_{i \in \mathcal{K}}]$ is proposed to model the spectrum sharing problem where one of the two operators is chosen to be the leader (primary operator). The *Stackelberg* Equilibrium (SE) [22] is the best response where a hierarchy of actions exists between players. *Backward induction* [27] is applied assuming that players can reliably forecast the behavior of other players and that they believe that the other player can do the same. For this reason, the key point in this setup is the capability of the follower of sensing the environment and, therefore, the power level of operator 1 (the leader).

5.1. Problem Formulation. Without loss of generality, we assume that primary operator 1 is the leader and secondary operator 2 is the follower. First, we give a definition of the Stackelberg equilibrium as follows.

Definition 5.1 (Stackelberg Equilibrium [27]). A strategy profile $(p_1^{\text{SE}}, p_2^{\text{SE}})$ is called a Stackelberg Equilibrium if p_1^{SE} maximizes the utility of the leader (operator 1) and p_2^{SE} is the best response of operator 2 to operator 1.

The Stackelberg spectrum sharing game can be formulated as follows. First, in the high-level problem (9), primary operator 1 maximizes his own utility function. Then, in the low-level problem (10), secondary operator 2 (follower) maximizes his own utility taking into account the optimal power allocation of operator 1 (p_1^{SE}). By denoting $(p_1^{\text{SE}}, p_2^{\text{SE}})$

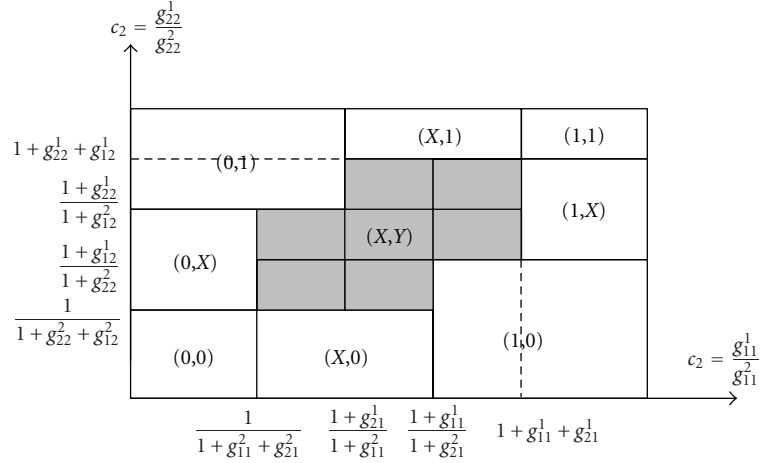
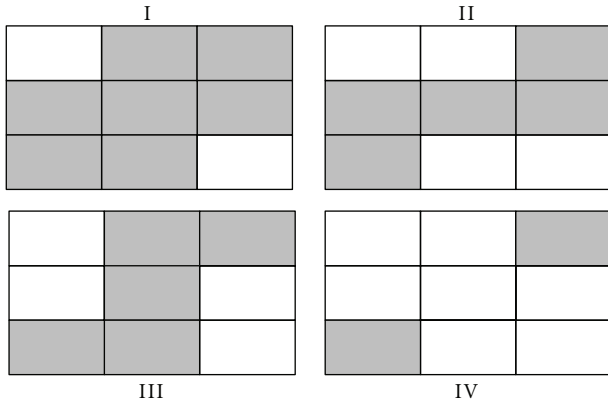
FIGURE 3: Characterization of the Nash equilibria regions given a set of channel realizations \bar{h} .

FIGURE 4: All four cases are depicted: (I) when $(1+g_{11}^1)/(1+g_{21}^2) > (1+g_{21}^1)/(1+g_{21}^2)$ and $(1+g_{22}^1)/(1+g_{22}^2) > (1+g_{12}^1)/(1+g_{22}^2)$, (II) when $(1+g_{11}^1)/(1+g_{21}^2) < (1+g_{21}^1)/(1+g_{21}^2)$ and $(1+g_{22}^1)/(1+g_{22}^2) > (1+g_{12}^1)/(1+g_{22}^2)$, (III) when $(1+g_{11}^1)/(1+g_{21}^2) > (1+g_{21}^1)/(1+g_{21}^2)$ and $(1+g_{22}^1)/(1+g_{22}^2) < (1+g_{12}^1)/(1+g_{22}^2)$, and finally when (IV) $(1+g_{11}^1)/(1+g_{21}^2) < (1+g_{21}^1)/(1+g_{21}^2)$ and $(1+g_{22}^1)/(1+g_{22}^2) < (1+g_{12}^1)/(1+g_{22}^2)$.

as the Stackelberg Equilibrium, the rate optimization problem for operator 1 (leader) is written as

$$\begin{aligned} \max_{p_1^n} \sum_{n=1}^N \log_2 \left(1 + \frac{|h_{11}^n|^2 p_1^n}{\sigma_n^2 + |h_{21}^n|^2 p_2^n(p_1^n)} \right) \\ \sum_{n=1}^N p_1^n \leq \bar{P}_1, \\ p_1^n \geq 0. \end{aligned} \quad (9)$$

The rate optimization problem for operator 2 (follower) is written as

$$\begin{aligned} \max_{p_2^n} \sum_{n=1}^N \log_2 \left(1 + \frac{|h_{22}^n|^2 p_2^n}{\sigma_n^2 + |h_{12}^n|^2 (p_1^n)^{\text{SE}}} \right) \\ \sum_{n=1}^N p_2^n \leq \bar{P}_2, \\ p_2^n \geq 0, \end{aligned} \quad (10)$$

where $p_2^{\text{SE}} = \text{BR}_2(p_1^{\text{SE}})$.

Using backward induction and given the best response of operator 2 (the follower), (10) can be rewritten as

$$\begin{aligned} \max_{p_1^n} \sum_{n=1}^N \log_2 \left(1 + \frac{|h_{11}^n|^2 p_1^n}{\sigma_n^2 + |h_{21}^n|^2 (1/\mu_2 - (\sigma_n^2 + |h_{12}^n|^2 p_1^n)/|h_{22}^n|^2)^+} \right) \\ \sum_{n=1}^N p_1^n \leq \bar{P}_1, \\ p_1^n \geq 0. \end{aligned} \quad (11)$$

The Stackelberg sharing game, therefore, boils down to solving (11) where several cases are considered. In our spectrum sharing problem ($K = N = 2$), the power strategies of operator 2 take 3 values. In the first case, operator 2 transmits with maximum power \bar{P}_2 in carrier 1 ($p_2^1 = \bar{P}_2, p_2^2 = 0$). In the second case, operator 2 transmits with \bar{P}_2 in carrier 2 ($p_2^1 = 0, p_2^2 = \bar{P}_2$), and finally in the third case, operator 2 transmits on both carriers (full spread) with ($p_2^1 = x, p_2^2 = \bar{P}_2 - x$), $0 < x < \bar{P}_2$. Therefore, the leader maximizes his utility function given the best response of the follower. In the following, all of the *three* cases are investigated. For simplicity sake, we assume $\bar{P}_1 = \bar{P}_2 = 1$.

5.1.1. *Operator 2 Transmits Only in Carrier 2* ($p_2^1 = 0, p_2^2 = 1$). Under this setup, $p_2^2 > 0 \Rightarrow p_1^1 \geq \beta_1$ where

$$\beta_1 = \frac{\sigma_2^2/|h_{22}^2|^2 + |h_{12}^2|^2/|h_{22}^2|^2 - \sigma_1^2/|h_{22}^1|^2 + 1}{|h_{12}^1|^2/|h_{22}^1|^2 + |h_{12}^2|^2/|h_{22}^2|^2}, \quad (12)$$

where β_1 depends on the set of channel realizations.

Furthermore, the maximization problem for the leader is written as

$$\begin{aligned} \max_{p_1^1} \log_2 \left(1 + \frac{|h_{11}^1|^2 p_1^1}{\sigma_1^2} \right) + \log_2 \left(1 + \frac{|h_{11}^1|^2 (\bar{P}_1 - p_1^1)}{\sigma_2^2 + |h_{21}^2|^2} \right), \\ \max(\beta_1, 0) \leq p_1^1 \leq \bar{P}_1, \end{aligned} \quad (13)$$

the KKT conditions are given such that

$$\begin{aligned} \lambda_1^* ((p_1^1)^* - \max(\beta_1, 0)) &= 0, \quad \lambda_1^* \geq 0, \\ \lambda_2^* ((p_1^1)^* - \bar{P}_1) &= 0, \quad \lambda_2^* \geq 0, \\ \frac{\partial \mathcal{L}_1}{\partial p_1^1} &= 0, \\ (p_1^1)^* &\leq \bar{P}_1, \\ (p_1^1)^* &\geq \max(\beta_1, 0), \end{aligned} \quad (14)$$

where λ_1^*, λ_2^* are the Lagrangian multipliers associated with the constraints

$$\begin{aligned} \mathcal{L}_1 &= \log_2 \left(1 + \frac{|h_{11}^1|^2 p_1^1}{\sigma_1^2} \right) + \log_2 \left(1 + \frac{|h_{11}^1|^2 (\bar{P}_1 - p_1^1)}{\sigma_2^2 + |h_{21}^2|^2} \right) \\ &\quad - \lambda_1 (p_1^1 - \max(\beta_1, 0)) + \lambda_2 (p_1^1 - \bar{P}_1), \\ \frac{\partial \mathcal{L}_1}{\partial p_1^1} &= 0 \\ \Rightarrow \frac{h_{11}^1}{\sigma_1^2 + h_{11}^1 p_1^1} - \frac{h_{11}^1}{\sigma_2^2 + h_{21}^2 + h_{11}^2 (\bar{P}_1 - p_1^1)} &= \lambda_1 - \lambda_2. \end{aligned} \quad (15)$$

Assume that $p_1^1 = \max(\beta_1, 0)$, then, $\lambda_1^* \geq 0, \lambda_2^* = 0$ and

$$\frac{h_{11}^1}{\sigma_1^2 + h_{11}^1 \max(\beta_1, 0)} - \frac{h_{11}^1}{\sigma_2^2 + h_{21}^2 + h_{11}^2 (\bar{P}_1 - \max(\beta_1, 0))} \geq 0. \quad (16)$$

Now assuming that $p_1^1 = \bar{P}_1$, then $\lambda_1^* = 0, \lambda_2^* \geq 0$ and

$$\frac{h_{11}^1}{\sigma_1^2 + h_{11}^1 \bar{P}_1} - \frac{h_{11}^1}{\sigma_2^2 + h_{21}^2} \leq 0. \quad (17)$$

Finally, assuming that $\max(\beta_1, 0) < p_1^1 < \bar{P}_1$, then $\lambda_1^* = \lambda_2^* = 0$ and

$$p_1^1 = \frac{\sigma_1^2 |h_{11}^1|^2 - \sigma_2^2 |h_{11}^2|^2 + |h_{11}^1|^2 (|h_{21}^2|^2 + |h_{11}^2|^2 \bar{P}_1)}{2 |h_{11}^1|^2 |h_{11}^2|^2}. \quad (18)$$

5.1.2. *Operator 2 Transmits Only in Carrier 1* ($p_2^1 = 1, p_2^2 = 0$). Under this setup, $p_2^1 > 0 \Rightarrow p_1^1 \leq \beta_2$, where

$$\beta_2 = \frac{\sigma_2^2/|h_{22}^2|^2 + |h_{12}^2|^2/|h_{22}^2|^2 - \sigma_1^2/|h_{22}^1|^2 - 1}{|h_{12}^1|^2/|h_{22}^1|^2 + |h_{12}^2|^2/|h_{22}^2|^2}. \quad (19)$$

Furthermore, the maximization problem for the leader is written as

$$\begin{aligned} \max_{p_1^1} \log_2 \left(1 + \frac{|h_{11}^1|^2 p_1^1}{\sigma_1^2 + |h_{21}^1|^2} \right) + \log_2 \left(1 + \frac{|h_{11}^1|^2 (\bar{P}_1 - p_1^1)}{\sigma_2^2} \right) \\ 0 \leq p_1^1 \leq \min(\beta_2, \bar{P}_1). \end{aligned} \quad (20)$$

Likewise, to derive the KKT conditions, form the Lagrangian denoted as \mathcal{L}_2 :

$$\begin{aligned} \mathcal{L}_2 &= \log_2 \left(1 + \frac{|h_{11}^1|^2 p_1^1}{\sigma_1^2 + |h_{21}^1|^2} \right) \\ &\quad + \log_2 \left(1 + \frac{|h_{11}^1|^2 (\bar{P}_1 - p_1^1)}{\sigma_2^2} \right) \\ &\quad - \lambda_1 p_1^1 + \lambda_2 (p_1^1 - \min(\beta_2, \bar{P}_1)) \end{aligned} \quad (21)$$

the KKT conditions are

$$\begin{aligned} \lambda_2^* ((p_1^1)^* - \min(\beta_2, \bar{P}_1)) &= 0, \quad \lambda_2^* \geq 0, \\ \lambda_1^* ((p_1^1)^*) &= 0, \quad \lambda_1^* \geq 0, \\ \frac{\partial \mathcal{L}_2}{\partial p_1^1} &= 0, \\ (p_1^1)^* &\geq 0, \\ (p_1^1)^* &\leq \min(\beta_2, \bar{P}_1), \end{aligned} \quad (22)$$

where λ_1^*, λ_2^* are the Lagrangian multipliers associated with the constraints

$$\begin{aligned} \frac{\partial \mathcal{L}_2}{\partial p_1^1} &= 0 \\ \Rightarrow \frac{h_{11}^1}{\sigma_1^2 + h_{21}^2 + h_{11}^1 p_1^1} - \frac{h_{11}^1}{\sigma_2^2 + h_{11}^2 (\bar{P}_1 - p_1^1)} &= \lambda_1^* - \lambda_2^*. \end{aligned} \quad (23)$$

Assume that $p_1^1 = 0, \lambda_1^* \geq 0$, then $\lambda_2^* = 0$ and, furthermore,

$$\frac{h_{11}^1}{\sigma_1^2 + h_{21}^2} - \frac{h_{11}^1}{\sigma_1^2 + h_{11}^2 \bar{P}_1} \geq 0. \quad (24)$$

Assuming that $p_1^1 = \min(\beta_2, \bar{P}_1), \lambda_2^* \geq 0$, then $\lambda_1^* = 0$ and, furthermore,

$$\begin{aligned} \frac{h_{11}^1}{\sigma_1^2 + h_{21}^2 + h_{11}^1 \min(\beta_2, \bar{P}_1)} - \frac{h_{11}^1}{\sigma_2^2 + h_{11}^2 (\bar{P}_1 - \min(\beta_2, \bar{P}_1))} \\ = \lambda_1^* - \lambda_2^*. \end{aligned} \quad (25)$$

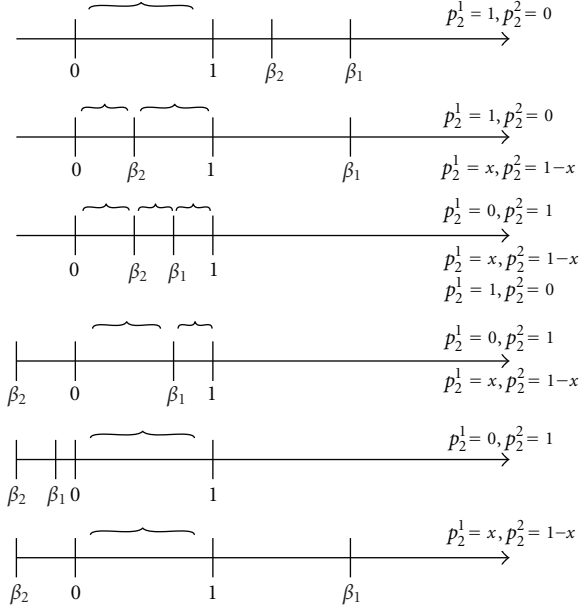


FIGURE 5: Power allocation strategies of the Stackelberg game in which 6 cases exist depending on the variables β_1 and β_2 . The X-axis depicts the strategy space for the leader (p_1^1).

Finally, assume that $0 < p_1^1 < \min(\beta_2, \bar{P}_1)$, then $\lambda_1^* = \lambda_2^* = 0$ and

$$p_1^1 = \frac{h_{11}^1(\sigma_1^2 + h_{11}^1 \bar{P}_1) - h_{11}^2(\sigma_2^2 + h_{21}^1)}{2h_{11}^1 h_{11}^2}. \quad (26)$$

5.1.3. *Operator 2 Transmits in Both Carriers* ($p_2^1 = x, p_2^2 = 1 - x$).

$$\begin{aligned} \max_{p_1^1} \log_2 \left(1 + \frac{|h_{11}^1|^2 p_1^1}{\sigma_1^2 + |h_{21}^1|^2 x} \right) \\ + \log_2 \left(1 + \frac{|h_{11}^2|^2 p_1^2}{\sigma_2^2 + |h_{22}^2|^2 (1-x)} \right), \quad (27) \\ p_1^1 + p_1^2 \leq \bar{P}_1, \\ p_1^1, p_1^2 \geq 0, \\ \beta_2 < p_1^1 < \beta_1. \end{aligned}$$

Since $p_2^1 = x = 1/\mu_2 - (\sigma_1^2 + |h_{12}^1|^2 p_1^1)/|h_{22}^1|^2 > 0$ depends on p_1^1 , the objective function (27) of operator 1 is nonconvex in p_1^1 (the KKT conditions can be written in the same way as done for the previous cases and the problem is solved numerically). Figure 5 depicts all the 6 different cases depending on the values of β_1 and β_2 (note that $\beta_2 < \beta_1$). The X-axis depicts the strategy space for the leader (p_1^1).

As can be seen, in the first case ($\beta_2 > 1, \beta_1 > 1$) the leader has to perform one maximization over the interval $[0, 1]$. In the *second* case ($0 < \beta_2 < 1, \beta_1 > 1$), the leader has to perform 2 maximizations ($[0, \beta_2]$ and $[\beta_2, 1]$) and pick the power allocation that maximizes his payoff. Similarly,

the leader has 3 maximizations to perform in the *third* case ($[0, \beta_2]$, $[\beta_2, \beta_1]$, and $[\beta_1, 1]$), where $0 < \beta_1 < 1, 0 < \beta_2 < 1$ and likewise for the remaining cases. In essence, in all these cases, the leader (operator 1) forces the follower to adopt a strategy that maximizes the leader's payoff. In this way, using backward induction, the Stackelberg equilibrium is *unique*, solving thereby the problem of nonuniqueness encountered in the noncooperative approach of Section 4. Additionally, one should note that there exist Stackelberg solutions that are *non-Nash* equilibria of the noncooperative game.

Another similar approach for solving the Stackelberg equilibrium is used in [31] where an explicit expression of p_2^n (function of p_1^n) needed to analytically find the SE is given as

$$p_2^n = \begin{cases} \left(\bar{P}_2 + \sum_{i=1}^k \left(\sigma_{\pi^{-1}(i)}^2 + |h_{12}^{\pi^{-1}(i)}|^2 p_1^{\pi^{-1}(i)} \right) \right) \\ - \sigma_2^2 - |h_{12}^n|^2 p_1^n, & \pi(i) \leq k, \\ 0, & \pi(i) > k, \end{cases} \quad (28)$$

where π is a permutation function ranking all channels according to their noise plus interference and k can be found from the following condition: $\varphi_k < \bar{P}_2 \leq \varphi_{k+1}$ and $\varphi_t = \sum_{i=1}^k (\sigma_{\pi^{-1}(i)}^2 + |h_{12}^{\pi^{-1}(i)}|^2 p_1^{\pi^{-1}(i)})$, $\forall t \in \{1, \dots, n\}$.

5.2. *Extension to the Multiple Operators' Case.* The inter-operator spectrum sharing in the context of two operators can be extended to the more general case with K operators sharing the same spectrum. The problem is formulated in the same way where the leader's optimization problem is written as

$$\begin{aligned} \max_{p_1^n} \sum_{n=1}^N \log_2 \left(1 + \frac{|h_{11}^n|^2 p_1^n}{\sigma_n^2 + \sum_{j \neq 1}^K |h_{j1}^n|^2 p_j^n(p_1^n)} \right), \\ \sum_{n=1}^N p_1^n \leq \bar{P}_1, \\ p_1^n \geq 0, \end{aligned} \quad (29)$$

and $p_j^{\text{SE}} = \text{BR}_j(p_1^{\text{SE}}, \dots, p_{-j}^{\text{SE}})$ is a function of p_1^n .

Solving (29) becomes much more involved in the general case in which the utility function of the primary operator is nonconvex (p_j^n is function of p_1^n). Nevertheless, there exist suboptimal and low-complexity methods to solve the problem. To this end and motivated by the work of [32], we use lagrangian duality theory wherein the *duality gap* [33] provides a nice tool for solving nonconvex optimization problem.


```

initialize  $\lambda, \bar{P}_1, \bar{P}_2, \dots, \bar{P}_K$ 
repeat
  for  $n = 1, \dots, N$ 
    set  $p_1^n = \operatorname{argmax}_{p_1^n} \sum_{n=1}^N \log_2 \left( 1 + \frac{|h_{11}^n|^2 p_1^n}{\sigma_n^2 + \sum_{j \neq 1}^K |h_{j1}^n|^2 p_j^n(p_1^n)} \right) + \lambda(\bar{P}_1 - \sum_{n=1}^N p_1^n)$ 
    by keeping  $p_1^1, \dots, p_1^{n-1}, p_1^{n+1}, p_1^N$  fixed.
  end
until  $(p_1^1, \dots, p_1^N)$  converges
update  $\lambda$  using subgradient [32] method until it converges.

```

ALGORITHM 1

The lagrangian of (29) is given by

$$\begin{aligned}
 g(\lambda) &= \max_{p_1^n} \mathcal{L}(p_1^n, \lambda) \\
 &= \max_{p_1^n} \sum_{n=1}^N \log_2 \left(1 + \frac{|h_{11}^n|^2 p_1^n}{\sigma_n^2 + \sum_{j \neq 1}^K |h_{j1}^n|^2 p_j^n} \right) \\
 &\quad + \lambda \left(\bar{P}_1 - \sum_{n=1}^N p_1^n \right), \quad (30)
 \end{aligned}$$

where λ is the lagrangian dual variable associated with the power constraint.

Consequently, solving the Stackelberg problem is done by locally optimizing the lagrangian function (30) via coordinate descent [33]. For each fixed set of λ , we find the optimal p_1^1 while keeping p_1^2, \dots, p_1^N fixed, then find the optimal p_1^2 keeping the other p_1^n ($n \neq 2$) fixed, and so on. Such process is guaranteed to converge because each iteration strictly increases the objective function. Finally, λ is found using subgradient [32] method. The algorithm is depicted in Algorithm 1.

Finally, in the case of multiple primary and secondary operators, the optimization problem in (29) becomes a multiobjective optimization problem. In this case, the issue of cooperation is at stake where primary operators i can operate at feasible points yielding rates that dominate the noncooperative equilibrium rate ($R_i \geq R^{NE}$). On the other hand, the secondary operators will adopt the same noncooperative and selfish approach (iterative water-filling).

6. Numerical Evaluation

In this section, numerical results are presented to validate the theoretical claims. Figure 6 depicts the average achievable rate of both operators for the Stackelberg approach. In the simulations, we let the individual power constraint $\bar{P}_1 = \bar{P}_2 = \bar{P} = 1$, $\text{SNR} = \bar{P}/\sigma^2$ and channel fading realizations are independent and identically distributed (i.i.d) Rayleigh distributed.

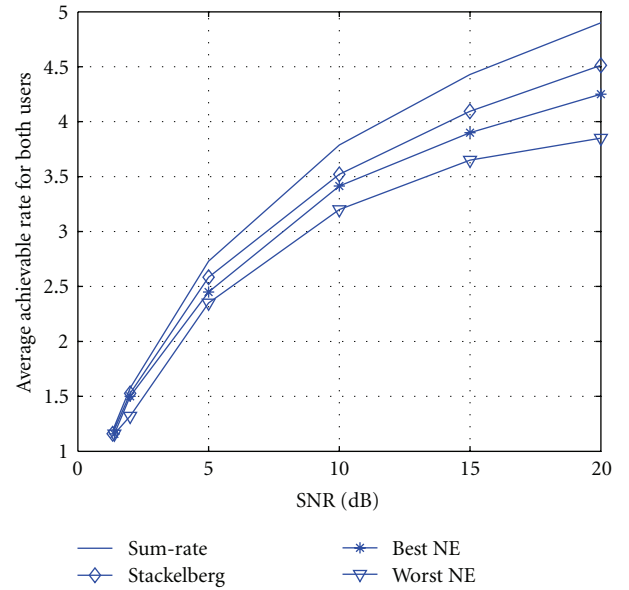


FIGURE 6: Average achievable rate for both users versus the signal-to-noise ratio for the centralized and Stackelberg approach. Moreover, the best and worst Nash equilibria for the non-cooperative game are illustrated.

It is important to quantify the performance loss from the optimal solution provided by the centralized strategy. To this end, we compare the Stackelberg rates with the rates obtained by sum-rate maximization (which are Pareto-optimal):

$$\begin{aligned}
 \max_{p_1^n, p_2^n} \sum_{i=1}^K \sum_{n=1}^N \log_2 \left(1 + \frac{|h_{i1}^n|^2 p_1^n}{\sigma_n^2 + |h_{i2}^n|^2 p_2^n} \right), \\
 \sum_{n=1}^N p_1^n \leq \bar{P}_1, \quad \sum_{n=1}^N p_2^n \leq \bar{P}_2, \\
 p_1^n \geq 0, \quad p_2^n \geq 0, \quad n = 1, \dots, N. \quad (31)
 \end{aligned}$$

The objective function is nonconvex in the power variables p_1^n and p_2^n . To solve (31), the maximization problem is transformed into a convex optimization problem using Geometric Programming [33]. Additionally, Figure 6 depicts the best and worst NE where the best NE refers to the equilibrium maximizing the sum-rate of both operators whereas the

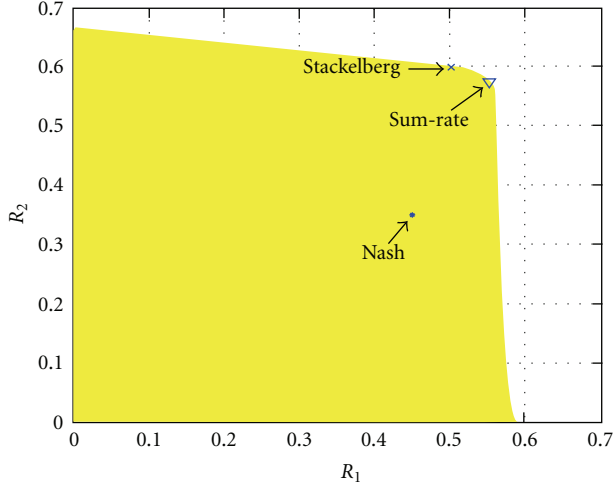


FIGURE 7: Achievable rate region for the inter-operator spectrum sharing game. Both operators achieve better payoffs when adopting the hierarchical (Stackelberg) approach.

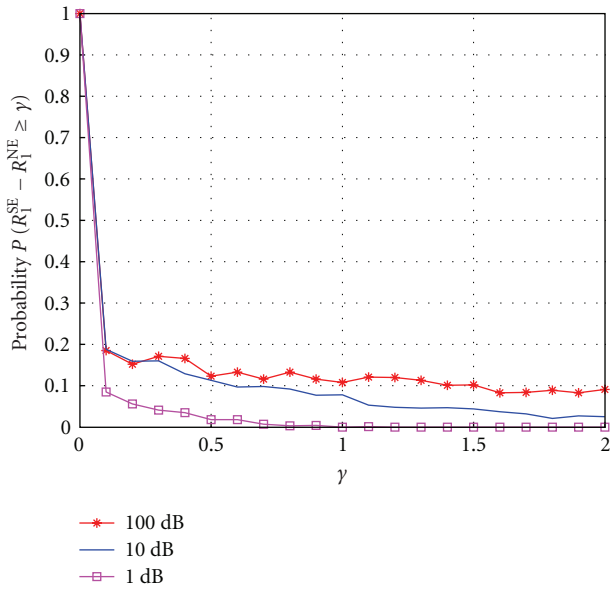


FIGURE 8: Probability when the NE approach is worse than the Stackelberg approach for the leader, for several SNR values.

worst NE case minimizes it. It is also worth noting that the worst Nash equilibrium acts like a lower-bound for the Nash equilibrium. Furthermore, the Stackelberg approach is closer to the centralized approach as compared to the selfish case. This is due to the fact that in the Stackelberg approach, operators take into account other operators' strategies whereas in the selfish case, operators behave carelessly by using water-filling. Figure 7 shows the achievable rate region for both operators in which the Nash and Stackelberg equilibria are illustrated. Since primary operator 1 is the leader, his rate is higher with the Stackelberg approach. Also, interestingly, the rate of operator 2 is also better off with the Stackelberg approach. As a result, both operators have strong incentives

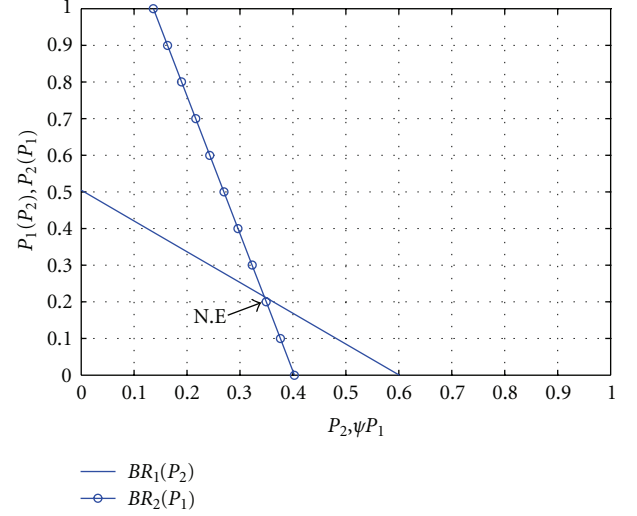


FIGURE 9: Best response functions illustrating the unique Nash equilibrium point where both operators transmit in both carriers.

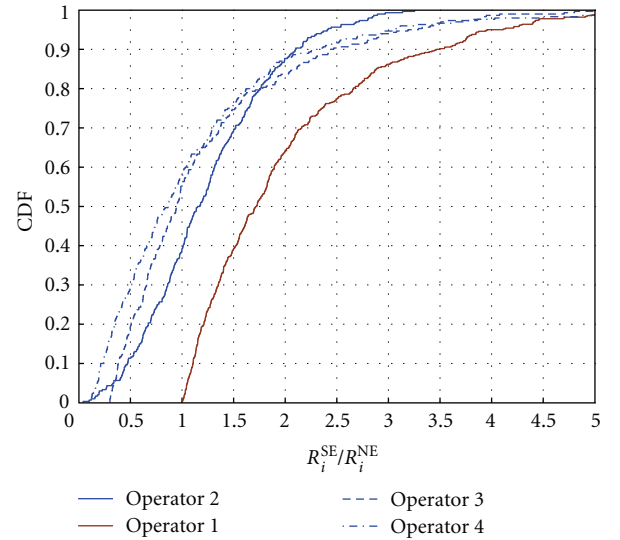


FIGURE 10: Cumulative distribution function (CDF) of the ratio of the rates between the hierarchical (Stackelberg) and noncooperative (selfish) approaches $R_i^{\text{SE}}/R_i^{\text{NE}}$ for operator i , $i = 1, 2, 3, 4$.

in adopting the hierarchical (Stackelberg) approach. Finally, Figure 8 depicts the probability $P(R_1^{\text{SE}} - R_1^{\text{NE}} \geq \gamma)$, that is, when the Nash equilibrium approach is worse than the Stackelberg's for operator 1, for several SNR values where it is seen that the Stackelberg approach outperforms the classical water-filling approach.

On the other hand, Figure 10 depicts the cumulative distribution function (cdf) of the ratio between the achievable rates of the hierarchical and noncooperative approaches. In this scenario, we assume $K = 4$ operators with 1 primary operator and 3 secondary wireless operators sharing the same spectrum composed of $N = 5$ carriers. As can be seen, the primary operator (operator 1) always improves

his achievable rate compared to the selfish approach. The cumulative distribution function of the secondary operators also provides insights on their achievable rates.

7. Conclusion

In this paper, we studied the problem of spectrum sharing between operators from two different perspectives. First, a *one-shot* game was studied where operators play simultaneously, operating at the NE point which exhibits different behaviors according to the set of channel realizations. Then, in a second approach, a *hierarchical* game is examined where the primary operator is the *leader* and the secondary operator is the *follower*, wherein their sum-rates are further improved compared to the water-filling approach. Simulation results show incentives for operators to behave cleverly by adopting the hierarchical (Stackelberg) approach. In our future work, we will use the concept of hierarchy to investigate power control schemes for femtocells networks [34].

Appendices

A.

We derive the a set of 9 inequalities for the Nash equilibria when $K = 2$ operators transmit over $N = 2$ carriers, for the noncooperative game Γ^{NC} .

Recall that $g_{ij}^n = \bar{P}_i |h_{ij}^n|^2 / \sigma^2$, $c_1 = g_{11}^1 / g_{11}^2$, and $c_2 = g_{22}^1 / g_{22}^2$.

It holds that $(\alpha_1, \alpha_2) = (0, 0)$ is a Nash Equilibrium

$$\Leftrightarrow \begin{cases} c_1 \leq \frac{1}{1 + g_{11}^2 + g_{21}^2}, \\ c_2 \leq \frac{1}{1 + g_{22}^2 + g_{12}^2}, \end{cases}$$

$(\alpha_1, \alpha_2) = (1, 0)$ is a Nash Equilibrium

$$\Leftrightarrow \begin{cases} c_1 \geq \frac{1 + g_{11}^1}{1 + g_{21}^2}, \\ c_2 \leq \frac{1 + g_{12}^1}{1 + g_{22}^2}, \end{cases}$$

$(\alpha_1, \alpha_2) = (0, 1)$ is a Nash Equilibrium

$$\Leftrightarrow \begin{cases} c_1 \leq \frac{1 + g_{21}^1}{1 + g_{11}^2}, \\ c_2 \geq \frac{1 + g_{12}^1}{1 + g_{22}^2}, \end{cases}$$

$(\alpha_1, \alpha_2) = (1, 1)$ is a Nash Equilibrium

$$\Leftrightarrow \begin{cases} c_1 \geq 1 + g_{11}^1 + g_{21}^1, \\ c_2 \geq 1 + g_{22}^1 + g_{12}^1, \end{cases}$$

$(\alpha_1, \alpha_2) = (x, 1)$ is a Nash Equilibrium

$$\Leftrightarrow \begin{cases} 2g_{22}^1 g_{11}^1 \leq c_2 g_{11}^1 (1 + g_{12}^2) - 2g_{11}^1 - (c_2 g_{12}^2 + g_{12}^1) \\ \quad \times (c_1 - 1 + c_1 g_{11}^2 - g_{21}^1), \\ \frac{1 + g_{21}^1}{1 + g_{11}^2} \leq c_1 \leq 1 + g_{11}^1 + g_{21}^1, \end{cases} \quad (\text{A.1})$$

where $x = 1/2 - 1/2g_{11}^1 - g_{21}^1/2g_{11}^1 + 1/2g_{11}^2$,

$(\alpha_1, \alpha_2) = (1, x)$ is a Nash Equilibrium

$$\Leftrightarrow \begin{cases} 2g_{11}^1 g_{22}^1 \leq c_1 g_{22}^1 (1 + g_{21}^2) - 2g_{22}^1 - (c_1 g_{21}^2 + g_{21}^1) \\ \quad \times (c_2 - 1 + c_2 g_{22}^2 - g_{12}^1), \\ \frac{1 + g_{12}^1}{1 + g_{22}^2} \leq c_2 \leq 1 + g_{22}^1 + g_{12}^1, \end{cases} \quad (\text{A.2})$$

where $x = 1/2 - 1/2g_{22}^1 - g_{12}^1/2g_{22}^1 + 1/2g_{22}^2$,

$(\alpha_1, \alpha_2) = (0, x)$ is a Nash Equilibrium

$$\Leftrightarrow \begin{cases} \frac{c_2 (1 + g_{12}^2) + g_{22}^1 - 1}{2g_{22}^1} \geq \frac{c_1 (1 + g_{21}^2 + g_{11}^2) - 1}{g_{21}^1 + c_1 g_{21}^2}, \\ \frac{1}{1 + g_{22}^2 + g_{12}^2} \leq c_2 \leq \frac{1 + g_{22}^1}{1 + g_{12}^2}, \end{cases} \quad (\text{A.3})$$

where $x = 1/2 + g_{12}^2/2g_{22}^2 - 1/2g_{22}^1 + 1/2g_{22}^2$,

$(\alpha_1, \alpha_2) = (x, 0)$ is a Nash Equilibrium

$$\Leftrightarrow \begin{cases} \frac{c_1 (1 + g_{21}^2) + g_{11}^1 - 1}{2g_{11}^1} \geq \frac{c_2 (1 + g_{12}^2 + g_{22}^2) - 1}{g_{12}^1 + c_2 g_{12}^2}, \\ \frac{1}{1 + g_{11}^2 + g_{21}^2} \leq c_1 \leq \frac{1 + g_{11}^1}{1 + g_{21}^2}, \end{cases} \quad (\text{A.4})$$

where $x = 1/2 + g_{21}^2/2g_{11}^2 - 1/2g_{11}^1 + 1/2g_{11}^2$,

$(\alpha_1, \alpha_2) = (x, y)$ is a Nash Equilibrium

$$\Leftrightarrow \begin{cases} 0 < x < 1, \\ 0 < y < 1, \\ x = \frac{(2g_{21}^1 (c_1 (1 + g_{12}^2 + g_{21}^2) - 1) - \mathfrak{A})}{(4g_{11}^1 g_{22}^1 - (c_2 g_{12}^2 + g_{12}^1) (c_1 g_{21}^2 + g_{21}^1))}, \\ y = \frac{(1/(2g_{22}^2)) - 1/(2g_{22}^1) + ((1 - x)g_{12}^2)}{(2g_{22}^1) - ((g_{12}^1 x)/2g_{22}^1 + 1/2)}, \end{cases} \quad (\text{A.5})$$

where \mathfrak{A} denotes $(c_1 g_{21}^2 + g_{21}^1)(c_2 (1 + g_{12}^2) + g_{22}^1 - 1)$.

B.

In this setup, the utility functions become

$$\begin{aligned}
 R_1(\alpha_1, \alpha_2) &= \log_2 \left(1 + \frac{|h_{1,1}^1|^2 \alpha_1}{\sigma_1^2 + |h_{2,1}^1|^2 \alpha_2} \right) \\
 &\quad + \log_2 \left(1 + \frac{|h_{1,1}^2|^2 (1 - \alpha_1)}{\sigma_2^2 + |h_{2,1}^2|^2 (1 - \alpha_2)} \right), \\
 R_2(\alpha_1, \alpha_2) &= \log_2 \left(1 + \frac{|h_{1,2}^1|^2 \alpha_2}{\sigma_1^2 + |h_{1,2}^1|^2 \alpha_1} \right) \\
 &\quad + \log_2 \left(1 + \frac{|h_{2,2}^2|^2 (1 - \alpha_2)}{\sigma_2^2 + |h_{1,2}^2|^2 (1 - \alpha_1)} \right).
 \end{aligned} \tag{B.1}$$

We will give now sufficient conditions that guarantee the uniqueness of the NE. By analyzing the first-order derivatives of the payoff functions, we can find explicit relations for the best response functions (BR):

$$\begin{aligned}
 BR_1(\alpha_2) &= \frac{-\mathfrak{B} - |h_{11}^2|^2 + |h_{11}^1|^2 (1 + |h_{21}^2|^2 + |h_{11}^2|^2)}{2 |h_{11}^1|^2 |h_{11}^2|^2}, \\
 BR_2(\alpha_1) &= \frac{-\mathfrak{C} - |h_{22}^2|^2 + |h_{22}^1|^2 (1 + |h_{12}^2|^2 + |h_{22}^2|^2)}{2 |h_{22}^1|^2 |h_{22}^2|^2},
 \end{aligned} \tag{B.2}$$

where \mathfrak{B} denotes $[|h_{11}^2|^2 |h_{21}^1|^2 + |h_{11}^1|^2 |h_{21}^2|^2] \alpha_2$ and \mathfrak{C} denotes $[|h_{22}^2|^2 |h_{12}^1|^2 + |h_{22}^1|^2 |h_{12}^2|^2] \alpha_1$.

We observe that the functions $BR_i(\alpha_{-i})$ are linear with respect to α_{-i} . Thus, the intersection of the BR functions is either a unique point or an infinity of points. Therefore, the sufficient conditions that ensure the uniqueness of the NE are the following:

$$\begin{aligned}
 &\frac{|h_{11}^2|^2 |h_{21}^1|^2 + |h_{11}^1|^2 |h_{21}^2|^2}{2 |h_{11}^1|^2 |h_{11}^2|^2} \\
 &\neq \frac{2 |h_{22}^1|^2 |h_{22}^2|^2}{|h_{22}^2|^2 |h_{12}^1|^2 + |h_{22}^1|^2 |h_{12}^2|^2}, \\
 &\frac{-|h_{11}^2|^2 + |h_{11}^1|^2 (1 + |h_{21}^2|^2 + |h_{11}^2|^2)}{2 |h_{11}^1|^2 |h_{11}^2|^2} \\
 &\neq \frac{-|h_{22}^2|^2 + |h_{22}^1|^2 (1 + |h_{12}^2|^2 + |h_{22}^2|^2)}{|h_{22}^2|^2 |h_{12}^1|^2 + |h_{22}^1|^2 |h_{12}^2|^2}.
 \end{aligned} \tag{B.3}$$

If these conditions are met, the unique point at the intersection of the BRs describes the Nash equilibrium. This is illustrated in Figure 9.

Acknowledgments

This work has been supported by the Finnish Funding Agency for Technology and Innovation (Tekes), Nokia,

Nokia Siemens Networks, Elektrobit and Tauno Tønning Foundation. This work has been performed in part in the framework of the CELTIC project CP5-026 WINNER+ and the french project TERROP. The valuable comments of Dr. Jorma Lilleberg are also very much appreciated.

References

- [1] Y. Xing, R. Chandramouli, S. Mangold, and S. S. N, "Dynamic spectrum access in open spectrum wireless networks," *IEEE Journal on Selected Areas in Communications*, vol. 24, no. 3, pp. 626–637, 2006.
- [2] Z. Ji and K. J. R. Liu, "Multi-stage pricing game for collusion-resistant dynamic spectrum allocation," *IEEE Journal on Selected Areas in Communications*, vol. 26, no. 1, pp. 182–191, 2008.
- [3] M. Bennis, C. Wijting, S. Abedi, S. Thilakawardana, and R. Tafazolli, "Performance evaluation of advanced spectrum functionalities for future radio networks," *Wireless Communications and Mobile Computing*. In press.
- [4] Z. Ji and K. J. R. Liu, "Dynamic spectrum sharing: a game theoretical overview," *IEEE Communications Magazine*, vol. 45, no. 5, pp. 88–94, 2007.
- [5] S. Sengupta, R. Chandramouli, S. Brahma, and M. Chatterjee, "A game theoretic frame-work for distributed self-coexistence among IEEE 802.22 networks," in *Proceedings of IEEE Global Telecommunications Conference (GLOBECOM '08)*, 2008.
- [6] Y. Xing, C. N. Mathur, M. A. Haleem, R. Chandramouli, and K. P. Subbalakshmi, "Dynamic spectrum access with QoS and interference temperature constraints," *IEEE Transactions on Mobile Computing*, vol. 6, no. 4, pp. 423–432, 2007.
- [7] D. Niyato and E. Hossain, "Competitive pricing for spectrum sharing in cognitive radio networks: dynamic game, inefficiency of nash equilibrium, and collusion," *IEEE Journal on Selected Areas in Communications*, vol. 26, no. 1, pp. 192–202, 2008.
- [8] M. Bennis and J. Lilleberg, "Inter base station resource sharing and improving the overall efficiency of B3G systems," in *Proceedings of IEEE Vehicular Technology Conference (VTC '07)*, pp. 1494–1498, Baltimore, Md, USA, September 2007.
- [9] M. Bennis, J. Lara, C. Wijting, and S. Thilakawardana, "WINNER spectrum sharing with fixed satellite services," in *Proceedings of IEEE Vehicular Technology Conference (VTC '09)*, Barcelona, Spain, April 2009.
- [10] M. Bennis and M. Debbah, "Opportunistic power allocation for point-to-point communication in self-organized networks," in *Proceedings of the IEEE Conference on Signals, Systems, and Computers (ASILOMAR '08)*, October 2008.
- [11] R. Etkin, A. Parekh, and D. Tse, "Spectrum sharing for unlicensed bands," *IEEE Journal on Selected Areas in Communications*, vol. 25, no. 3, pp. 517–528, 2007.
- [12] T. S. Han and K. Kobayashi, "A new achievable rate region for the interference channel," *IEEE Transactions on Information Theory*, vol. 27, no. 1, pp. 49–60, 1981.
- [13] W. Yu, G. Ginis, and J. M. Cioffi, "Distributed multiuser power control for digital subscriber lines," *IEEE Journal on Selected Areas in Communications*, vol. 20, no. 5, pp. 1105–1115, 2002.
- [14] A. Leshem and E. Zehavi, "Cooperative game theory and the Gaussian interference channel," *IEEE Journal on Selected Areas in Communications*, vol. 26, no. 7, pp. 1078–1088, 2008.
- [15] R. H. Etkin, D. N. C. Tse, and H. Wang, "Gaussian interference channel capacity to within one bit," *IEEE Transactions on Information Theory*, vol. 54, no. 12, pp. 5534–5562, 2008.

- [16] O. Popescu, D. C. Popescu, and C. Rose, "Simultaneous water filling in mutually interfering systems," *IEEE Transactions on Wireless Communications*, vol. 6, no. 3, pp. 1102–1113, 2007.
- [17] L. Gropop and D. N. C. Tse, "Spectrum sharing between wireless networks," submitted to *IEEE/ACM Transactions on Networking*, arxiv.org/abs/0809.2840.
- [18] E. G. Larsson and E. A. Jorswieck, "Competition versus cooperation on the MISO interference channel," *IEEE Journal on Selected Areas in Communications*, vol. 26, no. 7, pp. 1059–1069, 2008.
- [19] G. Scutari, D. P. Palomar, and S. Barbarossa, "Optimal linear precoding strategies for wideband noncooperative systems based on game theory—part I: nash equilibria," *IEEE Transactions on Signal Processing*, vol. 56, no. 3, pp. 1230–1249, 2008.
- [20] G. Scutari, D. P. Palomar, and S. Barbarossa, "Optimal linear precoding strategies for wideband non-cooperative systems based on game theory—part II: algorithms," *IEEE Transactions on Signal Processing*, vol. 56, no. 3, pp. 1250–1267, 2008.
- [21] G. Scutari, D. P. Palomar, and S. Barbarossa, "Competitive design of multiuser MIMO systems based on game theory: a unified view," *IEEE Journal on Selected Areas in Communications*, vol. 26, no. 7, pp. 1089–1103, 2008.
- [22] V. H. Stackelberg, *Marktform und Gleichgewicht*, Oxford University Press, Oxford, UK, 1934.
- [23] L. Lai and H. El Gamal, "The water-filling game in fading multiple-access channels," *IEEE Transactions on Information Theory*, vol. 54, no. 5, pp. 2110–2122, 2008.
- [24] B. Wang, Z. Han, and K. J. R. Liu, "Distributed relay selection and power control for multiuser cooperative communication networks using stackelberg game," *IEEE Transactions on Mobile Computing*, vol. 8, no. 7, pp. 975–990, 2009.
- [25] S. Lasaulce, Y. Hayel, R. El Azouzi, and M. Debbah, "Introducing hierarchy in energy games," *IEEE Transactions on Wireless Communications*, vol. 8, no. 7, pp. 3833–3843, 2009.
- [26] Y. Su and M. Schaar, "A new perspective on multi-user power control games in interference channels," *IEEE Transactions on Wireless Communications*, vol. 8, no. 6, pp. 2910–2919, 2009.
- [27] D. Fudenberg and J. Tirole, *Game Theory*, MIT Press, Cambridge, UK, 1991.
- [28] J. Nash, "Equilibrium points in n person game," *Proceedings of the National Academy of Sciences of the United States of America*, vol. 36, no. 1, pp. 48–49, 1950.
- [29] J. Rosen, "Existence and uniqueness of equilibrium points for concave n-person games," *Econometrica*, vol. 33, no. 3, pp. 520–534, 1965.
- [30] N. Liu and A. Goldsmith, "Capacity regions and bounds for a class of Z-interference channels," submitted to *IEEE transactions on information theory*, arxiv.org/abs/08.08.0876.
- [31] E. Altman, K. Avrachenkov, and A. Garnaev, "Closed form solutions for symmetric water filling games," in *Proceedings of IEEE International Conference on Computer Communications (INFOCOM '08)*, pp. 673–681, April 2008.
- [32] W. Yu and R. Lui, "Dual methods for nonconvex spectrum optimization of multicarrier systems," *IEEE Transactions on Communications*, vol. 54, no. 7, pp. 1310–1322, 2006.
- [33] S. Boyd and L. Vandenberghe, *Convex Optimization*, Cambridge University Press, Cambridge, UK, 2004.
- [34] V. Chandrasekhar, J. G. Andrews, and A. Gatherer, "Femtocell networks: a survey," *IEEE Communications Magazine*, vol. 46, no. 9, pp. 59–67, 2008.