

Ladder Metamodeling & PLC Program Validation through Time Petri Nets

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Motivation

- **Context:** PLC, Programmable Logical Controller
 - special purpose industrial computer used to automate industrial process
 - connected to inputs and outputs
 - controls the states of outputs according to inputs and internal state
 - programmed with Ladder Diagram (LD) and other languages [IEC 61131]
- **Problem:** Verification of Ladder Diagrams
 - actually, mainly achieved through exhaustive testing
 - purpose: experiment a model-checking approach
 - based on Model-Driven Engineering
- **Focus:** Race condition detection on Ladder diagrams
- **Approach**
 - Ladder translational semantics to Time Petri nets
 - use of Model-Driven Engineering (MDE)
 - use of Tina toolkit (model-checking)

Outline

- 1 Motivations and Approach
- 2 Validation of Ladder Diagrams using Time Petri Net
 - Ladder Diagrams
 - Time Petri Net
 - Translation of Ladder Diagram into Time Petri Net
 - Race Condition Formalisation
- 3 Implementation using MDE
 - General approach
 - Metamodels
 - Transformations
- 4 Conclusion & Future Works

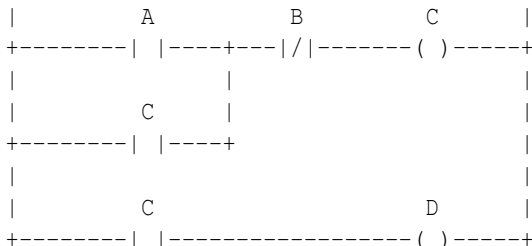
Plan

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Ladder Diagrams

Main concepts

■ An example:



■ **Main concepts:** Rail, Rung, Contact (|| and | / |), Coil (), Variable

■ **Interpretation:** from up to down.

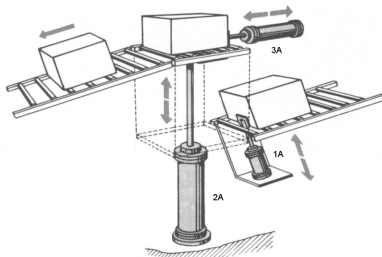
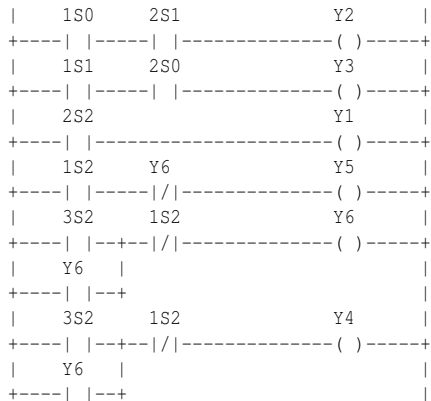
- $C = (A \vee C) \wedge \neg B = (A \wedge \neg B) \vee (C \wedge \neg B) = path1 \vee path2,$

- $D = C.$

■ **Other concepts:** (Not handled in this work)

- Function Blocks with inputs, outputs, internal states.

A more realistic example



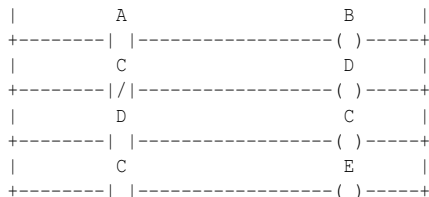
Ladder Diagrams

Race condition

■ Race condition:

Under fixed inputs (and function block states), one or more outputs keep changing their value.

■ An example



■ $B = A$

■ $D = \neg C$

■ $C = D$

■ $E = C$

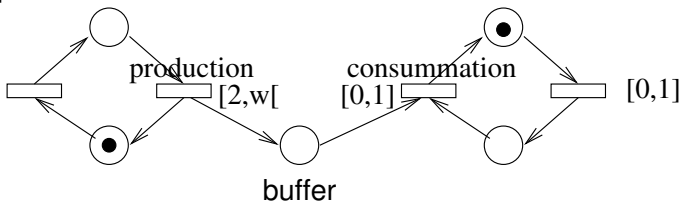
The values of D, C and E keep changing.

■ See later

Time Petri Net

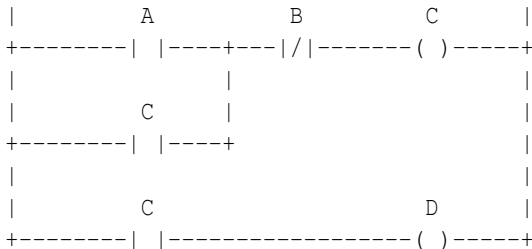
Main concepts

■ An example:

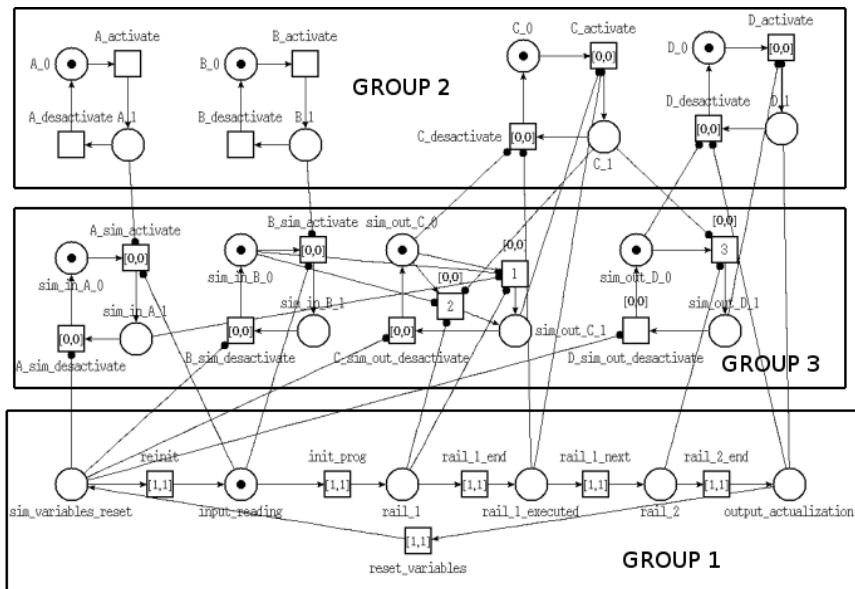


■ Main concepts: Place, Transition, Arc, Token, temporal constraint.

Translation of Ladder Diagram into Time Petri Net



Translation of Ladder Diagram into Time Petri Net



Translation of Ladder Diagram into Time Petri Net

Some explanations

- 1 **Three main parts:** Group 1, Group 2 and Group 3
- 2 **Group 2:** Variables
 - two places for one variable: set (*_1) and unset (*_0) places.
 - *memory* and *output* variables may be updated from simulation variables.
- 3 **Group 1:** Control one computation of all rungs, one at a time
 - initialize simulation variables from Group 2 Variables
 - compute the first rung
 - update Group 2 variables according to simulation variables
 - the same two steps for all other rungs
 - sequencing is done through [1,1] transition
- 4 **Group 3:** “simulation” variables used during a rung computation
 - the simulation variables
 - the way to update them from the Group 2 variables
 - the computation of their new values according to rungs :
 - rungs are decomposed into paths
 - $C = (A \vee C) \wedge \neg B = (A \wedge \neg B) \vee (C \wedge \neg B) = \text{path1} \vee \text{path2}$
 - each path becomes one transition

Race Condition Formalisation

Definition (race condition)

An LD program is free of race condition if

$$\Box (stable_inputs \Rightarrow \Diamond stable_outputs)$$

stable_inputs = logical AND between the stability condition for every input variable

stable_outputs = the same for every output and memory variable.

Definition (stable variable)

An LD variable called *x* is *stable* if $((\Box x_0) \vee (\Box x_1))$.

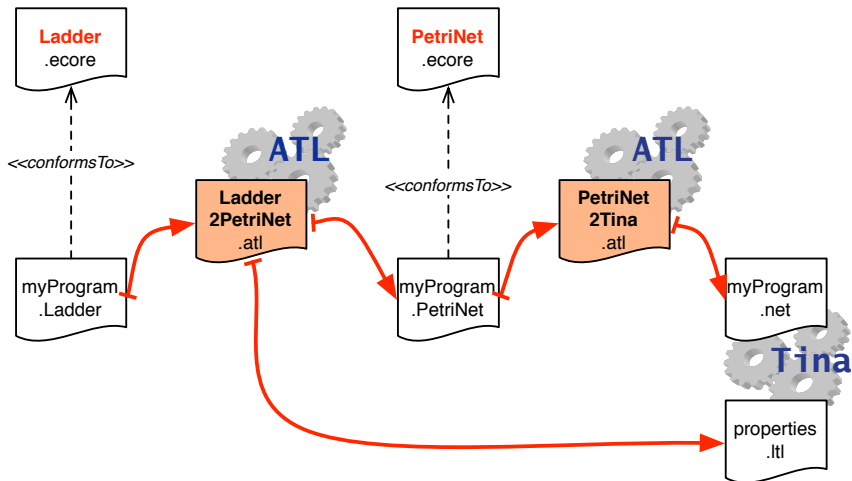
Example

$$\Box (((\Box A_0) \vee (\Box A_1)) \wedge ((\Box B_0) \vee (\Box B_1))) \\ \Rightarrow \Diamond (((\Box C_0) \vee (\Box C_1)) \wedge ((\Box D_0) \vee (\Box D_1)))$$

Plan

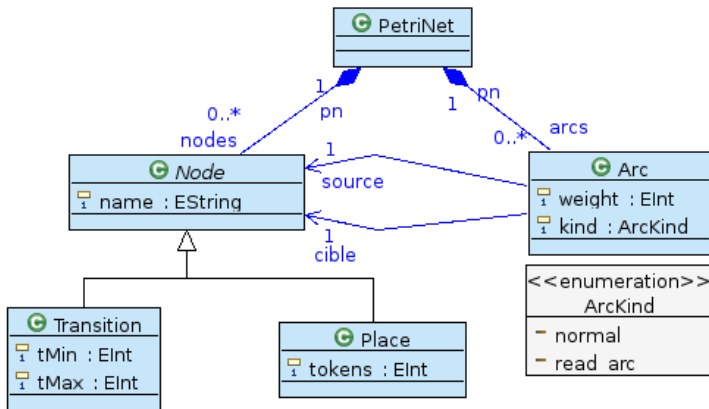
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Approach : Metamodels and Transformations



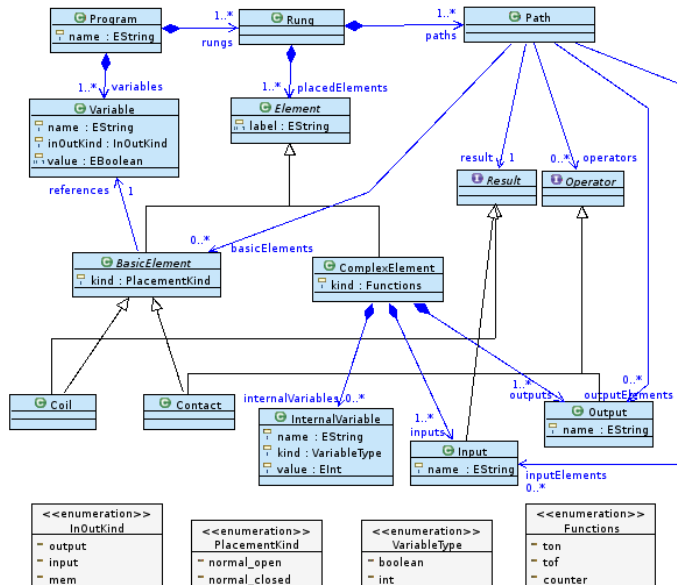
Time Petri Net Metamodel

■ A metamodel



■ and OCL constraints.

Ladder Metamodel



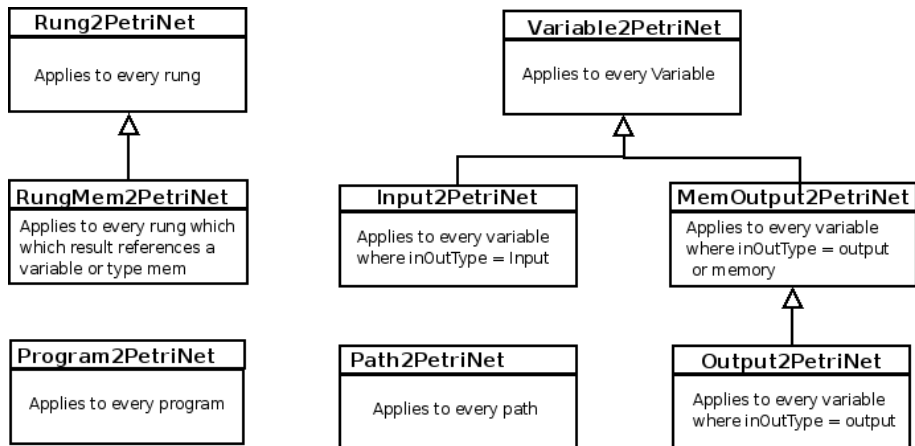
Ladder2PetriNet

Translational Semantics

- Transformation **Model to Model**
- Use of a hybrid transformation language : **ATL**
- Mainly relying on the **declarative style** (rules)
- **Rule inheritance** to structure and factorize code
Used to simulate a role of one source metamodel element (e.g. *Variable*)
- Use of the *resolveTemp* ATL operator (because of inheritance and to identify the right element generated by another rule)

Ladder2PetriNet

Rules architecture



Ladder2LTL

Properties generation

- Achieved through an ATL query (Model to Text)
- Requires to know the names used in the translational semantics

Tina Projector

Transformation from PetriNet model to Tina concrete syntax

- 1 Using an ATL query (Model2Text)
 - using helpers on metamodel elements
- 2 Using TCS (Textual Concrete Syntax) [Jouault *et al.* 2006]
 - Main characteristics:
 - defines textual concrete syntax for a DSL/metamodel
 - based on ANTLR
 - provides both an injector and a projector
 - Results:
 - wide gap between Tina syntax and PetriNet metamodel
 - able to write a TCS model with one warning
 - the injector works :-)
 - the projector does not work :(

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Conclusion

Contribution:

- Ladder metamodel,
- Ladder translational semantics through Petri nets,
- Race condition formalisation.

Conclusion:

- On the verification side:
 - Race conditions are indeed detected
 - Initial coding has been changed to avoid useless behaviours
 - \implies great improvement
- On the MDE side:
 - Should be able to handle large models
 - Easy to implement the evolutions of the Ladder2PetriNet mapping

Future Works

- Use FIACRE as an intermediate language to verification tools (Topcased),
- Handle other Ladder concepts like function blocks,
- Ladder injector to automate the building of the Ladder model,
- Find the changing variable in the counter-example provided by the model-checker.

Thank you
for your attention...

Questions?

Metamodels and transformations are available at

<http://combemale.svn.enseeiht.fr/proto/fr.irit.academie.ladder2tina/>