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► **To cite this version:**

Sylvain Gravier, Svante Janson, Tero Laihonen, Sanna Ranto. Graphs where every  $k$ -subset of vertices is an identifying set. *Discrete Mathematics and Theoretical Computer Science*, 2014, Vol. 16 no. 1 (1), pp.73-88. 10.46298/dmtcs.1253 . hal-00362184

**HAL Id: hal-00362184**

**<https://hal.science/hal-00362184>**

Submitted on 22 Jul 2015

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# Graphs where every $k$ -subset of vertices is an identifying set

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received 16<sup>th</sup> Feb. 2010, revised 7<sup>th</sup> Feb. 2014, accepted 10<sup>th</sup> Feb. 2014.

Let  $G = (V, E)$  be an undirected graph without loops and multiple edges. A subset  $C \subseteq V$  is called *identifying* if for every vertex  $x \in V$  the intersection of  $C$  and the closed neighbourhood of  $x$  is nonempty, and these intersections are different for different vertices  $x$ . Let  $k$  be a positive integer. We will consider graphs where every  $k$ -subset is identifying. We prove that for every  $k > 1$  the maximal order of such a graph is at most  $2k - 2$ . Constructions attaining the maximal order are given for infinitely many values of  $k$ . The corresponding problem of  $k$ -subsets identifying any at most  $\ell$  vertices is considered as well.

**Keywords:** identifying code, extremal graph, strongly regular graph, Plotkin bound

## 1 Introduction

Karpovsky *et al.* introduced identifying sets in [9] for locating faulty processors in multiprocessor systems. Since then identifying sets have been considered in many different graphs (see numerous references in [14]) and they find their motivations, for example, in sensor networks and environmental monitoring [10]. For recent developments see for instance [1, 2].

Let  $G = (V, E)$  be a simple undirected graph where  $V$  is the set of vertices and  $E$  is the set of edges. The adjacency between vertices  $x$  and  $y$  is denoted by  $x \sim y$ , and an edge between  $x$  and  $y$  is denoted by  $\{x, y\}$  or  $xy$ . Suppose  $x, y \in V$ . The (*graphical*) *distance* between  $x$  and  $y$  is the number of edges in any shortest path between these vertices and it is denoted by  $d(x, y)$ . If there is no such path, then  $d(x, y) = \infty$ . We denote by  $N(x)$  the set of vertices adjacent to  $x$  (*neighbourhood*) and the *closed neighbourhood* of a vertex  $x$  is  $N[x] = \{x\} \cup N(x)$ . The closed neighbourhood within radius  $r$  centered at  $x$  is denoted by  $N_r[x] = \{y \in V \mid d(x, y) \leq r\}$ . We denote further  $S_r(x) = \{y \in V \mid d(x, y) = r\}$ . Moreover, for  $X \subseteq V$ ,  $N_r[X] = \cup_{x \in X} N_r[x]$ . For  $C \subseteq V$ ,  $X \subseteq V$ , and  $x \in V$  we denote

$$I_r(C; x) = I_r(x) = N_r[x] \cap C$$

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and

$$I_r(C; X) = I_r(X) = N_r[X] \cap C = \bigcup_{x \in X} I_r(C; x).$$

If  $r = 1$ , we drop it from the notations. When necessary, we add a subscript  $G$ . We also write, for example,  $N[x, y]$  and  $I(C; x, y)$  for  $N[\{x, y\}]$  and  $I(C; \{x, y\})$ . The *symmetric difference* of two sets is

$$A \triangle B = (A \setminus B) \cup (B \setminus A).$$

The cardinality of a set  $X$  is denoted by  $|X|$ ; we will also write  $|G|$  for the order  $|V|$  of a graph  $G = (V, E)$ . The *degree* of a vertex  $x$  is  $\deg(x) = |N(x)|$ . Moreover,  $\delta_G = \delta = \min_{x \in V} \deg(x)$  and  $\Delta_G = \Delta = \max_{x \in V} \deg(x)$ . The *diameter* of a graph  $G = (V, E)$  is  $\text{diam}(G) = \max\{d(x, y) \mid x, y \in V\}$ .

We say that a vertex  $x \in V$  *dominates* a vertex  $y \in V$  if and only if  $y \in N[x]$ . As well we can say that a vertex  $y$  is *dominated* by  $x$  (or vice versa). A subset  $C$  of vertices  $V$  is called a *dominating set* (or *dominating*) if  $\cup_{c \in C} N[c] = V$ .

**Definition 1** A subset  $C$  of vertices of a graph  $G = (V, E)$  is called  $(r, \leq \ell)$ -identifying (or an  $(r, \leq \ell)$ -identifying set) if for all  $X, Y \subseteq V$  with  $|X| \leq \ell$ ,  $|Y| \leq \ell$ ,  $X \neq Y$  we have

$$I_r(C; X) \neq I_r(C; Y).$$

If  $r = 1$  and  $\ell = 1$ , then we speak about an identifying set.

The idea behind identification is that we can uniquely determine the subset  $X$  of vertices of a graph  $G = (V, E)$  by knowing only  $I_r(C; X)$  — provided that  $|X| \leq \ell$  and  $C \subseteq V$  is an  $(r, \leq \ell)$ -identifying set.

**Definition 2** Let, for  $n \geq k \geq 1$  and  $\ell \geq 1$ ,  $\mathfrak{G}\mathfrak{r}(n, k, \ell)$  be the set of graphs on  $n$  vertices such that every  $k$ -element set of vertices is  $(1, \leq \ell)$ -identifying. Moreover, we denote  $\mathfrak{G}\mathfrak{r}(n, k, 1) = \mathfrak{G}\mathfrak{r}(n, k)$  and  $\mathfrak{G}\mathfrak{r}(k) = \bigcup_{n \geq k} \mathfrak{G}\mathfrak{r}(n, k)$ .

In other words, in a sensor network which is modeled by a graph in the class  $\mathfrak{G}\mathfrak{r}(n, k, \ell)$  we can choose freely  $k$  sensors [10] i.e. vertices to locate any  $\ell$  objects in vertices.

**Example 3** (i) For every  $\ell \geq 1$ , an empty graph  $E_n = (\{1, \dots, n\}, \emptyset)$  belongs to  $\mathfrak{G}\mathfrak{r}(n, k, \ell)$  if and only if  $k = n$ .

(ii) A cycle  $C_n$  ( $n \geq 4$ ) belongs to  $\mathfrak{G}\mathfrak{r}(n, k)$  if and only if  $n - 1 \leq k \leq n$ . A cycle  $C_n$  with  $n \geq 7$  is in  $\mathfrak{G}\mathfrak{r}(n, n, 2)$ .

(iii) A path  $P_n$  of  $n$  vertices ( $n \geq 3$ ) belongs to  $\mathfrak{G}\mathfrak{r}(n, k)$  if and only if  $k = n$ .

(iv) A complete bipartite graph  $K_{n,m}$  ( $n+m \geq 4$ ) is in  $\mathfrak{G}\mathfrak{r}(n+m, k)$  if and only if  $n+m-1 \leq k \leq n+m$ .

(v) In particular, a star  $S_n = K_{1,n-1}$  ( $n \geq 4$ ) is in  $\mathfrak{G}\mathfrak{r}(n, k)$  if and only if  $n - 1 \leq k \leq n$ .

(vi) The complete graph  $K_n$  ( $n \geq 2$ ) is not in  $\mathfrak{G}\mathfrak{r}(n, k)$  for any  $k$ .

We are interested in the maximum number  $n$  of vertices which can be reached by a given  $k$ . We study mainly the case  $\ell = 1$  and define

$$\Xi(k) = \max\{n : \mathfrak{G}\tau(n, k) \neq \emptyset\}. \quad (1)$$

Conversely, the question is for a given graph on  $n$  vertices what is the smallest number  $k$  such that every  $k$ -subset of vertices is an identifying set (or a  $(1, \leq \ell)$ -identifying set). (Note that even if we take  $k = n$ , there are graphs on  $n$  vertices that do not belong to  $\mathfrak{G}\tau(n, n)$ , for example the complete graph  $K_n$ ,  $n \geq 2$ .) The ratio  $n/k$  is called the *rate*.

In particular, we are interested in the asymptotics as  $k \rightarrow \infty$ . Combining Theorem 17 and Corollary 26, we obtain the following, which in particular shows that the rate is always less than 2.

**Theorem 4**  $\Xi(k) \leq 2k - 2$  for all  $k \geq 2$ , and  $\lim_{k \rightarrow \infty} \frac{\Xi(k)}{k} = 2$ .

We will see in Section 4 that  $\Xi(k) = 2k - 2$  for infinitely many  $k$ .

We give some basic results in Section 2 and study small  $k$  in Section 2.1 where we give a complete description of the sets  $\mathfrak{G}\tau(k)$  for  $k \leq 4$ . In Section 3 we give an upper bound, which bases on a relation with error-correcting codes. We consider strongly regular graphs and some modifications of them in Section 4; this provides us with examples (e.g., Paley graphs) that attain or almost attain the upper bound in Theorem 4. In Section 5 we give results for the case  $\ell \geq 2$ .

## 2 Basic results

We begin with some simple consequences of the definition. We omit the simple proofs.

**Lemma 5** If  $G = (V, E) \in \mathfrak{G}\tau(n, k, \ell)$ , then every induced subgraph  $G[A]$ , where  $A \subseteq V$ , of order  $|A| = m \geq k$  belongs to  $\mathfrak{G}\tau(m, k, \ell)$ .

**Lemma 6** If  $G$  has connected components  $G_i$ ,  $i = 1, \dots, m$ , with  $|G| = n$  and  $|G_i| = n_i$ , then  $G \in \mathfrak{G}\tau(n, k, \ell)$  if and only if  $G_i \in \mathfrak{G}\tau(n_i, k + n_i - n, \ell)$  for every  $i$ . In other words,  $G_i \in \mathfrak{G}\tau(n_i, k_i, \ell)$  with  $n_i - k_i = n - k$ .

A graph  $G$  belongs to  $\mathfrak{G}\tau(n, k, \ell)$  if and only if every  $k$ -subset intersects every symmetric difference of the neighbourhoods of two sets that are of size at most  $\ell$ . Equivalently,  $G \in \mathfrak{G}\tau(n, k, \ell)$  if and only if the complement of every such symmetric difference of two neighbourhoods contains less than  $k$  vertices. We state this as a theorem.

**Theorem 7** Let  $G = (V, E)$  and  $|V| = n$ . A graph  $G$  belongs to  $\mathfrak{G}\tau(n, k, \ell)$  if and only if

$$n - \min_{\substack{X, Y \subseteq V \\ X \neq Y \\ |X|, |Y| \leq \ell}} \{|N[X] \Delta N[Y]|\} \leq k - 1. \quad (2)$$

Now take  $\ell = 1$ , and consider  $\mathfrak{G}\tau(n, k)$ . The characterization in Theorem 7 can be written as follows, since  $X$  and  $Y$  either are empty or singletons.

**Corollary 8** Let  $G = (V, E)$  and  $|V| = n$ . A graph  $G$  belongs to  $\mathfrak{G}\tau(n, k)$  if and only if

- (i)  $\delta_G \geq n - k$ , and

$$(ii) \max_{x,y \in V, x \neq y} \{|N[x] \cap N[y]| + |V \setminus (N[x] \cup N[y])|\} \leq k - 1.$$

In particular, if  $G \in \mathfrak{G}\tau(n, k)$  then every vertex is dominated by every choice of a  $k$ -subset, and for all distinct  $x, y \in V$  we have  $|N[x] \cap N[y]| \leq k - 1$ .

**Example 9** Let  $G$  be the 3-dimensional cube, with 8 vertices. Then  $|N[x]| = 4$  for every vertex  $x$ , and  $|N[x] \triangle N[y]|$  is 4 when  $d(x, y) = 1$ , 4 when  $d(x, y) = 2$ , and 8 when  $d(x, y) = 3$ . Hence, Theorem 7 shows that  $G \in \mathfrak{G}\tau(8, 5)$ .

**Lemma 10** Let  $G_0 = (V_0, E_0) \in \mathfrak{G}\tau(n_0, k_0)$  and let  $G = (V_0 \cup \{a\}, E_0 \cup \{\{a, x\} \mid x \in V_0\})$  for a new vertex  $a \notin V_0$ . In words, we add a vertex and connect it to all other vertices. Then  $G \in \mathfrak{G}\tau(n_0 + 1, k_0 + 1)$  if (and only if)  $|N_{G_0}[x]| \leq k_0 - 1$  for every  $x \in V_0$ , or, equivalently,  $\Delta_{G_0} \leq k_0 - 2$ .

**Proof:** An immediate consequence of Theorem 7 (or Corollary 8).  $\square$

**Example 11** If  $G_0$  is the 3-dimensional cube in Example 9, which belongs to  $\mathfrak{G}\tau(8, 5)$  and is regular with degree  $3 = 5 - 2$ , then Lemma 10 yields a graph  $G \in \mathfrak{G}\tau(9, 6)$ .  $G$  can be regarded as a cube with centre.

## 2.1 Small $k$

**Example 12** For  $k = 1$ , it is easily seen that  $\mathfrak{G}\tau(n, 1) = \emptyset$  for  $n \geq 2$ , and thus  $\mathfrak{G}\tau(1) = \{K_1\}$  and  $\Xi(1) = 1$ .

**Example 13** Let  $k = 2$ . If  $G \in \mathfrak{G}\tau(2)$ , then  $G$  cannot contain any edge  $xy$ , since then  $N[x] \cap \{x, y\} = \{x, y\} = N[y] \cap \{x, y\}$ , so  $\{x, y\}$  does not separate  $\{x\}$  and  $\{y\}$ . Consequently,  $G$  has to be an empty graph  $E_n$ , and then  $\delta_G = 0$  and Corollary 8(i) (or Example 3(i)) shows that  $n = k = 2$ . Thus  $\mathfrak{G}\tau(2) = \{E_2\}$  and  $\Xi(2) = 2$ .

**Example 14** Let  $k = 3$ . First, assume  $n = |G| = 3$ . There are only four graphs  $G$  with  $|G| = 3$ , and it is easily checked that  $E_3, P_3 \in \mathfrak{G}\tau(3, 3)$  (Example 3(i)(iii)), while  $C_3 = K_3 \notin \mathfrak{G}\tau(3, 3)$  (Example 3(vi)) and a disjoint union  $K_1 \cup K_2 \notin \mathfrak{G}\tau(3, 3)$ , for example by Lemma 6 since  $K_2 \notin \mathfrak{G}\tau(2, 2)$ . Hence  $\mathfrak{G}\tau(3, 3) = \{E_3, P_3\}$ .

Next, assume  $n \geq 4$ . Since there are no graphs in  $\mathfrak{G}\tau(n_1, k_1)$  if  $n_1 > k_1$  and  $k_1 \leq 2$ , it follows from Lemma 6 that there are no disconnected graphs in  $\mathfrak{G}\tau(n, 3)$  for  $n \geq 4$ . Furthermore, if  $G \in \mathfrak{G}\tau(n, 3)$ , then every induced subgraph with 3 vertices is in  $\mathfrak{G}\tau(3, 3)$  and is thus  $E_3$  or  $P_3$ ; in particular,  $G$  contains no triangle.

If  $G \in \mathfrak{G}\tau(4, 3)$ , it follows easily that  $G$  must be  $C_4$  or  $S_4$ , and indeed these belong to  $\mathfrak{G}\tau(4, 3)$  by Example 3(ii)(v). Hence  $\mathfrak{G}\tau(4, 3) = \{C_4, S_4\}$ .

Next, assume  $G \in \mathfrak{G}\tau(5, 3)$ . Then every induced subgraph with 4 vertices is in  $\mathfrak{G}\tau(4, 3)$  and is thus  $C_4$  or  $S_4$ . Moreover, by Corollary 8,  $\delta_G \geq 5 - 3 = 2$ . However, if we add a vertex to  $C_4$  or  $S_4$  such that the degree condition  $\delta_G \geq 2$  is satisfied and we do not create a triangle we get  $K_{2,3}$  – a complete bipartite graph, and we know already  $K_{2,3} \notin \mathfrak{G}\tau(5, 3)$  (Example 3(iv)). Consequently  $\mathfrak{G}\tau(5, 3) = \emptyset$ , and thus  $\mathfrak{G}\tau(n, 3) = \emptyset$  for all  $n \geq 5$ .

Consequently,  $\mathfrak{G}\tau(3) = \mathfrak{G}\tau(3, 3) \cup \mathfrak{G}\tau(4, 3) = \{E_3, P_3, S_4, C_4\}$  and  $\Xi(3) = 4$ .

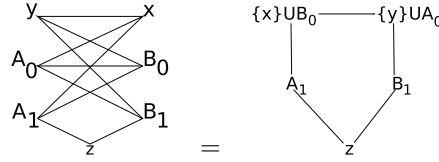
**Example 15** Let  $k = 4$ . First, it follows easily from Lemma 6 and the descriptions of  $\mathfrak{G}\tau(j)$  for  $j \leq 3$  above that the only disconnected graphs in  $\mathfrak{G}\tau(4)$  are  $E_4$  and the disjoint union  $P_3 \cup K_1$ ; in particular, every graph in  $\mathfrak{G}\tau(n, 4)$  with  $n \geq 5$  is connected.

Next, if  $G \in \mathfrak{G}\tau(n, 4)$ , there cannot be a triangle in  $G$  because otherwise if a 4-subset includes the vertices of a triangle, one more vertex cannot separate the vertices of the triangle from each other. (Cf. Lemma 19.)

For  $n = 4$ , the only connected graphs of order 4 that do not contain a triangle are  $C_4$ ,  $P_4$  and  $S_4$ , and these belong to  $\mathfrak{G}\tau(4, 4)$  by Example 3(ii)(iii)(v). Hence  $\mathfrak{G}\tau(4, 4) = \{C_4, P_4, S_4, E_4, P_3 \cup K_1\}$ .

Now assume that  $G \in \mathfrak{G}\tau(n, 4)$  with  $n \geq 5$ .

(i) Suppose first that a graph  $K_1 \cup K_2 = (\{x, y, z\}, \{\{x, y\}\})$  is an induced subgraph of  $G$ . Then all the other vertices of  $G$  are adjacent to either  $x$  or  $y$  but not both, since otherwise there would be an induced triangle or an induced  $E_2 \cup K_2$  or  $K_2 \cup K_2$ , and these do not belong to  $\mathfrak{G}\tau(4, 4)$ . Let  $A = N(x) \setminus \{y\}$  and  $B = N(y) \setminus \{x\}$ , so we have a partition of the vertex set as  $\{x, y, z\} \cup A \cup B$ . There can be further edges between  $A$  and  $B$ ,  $z$  and  $A$ ,  $z$  and  $B$  but not inside  $A$  and  $B$ . Let  $A = A_0 \cup A_1$  and  $B = B_0 \cup B_1$ , where  $A_1 = \{a \in A \mid a \sim z\}$ ,  $A_0 = A \setminus A_1$  and  $B_1 = \{b \in B \mid b \sim z\}$ ,  $B_0 = B \setminus B_1$ . If  $a \in A_0$  and  $b \in B$ , then the 4-subset  $\{a, b, x, z\}$  does not distinguish  $a$  and  $x$  unless  $a \sim b$ . Similarly, if  $a \in A$  and  $b \in B_0$ , then  $a \sim b$ . On the other hand, if  $a \in A_1$  and  $b \in B_1$ , then  $a \not\sim b$ , since otherwise  $abz$  would be a triangle. Thus, we have, where one or more of the sets  $A_0, A_1, B_0, B_1$  might be empty, where an edge



is a complete bipartite graph on sets incident to it, and there are no edges inside these sets.

If  $n \geq 6$ , then there are at least two elements in one of the sets  $\{x\} \cup B_0$ ,  $\{y\} \cup A_0$ ,  $A_1$  or  $B_1$ . However, these two vertices have the same neighbourhood and hence they cannot be separated by the other  $n - 2 \geq 4$  vertices. Thus,  $n = 5$ .

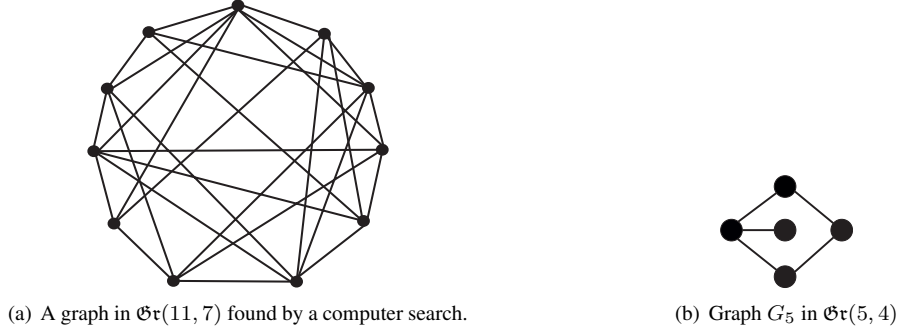
If  $n = 5$ , and both  $A_1$  and  $B_1$  are non-empty, we must have  $A_0 = B_0 = \emptyset$  and  $G = C_5$ , which is in  $\mathfrak{G}\tau(5, 4)$  by Example 3(ii).

Finally, assume  $n = 5$  and  $A_1 = \emptyset$  (the case  $B_1 = \emptyset$  is the same after relabelling). Then  $B_1$  is non-empty, since  $G$  is connected. If  $B_0$  is non-empty, let  $b_0 \in B_0$  and  $b_1 \in B_1$ , and observe that  $\{x, b_0, b_1, z\}$  does not separate  $z$  and  $b_1$ . Hence  $B_0 = \emptyset$ . We thus have either  $|A_0| = 1$  and  $|B_1| = 1$ , or  $|A_0| = 0$  and  $|B_1| = 2$ , and both cases yield the graph in Figure 1(b) which easily is seen to be in  $\mathfrak{G}\tau(5, 4)$ .

(ii) Suppose that there is no induced subgraph  $K_1 \cup K_2$ . Since  $G$  is connected, we can find an edge  $x \sim y$ . Let, as above,  $A = N(x) \setminus \{y\}$  and  $B = N(y) \setminus \{x\}$ . If  $a \in A$  and  $b \in B$  and  $a \not\sim b$ , then  $(\{a, x, b\}, \{\{a, x\}\})$  is an induced subgraph and we are back in case (i). Hence, all edges between sets  $A$  and  $B$  exist and thus, recalling that  $G$  has no triangles,  $G$  is the complete bipartite graph with bipartition  $(A \cup \{y\}, B \cup \{x\})$ . By Example 3(iv), then  $n \leq 5$ . If  $n = 5$ , we get  $G = K_{2,3}$  or  $G = K_{1,4} = S_5$ , which both belong to  $\mathfrak{G}\tau(5, 4)$  by Example 3(iv).

We summarize the result in a theorem.

**Theorem 16** We have  $\Xi(4) = 5$ . More precisely,  $\mathfrak{G}\tau(4) = \mathfrak{G}\tau(4, 4) \cup \mathfrak{G}\tau(5, 4)$ , where  $\mathfrak{G}\tau(4, 4) = \{C_4, P_4, S_4, E_4, P_3 \cup K_1\}$  and  $\mathfrak{G}\tau(5, 4) = \{S_5, C_5, K_{2,3}, G_5\}$  where  $G_5$  is the graph in Figure 1(b).



**Fig. 1:** Examples in  $\mathfrak{G}\mathfrak{r}(11, 7)$  and  $\mathfrak{G}\mathfrak{r}(5, 4)$ .

Upper and lower bounds for  $\Xi(k)$  for  $1 \leq k \leq 20$  are given in Table 1. Note that we have determined  $\Xi(k)$  exactly for  $k \leq 6$  and for 9, 19, but not for other values of  $k$  when  $k \leq 20$ .

### 3 Upper estimates on the order

In the next theorem we give an upper bound on  $\Xi(k)$ , which is obtained using knowledge on error-correcting codes.

**Theorem 17** *If  $k \geq 2$ , then  $\Xi(k) \leq 2k - 2$ .*

**Proof:** We begin by giving a construction from a graph in  $\mathfrak{G}\mathfrak{r}(n, k)$  to error-correcting codes. A non-existence result of error-correcting codes then yields the non-existence of  $\mathfrak{G}\mathfrak{r}(n, k)$  graphs of certain parameters. Let  $G = (V, E) \in \mathfrak{G}\mathfrak{r}(n, k)$ , where  $V = \{x_1, x_2, \dots, x_n\}$ . We construct  $n + 1$  binary strings  $y_i = (y_{i1}, \dots, y_{in})$  of length  $n$ , for  $i = 0, \dots, n$ , from the sets  $\emptyset = N[\emptyset]$  and  $N[x_i]$  for  $i = 1, \dots, n$  by defining  $y_{0j} = 0$  for all  $j$  and

$$y_{ij} = \begin{cases} 0 & \text{if } x_j \notin N[x_i] \\ 1 & \text{if } x_j \in N[x_i] \end{cases}, \quad 1 \leq i \leq n.$$

Let  $C$  denote the code which consists of these binary strings as codewords. Because  $G \in \mathfrak{G}\mathfrak{r}(n, k)$ , the symmetric difference of two closed neighbourhoods  $N[x_i]$  and  $N[x_j]$ , or of one neighbourhood  $N[x_i]$  and  $\emptyset$ , is at least  $n - k + 1$  by (2); in other words, the minimum Hamming distance  $d(C)$  of the code  $C$  is at least  $n - k + 1$ .

We first give a simple proof that  $\Xi(k) \leq 2k - 1$ . Thus, suppose that there is a  $G \in \mathfrak{G}\mathfrak{r}(n, k)$  such that  $n = 2k$ . In the corresponding error-correcting code  $C$ , the minimum distance is at least  $d = n - k + 1 = k + 1 > n/2$ . Let the maximum cardinality of the error-correcting codes of length  $n$  and minimum distance at least  $d$  be denoted by  $A(n, d)$ . We can apply the Plotkin bound (see for example [15, Chapter 2, §2]), which says  $A(n, d) \leq 2 \lfloor d/(2d - n) \rfloor$ , when  $2d > n$ . Thus, we have

$$A(n, d) \leq 2 \left\lfloor \frac{k+1}{2} \right\rfloor \leq k+1.$$

Because  $k + 1 < 2k = n < |C|$ , this contradicts the existence of  $C$ . Hence, there cannot exist a graph  $G \in \mathfrak{Gr}(2k, k)$ , and thus  $\mathfrak{Gr}(n, k) = \emptyset$  when  $n \geq 2k$ .

The Plotkin bound is not strong enough to imply  $\Xi(k) \leq 2k - 2$  in general, but we obtain this from the proof of the Plotkin bound as follows. (In fact, for odd  $k$ ,  $\Xi(k) \leq 2k - 2$  follows from the Plotkin bound for an odd minimum distance. We leave this to the reader since the argument below is more general.)

Suppose that  $G = (V, E) \in \mathfrak{Gr}(n, k)$  with  $n = 2k - 1$ . We thus have a corresponding error-correcting code  $C$  with  $|C| = n + 1 = 2k$  and minimum Hamming distance at least  $n - k + 1 = k$ . Hence, letting  $d$  denote the Hamming distance,

$$\sum_{0 \leq i < j \leq n} d(y_i, y_j) \geq \binom{n+1}{2} k = \frac{2k(2k-1)}{2} k = (2k-1)k^2. \quad (3)$$

On the other hand, if there are  $s_m$  strings  $y_i$  with  $y_{im} = 1$ , and thus  $|C| - s_m = 2k - s_m$  strings with  $y_{im} = 0$ , then the number of ordered pairs  $(i, j)$  such that  $y_{im} \neq y_{jm}$  is  $2s_m(2k - s_m)$  and this parabola gives  $2s_m(2k - s_m) \leq 2k^2$ . Hence each bit contributes at most  $k^2$  to the sum in (3), and summing over  $m$  we find

$$\sum_{0 \leq i < j \leq n} d(y_i, y_j) \leq nk^2 = (2k-1)k^2. \quad (4)$$

Consequently, we have equality in (3) and (4), and thus  $d(y_i, y_j) = k$  for all pairs  $(i, j)$  with  $i \neq j$ .

In particular,  $|N[x_i]| = d(y_i, y_0) = k$  for  $i = 1, \dots, n$ , and thus every vertex in  $G$  has degree  $k - 1$ , i.e.,  $G$  is  $(k - 1)$ -regular. Hence,  $2|E| = n(k - 1) = (2k - 1)(k - 1)$ , and  $k$  must be odd.

Further, if  $i \neq j$ , then  $|N[x_i] \triangle N[x_j]| = d(y_i, y_j) = k$ , and since  $N[x_i] \setminus N[x_j]$  and  $N[x_j] \setminus N[x_i]$  have the same size  $k - |N[x_i] \cap N[x_j]|$ , they have both the size  $k/2$  and  $k$  must be even.

This contradiction shows that  $\mathfrak{Gr}(2k - 1, k) = \emptyset$ , and thus  $\Xi(k) \leq 2k - 2$ .  $\square$

The next theorem (which does not use Theorem 17) will lead to another upper bound in Theorem 20. It can be seen as an improvement for the extreme case  $\mathfrak{Gr}(2k - 2, k)$  of Mantel's [16] theorem on existence of triangles in a graph. Note that this result fails for  $k = 5$  by Example 9.

**Theorem 18** *Suppose  $G \in \mathfrak{Gr}(n, k)$  and  $k \geq 6$ . If  $n \geq 2k - 2$ , then there is a triangle in  $G$ .*

**Proof:** Let  $G = (V, E) \in \mathfrak{Gr}(n, k)$ . Suppose to the contrary that there are no triangles in  $G$ . If there is a vertex  $x \in V$  such that  $\deg(x) \geq k + 1$ , then we select in  $N(x)$  a  $k$ -set  $X$  and a vertex  $y$  outside it; since  $X$  has to dominate  $y$ , it is clear that there exists a triangle  $xyz$ . Hence  $\deg(x) \leq k$  for every  $x$ . On the other hand, we know by Corollary 8(i) that for all  $x \in V$   $\deg(x) \geq n - k \geq k - 2$ .

Let  $x \in V$  be a vertex whose degree is minimum. We denote  $V \setminus N[x] = B$  and we use the fact that  $|B| \leq k - 1$ .

1) Suppose first  $\deg(x) = k$ . Because  $\deg(x)$  is minimum we know that for all  $a \in N(x)$ ,  $\deg(a) = k$ . This is possible if and only if  $|B| = k - 1$  and for all  $a \in N(x)$  we have  $B \cap N(a) = B$ . But then in the  $k$ -subset  $C = \{x\} \cup B$  we have  $I(C; a) = I(C; b)$  for all  $a, b \in N(x)$ . This is impossible.

2) Suppose then  $\deg(x) = k - 1$ . If now  $|B| \leq k - 2$  the graph is impossible as in the first case (choose  $C = N[x]$ ). Hence,  $|B| = k - 1$ . For every  $a \in N(x)$  there are at least  $k - 2$  adjacent vertices in  $B$ , and thus at most 1 non-adjacent. This implies that for all  $a, b \in N(x)$ ,  $a \neq b$ , we have  $|N(a) \cap N(b) \cap B| \geq k - 3 \geq 2$ , when  $k \geq 5$ . Hence, by choosing  $a, b \in N(x)$ ,  $a \neq b$ , we have



the  $k$ -subset  $C = \{x\} \cup (N(x) \setminus \{a, b\}) \cup \{c_1, c_2\}$ , where  $c_1, c_2 \in N(a) \cap N(b) \cap B$ . In this  $k$ -subset  $I(C; a) = I(C; b)$ , which is impossible.

3) Suppose finally  $\deg(x) = k - 2$ . Now  $|B| = k - 1$ , otherwise we cannot have  $n \geq 2k - 2$ . If there is  $b \in B$  such that  $|N(b) \cap N(x)| = k - 2$ , then because  $\deg(b) \leq k$  we have  $|N(b) \cap B| \geq 2$  and  $|B \setminus N[b]| \geq k - 4 \geq 2$ , when  $k \geq 6$ . Hence, there are  $c_1, c_2 \in B \setminus N[b]$ ,  $c_1 \neq c_2$ , and in the  $k$ -subset  $C = N(x) \cup \{c_1, c_2\}$  we have  $I(C; x) = I(C; b)$  which is impossible.

Thus, for all  $b \in B$  we have  $|N(b) \cap N(x)| \leq k - 3$ . On the other hand, each of the  $k - 2$  vertices in  $N(x)$  has at least  $k - 3$  adjacent vertices in  $B$ , so the vertices in  $B$  have on the average at least  $(k - 2)(k - 3)/(k - 1) > k - 4$  adjacent vertices in the set  $N(x)$ . Hence, we can find  $b \in B$  such that  $|N(b) \cap N(x)| = k - 3$ . Because  $\deg(b) \geq k - 2$  we have at least one  $b_0 \in B$  such that  $d(b, b_0) = 1$ . Because there are no triangles, each of the  $k - 3$  neighbours of  $b$  in  $N(x)$  is not adjacent with  $b_0$ , and therefore adjacent to at least  $k - 3$  of the  $k - 2$  vertices in  $B \setminus \{b_0\}$ . Hence, for all  $a_1, a_2 \in N(x) \cap N(b)$ ,  $a_1 \neq a_2$ , we have  $|N(a_1) \cap N(a_2) \cap B| \geq k - 4 \geq 2$  when  $k \geq 6$ . In the  $k$ -subset  $C = \{x, b, c_1, c_2\} \cup (N(x) \setminus \{a_1, a_2\})$ , where  $c_1, c_2 \in N(a_1) \cap N(a_2) \cap B$ , we have  $I(C; a_1) = I(C; a_2)$ , which is impossible.  $\square$

**Lemma 19** *If there is a graph  $G \in \mathfrak{Gr}(n, k)$  that contains a triangle, then  $n \leq 3k - 9$ . (In particular,  $k \geq 5$ .)*

**Proof:** Suppose that  $G = (V, E) \in \mathfrak{Gr}(n, k)$  and that there is a triangle  $\{x, y, z\}$  in  $G$ . Let, for  $v, w \in V$ ,  $J_w(v)$  denote the indicator function given by  $J_w(v) = 1$  if  $v \in N[w]$  and  $J_w(v) = 0$  if  $v \notin N[w]$ . Define the set  $M_{xy} = \{v \in V : J_x(v) = J_y(v)\}$ , and  $M'_{xy} = M_{xy} \setminus \{x, y, z\}$ . Since  $M_{xy}$  does not distinguish  $x$  and  $y$ , we have  $|M_{xy}| \leq k - 1$ . Further,  $\{x, y, z\} \subseteq M_{xy}$ , and thus  $|M'_{xy}| \leq k - 4$ . Define similarly  $M_{xz}, M_{yz}, M'_{xz}, M'_{yz}$ ; the same conclusion holds for these.

Since the indicator functions take only two values,  $M_{xy}, M_{xz}$  and  $M_{yz}$  cover  $V$ , and thus

$$n = |V| = |M'_{xy} \cup M'_{xz} \cup M'_{yz} \cup \{x, y, z\}| \leq 3(k - 4) + 3 = 3k - 9.$$

Since  $n \geq k$ , this entails  $3k - 9 \geq k$  and thus  $k \geq 5$ .  $\square$

The following upper bound is generally weaker than Theorem 17, but it gives the optimal result for  $k = 6$ .

**Theorem 20** *Suppose  $k \geq 6$ . Then  $\Xi(k) \leq 3k - 9$ .*

**Proof:** Suppose that  $G \in \mathfrak{Gr}(n, k)$ . If  $G$  does not contain any triangle, then Theorem 18 yields  $n \leq 2k - 3 \leq 3k - 9$ . If  $G$  does contain a triangle, then Lemma 19 yields  $n \leq 3k - 9$ .  $\square$

## 4 Strongly regular graphs

A graph  $G = (V, E)$  is called *strongly regular* with parameters  $(n, t, \lambda, \mu)$  if  $|V| = n$ ,  $\deg(x) = t$  for all  $x \in V$ , any two adjacent vertices have exactly  $\lambda$  common neighbours, and any two nonadjacent vertices have exactly  $\mu$  common neighbours; we then say that  $G$  is an  $(n, t, \lambda, \mu)$ -SRG. See [3] for more information. By [3, Proposition 1.4.1] we know that if  $G$  is an  $(n, t, \lambda, \mu)$ -SRG, then  $n = t + 1 + t(t - 1 - \lambda)/\mu$ .

We give two examples of strongly regular graphs that will be used below.

**Example 21** The well-known Paley graph  $P(q)$ , where  $q$  is a prime power with  $q \equiv 1 \pmod{4}$ , is a  $(q, (q-1)/2, (q-5)/4, (q-1)/4)$ -SRG, see for example [3]. The vertices of  $P(q)$  are the elements of the finite field  $F_q$ , with an edge  $ij$  if and only if  $i-j$  is a non-zero square in the field; when  $q$  is a prime, this means that the vertices are  $\{1, \dots, q\}$  with edges  $ij$  when  $i-j$  is a quadratic residue mod  $q$ .

**Example 22** Another construction of strongly regular graphs uses a regular symmetric Hadamard matrix with constant diagonal (RSHCD) [6], [4], [5]. In particular, in the case (denoted RSHCD+) of a regular symmetric  $n \times n$  Hadamard matrix  $H = (h_{ij})$  with diagonal entries  $+1$  and constant positive row sums  $2m$  (necessarily even when  $n > 1$ ), then  $n = (2m)^2 = 4m^2$  and the graph  $G$  with vertex set  $\{1, \dots, n\}$  and an edge  $ij$  (for  $i \neq j$ ) if and only if  $h_{ij} = +1$  is a  $(4m^2, 2m^2 + m - 1, m^2 + m - 2, m^2 + m)$ -SRG [4, §8D].

It is not known for which  $m$  such RSHCD+ exist (it has been conjectured that any  $m \geq 1$  is possible) but constructions for many  $m$  are known, see [6], [17, V.3] and [5, IV.24.2]. For example, starting with the  $4 \times 4$  RSHCD+

$$H_4 = \begin{pmatrix} 1 & 1 & 1 & -1 \\ 1 & 1 & -1 & 1 \\ 1 & -1 & 1 & 1 \\ -1 & 1 & 1 & 1 \end{pmatrix}$$

its tensor power  $H_4^{\otimes r}$  is an RSHCD+ with  $n = 4^r$ , and thus  $m = 2^{r-1}$ , for any  $r \geq 1$ . This yields a  $(2^{2r}, 2^{2r-1} + 2^{r-1} - 1, 2^{2r-2} + 2^{r-1} - 2, 2^{2r-2} + 2^{r-1})$ -SRG with vertex set  $\{1, 2, 3, 4\}^r$ , where two different vertices  $(i_1, \dots, i_r)$  and  $(j_1, \dots, j_r)$  are adjacent if and only if the number of coordinates  $\nu$  such that  $i_\nu + j_\nu = 5$  is even.

**Theorem 23** A strongly regular graph  $G = (V, E)$  with parameters  $(n, t, \lambda, \mu)$  belongs to  $\mathfrak{S}\mathfrak{r}(n, k)$  if and only if

$$k \geq \max\{n - t, n - 2t + 2\lambda + 3, n - 2t + 2\mu - 1\},$$

or, equivalently,  $t \geq n - k$  and  $2 \max\{\lambda + 1, \mu - 1\} \leq k + 2t - n - 1$ .

**Proof:** An immediate consequence of Theorem 7, since  $|N[x]| = t + 1$  for every vertex  $x$  and  $|N[x] \triangle N[y]|$  equals  $2(t - \lambda - 1)$  when  $x \sim y$  and  $2(t + 1 - \mu)$  when  $x \not\sim y, x \neq y$ .  $\square$

We can extend this construction to other values of  $n$  by modifying the strongly regular graph.

**Theorem 24** If there exists a strongly regular graph with parameters  $(n_0, t, \lambda, \mu)$ , then for every  $i = 0, \dots, n_0 + 1$  there exists a graph in  $\mathfrak{S}\mathfrak{r}(n_0 + i, k_0 + i)$ , where

$$k_0 = \max\{n_0 - t, t, n_0 - 2t + 2\lambda + 3, n_0 - 2t + 2\mu - 1, 2t - 2\lambda - 1, 2t - 2\mu + 2\},$$

provided  $k_0 \leq n_0$ .

**Proof:** For  $i = 0$ , this is a weaker form of Theorem 23. For  $i \geq 1$ , we suppose that  $G_0 = (V_0, E_0)$  is  $(n_0, t, \lambda, \mu)$ -SRG and build a graph  $G_i$  in  $\mathfrak{S}\mathfrak{r}(n_0 + i, k_0 + i)$  from  $G_0$  by adding suitable new vertices and edges.

If  $1 \leq i \leq n_0$ , choose  $i$  different vertices  $x_1, x_2, \dots, x_i$  in  $V_0$ . Construct a new graph  $G_i = (V_i, E_i)$  by taking  $G_0$  and adding to it new vertices  $x'_1, x'_2, \dots, x'_i$  and new edges  $x'_j y$  for  $j \leq i$  and all  $y \notin N_{G_0}(x_j)$ .

First,  $\deg_{G_i}(x) \geq \deg_{G_0}(x) = t$  for  $x \in V_0$  and  $\deg_{G_i}(x') = n_0 - t$  for  $x' \in V'_i = V_i \setminus V_0$ . We proceed to investigate  $N[x] \Delta N[y]$ , and separate several cases.

(i) If  $x, y \in V_0$ , with  $x \neq y$ , then

$$|N[x] \Delta N[y]| \geq |(N[x] \Delta N[y]) \cap V_0| = |N_{G_0}[x] \Delta N_{G_0}[y]|,$$

which equals  $2(t - \lambda - 1)$  if  $x \sim y$  and  $2(t - \mu + 1)$  if  $x \not\sim y$ .

(ii) If  $x \in V_0, y' \in V'_i$ , then, since  $\Delta$  is associative and commutative,

$$|(N[x] \Delta N[y']) \cap V_0| = |(N_{G_0}[x] \Delta (V_0 \Delta N_{G_0}(y)))| = n_0 - |N_{G_0}[x] \Delta N_{G_0}(y)|,$$

which equals  $n_0 - 1$  if  $x = y, n_0 - (2t - 2\lambda - 1)$  if  $x \sim y$ , and  $n_0 - (2t - 2\mu + 1)$  if  $x \not\sim y$  and  $x \neq y$ . If  $x \sim y$ , further,  $|(N[x] \Delta N[y']) \cap V'_i| \geq 1$ , since  $y' \notin N[x]$ .

(iii) If  $x', y' \in V'_i$ , with  $x' \neq y'$ , then

$$|(N[x'] \Delta N[y']) \cap V_0| = |(V_0 \setminus N_{G_0}(x)) \Delta (V_0 \setminus N_{G_0}(y))| = |N_{G_0}(x) \Delta N_{G_0}(y)|,$$

which equals  $2(t - \lambda)$  if  $x \sim y$  and  $2(t - \mu)$  if  $x \not\sim y$ . Further,  $|(N[x'] \Delta N[y']) \cap V'_i| = |\{x', y'\}| = 2$ .

Collecting these estimates, we see that  $G_i \in \mathfrak{Gr}(n_0 + i, k_0 + i)$  by Theorem 7 (or Corollary 8) with our choice of  $k_0$ . Note that  $2k_0 \geq (n_0 - 2t + 2\lambda + 3) + (2t - 2\lambda - 1) = n_0 + 2 \geq 3$ , so  $k_0 \geq 2$ .

Finally, for  $i = n_0 + 1$ , we construct  $G_{n_0+1}$  by adding a new vertex to  $G_{n_0}$  and connecting it to all other vertices. The graph  $G_{n_0}$  has by construction maximum degree  $\Delta_{G_{n_0}} = n_0 \leq k_0 + n_0 - 2$ . Hence, Lemma 10 shows that  $G_{n_0+1} \in \mathfrak{Gr}(2n_0 + 1, k_0 + n_0 + 1)$ .  $\square$

We specialize to the Paley graphs, and obtain from Example 21 and Theorems 23–24 the following.

**Theorem 25** *Let  $q$  be an odd prime power such that  $q \equiv 1 \pmod{4}$ .*

(i) *The Paley graph  $P(q) \in \mathfrak{Gr}(q, (q+3)/2)$ .*

(ii) *There exists a graph in  $\mathfrak{Gr}(q+i, (q+3)/2+i)$  for all  $i = 0, 1, \dots, q+1$ .*

Note that the rate  $2q/(q+3)$  for the Paley graphs approaches 2 as  $q \rightarrow \infty$ ; in fact, with  $n = q$  and  $k = (q+3)/2$  we have  $n = 2k - 3$ , almost attaining the bound  $2k - 2$  in Theorem 17. (The Paley graphs thus almost attain the bound in Theorem 17, but never attain it exactly.)

**Corollary 26**  $\Xi(k) \geq 2k - o(k)$  as  $k \rightarrow \infty$ .

**Proof:** Let  $q = p^2$  where (for  $k \geq 6$ )  $p$  is the largest prime such that  $p \leq \sqrt{2k - 3}$ . It follows from the prime number theorem that  $p/\sqrt{2k - 3} \rightarrow 1$  as  $k \rightarrow \infty$ , and thus  $q = 2k - o(k)$ . Hence, if  $k$  is large enough, then  $k \leq q \leq 2k - 3$ , and Theorem 25 shows that  $P(q) \in \mathfrak{Gr}(q, (q+3)/2) \subseteq \mathfrak{Gr}(q, k)$ , so  $\Xi(k) \geq q = 2k - o(k)$ . (Alternatively, we may let  $q$  be the largest prime such that  $q \leq 2k - 3$  and  $q \equiv 1 \pmod{4}$  and use the prime number theorem for arithmetic progressions [8, Chapter 17] to see that then  $q = 2k - o(k)$ .)  $\square$

We turn to the strongly regular graphs constructed in Example 22 and find from Theorem 23 that they are in  $\mathfrak{Gr}(4m^2, 2m^2 + 1)$ , thus attaining the bound in Theorem 17. We state that as a theorem.

**Theorem 27** *The strongly regular graph constructed in Example 22 from an  $n \times n$  RSHCD+ belongs to  $\mathfrak{Gr}(n, n/2 + 1)$ .*

**Corollary 28** *There exist infinitely many integers  $k$  such that  $\Xi(k) = 2k - 2$ .*

**Proof:** If  $k = n/2 + 1$  for an even  $n$  such that there exists an  $n \times n$  RSHCD+, then  $\Xi(k) \geq n = 2k - 2$  by Theorem 27. The opposite inequality is given by Theorem 17. By Example 22, this holds at least for  $k = 2^{2r-1} + 1$  for any  $r \geq 1$ .  $\square$

## 5 On $\mathfrak{Gr}(n, k, \ell)$

In this section we consider  $\mathfrak{Gr}(n, k, \ell)$  for  $\ell \geq 2$ . Let us denote

$$\Xi(k, \ell) = \max\{n : \mathfrak{Gr}(n, k, \ell) \neq \emptyset\}.$$

Trivially, the empty graph  $E_k \in \mathfrak{Gr}(k, k, \ell)$  for any  $\ell \geq 1$ ; thus  $\Xi(k, \ell) \geq k$ .

Note that a graph  $G = (V, E)$  with  $|V| = n$  admits a  $(1, \leq \ell)$ -identifying set  $\iff V$  is  $(1, \leq \ell)$ -identifying  $\iff G \in \mathfrak{Gr}(n, n, \ell)$ .

**Theorem 29** *Suppose that  $G = (V, E) \in \mathfrak{Gr}(n, k, \ell)$ , where  $n > k$  and  $\ell \geq 2$ . Then the following conditions hold:*

- (i) *For all  $x \in V$  we have  $\ell + 1 < n - k + \ell + 1 \leq |N[x]| \leq k - \ell$ . In other words,  $\delta_G \geq n - k + \ell$  and  $\Delta_G \leq k - \ell - 1$ .*
- (ii) *For all  $x, y \in V, x \neq y, |N[x] \cap N[y]| \leq k - 2\ell + 1$ .*
- (iii)  *$n \leq 2k - 2\ell - 1$  and  $k \geq 2\ell + 2$ .*

**Proof:** (i) Suppose first that there is a vertex  $x \in V$  such that  $|N[x]| \leq n - k + \ell$ . By removing  $n - k$  vertices from  $V$ , starting in  $N[x]$ , we find a  $k$ -subset  $C$  with  $I(C; x) = \{c_1, \dots, c_m\}$  for some  $m \leq \ell$ . If  $m = 0$ , then  $I(C; x) = I(C; \emptyset)$ , which is impossible. If  $1 \leq m < \ell$ , we can arrange (by removing  $x$  first) so that  $x \notin C$ , and thus  $x \notin Y = \{c_1, \dots, c_m\}$ . Then  $I(C; \{x\} \cup Y) = I(C; Y)$ , a contradiction. If  $m = \ell \geq 2$ , we can conversely arrange so that  $x \in C$ , and thus  $x \in I(C; x)$ , say  $c_1 = x$ . Then  $I(C; c_2, \dots, c_m) = I(C; c_1, \dots, c_m)$ , another contradiction. Consequently,  $|N[x]| \geq n - k + \ell + 1$ .

Suppose then  $|N[x]| \geq k - \ell + 1$ . If  $|N[x]| \geq k$ , we can choose a  $k$ -subset  $C$  of  $N[x]$ ; then  $I(C; x) = C = I(C; x, y)$  for any  $y$ , which is impossible. If  $k > |N[x]| \geq k - \ell + 1$ , we can choose a  $k$ -subset  $C = N[x] \cup \{c_1, \dots, c_{k-|N[x]|}\}$ . Choose also  $a \in N(c_1)$  (which is possible because  $\deg(c_1) \geq 1$  by (i)). Now  $I(C; x, c_1, \dots, c_{k-|N[x]|}) = C = I(C; x, a, c_2, \dots, c_{k-|N[x]|})$ , which is impossible.

(ii) Suppose to the contrary that there are  $x, y \in V, x \neq y$ , such that  $|N[x] \cap N[y]| \geq k - 2\ell + 2$ . Let  $A = N(y) \setminus N[x]$ . Then, according to (i),  $|A| \leq |N[y] \setminus N[x]| = |N[y]| - |N[x] \cap N[y]| \leq k - \ell - (k - 2\ell + 2) = \ell - 2$ . Since  $k > \ell - 2$  by (i), there is a  $k$ -subset  $C \subseteq V \setminus \{y\}$  such that  $A \subset C$ . Then  $I(C; A \cup \{x, y\}) = I(C; A \cup \{x\})$ , a contradiction.

(iii) An immediate consequence of (i), which implies  $n - k + \ell + 1 \leq k - \ell$  and  $\ell + 1 < k - \ell$ .  $\square$

**Theorem 30** *For  $\ell \geq 2, \Xi(k, \ell) \leq \max\{\frac{\ell}{\ell-1}(k-2), k\}$ .*

**Proof:** If  $\Xi(k, \ell) = k$ , there is nothing to prove. Assume then that there exists a graph  $G = (V, E) \in \mathfrak{G}\mathfrak{r}(n, k, \ell)$ , where  $n > k$ . By Theorem 29(iii),  $\ell < k/2 < n$ . Let us consider any set of vertices  $Z = \{z_1, z_2, \dots, z_\ell\}$  of size  $\ell$ . We will estimate  $|N[Z]|$  as follows. By Theorem 29(i) we know  $|N[z_1]| \geq n - k + \ell + 1$ . Now  $N[z_1, z_2]$  must contain at least  $n - k + 1$  vertices, which *do not* belong to  $N[z_1]$  due to Theorem 7 which says that  $|N[X] \triangle N[Y]| \geq n - k + 1$ , where we take  $X = \{z_1\}$  and  $Y = \{z_1, z_2\}$ . Analogously, each set  $N[z_1, \dots, z_i]$  ( $i = 2, \dots, \ell$ ) must contain at least  $n - k + 1$  vertices which are not in  $N[z_1, \dots, z_{i-1}]$ . Hence, for the set  $Z$  we have  $|N[Z]| \geq n - k + \ell + 1 + (\ell - 1)(n - k + 1) = \ell(n - k + 2)$ . Since trivially  $|N[Z]| \leq n$ , we have  $(\ell - 1)n \leq \ell(k - 2)$ , and the claim follows.  $\square$

**Corollary 31** For  $\ell \geq 2$ , we have  $\frac{\Xi(k, \ell)}{k} \leq 1 + \frac{1}{\ell - 1}$ .

The next results improve the result of Theorem 30 for  $\ell = 2$ .

**Lemma 32** Assume that  $n > k$ . Let  $G = (V, E)$  belong to  $\mathfrak{G}\mathfrak{r}(n, k, 2)$ . Then

$$n + \frac{n - k + 2}{n - 1}(n - k + 3) \leq 2k - 3$$

**Proof:** Suppose  $x \in V$ . Let

$$f(n, k) = \frac{n - k + 2}{n - 1}(n - k + 3).$$

Our aim is first to show that there exists a vertex in  $N(x)$  or in  $S_2(x)$  which dominates at least  $f(n, k)$  vertices of  $N[x]$ . Let

$$\lambda_x = \max\{|N[x] \cap N[a]| \mid a \in N(x)\}.$$

If  $\lambda_x \geq f(n, k)$ , we are already done. But if  $\lambda_x < f(n, k)$ , then we show that there is a vertex in  $S_2(x)$  that dominates at least  $f(n, k)$  vertices of  $N[x]$ . Let us estimate the number of edges between the vertices in  $N(x)$  and in  $S_2(x)$  — we denote this number by  $M$ . By Theorem 29(i), every vertex  $y \in N(x)$  yields at least  $|N[y]| - \lambda_x \geq n - k + 3 - \lambda_x$  such edges and there are at least  $n - k + 2$  vertices in  $N(x)$ . Consequently,  $M \geq (n - k + 2)(n - k + 3 - \lambda_x)$ . On the other hand, again by Theorem 29(i),  $|S_2(x)| \leq n - |N[x]| \leq k - 3$ . Hence, there must exist a vertex in  $S_2(x)$  incident with at least  $M/(k - 3)$  edges whose other endpoint is in  $N(x)$ . Now, if  $\lambda_x < f(n, k)$ , then

$$\frac{M}{k - 3} > \frac{(n - k + 2)(n - k + 3 - f(n, k))}{k - 3} = f(n, k).$$

Hence there exists in this case a vertex in  $S_2(x)$  that is incident to at least  $f(n, k)$  such edges, i.e., it dominates at least  $f(n, k)$  vertices in  $N(x)$ .

In any case there thus exists  $z \neq x$  such that  $|N[x] \cap N[z]| \geq f(n, k)$ . Let  $C = (N[x] \cap N[z]) \cup (V \setminus N[x])$ . Then  $I(C; x, z) = I(C; z)$ , so  $C$  is not  $(1, \leq 2)$ -identifying and thus  $|C| < k$ . Hence, using Theorem 29(i),

$$k - 1 \geq |C| \geq f(n, k) + n - |N[x]| \geq f(n, k) + n - (k - 2),$$

and thus  $n + f(n, k) \leq 2k - 3$  as asserted.  $\square$

**Theorem 33** *If  $k \leq 5$ , then  $\Xi(k, 2) = k$ . If  $k \geq 6$ , then*

$$\Xi(k, 2) < \left(1 + \frac{1}{\sqrt{2}}\right)(k - 2) + \frac{1}{4}.$$

**Proof:** Let  $n = \Xi(k, 2)$ , and let  $m = k - 2$ . If  $n > k$ , then  $k \geq 6$  by Theorem 29(iii); hence  $n = k$  when  $k \leq 5$ . Further, still assuming  $n > k$ , Lemma 32 yields

$$n + \frac{(n - m)(n - m + 1)}{n - 1} \leq 2m + 1$$

or

$$0 \geq n(n - 1) + (n - m)^2 + n - m - (2m + 1)(n - 1) = 2\left(n - \left(m + \frac{1}{4}\right)\right)^2 - m^2 + \frac{7}{8}.$$

Hence,  $n - \left(m + \frac{1}{4}\right) < m/\sqrt{2}$ . □

**Corollary 34** *For  $\ell = 2$ , we have  $\Xi(k, 2)/k \leq 1 + \frac{1}{\sqrt{2}}$ .*

**Problem 35** *What is  $\limsup_{k \rightarrow \infty} \Xi(k, \ell)/k$  for  $\ell \geq 2$ ? In particular, is  $\limsup_{k \rightarrow \infty} \Xi(k, \ell)/k > 1$ ?*

The following theorem implies that for any  $\ell \geq 2$  there exist graphs in  $\mathfrak{Gr}(n, k, \ell)$  for  $n \approx k + \log_2 k$ . In particular, we have such graphs with  $n > k$ .

**Theorem 36** *Let  $\ell \geq 2$  and  $m \geq \max\{2\ell - 2, 4\}$ . A binary hypercube of dimension  $m$  belongs to  $\mathfrak{Gr}(2^m, 2^m - m + 2\ell - 2, \ell)$ .*

**Proof:** Suppose first  $\ell \geq 3$ . By [11, Theorem 2] we know that then a set in a binary hypercube is  $(1, \leq \ell)$ -identifying if and only if every vertex is dominated by at least  $2\ell - 1$  different vertices belonging to the set. Hence, we can remove any  $m + 1 - (2\ell - 1)$  vertices from the set of vertices, and there will still be a big enough multiple domination to assure that the remaining set is  $(1, \leq \ell)$ -identifying.

Suppose then that  $\ell = 2$  and  $G = (V, E)$  is the binary  $m$ -dimensional hypercube. Let us denote by  $C \subseteq V$  a  $(2^m - m + 2)$ -subset. Every vertex is dominated by at least  $m + 1 - (m - 2) = 3$  vertices of  $C$ . For all  $x, y \in V$ ,  $x \neq y$  we have  $|N[x] \cap N[y]| = 2$  if and only if  $1 \leq d(x, y) \leq 2$  and otherwise  $|N[x] \cap N[y]| = 0$ . Hence, for all  $x, y, z \in V$  with  $x \neq y$ ,  $I(y) = N[y] \cap C$  contains at least 3 vertices, and these cannot all be dominated by  $x$ ; thus, we have  $I(x) \neq I(y)$  and  $I(x) \neq I(y, z)$ .

We still need to show that  $I(x, y) \neq I(z, w)$  for all  $x, y, z, w \in V$ ,  $x \neq y$ ,  $z \neq w$ ,  $\{x, y\} \neq \{z, w\}$ . By symmetry we may assume that  $x \notin \{z, w\}$ . Suppose  $I(x, y) = I(z, w)$ .

If  $|I(x)| \geq 5$ , then any two vertices  $z, w \neq x$  cannot dominate  $I(x)$ , a contradiction.

If  $|I(x)| = 4$ , then  $|I(z) \cap I(x)| = |I(w) \cap I(x)| = 2$  and  $I(x) \cap I(z) \cap I(w) = \emptyset$ . It follows that  $3 \leq d(z, w) \leq 4$  which implies  $I(z) \cap I(w) = \emptyset$ . Since  $|N[x] \setminus C| = |N[x]| - |I(x)| = m - 3$ , all except one vertex, say  $v$ , of  $V \setminus C$  belong to  $N[x]$ , so  $V \setminus N[x] \subseteq C \cup \{v\}$ ; the vertex  $v$  cannot belong to both  $N[z]$  and  $N[w]$  since these are disjoint, so we may (w.l.o.g.) assume that  $v \notin N[z]$ , and thus  $N[z] \setminus N[x] \subseteq C$ , whence  $N[z] \setminus N[x] \subseteq I(z) \setminus I(x)$ . Hence,  $|I(z) \cap I(y)| \geq |I(z) \setminus I(x)| \geq |N[z] \setminus N[x]| = |N[z]| - |N[z] \cap N[x]| = m + 1 - 2 \geq 3$ . Thus  $y = z$ ; however, then  $I(y) \cap I(w) = I(z) \cap I(w) = \emptyset$  and since  $I(w) \not\subseteq I(x)$ , we have  $I(w) \not\subseteq I(x, y)$ .

Suppose finally that  $|I(x)| = 3$ ; w.l.o.g. we may assume  $|I(z) \cap I(x)| = 2$ . Now  $|N[x] \setminus C| = |N[x]| - |I(x)| = m - 2 = |V \setminus C|$ , and thus  $V \setminus C = N[x] \setminus C \subseteq N[x]$ ; hence,  $V \setminus N[x] \subseteq C$  and thus

$N[z] \setminus N[x] \subseteq I(z) \setminus I(x)$ . Consequently,  $|I(z) \cap I(y)| \geq |I(z) \setminus I(x)| \geq |N[z] \setminus N[x]| \geq m+1-2 \geq 3$ , and thus  $z = y$ . But similarly  $N[w] \setminus N[x] \subseteq I(w) \setminus I(x)$  and the same argument shows  $w = y$ , and thus  $w = z$ , a contradiction.  $\square$

We finally consider graphs without isolated vertices (i.e., no vertices with degree zero), and in particular connected graphs.

By [13, Theorem 8] a graph with no isolated vertices admitting a  $(1, \leq \ell)$ -identifying set has minimum degree at least  $\ell$ . Hence, always  $n \geq \ell + 1$ .

In [7] and [12] it has been proven that there exist connected graphs which admit  $(1, \leq \ell)$ -identifying sets. For example, the smallest known connected graph admitting a  $(1, \leq 3)$ -identifying set has 16 vertices [12]. It is unknown whether there are such graphs with smaller order. In the next theorem we solve the case of graphs admitting  $(1, \leq 2)$ -identifying sets.

**Theorem 37** *The smallest  $n \geq 3$  such that there exists a connected graph (or a graph without isolated vertices) in  $\mathfrak{Gr}(n, n, 2)$  is  $n = 7$ .*

(If we allow isolated vertices, we can trivially take the empty graph  $E_n$  for any  $n \geq 2$ .)

**Proof:** The cycle  $C_n \in \mathfrak{Gr}(n, n, 2)$  for  $n \geq 7$  by Example 3(ii) (see also [12]).

Assume that  $G = (V, E) \in \mathfrak{Gr}(n, n, 2)$  is a graph of order  $n \leq 6$  without isolated vertices; we will show that this leads to a contradiction. By [13], we know that  $\deg(v) \geq 2$  for all  $v \in V$ . We will use this fact frequently in the sequel.

If  $G$  is disconnected, the only possibility is that  $n = 6$  and that  $G$  consists of two disjoint triangles, but this graph is not even in  $\mathfrak{Gr}(n, n, 1)$ .

Hence,  $G$  is connected. Let  $x, y \in V$  be such that  $d(x, y) = \text{diam}(G)$ .

(i) Suppose that  $\text{diam}(G) = 1$ , or more generally that there exists a dominating vertex  $x$ . Then  $N[x, y] = N[x]$  for any  $y \in V$ , which is a contradiction.

(ii) Suppose next  $\text{diam}(G) = 2$ . Moreover, by the previous case we can assume that for any  $v \in V$  there is  $w \in V$  such that  $d(v, w) = 2$ .

Assume first  $|N(x)| = 4$ . Then  $S_2(x) = \{y\}$ . Since  $\deg(y) \geq 2$ , there exist two vertices  $w_1, w_2 \in N(y) \cap N(x)$ , but then  $N[x, w_1] = N[x, w_2]$ .

Assume next  $|N(x)| = 3$ , say  $N(x) = \{u_1, u_2, u_3\}$ . Then  $|S_2(x)| = n - |N[x]| \leq 2$ . Since the four sets  $N[x]$  and  $N[x, u_i]$ ,  $i = 1, 2, 3$ , must be distinct, we can assume without loss of generality that  $|S_2(x)| = 2$ , say  $S_2(x) = \{y, w\}$ , and that the only edges between the elements in  $S_2(x)$  and  $N(x)$  are  $u_1y, u_2w, u_3y$  and  $u_3w$ . Then  $N[x, u_3] = N[y, u_2]$ .

Assume finally that  $|N(x)| = 2$ . By the previous discussion we may assume that  $|N(v)| = 2$  for all  $v \in V$ . Then  $G$  must be a cycle  $C_n$ , but it can easily be seen that  $C_n \notin \mathfrak{Gr}(n, n, 2)$  for  $3 \leq n \leq 6$ .

(iii) Suppose that  $\text{diam}(G) = 3$ . Clearly  $|N(x)| \geq 2$  and  $|S_2(x)| \geq 1$ . If  $|S_2(x)| = 1$ , say  $S_2(x) = \{w\}$ , then  $N[w, y] = N[w]$ , which is not allowed. Since  $n \leq 6$ , we thus have  $|N(x)| = 2$  and  $|S_2(x)| = 2$ , say  $N(x) = \{u_1, u_2\}$  and  $S_2(x) = \{w_1, w_2\}$ . We can assume without loss of generality that  $u_1w_1 \in E$ . If  $w_2u_2 \in E$ , then  $N[w_1, u_2] = N[x, y]$ . If  $w_2u_2 \notin E$ , then  $N[w_1, w_2] = N[w_1]$ .

(iv) Suppose that  $\text{diam}(x, y) \geq 4$ . Then  $G$  contains an induced path  $P_5$ . There is at most one additional vertex, but it is impossible to add it to  $P_5$  and obtain  $\delta_G \geq 2$  and  $\text{diam}(G) \geq 4$ .

This completes the proof.  $\square$

## Acknowledgements

Part of this research was done during the Workshop on Codes and Discrete Probability in Grenoble, France. We would like to thank the referees for their comments.

**Tab. 1:** Lower and upper bounds for  $\Xi(k)$  for  $1 \leq k \leq 20$ . The lower bounds come from the examples given in the last column; for  $n \geq 8$  using Theorem 23, 25 or 27 or Lemma 10. The strongly regular graphs used here can be found from [5]. The upper bounds for  $k \geq 7$  come from Theorem 17.

k	lower bound	upper bound	example
1	1	1 (Ex. 12)	$E_1$
2	2	2 (Ex. 13)	$E_2$
3	4	4 (Ex. 14, Th.17)	$C_4, S_4$
4	5	5 (Th. 16)	Figure 1(b)
5	8	8 (Th. 17)	Example 9
6	9	9 (Th. 20)	Example 11, $P(9)$
7	11	12 (Th. 17, Th. 20)	Figure 1(a)
8	13	14	$P(13)$
9	16	16	RSHCD+
10	17	18	$P(17)$
11	18	20	Th. 25(ii)
12	21	22	(21,10,3,6)-SRG
13	22	24	Lemma 10
14	25	26	$P(25)$
15	26	28	(26,15,8,9)-SRG
16	29	30	$P(29)$
17	30	32	Th. 25(ii)
18	31	34	Th. 25(ii)
19	36	36	RSHCD+
20	37	38	$P(37)$



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