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When are Swing options bang-bang and how to use it?

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Abstract

In this paper we investigate a class of swing options with firm constraints in view of the modeling of supply agreements. We show, for a fully general payoff process, that the premium, solution to a stochastic control problem, is concave and piecewise affine as a function of the global constraints of the contract. The existence of bang-bang optimal controls is established for a set of constraints which generates by affinity the whole premium function.

When the payoff process is driven by an underlying Markov process, we propose a quantization based recursive backward procedure to price these contracts. A priori error bounds are established, uniformly with respect to the global constraints.

Key words: Swing option, stochastic control, optimal quantization, energy.

1 Introduction

The deregulation of energy markets has given rise to various families of contracts. Many of them appear as some derivative products whose underlying is some tradable futures (day-ahead, etc) on gas or electricity (see [12] for an introduction). The class of swing options has been paid a special attention in the literature, because it includes many of these derivative products. A common feature to all these options is that they introduce some risk sharing between a producer and a trader, of gas or electricity for example. From a probabilistic viewpoint, they appear as some stochastic control problems modeling multiple optimal stopping problems (the control variable is the purchased quantity of energy); see *e.g.* [10, 9] in a continuous time setting. Gas storage contracts (see [6], [8]) or electricity supply agreements (see [18], [7]) are examples of such swing options. Indeed, energy supply contracts are one simple and important example of such swing options that will be deeply investigated in this paper (see below, see also [12] for an introduction). It is worth mentioning that this kind of contracts are slightly different from multiple exercises American options as considered in [10] for example. In our setting the volumetric constraints play a key role and thus, the flexibility is not restricted to time decisions, but also has to take into account volumes management.

Designing efficient numerical procedures for the pricing of swing option contracts remains a very challenging question as can be expected from a possibly multi-dimensional stochastic control problem subject to various constraints (due to the physical properties of the assets like in storage

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contracts). Most recent approaches developed in mathematical finance, especially for the pricing of American options, have been adapted and transposed to the swing framework: tree (or “forest”) algorithms in the pioneering work [17], Least squares regression MC methods (see [6]), PDE’s numerical methods (finite elements, see [29]).

The aim of this paper is to deeply investigate an old question, namely to elucidate the structure of the optimal control in supply contracts (with firm constraints) and how it impacts the numerical methods of pricing. We will provide in a quite general (and abstract) setting some “natural” (and simple) conditions involving the local and global purchased volume constraints to ensure the existence of *bang-bang optimal strategy* (such controls usually do not exist). It is possible to design *a priori* the contract so that their parameters satisfy these conditions. To our knowledge very few theoretical results have been established so far on this problem (see however [6] in a Markovian framework for contracts with penalized constraints and [27], also in a Markovian framework).

This first result of the paper not only enlightens the understanding of the management of a swing contract: it also has some deep repercussions on the numerical methods to price it. As a matter of fact, taking advantage of the existence of a bang-bang optimal strategy, we propose and analyze in details (when the underlying asset has a Markovian dynamics) a quantized Dynamic Programming procedure to price any swing options whose volume constraints satisfy the “bang-bang” assumption. Furthermore some *a priori* error bounds are established. This procedure turns out to be dramatically efficient, as emphasized in the companion paper [5] where the method is extensively tested with assets having multi-factor Gaussian underlying dynamics and compared to the least squares regression method.

The abstract swing contract with firm constraints The holder of a supply contract has the right to purchase periodically (daily, monthly, etc) an amount of energy at a given unitary price. This amount of energy is subject to some lower and upper “local” constraints. The total amount of energy purchased at the end of the contract is also subject to a “global” constraint. Given dynamics on the energy price process, the problem is to evaluate the price of such a contract, at time $t = 0$ when it is emitted and during its whole life up to its maturity.

To be precise, the owner of the contract is allowed to purchase at times t_i , $i = 0, \dots, n - 1$ a quantity q_i of energy at a unitary *strike price* $K_i := K(t_i)$. At every date t_i , the purchased quantity q_i is subject to the firm “local” constraint,

$$q_{\min} \leq q_i \leq q_{\max}, \quad i = 0, \dots, n - 1,$$

whereas the global purchased quantity $\bar{q}_n := \sum_{i=0}^{n-1} q_i$ is subject to the (firm) global constraint

$$\bar{q}_n \in [Q_{\min}, Q_{\max}] \quad (0 < Q_{\min} \leq Q_{\max} < +\infty).$$

The strike price process $(K_i)_{0 \leq i \leq n-1}$ can be either deterministic (even constant) or stochastic, *e.g.* indexed on past values of other commodities (oil, etc). Usually, on energy markets the price is known through future contracts $(F_{s,t})_{0 \leq s \leq t}$ where $F_{s,t}$ denotes the price at time s of the forward contract delivered at maturity t . The available data at time 0 are $(F_{0,t})_{0 \leq t \leq T}$ (in real markets this is of course not a continuum).

The underlying asset price process, temporarily denoted $(S_{t_i})_{0 \leq i \leq n-1}$, is often the so-called “day-ahead” contract $F_{t,t+1}$ which is a tradable instrument or the spot price $F_{t,t}$ which is not. All the decisions about the contract need to be adapted to the filtration of (S_{t_i}) *i.e.* $\mathcal{F}_i := \sigma(S_{t_j}, j = 0, \dots, i)$, $i = 0, \dots, n - 1$ (with $\mathcal{F}_0 = \{\emptyset, \Omega\}$). This means that the price of such a contract is given

at any time t_k , by

$$P_k^n(Q_{\min}^k, Q_{\max}^k) := \operatorname{esssup} \left\{ \mathbb{E} \left(\sum_{j=k}^{n-1} q_j e^{-r(t_j - t_k)} (S_{t_j} - K_j) \mid \mathcal{F}_k \right), \right. \\ \left. q_j : (\Omega, \mathcal{F}_j) \rightarrow [q_{\min}, q_{\max}], j = k, \dots, n-1, \sum_{j=k}^{n-1} q_j \in [Q_{\min}^k, Q_{\max}^k] \right\}$$

where $Q_{\min}^k = Q_{\min} - \bar{q}_k$, $Q_{\max}^k = Q_{\max} - \bar{q}_k$ denote the *residual global constraints* and r denotes the (deterministic) interest rate. This pricing problem clearly appears as a stochastic control problem.

In the pioneering work by [17], this type of contract was computed by using some forests of (multinomial) trees. A natural variant, at least for numerical purpose, is to consider a penalized version of this stochastic control. Thus, in [6], a penalization $Q_n(V_n, \bar{q}_n)$ with $Q_n(x, q) = -((x - Q_{\max})^+ + (x - Q_{\min})^+)/\varepsilon$ is added (Q_n is negative outside $[Q_{\min}, Q_{\max}]$ and zero inside).

As concerns more sophisticated contracts (like storages), the holder of the contract receive a quantity $\Psi(t_i, S_{t_i}, q_i)$ when deciding q_i . When dealing with gas this is due to the storing constraints since injecting or withdrawing gas from its storing units induce fixed costs (and physical constraints (pressure, etc)).

As concerns the underlying asset dynamics, it is commonly shared in finance to assume that the traded asset has a Markovian dynamics (or is a component of a Markov process like with stochastic volatility models). The dynamics of physical assets for many reasons (some of them simply coming from history) are often modeled using some more deeply non-Markovian models like long memory processes, etc.

All these specific features of energy derivatives suggest to tackle the above pricing problem in a rather general framework, trying to avoid as long as possible to call upon Markov properties. This is what we do in the first part of the paper where the general setting of a swing option defined by an abstract sequence of \mathcal{F}_k -adapted payoffs is deeply investigated as a function of its global constraints (Q_{\min}, Q_{\max}) (when the local constraints are normalized *i.e.* q_i is $[0, 1]$ -valued for every $i \in \{0, \dots, n-1\}$). We show that this premium is a concave, piecewise affine, function of the global constraints, affine on triangles of the $(m, M) + \{(u, v), 0 \leq u \leq v \leq 1\}$, $m, M \in \mathbb{N}^2$, $m \leq M \leq n$ and $(m, M) + \{(u, v), 0 \leq v \leq u \leq 1\}$, $m, M \in \mathbb{N}^2$, $m \leq M - 1 \leq n - 1$. We also show that for integral valued global constraints, the optimal controls are always bang-bang *i.e.* the *a priori* $[0, 1]$ -valued optimal purchased quantities q_i^* are in fact always equal to 0 or 1. Such a result can be extended in some way to any couple of global constraints when all the payoffs are nonnegative.

Then, when there is an underlying Markov “structure process”, we propose an optimal quantization based on numerical approach to price efficiently swing options. This Markov “structure process” can be the underlying traded asset itself or a higher dimensional hidden Markov process: such a framework comes out in case of multi-factor processes having some long-memory properties.

Optimal Quantization was first introduced as a numerical method to solve nonlinear problem arising in Mathematical Finance in a series of papers [1, 2, 3, 4] devoted to the pricing and hedging of American style multi-asset options. It has also been applied to stochastic control problem, namely portfolio optimization in [24]. The purely numerical aspects induced by optimal quantization, with a special emphasis on the Gaussian distribution, have been investigated in [26]. See [25] for a survey on numerical application of optimal quantization to Finance. For other applications (to automatic classification, clustering, etc), see also [14]. In this paper, we propose a quantized backward dynamic programming to approximate the premium of a swing contract. We analyze the rate of convergence of this algorithm and provide some *a priori* error bounds in terms of quantization errors.

We illustrate the method by computing the whole graph of the premium viewed as a function of the global constraints, combining the affine property of the premium and the quantized algorithm

in “toy model”: the future prices of gas are modeled by a two factor Gaussian model. An extensive study of the pricing method by optimal quantization is carried out from both a theoretical and numerical point of view in [5].

The paper is organized as follows. In the section below we detail the decomposition of swing options into a swap contract and a normalized swing option. In Section 3, we precisely describe our abstract setting for normalized swing options with firm constraints and the variable of interest (global constraints, local constraints, etc). In Section 4, we establish the dynamic programming formula satisfied in full generality by the premium as a function of the global constraints (this unifies the similar results obtained in Markov settings, see [17], [6], etc) and we show this is a concave function with respect to the global constraints. Then, in Section 5, we prove in our abstract framework that the premium function is piecewise affine and that the optimal purchased quantities satisfy a “0-1” or bang-bang principle (Theorem 2). A special attention is paid to the 2-period model which provides an intuitive interpretation of the results. In Section 6, after some short background on quantization and its optimization, we propose a quantization based backward dynamic programming formula as a numerical method to solve the swing pricing problem. Then we provide some error bounds for the procedure depending on the quantization error induced by the quantization of the Markov structure process.

NOTATIONS. • The Lipschitz coefficient of a function $f : \mathbb{R}^d \rightarrow \mathbb{R}$ is defined by $[f]_{\text{Lip}} := \sup_{x \neq y} \frac{|f(x) - f(y)|}{|x - y|} \leq +\infty$. The coefficient $[f]_{\text{Lip}}$ is finite if and only if f Lipschitz continuous.

- The canonical Euclidean norm on \mathbb{R}^d will be denoted $|\cdot|$.

2 Canonical decomposition, normalized swing option

As a first step we need to normalize this contract to reduce some useless technical aspects. In practice this normalization, in fact decomposition, corresponds to the splitting of the contract into a swap and a normalized swing. The decomposition can be derived from the fact that an \mathcal{A} -measurable random variable q is $[q_{\min}, q_{\max}]$ -valued if and only if there exists a $[0, 1]$ -valued \mathcal{A} -measurable random variable q' such that $q = q_{\min} + (q_{\max} - q_{\min})q'$. Then, for every $k \in \{0, \dots, n-1\}$,

$$P_k^n(Q_{\min}^k, Q_{\max}^k) = \underbrace{q_{\min} \sum_{j=k}^{n-1} e^{-r(t_j - t_k)} \mathbb{E}(S_{t_j} - K_j | \mathcal{F}_k)}_{\text{swap contract}} + (q_{\max} - q_{\min}) \underbrace{P_k^{[0,1],n}(\tilde{Q}_{\min}^k, \tilde{Q}_{\max}^k)}_{\text{normalized contract}}$$

where

$$\tilde{Q}_{\min}^k = \frac{Q_{\min}^k - (n-k)q_{\min}}{q_{\max} - q_{\min}} \quad \text{and} \quad \tilde{Q}_{\max}^k = \frac{Q_{\max}^k - (n-k)q_{\min}}{q_{\max} - q_{\min}}$$

and $P_k^{[0,1],n}(\tilde{Q}_{\min}^k, \tilde{Q}_{\max}^k)$ is a *normalized swing contract* in which the local constraints are $[0, 1]$ -valued.

3 An abstract model for swing options with firm constraints

One considers a sequence $(V_k)_{0 \leq k \leq n-1}$ of integrable random variables defined on a probability space $(\Omega, \mathcal{A}, \mathbb{P})$. Let $\mathcal{F}_k^V := \sigma(V_0, V_1, \dots, V_k)$, $k = 0, \dots, n-1$ denote its natural filtration. For convenience we introduce a more general discrete time filtration $\mathcal{F} := (\mathcal{F}_k)_{0 \leq k \leq n-1}$ to which $V =$

$(V_k)_{0 \leq k \leq n-1}$ is adapted *i.e.* satisfying $\mathcal{F}_k^V \subset \mathcal{F}_k$, $k = 0, \dots, n-1$ *i.e.* such that the sequence $(V_k)_{0 \leq k \leq n-1}$ is \mathcal{F} -adapted.

We aim to solve the following abstract stochastic control problem (with *maturity* n)

$$(\mathcal{S})_n^{\mathcal{F}} \equiv \text{esssup} \left\{ \mathbb{E} \left(\sum_{k=0}^{n-1} q_k V_k \mid \mathcal{F}_0 \right), q_k : (\Omega, \mathcal{F}_k) \rightarrow [0, 1], 0 \leq k \leq \dots, n-1, \sum_{k=0}^{n-1} q_k \in [Q_{\min}, Q_{\max}] \right\} \quad (3.1)$$

where Q_{\min} and Q_{\max} are two non-negative \mathcal{F}_0 -measurable random variables satisfying

$$0 \leq Q_{\min} \leq Q_{\max} \leq n. \quad (3.2)$$

(The inequality $Q_{\max} \leq n$ induces no loss of generality: one can always replace Q_{\max} by $Q_{\max} \wedge n$ in (3.1) since in any case $q_0 + \dots + q_n \leq n$).

Note that no assumption is made on the dynamics of the state process $V = (V_k)_{0 \leq k \leq n-1}$.

We need to introduce the following notations and terminology:

– Throughout the paper, esssup will always be taken with respect to the probability \mathbb{P} so \mathbb{P} will be dropped from now on.

– A couple $Q := (Q_{\min}, Q_{\max})$ of non-negative \mathcal{F}_k -measurable random variables satisfying $0 \leq Q_{\min} \leq Q_{\max} \leq n-k$ is called a couple of *global constraints* at time k .

– An \mathcal{F} -adapted sequence $q = (q_k)_{0 \leq k \leq n-1}$ of $[0, 1]$ -valued r.v. is called a *locally admissible control*. For any locally admissible control, one defines the cumulative purchase process by

$$\bar{q}_0 := 0, \quad \bar{q}_k := q_0 + \dots + q_{k-1}, \quad k = 1, \dots, n.$$

If $\bar{q}_n \in [Q_{\min}, Q_{\max}]$, q is called an (\mathcal{F}, Q) -*admissible control*.

– For every $k \in \{0, \dots, n\}$ and every couple of global constraints $Q := (Q_{\min}, Q_{\max})$ at time k , set

$$P_k^n(Q, (V, \mathcal{F})) := \text{esssup} \left\{ \mathbb{E} \left(\sum_{\ell=k}^{n-1} q_\ell V_\ell \mid \mathcal{F}_k \right), q \text{ } (\mathcal{F}, Q)\text{-admissible control} \right\} \quad (3.3)$$

so that $P_0^n(Q, (V, \mathcal{F}))$ is the value function of the stochastic control problem (3.1) when the global constraints Q_{\min} and Q_{\max} (at time 0) satisfy (3.2). Note that the standard convention $\text{esssup}(\emptyset) = 0$ yields $P_n^n \equiv 0$. To alleviate notations, the payoff process V the filtration \mathcal{F} will be often dropped in $P_k^n(Q, (V, \mathcal{F}))$.

To be precise we will answer the following questions:

- Existence of an optimal control $q^* = (q_k^*)_{0 \leq k \leq n-1}$.
- Regularity of the value function $Q \mapsto P_k^n(Q, (V, \mathcal{F}))$.
- Existence of a *bang-bang* optimal control $(q_k^*)_{0 \leq k \leq n-1}$ for certain values of the global constraints Q (namely when Q has integral components) ?

By *bang-bang* we mean that $\mathbb{P}(d\omega)$ -*a.s.*

– all the *local* constraints on the $q_k(\omega)$ are saturated *i.e.* for every $k \in \{0, \dots, n-1\}$, $q_k(\omega) \in \{0, 1 \wedge Q_{\max}\}$

or

– there exists at most one instant $k_0(\omega)$ such that $q_{k_0(\omega)}(\omega) \in (0, 1 \wedge Q_{\max})$ and one global constraint is saturated.

Note that if $Q \in \mathbb{N}^2$, then a *bang-bang* Q -admissible control necessarily satisfies $\mathbb{P}(d\omega)$ -*a.s.* $q_k(\omega) \in \{0, 1\}$, $k = 0, \dots, n-1$. The existence of *bang-bang* optimal controls combined with the piecewise affinity of $P_k^n(Q)$ will be the key in the design of a numerical.

- When there is an underlying structure Markov process ($V_k = v_k(Y_k)$), we will show that the optimal control turns out to be a function of Y_k at every time k as well.

4 Abstract dynamical programming principle and first properties

4.1 Basic properties

As a first step, we need to establish the following easy properties of P_k^n as a function of the global constraints $Q = (Q_{\min}, Q_{\max})$.

P1. For every $k \in \{0, \dots, n\}$, and every couple of global constraints (at time k),

$$P_k^n(Q, (V_k, \mathcal{F}_k)_{0 \leq k \leq n-1}) = P_0^{n-k}(Q, (V_{k+\ell}, \mathcal{F}_{k+\ell})_{0 \leq \ell \leq n-k-1}).$$

This is obvious from (3.3).

P2. If Q and Q' are two admissible global constraints (at time k) then

$$P_k^n(Q) = P_k^n(Q') \quad \text{on the event} \quad \{Q = Q'\}.$$

Proof. Owing to **P1** one may assume without loss of generality that $k = 0$. Let q and q' be two admissible controls with respect to Q and Q' respectively. Set $\tilde{q}_\ell = q_\ell \mathbf{1}_{\{Q=Q'\}} + q'_\ell \mathbf{1}_{\{Q \neq Q'\}}$, $\ell = 0, \dots, n-1$. The control \tilde{q} is admissible with respect to Q' since $\{Q = Q'\} \in \mathcal{F}_0$. Furthermore,

$$\begin{aligned} \mathbf{1}_{\{Q=Q'\}} \mathbb{E} \left(\sum_{\ell=0}^{n-1} q_\ell V_\ell \mid \mathcal{F}_0 \right) &= \mathbb{E} \left(\sum_{\ell=0}^{n-1} \mathbf{1}_{\{Q=Q'\}} q_\ell V_\ell \mid \mathcal{F}_0 \right) = \mathbf{1}_{\{Q=Q'\}} \mathbb{E} \left(\sum_{\ell=0}^{n-1} \tilde{q}_\ell V_\ell \mid \mathcal{F}_0 \right). \\ &\leq \mathbf{1}_{\{Q=Q'\}} P_0^n(Q') \quad \text{a.s.} \end{aligned}$$

Hence,

$$\mathbf{1}_{\{Q=Q'\}} P_0^n(Q) = \text{esssup} \left\{ \mathbf{1}_{\{Q=Q'\}} \mathbb{E} \left(\sum_{\ell=0}^{n-1} q_\ell V_\ell \mid \mathcal{F}_0 \right), q \text{ } Q\text{-admissible} \right\} \leq \mathbf{1}_{\{Q=Q'\}} P_0^n(Q') \quad \text{a.s.}$$

The equality follows by symmetry. \diamond

P3. Let $k \in \{0, \dots, n-1\}$. The set of admissible global constraints $Q = (Q_{\min}, Q_{\max})$ at time k is convex and the mapping $Q \mapsto P_k^n(Q)$ is concave in the following sense: if Q and Q' are two couples of admissible constraints, then for every random variable $\lambda : (\Omega, \mathcal{F}_k) \rightarrow [0, 1]$, $\lambda Q + (1 - \lambda)Q'$ is an admissible couple of constraints and

$$P_k^n(\lambda Q + (1 - \lambda)Q') \geq \lambda P_k^n(Q) + (1 - \lambda)P_k^n(Q') \quad \text{a.s.}$$

Furthermore $Q_{\min} \mapsto P_k^n(Q_{\min}, Q_{\max})$ is non-increasing and $Q_{\max} \mapsto P_k^n(Q_{\min}, Q_{\max})$ is non-decreasing i.e.

$$P_k^n(Q) \leq P_k^n(Q') \quad \text{a.s.} \quad \text{on} \quad \{Q'_{\min} \leq Q_{\min} \leq Q_{\max} \leq Q'_{\max}\}.$$

Proof. One may assume by **P1** that $k = 0$. The convexity of admissible global constraints is obvious. As concerns the concavity of the value function, note that if q and q' are locally admissible controls then $\lambda q + (1 - \lambda)q' := (\lambda q_k + (1 - \lambda)q'_k)_{0 \leq k \leq n-1}$ is still locally admissible. If q and q' satisfy

the global constraints induced by Q and Q' respectively, then $\lambda q + (1 - \lambda)q'$ always satisfies that induced by $\lambda Q + (1 - \lambda)Q'$. Consequently, using that λ is \mathcal{F}_0 -measurable,

$$\begin{aligned}
& P_0^n(\lambda Q + (1 - \lambda)Q') \\
\geq & \operatorname{esssup} \left\{ \mathbb{E} \left(\sum_{k=0}^{n-1} (\lambda q_k + (1 - \lambda)q'_k) V_k \mid \mathcal{F}_0 \right), q, q' \text{ locally admissible, } \bar{q}_n \in [Q_{\min}, Q_{\max}], \bar{q}'_n \in [Q'_{\min}, Q'_{\max}] \right\} \\
= & \lambda \operatorname{esssup} \left\{ \mathbb{E} \left(\sum_{k=0}^{n-1} q_k V_k \mid \mathcal{F}_0 \right), q \text{ locally admissible, } \bar{q}_n \in [Q_{\min}, Q_{\max}] \right\} \\
& + (1 - \lambda) \operatorname{esssup} \left\{ \mathbb{E} \left(\sum_{k=0}^{n-1} q'_k V_k \mid \mathcal{F}_0 \right), q' \text{ locally admissible, } \bar{q}'_n \in [Q'_{\min}, Q'_{\max}] \right\} \\
= & \lambda P_0^n(Q) + (1 - \lambda) P_0^n(Q').
\end{aligned}$$

The monotony property is as follows: let $A := \{Q'_{\min} \leq Q_{\min} \leq Q_{\max} \leq Q'_{\max}\} \in \mathcal{F}_0$ and let \tilde{q}' be a fixed Q' -admissible control. Then, for every Q -admissible control q , set

$$q' := q \mathbf{1}_A + \tilde{q}' \mathbf{1}_{c_A}.$$

Then q' is clearly Q' -admissible and

$$\begin{aligned}
\mathbf{1}_A \mathbb{E} \left(\sum_{k=0}^{n-1} q_k V_k \mid \mathcal{F}_0 \right) &= \mathbb{E} \left(\sum_{k=0}^{n-1} \mathbf{1}_A q_k V_k \mid \mathcal{F}_0 \right) \\
&= \mathbb{E} \left(\mathbf{1}_A \sum_{k=0}^{n-1} q'_k V_k \mid \mathcal{F}_0 \right) \\
&\leq \mathbf{1}_A P_0^n(Q') \quad a.s.
\end{aligned}$$

so that

$$\mathbf{1}_A P_0^n(Q) \leq \mathbf{1}_A P_0^n(Q') \quad a.s. \quad \diamond$$

P4. Let $k \in \{0, \dots, n-1\}$. Let $Q^{(n)}$, $n \geq 1$, be a sequence of admissible global constraints such that $Q_{\max}^{(n)} \rightarrow 0$ as $n \rightarrow \infty$. Then

$$P_k^n(Q^{(n)}) \longrightarrow 0 \quad \text{as} \quad n \rightarrow \infty.$$

Proof. Owing to **P1**, one may assume $k = 0$ without loss of generality. The result is straightforward once noticed that for every Q -admissible control

$$\left| \mathbb{E} \left(\sum_{k=0}^{n-1} q_k V_k \mid \mathcal{F}_0 \right) \right| \leq Q_{\max} \mathbb{E} \left(\max_{0 \leq k \leq n-1} |V_k| \mid \mathcal{F}_0 \right). \quad \diamond$$

P5. Let $T^+(n) := \{(u, v) \in \mathbb{R}_+^2, 0 \leq u \leq v \leq n\}$ and let $k \in \{0, \dots, n-1\}$. There is process $\Pi_k^n : (T^+(n) \times \Omega, \mathcal{B}or(T^+(n)) \otimes \mathcal{F}_k) \rightarrow \mathbb{R}$ such that

- (i) $(u, v) \mapsto \Pi_k^n(u, v, \omega)$ is concave and continuous on $T^+(n)$ for every $\omega \in \Omega$,
- (ii) $u \mapsto \Pi_k^n(u, v, \omega)$ is non-increasing on $[0, v]$ and $v \mapsto \Pi_k^n(u, v, \omega)$ is nondecreasing on $[u, n]$ for every $\omega \in \Omega$.

(iii) For every admissible constraint $Q = (Q_{\min}, Q_{\max})$ at time k , $P_k^n(Q)(\omega) = \Pi_k^n(Q(\omega), \omega) \mathbb{P}(d\omega)$ -a.s.

Proof. Classical consequence of **P3**: for every $(r, s) \in T^+(n) \cap \mathbb{Q}^2$ set

$$\Pi_k^n((r, s), \omega) := P_k^n((r, s))(\omega).$$

Then, set for every $(u, v) \in T^+(n)$

$$\Pi_k^n((u, v), \omega) := a.s. \lim_{(r,s) \rightarrow (u,v), (r,s) \in T^+(n) \cap \mathbb{Q}^2, r \leq u \leq v \leq s} P_k^n((r, s))(\omega).$$

One shows using the concavity and monotony properties established in **P3** that the above limit does exist and that $\mathbb{P}(d\omega)$ -a.s. $(u, v) \mapsto \Pi_k^n((u, v), \omega)$ is continuous on $T^+(n)$ and that $P_k^n(Q)(\omega) = \Pi_k^n(Q(\omega), \omega)$ \mathbb{P} -a.s..

P6. Let $Q = (Q_{\min}, Q_{\max})$ be a couple of global constraints (at time 0).

$$P_0^n(Q) = \text{esssup} \left\{ \mathbb{E} \left(\sum_{\ell=0}^{n-1} q_\ell (V_\ell)^+ \mid \mathcal{F}_0 \right), q \text{ } (\mathcal{F}, Q)\text{-admissible} \right\} \text{ on the event } \{Q_{\min} = 0\} \in \mathcal{F}_0.$$

Proof. This follows from the simple remark that one can define from any (\mathcal{F}, Q) -admissible control q a new (\mathcal{F}, Q) -admissible control \tilde{q} by $\tilde{q}_\ell := q_\ell \mathbf{1}_{\{V_\ell \geq 0\} \cap \{Q_{\min} = 0\}} + q_\ell \mathbf{1}_{\{Q_{\min} > 0\}}$ and

$$\mathbf{1}_{\{Q_{\min} = 0\}} \sum_{k=0}^{n-1} q_k V_k \leq \mathbf{1}_{\{Q_{\min} = 0\}} \sum_{k=0}^{n-1} q_k (V_k)^+ = \mathbf{1}_{\{Q_{\min} = 0\}} \sum_{k=0}^{n-1} \tilde{q}_k (V_k)^+. \quad \diamond$$

4.2 Dynamic programming principle

The main consequence of **P5** is that at every time k , one may assume without loss of generality that the couple of admissible global constraints $Q = (Q_{\min}, Q_{\max})$ is deterministic since, for any possibly random admissible constraint Q (at time k) and every $\omega \in \Omega$, $P_k^n(Q)(\omega) = \Pi_k^n(\omega, Q(\omega))$.

As a consequence, for notational convenience, we will still denote P_k^n instead of Π_k^n so that for any admissible global constraints Q at time k

$$P_k^n(Q)(\omega) = P_k^n(\omega, x)|_{x=Q(\omega)}.$$

Theorem 1 (*Backward Dynamic Programming Principle*) Set $P_n^n \equiv 0$.

(a) LOCAL DYNAMIC PROGRAMMING FORMULA. For every $k \in \{0, \dots, n-1\}$ and every couple $Q = (Q_{\min}, Q_{\max})$ of deterministic admissible global constraints at time k

$$P_k^n(Q) = \sup \left\{ xV_k + \mathbb{E} \left(P_{k+1}^n(\chi^{n-k-1}(Q, x)) \mid \mathcal{F}_k \right), x \in I_Q^{n-1-k} \right\} \quad (4.1)$$

where $\chi^M(Q, x) = ((Q_{\min} - x)^+, (Q_{\max} - x) \wedge M)$ and $I_Q^M := [(Q_{\min} - M)^+ \wedge 1, Q_{\max} \wedge 1]$.

(b) GLOBAL DYNAMIC PROGRAMMING FORMULA. For every couple $Q = (Q_{\min}, Q_{\max})$ of admissible global constraints at time 0, the price of the contract at time $k \in \{0, \dots, n-1\}$ is given by $P_k^n(Q^{k,*})$ where

$$Q^{k,*} := (Q_{\max} - \bar{q}_k^*, Q_{\min} - \bar{q}_k^*) \quad (\text{residual global constraints}) \quad (4.2)$$

$$\text{with } q_k^* = q_k^*(Q^{k,*})$$

$$q_k^*(Q) := \operatorname{argmax}_{x \in I_Q^{n-1-k}} \left(xV_k + \mathbb{E} \left(P_{k+1}^n(\chi^{n-k-1}(Q, x)) \mid \mathcal{F}_k \right) \right), k = 0, \dots, n-1. \quad (4.3)$$

Furthermore,

$$P_k^n(Q^{k,*}) = \mathbb{E} \left(\sum_{\ell=k}^{n-1} q_\ell^* V_\ell \mid \mathcal{F}_k \right).$$

Remark. The definition (4.3) may be ambiguous when argmax is not reduced to a single point. Then, one considers $\min \operatorname{argmax}$ to define $q_k^*(Q)$.

Proof. (a) It is clear that owing to **P1** the case $k = 0$ is the only one to be proved. As a first step, we prove that

$$P_0^n(Q) = \operatorname{esssup} \left\{ q_0 V_0 + \mathbb{E} \left(P_1^n(\chi^{n-1}(Q, q_0)) \mid \mathcal{F}_0 \right), q_0 : (\Omega, \mathcal{F}_0) \rightarrow I_Q^{n-1} \right\} \quad (4.4)$$

$\boxed{\leq}$: Let $q = (q_k)_{0 \leq k \leq n-1}$ be a Q -admissible control. Then, q_0 is $[0, 1]$ -valued (as well as the q_k 's) and

$$Q_{\min} - (n-1) \leq Q_{\min} - (q_1 + \dots + q_{n-1}) \leq q_0 \leq Q_{\max} - (q_1 + \dots + q_{n-1}) \leq Q_{\max}$$

so that q_0 is I_Q^{n-1} -valued. Furthermore,

$$\max(0, Q_{\min} - q_0) \leq q_1 + \dots + q_{n-1} \leq (Q_{\max} - q_0) \wedge (n-1). \quad (4.5)$$

note that $\chi^{n-1}(Q, q_0)$ is an admissible couple of (\mathcal{F}_0 -measurable) constraints at time 1. Consequently,

$$\begin{aligned} \mathbb{E} \left(\sum_{1 \leq \ell \leq n-1} q_\ell V_\ell \mid \mathcal{F}_0 \right) &= \mathbb{E} \left(\mathbb{E} \left(\sum_{1 \leq \ell \leq n-1} q_\ell V_\ell \mid \mathcal{F}_1 \right) \mid \mathcal{F}_0 \right) \\ &\leq \mathbb{E} \left(P_1^n(\chi^{n-1}(Q, q_0)) \mid \mathcal{F}_0 \right) \quad a.s. \end{aligned}$$

where the last inequality follows from the definition of P_1^n , χ^{n-1} , (4.5) and the monotony of conditional expectation. Then,

$$\mathbb{E} \left(\sum_{0 \leq \ell \leq n-1} q_\ell V_\ell \mid \mathcal{F}_0 \right) \leq \mathbb{E} \left(q_0 V_0 + P_1^n(\chi^{n-1}(Q, q_0)) \mid \mathcal{F}_0 \right) \quad a.s.$$

One concludes that

$$P_0^n(Q) \leq \operatorname{esssup} \left\{ \mathbb{E} \left(q_0 V_0 + P_1^n(\chi^{n-1}(Q, q_0)) \mid \mathcal{F}_0 \right), q_0 : (\Omega, \mathcal{F}_0) \rightarrow I_Q^{n-1} \right\}.$$

$\boxed{\geq}$: We proceed as usual by proving a bifurcation property for the controls. Let q_0 and q'_0 be two I_Q^{n-1} -valued \mathcal{F}_0 -measurable random variables. Set

$$A_0 := \{ q_0 V_0 + \mathbb{E} \left(P_1^n(\chi^{n-1}(Q, q_0)) \mid \mathcal{F}_0 \right) > q'_0 V_0 + \mathbb{E} \left(P_1^n(\chi^{n-1}(Q, q'_0)) \mid \mathcal{F}_0 \right) \} \in \mathcal{F}_0.$$

and

$$\tilde{q}_0 = q_0 \mathbf{1}_{A_0} + q'_0 \mathbf{1}_{c_{A_0}} : (\Omega, \mathcal{F}_0) \rightarrow I_Q^{n-1}.$$

Then, one checks that

$$\tilde{q}_0 V_0 + \mathbb{E} \left(P_1^n(\chi^{n-1}(Q, \tilde{q}_0)) \mid \mathcal{F}_0 \right) = \max_{y=q_0, q'_0} \left(y V_0 + \mathbb{E} \left(P_1^n(\chi^{n-1}(Q, y)) \mid \mathcal{F}_0 \right) \right).$$

Consequently there exists a sequence $q_0^{(n)}$ of $[0, Q_{\max} \wedge 1]$ -valued random variables such that

$$\begin{aligned} \operatorname{esssup} \left\{ q_0 V_0 + \mathbb{E} \left(P_1^n(\chi^{n-1}(Q, q_0)) \mid \mathcal{F}_0 \right), q_0 : (\Omega, \mathcal{F}_0) \rightarrow I_Q^{n-1} \right\} \\ &= \sup_n q_0^{(n)} V_0 + \mathbb{E} \left(P_1^n(\chi^{n-1}(Q, q_0^{(n)})) \mid \mathcal{F}_0 \right) \\ &= \lim_n^{\uparrow} q_0^{(n)} V_0 + \mathbb{E} \left(P_1^n(\chi^{n-1}(Q, q_0^{(n)})) \mid \mathcal{F}_0 \right). \end{aligned}$$

One may assume by applying the above bifurcation property that the above supremum holds as a nondecreasing limit as $n \rightarrow \infty$.

Now, for every fixed $n \geq 1$, there exists a sequence of $[0, 1]^{n-1}$ -valued random vectors $(q_k^{(n,m)})_{1 \leq k \leq n-1}$ such that $\sum_{k=1}^{n-1} q_k^{(n,m)} \in [(Q_{\min} - q_0^{(n)})^+, Q_{\max} - q_0^{(n)}]$ and

$$P_1^n(\chi^{n-1}(Q, q_0^{(n)})) = \lim_m^\uparrow \left(\sum_{k=1}^{n-1} q_k^{(n,m)} V_k \mid \mathcal{F}_1 \right) \quad a.s.$$

where we used that the admissible sequences $(q_k)_{1 \leq k \leq n-1}$ for the problem starting at 1 clearly satisfy the bifurcation principle due to the homogeneity of conditional expectation $\mathbb{E}(\cdot \mid \mathcal{F}_0)$ with respect to \mathcal{F}_0 -measurable r.v.. Consequently,

$$\begin{aligned} q_0^{(n)} V_0 + P_1^n(\chi^{n-1}(Q, q_0^{(n)})) &= \lim_m^\uparrow q_0^{(n)} V_0 + \mathbb{E} \left(\sum_{k=1}^{n-1} q_k^{(n,m)} V_k \mid \mathcal{F}_1 \right) \quad a.s. \\ \mathbb{E} \left(q_0^{(n)} V_0 + P_1^n(\chi^{n-1}(Q, q_0^{(n)})) \mid \mathcal{F}_0 \right) &= \lim_m^\uparrow \mathbb{E} \left(q_0^{(n)} V_0 + \sum_{k=1}^{n-1} q_k^{(n,m)} V_k \mid \mathcal{F}_0 \right) \quad a.s. \end{aligned}$$

where we used the conditional Beppo Levi Theorem. Note that for every $n, m \geq 1$, $(q_0^{(n)}, q_k^{(n,m)}, k = 1, \dots, n-1)$ is an admissible control with respect to Q since $x + (Q_{\min} - x)^+ \geq Q_{\min}$. Hence

$$\begin{aligned} \mathbb{E} \left(q_0^{(n)} V_0 + P_1^n(\chi^{n-1}(Q, q_0^{(n)})) \mid \mathcal{F}_0 \right) &\leq \text{esssup} \left\{ \mathbb{E} \left(q_0 V_0 + \sum_{k=1}^{n-1} q_k V_k \mid \mathcal{F}_0 \right), q \text{ } Q\text{-admissible} \right\} \quad a.s. \\ &= P_0^n(Q) \quad a.s. \end{aligned}$$

To pass from (4.4) to (4.1) is standard using **P5**. Let $q_0 : (\Omega, \mathcal{F}_0) \rightarrow I_Q^{n-1}$.

$$\begin{aligned} q_0 V_0 + \mathbb{E}(P_1^n(\chi^{n-1}(Q, q_0)) \mid \mathcal{F}_0) &= q_0 V_0 + \mathbb{E}(P_1^n(y) \mid \mathcal{F}_0)_{|y=\chi^{n-1}(Q, q_0)} \\ &\leq \sup_{x \in I_Q^{n-1}} (x V_0 + \mathbb{E}(P_1^n(y) \mid \mathcal{F}_0)_{|y=\chi^{n-1}(Q, x)}) \\ &= \sup_{x \in I_Q^{n-1}} (x V_0 + \mathbb{E}(P_1^n(\chi^{n-1}(Q, x)) \mid \mathcal{F}_0)). \end{aligned}$$

Conversely, setting $q_0^x(\omega) := x \in I_Q^{n-1}$ yields the reverse inequality.

(b) This item follows from **P5**. \diamond

5 Affine value function with bang-bang optimal controls

5.1 The main result

We recall that, for every integer $n \geq 1$, the triangular set of admissible values for a couple of (deterministic) global constraint (at time 0), as introduced in **P5**, is defined as:

$$T^+(n) := \{(u, v), 0 \leq u \leq v \leq n\}.$$

Then we will define a triangular tiling of $T^+(n)$ as follows: for every couple of integers (i, j) , $0 \leq i \leq j \leq n-1$,

$$T_{ij}^+ := \{(u, v) \in [i, i+1] \times [j, j+1], v \geq u+j-i\} \quad \text{and} \quad T_{ij}^- := \{(u, v) \in [i, i+1] \times [j, j+1], v \leq u+j-i\}.$$

One checks that

$$T^+(n) = \left(\bigcup_{0 \leq i < j \leq n-1} T_{ij}^+ \right) \cup \left(\bigcup_{0 \leq i < j \leq n-1} T_{ij}^- \right).$$

Theorem 2 *The multi-period swing option premium with deterministic global constraints $Q := (Q_{\min}, Q_{\max}) \in T^+(n)$ as defined by (3.3) is always obtained as the result of an optimal strategy.*

(a) *The value function (premium):*

– *the mapping $Q \mapsto P_0(Q, \mathcal{F})$ is a concave, continuous, piecewise affine process, affine on every triangle $T_{i,j}^\pm$ of the tiling of $T^+(n)$. Furthermore,*

$$P_0^n(0, 0) = 0, \quad P_0^n(0, n) = \mathbb{E} \left(\sum_{k=0}^{n-1} V_k^+ \mid \mathcal{F}_0 \right) \quad \text{and} \quad P_0^n(n, n) = \mathbb{E} \left(\sum_{k=0}^{n-1} V_k \mid \mathcal{F}_0 \right).$$

– *If $V_i \geq 0$ a.s. for every $i \in \{0, \dots, n-1\}$, then, for every $Q = (Q_{\min}, Q_{\max}) \in T^+(n)$, $P_0(Q, \mathcal{F}) = P_0((0, Q_{\max}), \mathcal{F}) = P_0((Q_{\max}, Q_{\max}), \mathcal{F})$ a.s.. (in particular $Q_{\max} \mapsto P_0((Q_{\max}, Q_{\max}), \mathcal{F})$ is a.s. non-decreasing).*

(b) *The optimal control:*

– *If the global constraint $Q = (Q_{\min}, Q_{\max}) \in \mathbb{N}^2 \cap T^+(n)$, then there always exists a bang-bang optimal control $q^* = (q_k^*)_{0 \leq k \leq n-1}$ with q_k^* is $\{0, 1\}$ -valued for every $k = 0, \dots, n-1$.*

– *If $V_i \geq 0$ a.s. for every $i \in \{0, \dots, n-1\}$, then there always exists a bang-bang optimal control which satisfies $\sum_{0 \leq k \leq n-1} q_k^* = Q_{\max}$.*

– *Otherwise the optimal control is not bang-bang as emphasized by the case $n = 2$ (see proposition 2 below)*

We will first inspect the case of a two period swing contract. It will illustrate in a simpler setting the approach developed in the general case. Furthermore, we will obtain a slightly more precise result about the optimal controls.

5.2 The two period option

We assume $n = 2$ throughout this section. The first result is the following

Proposition 1 *Let $Q = (Q_{\min}, Q_{\max}) \in T^+(2)$ denote an admissible global constraint and $I_Q^1 := [(Q_{\min} - 1)^+, Q_{\max} \wedge 1]$. There is an optimal control $q^* = (q_0^*, q_1^*)$ given by*

$$q_0^* = \operatorname{argmax}_{x \in I_Q^1} \{xV_0 + 1 \wedge (Q_{\max} - x) \mathbb{E}(V_1^+ | \mathcal{F}_0) - (Q_{\min} - x)^+ \mathbb{E}(V_1^- | \mathcal{F}_0)\} \quad (5.1)$$

$$q_1^* = 1 \wedge (Q_{\max} - q_0^*) \mathbf{1}_{\{V_1 \geq 0\}} + (Q_{\min} - q_0^*)^+ \mathbf{1}_{\{V_1 < 0\}} \quad (5.2)$$

so that

$$P_0^2(Q, \mathcal{F}) = \mathbb{E}(q_0^* V_0 + q_1^* V_1 \mid \mathcal{F}_0).$$

Proof. Let $q = (q_0, q_1)$ be an admissible control: $q_0 + q_1 \in [Q_{\min}, Q_{\max}]$ and q_i are $[0, 1]$ -valued \mathcal{F}_i -measurable, $i = 0, 1$. Consequently q_0 is I_Q^1 -valued and q_1 is $[(Q_{\min} - q_0)^+, (Q_{\max} - q_0) \wedge 1]$ -valued. Hence

$$q_0 V_0 + q_1 V_1 \leq q_0 V_0 + 1 \wedge (Q_{\max} - q_0) V_1^+ - (Q_{\min} - q_0)^+ V_1^-. \quad (5.3)$$

On the other hand

$$\begin{aligned} \mathbb{E}(q_0 V_0 + 1 \wedge (Q_{\max} - q_0) V_1^+ - (Q_{\min} - q_0)^+ V_1^- \mid \mathcal{F}_0) \\ = q_0 V_0 + 1 \wedge (Q_{\max} - q_0) \mathbb{E}(V_1^+ | \mathcal{F}_0) - (Q_{\min} - q_0)^+ \mathbb{E}(V_1^- | \mathcal{F}_0). \end{aligned}$$

The mapping $x \mapsto x V_0 + 1 \wedge (Q_{\max} - x) \mathbb{E}(V_1^+ | \mathcal{F}_0) - (Q_{\min} - x)^+ \mathbb{E}(V_1^- | \mathcal{F}_0)$ (called the *objective variable* from now on) is piecewise affine on I_Q^1 with \mathcal{F}_0 -measurable coefficients so the above definition of q_0^* defines an \mathcal{F}_0 -measurable I_Q^1 -valued random variable. Now, combining the above inequalities yields

$$\begin{aligned} \mathbb{E}(q_0 V_0 + q_1 V_1 | \mathcal{F}_0) &\leq q_0 V_0 + 1 \wedge (Q_{\max} - q_0) \mathbb{E}(V_1^+ | \mathcal{F}_0) - (Q_{\min} - q_0)^+ \mathbb{E}(V_1^- | \mathcal{F}_0) \\ &\leq \sup_{q_0 \in I_Q^1} q_0 V_0 + 1 \wedge (Q_{\max} - q_0) \mathbb{E}(V_1^+ | \mathcal{F}_0) - (Q_{\min} - q_0)^+ \mathbb{E}(V_1^- | \mathcal{F}_0) \\ &= \mathbb{E}(q_0^* V_0 + q_1^* V_1 | \mathcal{F}_0). \quad \diamond \end{aligned}$$

In the proposition below we investigate in full details the case $n = 2$.

Proposition 2 *Let $n = 2$. The two period swing option premium with admissible global constraints $Q = (Q_{\min}, Q_{\max}) \in T^+(2)$ as defined by (3.1) is always obtained as the result of an optimal strategy.*

(a) *The optimal control:*

– *If the global constraints Q_{\min}, Q_{\max} only take integral values (in $\{0, 1, 2\}$) then there always exists a $\{0, 1\}$ -valued bang-bang optimal control. When Q simply satisfies $Q_{\max} - Q_{\min} \in \{0, 1, 2\}$, there always exists a bang-bang optimal control.*

– *If $V_0, V_1 \geq 0$ a.s., then any optimal control q^* is a.s. bang-bang and satisfies $q_0^* + q_1^* = Q_{\max}$ on $\{V_i > 0, i = 1, 2\}$.*

– *Otherwise the optimal control is generally not bang-bang.*

(b) *The value function (premium):*

– *the mapping $Q = (Q_{\min}, Q_{\max}) \mapsto P_0(Q, \mathcal{F})$ is affine on the four triangles $T_{i,j}^\pm$ that tile $T^+(2)$.*

– *Furthermore, when V_0 and V_1 are a.s. non negative,*

$$P_0^2(Q, \mathcal{F}) = (Q_{\max} - 1)^+ \wedge 1 (V_0 \wedge \mathbb{E}(V_1 | \mathcal{F}_0)) + (Q_{\max} \wedge 1) (V_0 \vee \mathbb{E}(V_1 | \mathcal{F}_0)).$$

The objective variable being piecewise affine on I_Q^1 , q_0^* is equal either to one of its monotony breaks or to the endpoints of I_Q^1 . Consequently, a careful inspection of all possible situations for the global constraints yields the complete set of explicit optimal rules for the optimal exercise of the swing option involving the values V_0 and $\mathbb{E}(V_1^\pm | \mathcal{F}_0)$ (expected gain or loss at time 0) at time 0 and V_1 at time 1.

$Q \in T_{00}^+$ i.e. $Q_{\min} \leq Q_{\max} \leq 1$: $I_Q^1 = [0, Q_{\max}]$ and the objective variable reads

$$q_0 V_0 + (Q_{\max} - q_0) \mathbb{E}(V_1^+ | \mathcal{F}_0) - (Q_{\min} - q_0)^+ \mathbb{E}(V_1^- | \mathcal{F}_0)$$

with one monotony break at Q_{\min} . One checks that

$$\begin{array}{lll} q_0^* = Q_{\max}, & q_1^* = 0 & \text{on } \{V_0 \geq \mathbb{E}(V_1^+ | \mathcal{F}_0)\}, \\ q_0^* = Q_{\min}, & q_1^* = (Q_{\max} - Q_{\min}) \mathbf{1}_{\{V_1 \geq 0\}} & \text{on } \{\mathbb{E}(V_1 | \mathcal{F}_0) \leq V_0 < \mathbb{E}(V_1^+ | \mathcal{F}_0)\}, \\ q_0^* = 0, & q_1^* = Q_{\max} \mathbf{1}_{\{V_1 \geq 0\}} + Q_{\min} \mathbf{1}_{\{V_1 < 0\}} & \text{on } \{V_0 < \mathbb{E}(V_1 | \mathcal{F}_0)\}. \end{array}$$

Note that

$$q_0^* = Q_{\min}, \quad q_1^* = Q_{\max} - Q_{\min} \quad \text{on } \{\mathbb{E}(V_1 | \mathcal{F}_0) \leq V_0 < \mathbb{E}(V_1^+ | \mathcal{F}_0)\} \cap \{V_1 \geq 0\}$$

so that the above optimal control is not bang-bang on this event except if $Q_{\min} \in \{0, 1\}$ or $Q_{\max} = Q_{\min}$.

$Q \in T_{01}^+$ i.e. $Q_{\min} \leq 1 \leq Q_{\max} \leq 2$, $Q_{\min} \leq Q_{\max} - 1$: $I_Q^1 = [0, 1]$ and the objective variable reads

$$q_0 V_0 + 1 \wedge (Q_{\max} - q_0) \mathbb{E}(V_1^+ | \mathcal{F}_0) - (Q_{\min} - q_0)^+ \mathbb{E}(V_1^- | \mathcal{F}_0)$$

with monotony breaks at Q_{\min} , $Q_{\max} - 1$. One checks that

$$\begin{array}{lll} q_0^* = 1, & q_1^* = (Q_{\max} - 1) \mathbf{1}_{\{V_1 \geq 0\}} & \text{on } \{V_0 \geq \mathbb{E}(V_1^+ | \mathcal{F}_0)\}, \\ q_0^* = Q_{\max} - 1, & q_1^* = \mathbf{1}_{\{V_1 \geq 0\}} & \text{on } \{0 \leq V_0 < \mathbb{E}(V_1^+ | \mathcal{F}_0)\}, \\ q_0^* = Q_{\min}, & q_1^* = \mathbf{1}_{\{V_1 \geq 0\}} & \text{on } \{-\mathbb{E}(V_1^- | \mathcal{F}_0) \leq V_0 < 0\} \\ q_0^* = 0, & q_1^* = \mathbf{1}_{\{V_1 \geq 0\}} + Q_{\min} \mathbf{1}_{\{V_1 < 0\}} & \text{on } \{V_0 < -\mathbb{E}(V_1^- | \mathcal{F}_0)\}. \end{array}$$

Note that

$$q_0^* + q_1^* = Q_{\max} - 1 \quad \text{on} \quad \{0 \leq V_0 < \mathbb{E}(V_1^+ | \mathcal{F}_0)\} \cap \{V_1 < 0\}$$

so that the control is not bang-bang on this event, except if $Q_{\max} \in \{1, 2\}$ or $Q_{\max} = 1 + Q_{\min}$, since the local control q_0^* and the global constraint are not saturated. Likewise

$$q_0^* + q_1^* = 1 + Q_{\min} \quad \text{on} \quad \{-\mathbb{E}(V_1^- | \mathcal{F}_0) \leq V_0 < 0\} \cap \{V_1 > 0\}$$

and the optimal control is not bang-bang on this event, except when $Q_{\min} \in \{0, 1\}$ or $Q_{\max} = 1 + Q_{\min}$.

Note that *both events correspond to prediction errors*: V_1 has not the predicted sign. Moreover, these events are *a.s.* empty when $V_i \geq 0$ *a.s.*, $i = 1, 2$. On all other events the optimal control is bang-bang.

$Q \in T_{01}^-$ i.e. $Q_{\min} \leq 1 \leq Q_{\max} \leq 2$, $Q_{\min} \geq Q_{\max} - 1$: Then the monotony breaks of the objective process (with the same expression as in the former case) still are Q_{\min} , $Q_{\max} - 1$. A careful inspection of the four possible cases leads to

$$\begin{array}{lll} q_0^* = 1, & q_1^* = (Q_{\max} - 1) \mathbf{1}_{\{V_1 \geq 0\}} & \text{on } \{V_0 \geq \mathbb{E}(V_1^+ | \mathcal{F}_0)\}, \\ q_0^* = Q_{\min}, & q_1^* = (Q_{\max} - Q_{\min}) \mathbf{1}_{\{V_1 \geq 0\}} & \text{on } \{\mathbb{E}(V_1 | \mathcal{F}_0) \leq V_0 < \mathbb{E}(V_1^+ | \mathcal{F}_0)\}, \\ q_0^* = Q_{\max} - 1, & q_1^* = \mathbf{1}_{\{V_1 \geq 0\}} + (Q_{\min} - Q_{\max} + 1) \mathbf{1}_{\{V_1 < 0\}} & \text{on } \{-\mathbb{E}(V_1^- | \mathcal{F}_0) \leq V_0 < \mathbb{E}(V_1 | \mathcal{F}_0)\} \\ q_0^* = 0, & q_1^* = \mathbf{1}_{\{V_1 \geq 0\}} + Q_{\min} \mathbf{1}_{\{V_1 < 0\}} & \text{on } \{V_0 < -\mathbb{E}(V_1^- | \mathcal{F}_0)\}. \end{array}$$

Note that on the event

$$\{-\mathbb{E}(V_1^- | \mathcal{F}_0) \leq V_0 < \mathbb{E}(V_1 | \mathcal{F}_0)\} \cap \{V_1 < 0\}$$

the optimal control is not bang-bang, except if $Q_{\max} \in \{1, 2\}$ or $Q_{\max} = Q_{\min}$ (both q_0^* and q_1^* are $(0, 1)$ -valued) or $Q_{\max} = Q_{\min} + 1$ ($q_0^* = Q_{\min}$, $q_1^* = 0$); and on the event

$$\{\mathbb{E}(V_1 | \mathcal{F}_0) \leq V_0 < \mathbb{E}(V_1^+ | \mathcal{F}_0)\} \cap \{V_1 > 0\}$$

the optimal control is not bang-bang either (except if $Q_{\min} \in \{0, 1\}$ or $Q_{\max} = Q_{\min}$ or $Q_{\max} = Q_{\min} + 1$) by similar arguments.

Note that these events *do not correspond to an error of prediction*. On all other events the optimal control is bang-bang.

$Q \in T_{11}^+$ i.e. $1 < Q_{\min} \leq Q_{\max} \leq 2$: The objective variable is defined on $I_Q^1 = [Q_{\min} - 1, 1]$ by

$$q_0 V_0 + 1 \wedge (Q_{\max} - q_0) \mathbb{E}(V_1^+ | \mathcal{F}_0) - (Q_{\min} - q_0) \mathbb{E}(V_1^- | \mathcal{F}_0)$$

with only one breakpoint at $Q_{\max} - 1$. One checks that

$$\begin{aligned} q_0^* &= 1, & q_1^* &= (Q_{\max} - 1)\mathbf{1}_{\{V_1 \geq 0\}} + (Q_{\min} - 1)\mathbf{1}_{\{V_1 < 0\}} && \text{on } \{\mathbb{E}(V_1 | \mathcal{F}_0) \leq V_0\}, \\ q_0^* &= Q_{\max} - 1, & q_1^* &= \mathbf{1}_{\{V_1 \geq 0\}} + (Q_{\min} - Q_{\max} + 1)\mathbf{1}_{\{V_1 < 0\}} && \text{on } \{-\mathbb{E}(V_1^- | \mathcal{F}_0) \leq V_0 < \mathbb{E}(V_1 | \mathcal{F}_0)\} \\ q_0^* &= Q_{\min} - 1, & q_1^* &= 1 && \text{on } \{V_0 < -\mathbb{E}(V_1^- | \mathcal{F}_0)\}. \end{aligned}$$

Once again on the event

$$\{-\mathbb{E}(V_1^- | \mathcal{F}_0) \leq V_0 < \mathbb{E}(V_1 | \mathcal{F}_0)\} \cap \{V_1 < 0\}$$

the optimal control is not bang-bang, except if $Q_{\max} \in \{1, 2\}$ or $Q_{\max} = Q_{\min}$ or $Q_{\max} = Q_{\min} + 1$.

Finally, note that when $V_0, V_1 \geq 0$, the events on which the optimal controls are not bang-bang are empty. \diamond

5.3 Proof of Theorem 2

(a) We proceed by induction on n . For $n = 1$ the result is trivial since $T^+(1) = T_{00}^+$ and $P_0(Q) = Q_{\max}\mathbf{1}_{\{V_0 \geq 0\}} + Q_{\min}\mathbf{1}_{\{V_0 < 0\}}$. (When $n = 2$ this follows from Proposition 2.)

Now, we pass from n to $n + 1$. Note that combining the backward programming principle and **P1** yields

$$P_0^{n+1}(Q, \mathcal{F}) = \sup \{xV_0 + \mathbb{E}(P_0^n(\chi^n(Q, x), (\mathcal{F}_{1+\ell})_{0 \leq \ell \leq n-1}) | \mathcal{F}_0), x \in I_Q^n\}. \quad (5.4)$$

We inspect successively all the triangles of the tiling of $T^+(n + 1)$ as follows: the upper and lower triangles which lie strictly inside the tiling, then the triangles which lie on the boundary of the tiling.

$Q \in T_{ij}^+, 1 \leq i \leq j \leq n - 1$ Then, $\chi^n(Q, x) = Q - x(1, 1)$ and $I_Q^n = [0, 1]$. One checks that $\chi^n(Q, x) \in T_{ij}^+$ if $x \in [0, Q_{\min} - i]$, $\chi^n(Q, x) \in T_{i-1, j}^-$ if $x \in [Q_{\min} - i, Q_{\max} - j]$ and $\chi^n(Q, x) \in T_{i-1, j-1}^+$ if $x \in [Q_{\max} - j, 1]$ (see Figure 1). These three triangles $T_{ij}^+, T_{i-1, j}^-$ and $T_{i-1, j-1}^+$ are included in $T^+(n)$. It follows from the induction assumption that $(u, v) \mapsto P_0^n((u, v), (\mathcal{F}_{1+}))$ is *a.s.* affine on them. Hence there exists three triplets of \mathcal{F}_1 -measurable random variables (A^m, B^m, C^m) , $m = 1, 2, 3$, such that, for every $Q \in T_{ij}^+$,

$$P_0^n(\chi^n(Q, x), (\mathcal{F}_{1+})) = \sum_{m=1}^3 \mathbf{1}_{J_Q^m} (A^m(Q_{\min} - x) + B^m(Q_{\max} - x) + C^m)$$

where $J_Q^1 = [0, Q_{\min} - i]$, $J_Q^2 = [Q_{\min} - i, Q_{\max} - j]$ and $J_Q^3 = [Q_{\max} - j, 1]$. Note these random coefficients satisfy some compatibility constraints to ensure concavity (and continuity). Consequently

$$xV_0 + \mathbb{E}(P_0^n(\chi^n(Q, x), (\mathcal{F}_{1+})) | \mathcal{F}_0) = \sum_{m=1}^3 \mathbf{1}_{J_Q^m} (xV_0 + A_0^m(Q_{\min} - x) + B_0^m(Q_{\max} - x) + C_0^m)$$

where $A_0^m = \mathbb{E}(A^m | \mathcal{F}_0)$, etc. A piecewise affine function reaches its maximum on a compact interval either at its endpoint or at its monotony breakpoints $x_1 = 0$, $x_2 = Q_{\min} - i$, $x_3 = Q_{\max} - j$, $x_4 = 1$. Hence,

$$\sup_{x \in I_Q^n} (xV_0 + \mathbb{E}(P_0^n(\chi^n(Q, x), (\mathcal{F}_{1+})) | \mathcal{F}_0)) = \max \{x_\ell V_0 + A_0^m(Q_{\min} - x_\ell) + B_0^m(Q_{\max} - x_\ell) + C_0^m,$$

$$(\ell, m) = (1, 1), (2, 1), (3, 2), (4, 3)\}.$$

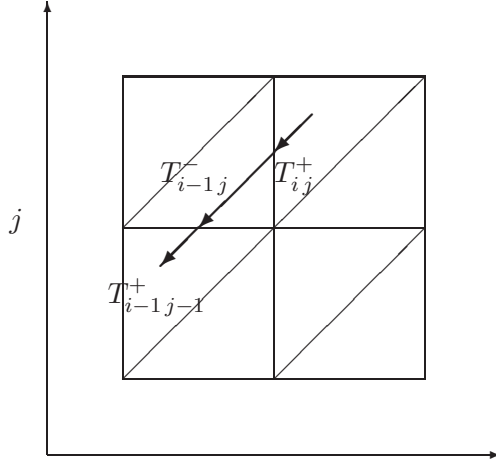


Figure 1: $x \mapsto \chi^n(x)$ for $Q \in T_{i,j}^+$, $1 \leq i \leq j \leq n-1$

It is clear that the right hand side of the previous equality stands as the maximum of four affine functions of Q . One derives that $Q \mapsto P_0^{n+1}(Q, \mathcal{F})$ is a convex function on $T_{i,j}^+$ as the maximum of affine functions. Hence it is affine since we know that it is also concave.

$Q \in T_{i,j}^-$, $1 \leq i < j \leq n-1$: This case can be treated likewise.

$Q \in T_{0,j}^\pm$, $1 \leq j \leq n-1$: In that case $I_Q^n = [0, 1]$, $\chi^n(Q, x) = ((Q_{\min} - x)^+, Q_{\max} - x)$, $x \in I_Q^n$.

- If $Q \in T_{0,j}^+$, $\chi^n(Q, x) = Q - x(1, 1) \in T_{0,j}^+$, $x \in [0, Q_{\min}]$, $\chi^n(Q, x) = (0, Q_{\max} - x) \in T_{0,j}^+$, $x \in [Q_{\min}, Q_{\max} - j]$, $\chi^n(Q, x) = (0, Q_{\max} - x) \in T_{0,j-1}^+$, $x \in [Q_{\max} - j, 1]$ (see Figure 2). The induction assumption implies that $x \mapsto P_0^n(\chi^n(Q, x), (\mathcal{F}_{1+}))$ is piecewise affine with monotony breaks at Q_{\min} and $Q_{\max} - j$.

- If $Q \in T_{0,j}^-$, $\chi^n(Q, x)$ crosses the upper (horizontal) edge of $T_{0,j-1}^+$ at $x = Q_{\max} - j$ and the left (vertical) edge of $T^+(n)$ at $x = Q_{\min}$. Hence $x \mapsto P_0^n(\chi^n(Q, x), (\mathcal{F}_{1+}))$ is again piecewise affine with monotony breaks at Q_{\min} and $Q_{\max} - j$.

In both cases one concludes as above.

$Q \in T_{0,0}^\pm$: On proceeds like with $T_{0,j}^+$ except that $I_Q^n = [0, Q_{\max}]$ which yields only one monotony break at Q_{\min} .

$Q \in T_{i,n}^\pm$, $1 \leq i \leq n-1$: Assume first $Q \in T_{i,n}^+$. $I_Q^n = [0, 1]$ and $\chi^n(Q, x) = (Q_{\min} - x, n)$ if $x \in [0, Q_{\max} - n]$. Otherwise $\chi^n(Q, x) = (Q_{\min} - x, Q_{\max} - x)$. It follows (see Figure 3) that $\chi^n(Q, x) \in T_{i,n-1}^+$ if $x \in [0, Q_{\min} - i]$ and $\chi^n(Q, x) \in T_{i-1,n-1}^+$ if $x \in [Q_{\min} - i, 1]$. Both $T_{i,n-1}^+$ and $T_{i-1,n-1}^+$ are included in $T^+(n)$. Hence $(u, v) \mapsto P_0^n((u, v), (\mathcal{F}_{1+}))$ is affine on both triangles, one derives that

$$\begin{aligned} & P_0^n(\chi^n(Q, x), (\mathcal{F}_{1+})) \\ &= \mathbf{1}_{x \in [0, Q_{\max} - n]}(A^1(Q_{\min} - x) + B^1 n + C^1) + \mathbf{1}_{x \in [Q_{\max} - n, 1]}(A^2(Q_{\min} - x) + B^2(Q_{\max} - x) + C^2). \end{aligned}$$

where A^m, B^m, C^m , $m = 1, 2$ are \mathcal{F}_1 -measurable r.v.. Then, one concludes like in the first case.

If $Q \in T_{i,n}^-$, one proceeds likewise except that the two “visited” triangles of $T^+(n)$ are $T_{i-1,n-1}^\pm$.

$Q \in T_{0,n}^\pm$: $I_Q^n = [0, 1]$ and

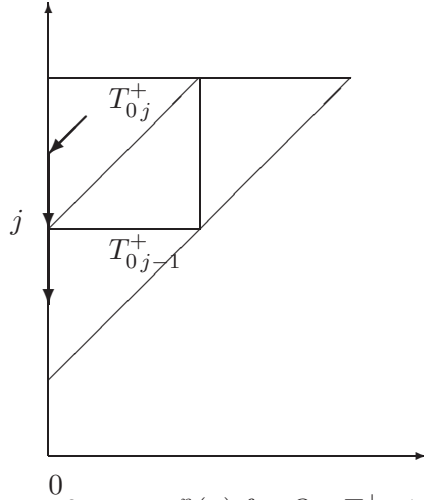


Figure 2: $x \mapsto \chi^n(x)$ for $Q \in T_{0j}^+$, $1 \leq j \leq n-1$

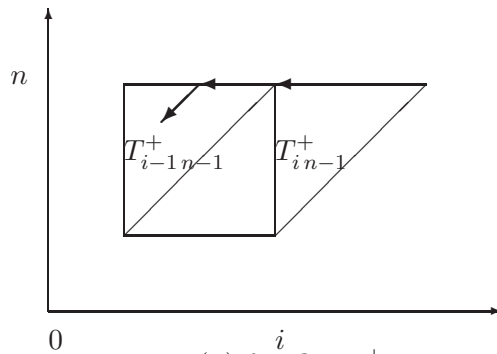


Figure 3: $x \mapsto \chi^n(x)$ for $Q \in T_{in}^+$, $1 \leq i \leq n-1$

- $\chi^n(Q, x) = ((Q_{\min} - x)^+, n)$, $x \in I_Q^n$ if $Q \in T_{0n}^+$,
- $\chi^n(Q, x) = (Q_{\min} - x, n)$, $x \in [0, Q_{\max} - n]$, $\chi^n(Q, x) = (Q_{\min} - x, Q_{\max} - x)$, $x \in [Q_{\max} - n, Q_{\min}]$ if $Q \in T_{0n}^-$.

In both cases the only “visited” triangle is $T_{0n-1}^+ \subset T^+(n)$ and one concludes as usual.

$Q \in T_{nn}^+$: $I_Q^n = [Q_{\min} - n, 1]$ and $\chi^n(Q, x) = (Q_{\min} - x, n)$ if $x \in [Q_{\min} - n, Q_{\max} - n]$, $\chi^n(Q, x) = (Q_{\min} - x, Q_{\max} - x)$ otherwise. Hence $\chi^n(Q, x)$ takes its values in T_{n-1n-1}^+ on which $(u, v) \mapsto P_0^n((u, v), (\mathcal{F}_{1+}))$ is affine. The conclusion follows.

The inspection of all these cases completes the proof of the induction.

The values of $P_0^n(Q)$ when $Q \in \{(0, 0), (0, n), (n, n)\}$ are obvious consequences of the degeneracy of the global constraints.

(b) We deal successively with the two announced settings.

- GLOBAL CONSTRAINTS IN \mathbb{N}^2 : Let $n \geq 1$. We rely on the characterization (4.3) of q_0^*

$$q_0^* = \operatorname{argmax}_{x \in I_Q^{n-1}} (xV_0 + \mathbb{E}(P_0^{n-1}(\chi^{n-1}(Q, x), (\mathcal{F}_{1+\ell})_{0 \leq \ell \leq n-2}) | \mathcal{F}_0)).$$

We know from item (a) that $(u, v) \mapsto P^{n-1}((u, v), \mathcal{F}_{1+})$ is affine on every tile T_{ij}^\pm of $T^+(n)$.

If $Q \in T^+(n) \cap \mathbb{N}^2$ then one checks that $x \mapsto \chi^{n-1}(Q, x)$, $x \in I_Q^{n-1}$, is always affine with I_Q^{n-1} having 0 and/or 1 as endpoints. To be precise

- if $Q = (i, j)$, $1 \leq i \leq j \leq n-1$, $\chi^{n-1}(Q, x) = (i-x, j-x) \in \partial T_{i-1j-1}^+ \cap \partial T_{i-1j-1}^-$, $I_Q^{n-1} = [0, 1]$,
- if $Q = (0, j)$, $1 \leq j \leq n-1$, $\chi^{n-1}(Q, x) = (0, j-x) \in \partial T_{0j-1}^+$, $I_Q^{n-1} = [0, 1]$,
- if $Q = (0, 0)$, $\chi^{n-1}(Q, x) = (0, 0)$, $I_Q^{n-1} = \{0\}$,
- if $Q = (i, n)$, $1 \leq i \leq n-1$, $\chi^{n-1}(Q, x) = (i-x, n-1) \in \partial T_{i-1n-1}^-$, $I_Q^{n-1} = [0, 1]$,
- if $Q = (0, n)$, $\chi^{n-1}(Q, x) = (0, n-1)$, $I_Q^{n-1} = [0, 1]$,
- if $Q = (n, n)$, $\chi^{n-1}(Q, x) = (n-1, n-1)$, $I_Q^{n-1} = \{1\}$.

As a consequence, affinity being stable by composition, $x \mapsto P_0^{n-1}(\chi^{n-1}(Q, x), \mathcal{F}_{1+})$ is affine on $I_Q^{n-1} \in \{[0, 1], \{0\}, \{1\}\}$. In turn,

$x \mapsto xV_0 + \mathbb{E}(P_0^{n-1}(\chi^{n-1}(Q, x), \mathcal{F}_{1+}) | \mathcal{F}_0)$ is affine and reaches its maximum at some endpoint of I_Q^{n-1} i.e. $q_0^* = 0$ or at $q_0^* = 1$. Then, inspecting the above cases shows that $Q^{1,*} = Q - q_0^*(1, 1) \in T^+(n-1) \cap \mathbb{N}^2$. Using (4.2) and (4.3), one shows by induction on k that q_k^* is always $\{0, 1\}$ -valued.

- NON NEGATIVE V_i :

Step 1: Global constraint saturated. Let $n \geq 1$. Let $q^* = (q_k^*)_{0 \leq k \leq n-1}$ be an optimal Q -admissible control. We introduce the \mathcal{F} -stopping time

$$\tau(q^*) := \min \{k \mid q_0^* + \dots + q_k^* < Q_{\max} - (n-1) + k\}$$

with the convention $\min \emptyset = +\infty$.

On $\tau(q^*) = +\infty$, $q_0^* + \dots + q_k^* \geq Q_{\max} - (n-1) + k$ for every $k = 0, \dots, n-1$. In particular the global constraint is saturated at time $n-1$, i.e.

$$q_0^* + \dots + q_{n-1}^* = Q_{\max}.$$

Set

$$\tilde{q}_k = q_k^* \mathbf{1}_{\{k \leq \tau(q^*)-1\}} + (Q_{\max} - (n-1) + \tau(q^*) - (q_0^* + \dots + q_{\tau(q^*)-1}^*)) \mathbf{1}_{\{k = \tau(q^*)\}}.$$

One check that \tilde{q} is a Q -admissible control: this follows from the fact that $\tau(q^*)$ is an \mathcal{F} -stopping time (note that q_τ^* is $[0, 1]$ -valued on $\{\tau(q^*) < +\infty\}$ since then $q_0^* + \dots + q_{\tau-1}^* \geq Q_{\max} - n + \tau$ and $q_0^* + \dots + q_\tau^* < Q_{\max} - (n-1) + \tau$). Likewise one shows that $\tilde{q}_k \geq q_k^*$ for every $k = 0, \dots, n-1$.

Furthermore, note that if Q_{\max} is an integer and q^* is $\{0, 1\}$ -valued then \tilde{q} is still $\{0, 1\}$ -valued. The V_k being non negative

$$\sum_{k=0}^{n-1} \tilde{q}_k V_k \geq \sum_{k=0}^{n-1} q_k^* V_k$$

hence \tilde{q} is still an optimal control. Furthermore $\tilde{q}_0 + \dots + \tilde{q}_k \geq Q_{\max} - (n-1) + k$ on $\{k \geq \tau(q^*)\}$ so that the stopping time $\tau(\tilde{q})$ satisfies by construction

$$\tau(\tilde{q}) \leq \tau(q^*) - 1 \quad \text{on } \{1 \leq \tau(q^*) < +\infty\} \quad \text{and} \quad \tau(\tilde{q}) = +\infty \quad \text{on } \{\tau(q^*) = 0\} \cup \{\tau(q^*) = +\infty\}.$$

Note that if the control q^* is bang-bang, iterating the above construction at most n times yields an optimal control q^{opt} for which $\tau(q^{opt}) = +\infty$ *a.s.*. Such a control q^{opt} saturates the global constraint.

As a consequence, this shows that $P_0(Q, \mathcal{F}) = P_0((0, Q_{\max}), \mathcal{F}) = P_0((Q_{\max}, Q_{\max}), \mathcal{F})$ *a.s.* so that $Q_{\max} \mapsto P_0((Q_{\max}, Q_{\max}), \mathcal{F})$ is *a.s.* non-decreasing and concave.

Step 2: Local constraints. Since there is an optimal control q^* which saturates the global constraint, one may assume without loss of generality that $Q_{\min} = Q_{\max}$. We proceed again by induction on n based on the dynamic programming formula (5.4). When $n = 1$ the result is obvious (and true when $n = 2$ as well).

Assume now the announced result is true for $n \geq 1$.

Let $j \in \{0, \dots, n-1\}$ and $Q_{\max} \in [j, j+1]$. Then, $I_Q^n = [0, Q_{\max} \wedge 1]$ and $\chi^n((Q_{\max}, Q_{\max}), x) = (Q_{\max} - x, Q_{\max} - x)$, $x \in I_Q^n$. Hence $\chi^n((Q_{\max}, Q_{\max}), x) \in T_{jj}^+$, $x \in [0, Q_{\max} - j]$ and $\chi^n((Q_{\max}, Q_{\max}), x) \in T_{j-1, j-1}^+$, $x \in [Q_{\max} - j, Q_{\max} \wedge 1]$. Now $v \mapsto P_0^n((v, v), \mathcal{F}_{1+})$ is *a.s.* concave, non-decreasing, affine on $[j-1, j]$ and on $[j, j+1]$ and non-decreasing. Consequently, there exists B^m, C^m , $m = 1, 2$, \mathcal{F}_1 -measurable random variables satisfying

$$\begin{aligned} P_0(\chi^n((Q_{\max}, Q_{\max}), x), \mathcal{F}_{1+}) &= B^1(Q_{\max} - x) + C^1, \quad x \in [0, Q_{\max} - j], \\ P_0(\chi^n((Q_{\max}, Q_{\max}), x), \mathcal{F}_{1+}) &= B^2(Q_{\max} - x) + C^2, \quad x \in [Q_{\max} - j, Q_{\max} \wedge 1]. \end{aligned}$$

with $0 \leq B^1 \leq B^2$ and $B^2 j + C^2 = B^1 j + C^1$ *a.s.*. Set temporarily

$$\Psi(x) := xV_0 + \mathbb{E}(P_0(\chi^n((Q_{\max}, Q_{\max}), x), \mathcal{F}_{1+}) | \mathcal{F}_0).$$

Hence,

$$\sup_{x \in I_Q^n} \Psi(x) = \max(\Psi(0), \Psi(Q_{\max} - j), \Psi(Q_{\max} \wedge 1)).$$

Set $B_0^m := \mathbb{E}(B^m | \mathcal{F}_0)$ and $C_0^m := \mathbb{E}(C^m | \mathcal{F}_0)$ and note that $B_0^1 \leq B_0^2$ and $B_0^2 j + C_0^2 = B_0^1 j + C_0^1$ *a.s.*. Elementary computations show that:

- $\Psi(0) \leq \Psi(Q_{\max} - j)$ on $\{V_0 \geq B_0^1\}$ and $\Psi(0) \geq \Psi(Q_{\max} - j)$ on $\{V_0 \leq B_0^1\}$,
- $\Psi(Q_{\max} - j) \leq \Psi(Q_{\max} \wedge 1)$ on $\{V_0 \geq B_0^2\}$ and $\Psi(Q_{\max} - j) \geq \Psi(Q_{\max} \wedge 1)$ on $\{V_0 \leq B_0^2\}$.

Consequently q_0^* can be chosen $\{0, Q_{\max} \wedge 1\}$ -valued on $E_0 := \{B_0^2 < V_0 \leq B_0^1\} \in \mathcal{F}_0$ and equal to $Q_{\max} - j$ on ${}^c E_0 := \{B_0^1 < V_0 \leq B_0^2\} \in \mathcal{F}_0$.

On $E_0^1 = E_0 \cap \{q_0^* = Q_{\max} \wedge 1\} \in \mathcal{F}_0$, one has $P_0^{n+1}((Q_{\max}, Q_{\max}), \mathcal{F}) = P_0^{n+1}((Q_{\max}, Q_{\max}), \mathcal{F} \cap E_0^1)$.

Then, the dynamic programming formula shows that the other components $(q_k^*)_{1 \leq k \leq n}$ of the optimal control on E_0 can be obtained as the optimal control of the pricing problem $P_0^n(((Q_{\max} - 1)^+, (Q_{\max} - 1)^+), (\mathcal{F}_{1+k} \cap E_0^1)_{0 \leq k \leq n-1})$. One derives from the induction assumption at time n that $(q_k^*)_{1 \leq k \leq n}$ can be chosen bang-bang and $((Q_{\max} - 1)^+, (Q_{\max} - 1)^+)$ -admissible which implies that q is Q -admissible and bang-bang since $q_0^* = Q_{\max} \wedge 1$ (on E_0^1). A similar proof holds on $E_0^0 = E_0 \cap \{q_0^* = 0\}$.

On ${}^c E_0$, one has $P_0^{n+1}((Q_{\max}, Q_{\max}), \mathcal{F}) = P_0^{n+1}((Q_{\max}, Q_{\max}), \mathcal{F} \cap {}^c E_0)$. Then, the dynamic programming formula shows that the other components $(q_{1+k}^*)_{0 \leq k \leq n-1}$ of the optimal control on ${}^c E_0$ can be obtained as the optimal control of the pricing problem $P_0^n((j, j), (\mathcal{F}_{1+k} \cap {}^c E_0)_{0 \leq k \leq n-1})$. As $(j, j) \in \mathbb{N}^2$ there exists a (j, j) -admissible bang-bang optimal control $(q_k^*)_{0 \leq k \leq n-1}$ (with respect to $(\mathcal{F}_{1+k} \cap {}^c E_0)_{0 \leq k \leq n-1}$) on ${}^c E_0$. Then q_{1+k}^* is $\{0, 1\}$ -valued for every $k = 0, \dots, n-1$ (in fact identically 0 if $j = 0$). At this stage one can recursively modify $(q_{1+k}^*)_{0 \leq k \leq n-1}$ using the procedure described in Step 1 to saturate the upper global constraint. Finally one may assume that $\sum_{0 \leq k \leq n-1} q_k^* = j$ which in turn implies that $(q_k^*)_{0 \leq k \leq n}$ is a bang-bang (Q_{\max}, Q_{\max}) -admissible optimal control. \diamond

APPLICATION. When a global constraint Q belongs to the interior of a triangle T_{ij}^\pm , *one only needs to compute the value of $P_0(\cdot, \mathcal{F})$ at the vertices of this triangle to derive the value of the premium at every $Q \in T_{ij}^\pm$* . When Q is itself an integral valued couple, at most six further points allow to compute the premium in a neighborhood of Q . We will use this result extensively when designing our quantization based numerical procedure in Section 6.

AN ADDITIONAL RESULT. Proposition 2 shows that it is hopeless to produce in full generality some bang-bang optimal control when $Q_{\max} - Q_{\min} \in \mathbb{N}$. This comes from the fact that at integral valued global constraints the bang-bang optimal control may saturate none of the global constraints (indeed, so is the case at $(2, 2)$ when $n = 2$). However, using the same approach as that developed in that in the proof of Theorem 2, one can show the following result, whose details of proof are left to the reader.

Corollary 1 *Assume the assumptions of Theorem 2 hold. If a couple of admissible constraint $(Q_{\min}, Q_{\max}) \in T^+(n)$ satisfies*

$$Q_{\max} - Q_{\min} \in \{0, \dots, n\}$$

then there exist a quasi-bang-bang control in the following sense: \mathbb{P} -a.s., q_k^ is $\{0, 1\}$ -valued except for at most one local constraint $q_{k_0}^*$.*

5.4 The Markov setting

By Markov setting we simply mean that the payoffs V_k are function of an \mathbb{R}^d -valued underlying \mathcal{F} -Markov structure process $(Y_k)_{0 \leq k \leq n-1}$ *i.e.*

$$V_k = v_k(Y_k), \quad k = 0, \dots, n-1.$$

The Markovian dynamics of Y reads on Borel functions $g : \mathbb{R} \rightarrow \mathbb{R}$

$$\mathbb{E}(g(Y_{k+1}) | \mathcal{F}_k) = \mathbb{E}(g(Y_{k+1}) | Y_k) = \Theta_k(g)(Y_k)$$

where $(\Theta_k)_{0 \leq k \leq n-1}$ is sequence of Borel probability transitions on $(\mathbb{R}^d, \mathbb{B}(\mathbb{R}^d))$.

Then the backward dynamic programming principle (4.1) can be rewritten as follows

$$P_k^n(Q) = p_k^n(Q, Y_k), \quad k = 0, \dots, n$$

with $p_n^n(\cdot, \cdot) \equiv 0$ and for every $k = 0, \dots, n-1$ and every $Q \in T^+(n-k)$,

$$p_k^n(Q, y) = \sup \left\{ x v_k(y) + \Theta_k(p_{k+1}^n(\chi^{n-k-1}(Q, x, \cdot)))(y), x \in I_Q^{n-k-1} \right\}, \quad (5.1)$$

where $\chi^M(Q, x) = ((Q_{\min} - x)^+, (Q_{\max} - x) \wedge M)$ and $I_Q^M := [(Q_{\min} - M)^+ \wedge 1, Q_{\max} \wedge 1]$.

Pointwise estimation of $P_0^n(Q^0)$. As established in Theorem 2, one only needs to compute the value function $P_0^n(Q)$ at global constraints $Q = (Q_{\min}, Q_{\max}) \in T^+(n) \cap \mathbb{N}^2$ *i.e.* with integral components. Moreover, for these constraints, the local optimal control q_k^* is always bang-bang *i.e.* $q_k^* \in \{0, 1\}$.

Let $Q^0 = (Q_{\min}^0, Q_{\max}^0) \in T^+(n) \cap \mathbb{N}^2$. For every $k = 0, \dots, n-1$, one defines the set of attainable residual global constraints at time k , namely

$$\mathcal{Q}_k^n(Q^0) := \{((Q_{\min}^0 - \ell)^+, (Q_{\max}^0 - \ell)^+ \wedge (n-k)), \ell = 0, \dots, k\}. \quad (5.2)$$

(thus $\mathcal{Q}_0^n(Q^0) = \{Q^0\}$). Note that the running parameter ℓ represents the possible values of the cumulative purchase process \tilde{q}_k^* .

One checks that for every $Q \in \mathcal{Q}_k^n(Q^0)$, $\chi^{n-k-1}(Q, 1)\chi^{n-k-1}(Q, 0) \in \mathcal{Q}_{k+1}^n(Q^0)$ since

$$\chi^{n-k-1}(Q, 1) = ((Q_{\min} - 1)^+, (Q_{\max} - 1)^+) \quad \text{and} \quad \chi^{n-k-1}(Q, 0) = (Q_{\min}, Q_{\max} \wedge (n-k-1)).$$

Consequently the backward dynamic programming formula having $p_0^n(Q^0, y)$ as a result reads:

$$p_k^n(Q, y) = \max \left\{ x v_k(y) + \Theta_k(p_{k+1}^n(\chi^{n-k-1}(Q, x, \cdot)))(y), x \in \{0, 1\} \cap I_Q^{n-1-k} \right\}, Q \in \mathcal{Q}_k^n(Q^0), k = 0, \dots, n-1.$$

At this stage no numerical computation is possible yet since no space discretization has been achieved. This is the aim of the next section where we will approximate the above dynamic programming principle by (optimal) quantization of the state process Y .

6 Computing swing contracts by (optimal) quantization

6.1 The abstract quantization tree approach

Abstract quantization In this section, we propose a quantization based approach to compute the premium of the swing contracts with firm constraints. Quantization has been originally introduced and developed in the early 1950' in Signal processing (see [13]). The starting idea is simply to replace every random vector $Y : (\Omega, \mathcal{A}) \rightarrow \mathbb{R}^d$ by a random vector $\hat{Y} = g(Y)$ taking finitely many values in a *grid* (or *codebook*) $\Gamma := \{y_1, \dots, y_N\}$ (of size N). The grid Γ is also called an N -quantizer of Y . When the Borel function g satisfies

$$|Y - \hat{Y}| = d(Y, \Gamma) = \min_{1 \leq i \leq N} |Y - y^i|, \quad (6.3)$$

\hat{Y} is called a *Voronoi quantization* of Y (and g as well). One easily checks that g is necessarily a nearest neighbor projection on Γ *i.e.* satisfies

$$\forall i \in \{1, \dots, N\}, \quad \{g = y^i\} \subset \{u \in \mathbb{R}^d : |u - y^i| = \min_{1 \leq j \leq N} |u - y^j|\}.$$

The so-called *Voronoi cells* $\{g = y^i\}$, $1 \leq i \leq N$, make up a *Voronoi tessellation* or partition of \mathbb{R}^d (induced by Γ . Note that when the distribution \mathbb{P}_Y of Y weights no hyperplanes the boundary of the Voronoi tessellation of Γ are \mathbb{P}_Y -negligible so that the \mathbb{P}_Y -weights of the Voronoi cells entirely characterize the distribution of \hat{Y}^Γ .

When $p \in [1, \infty)$, the L^p -mean error induced by replacing Y by \widehat{Y} , namely

$$\|Y - \widehat{Y}\|_p = \left(\mathbb{E} \left(\min_{1 \leq i \leq N} |Y - y^i|^p \right) \right)^{\frac{1}{p}}$$

is called the L^p -mean quantization error induced by Γ_k and its p^{th} power is known as the L^p -distortion. We will see in the next section that the codebook Γ can be optimized so as to minimize the L^p -quantization error with respect to the (distribution of) Y .

Our aim in this section is to design an algorithm based on the quantization of the Markov chain (Y_k) at every time k to approximate the premium of the swing contract with firm constraints and to provide some *a priori* error estimates in terms of quantization errors.

Quantized tree for pricing swing options. As a first step we consider at every time k a grid $\Gamma_k := \{y_k^1, \dots, y_k^{N_k}\}$ (of size N_k). Then, we design the quantized tree algorithm to price swing contracts by simply mimicking the original dynamic programming formula (4.1). This means in particular that we force in some way the Markov property on $(\widehat{Y}_k)_{0 \leq k \leq n-1}$ by considering the quantized transition operator

$$\widehat{\Theta}_k(g)(y_k^i) = \sum_{j=1}^{N_{k+1}} g(y_{k+1}^j) p_k^{ij}, \quad p_k^{ij} := \mathbb{P}(\widehat{Y}_{k+1} = y_{k+1}^j \mid \widehat{Y}_k = y_k^i)$$

so that

$$\widehat{\Theta}_k(g)(\widehat{Y}_k) = \mathbb{E} \left(g(\widehat{Y}_{k+1}) \mid \widehat{Y}_k \right), \quad k = 0, \dots, n-1.$$

▷ Let $Q^0 \in T^+(n) \cap \mathbb{N}^2$ be a couple of (deterministic) global constraints (at time 0). The quantized dynamic programming principle is defined by

$$\begin{aligned} \widehat{P}_n^n(Q) &:= 0, \quad Q \in T^+(n) \cap \mathbb{N}^2, \\ \widehat{P}_k^n(Q) &:= \max \left(x v_k(\widehat{Y}_k) + \mathbb{E} \left(\widehat{P}_{k+1}^n(\chi^{n-k-1}(Q, x)) \mid \widehat{Y}_k \right), x \in \partial I_Q^{n-1-k} \right), \quad Q \in \mathcal{Q}_k^n(Q^0), \quad k = 0, \dots, n-1. \end{aligned} \quad (6.4)$$

One easily shows by induction that, for every $Q \in \mathcal{Q}_k^n(Q^0)$ (residual global constraint at time k),

$$\widehat{P}_k^n(Q) = \widehat{p}_k^n(Q, \widehat{Y}_k)$$

where $\widehat{p}_n^n(Q, y) = 0$, $Q \in \mathcal{Q}_n^n(Q^0)$, $y \in \mathbb{R}^d$,

$$\begin{aligned} \widehat{p}_k^n(Q, y_k^i) &= \sup \left\{ x v_k(y_k^i) + \widehat{\Theta}_k(\widehat{p}_{k+1}^n(\chi^{n-k-1}(Q, x, \cdot)))(y_k^i), x \in \partial I_Q^{n-1-k} \right\} \\ & \quad i = 1, \dots, N_k, \quad Q \in \mathcal{Q}_k^n(Q^0), \quad k = 0, \dots, n-1. \end{aligned} \quad (6.5)$$

▷ When $Q^0 \in T^+(n) \setminus \mathbb{N}^2$, one defines $\widehat{P}_0^n(Q)$ (and $\widehat{p}_0^n(Q, \cdot)$) by affinity on each elementary triangle T_{ij}^\pm that tiles $T^+(n)$.

Complexity Let us briefly discuss the complexity of this quantized backward dynamic procedure. Let $k \in \{0, \dots, n-1\}$. At every “nod” y_k^i the computation of $\widehat{\Theta}_k(\widehat{p}_{k+1}^n(\chi^{n-k-1}(Q, x, \cdot)))(y_k^i)$ requires N_{k+1} products (up to a constant), so that for a given residual global constraint the complexity at time k in the dynamic programming is proportional to $N_k N_{k+1}$. On the other hand, one checks that

$$\text{card}(\mathcal{Q}_k^n(Q^0)) = (Q_{\max}^0 \wedge k) + 1 - (Q_{\max}^0 - Q_{\min}^0 - (n - k) - 1)^+.$$

Consequently, the complexity of the computation of $\widehat{p}_0^n(Q^0, \widehat{Y}_0)$ is proportional to

$$\sum_{k=0}^{n-1} \text{card}(\mathcal{Q}_k^n(Q^0)) N_k N_{k+1}.$$

A simple upper-bound is provided by

$$\sum_{k=0}^{n-1} ((Q_{\max}^0 \wedge k) + 1) N_k N_{k+1}$$

and a uniform one by

$$\sum_{k=0}^{n-1} (k + 1) N_k N_{k+1}.$$

Note that this last upper bound corresponds to the complexity of the quantized version of the algorithm based on some *penalized global volume constraints* (see the companion paper [5]).

A priori error bounds for the quantized procedure

Theorem 3 *Assume that the Markov process $(Y_k)_{0 \leq k \leq n-1}$ is Lipschitz Feller in the following sense: for every bounded Lipschitz continuous function $g : \mathbb{R}^d \rightarrow \mathbb{R}$ and every $k \in \{0, \dots, n-1\}$, $\Theta_k(g)$ is a Lipschitz function satisfying $[\Theta_k(g)]_{\text{Lip}} \leq [\Theta_k]_{\text{Lip}} [g]_{\text{Lip}}$. Assume that every function $v_k : \mathbb{R}^d \rightarrow \mathbb{R}$ is Lipschitz continuous with Lipschitz coefficient $[v_k]_{\text{Lip}}$. Let $p \in [1, \infty)$ such that $\max_{0 \leq k \leq n-1} |Y_k| \in L^p(\mathbb{P})$. Then, there exists a real constant $C_p > 0$ such that*

$$\| \sup_{Q \in T_{\mathbb{N}}^+(n)} |\widehat{P}_0^n(Q) - P_0^n(Q)| \|_p \leq C_p \sum_{k=0}^{n-1} \|Y_k - \widehat{Y}_k\|_p \quad (6.6)$$

Remark. In most situations $\mathcal{F}_0 = \{\emptyset, \Omega\}$ so that the error term $|\widehat{P}_0^n(Q) - P_0^n(Q)|$ is deterministic. When \mathcal{F}_0 is not trivial, it is straightforward from (6.6) (with $p = 1$) that

$$\sup_{Q \in T^+(n)} |\mathbb{E}(\widehat{P}_0^n(Q)) - \mathbb{E}(P_0^n(Q))| \leq C_1 \sum_{k=0}^{n-1} \|Y_k - \widehat{Y}_k\|_1.$$

We first need a lemma about the Lipschitz regularity of the p_k^n functions.

Lemma 6.1 *For every $k \in \{0, \dots, n-1\}$, the function $y \mapsto p_k^n(Q, y)$ is Lipschitz on \mathbb{R}^d , uniformly with respect to $Q \in T^+(n-k)$ and its Lipschitz coefficient $[p_k^n]_{\text{Lip}, y} := \sup_{Q \in T^+(n-k)} [p_k^n(Q, \cdot)]_{\text{Lip}}$ satisfies for every $k \in \{0, \dots, n-1\}$,*

$$[p_{n-1}^n]_{\text{Lip}, y} \leq [v_{n-1}]_{\text{Lip}}, \quad [p_k^n]_{\text{Lip}, y} \leq [v_k]_{\text{Lip}} + [\Theta_k]_{\text{Lip}} [p_{k+1}^n]_{\text{Lip}, y}.$$

Proof. This follows easily by a backward induction on k , based on the dynamic programming formula (5.1) and the elementary inequality $|\sup_{i \in I} a_i - \sup_{i \in I} b_i| \leq \sup_{i \in I} |a_i - b_i|$. \diamond

Proof of Theorem 3. First note that, by piecewise affinity of \widehat{P}_0^n and P_0^n , one has

$$\sup_{Q \in T_{\mathbb{N}}^+(n)} |\widehat{P}_0^n(Q) - P_0^n(Q)| = \sup_{Q \in T_{\mathbb{N}}^+(n) \cap \mathbb{N}^2} |\widehat{P}_0^n(Q) - P_0^n(Q)|.$$

Temporarily set $T_{\mathbb{N}}^+(n) := T^+(n) \cap \mathbb{N}$. Let $k \in \{0, \dots, n-1\}$. Now

$$\begin{aligned} \sup_{Q \in T_{\mathbb{N}}^+(n-k)} |p_k^n(Q, Y_k) - \widehat{p}_k^n(Q, \widehat{Y}_k)| &\leq |v_k(Y_k) - v_k(\widehat{Y}_k)| \\ &+ \sup_{Q \in T_{\mathbb{N}}^+(n-k), x \in \partial I_Q^{n-1-k}} \left| \mathbb{E}(p_{k+1}^n(\chi^{n-1-k}(Q, x), Y_{k+1}) | \mathcal{F}_k) - \mathbb{E}(\widehat{p}_{k+1}^n(\chi^{n-1-k}(Q, x), \widehat{Y}_{k+1}) | \widehat{Y}_k) \right|. \end{aligned} \quad (6.7)$$

Now, using that Θ_k is a Markov transition and that $\sigma(\widehat{Y}_k) \subset \sigma(Y_k)$ -measurable, one gets

$$\begin{aligned} \mathbb{E} \left(p_{k+1}^n(\chi^{n-1-k}(Q, x), Y_{k+1}) | \mathcal{F}_k \right) &- \mathbb{E} \left(\widehat{p}_{k+1}^n(\chi^{n-1-k}(Q, x), \widehat{Y}_{k+1}) | \widehat{Y}_k \right) \\ &= \Theta_k(p_{k+1}^n(\chi^{n-1-k}(Q, x), \cdot))(Y_k) - \mathbb{E} \left(\Theta_k(p_{k+1}^n(\chi^{n-1-k}(Q, x), \cdot))(Y_k) | \widehat{Y}_k \right) \\ &+ \mathbb{E} \left(p_{k+1}^n(\chi^{n-1-k}(Q, x), Y_{k+1}) - \widehat{p}_{k+1}^n(\chi^{n-1-k}(Q, x), \widehat{Y}_{k+1}) | \widehat{Y}_k \right) \\ &= \Theta_k(p_{k+1}^n(\chi^{n-1-k}(Q, x), \cdot))(Y_k) - \Theta_k(p_{k+1}^n(\chi^{n-1-k}(Q, x), \cdot))(\widehat{Y}_k) \\ &+ \mathbb{E} \left(\Theta_k(p_{k+1}^n(\chi^{n-1-k}(Q, x), \cdot))(\widehat{Y}_k) - \Theta_k(p_{k+1}^n(\chi^{n-1-k}(Q, x), \cdot))(Y_k) | \widehat{Y}_k \right) \\ &+ \mathbb{E} \left(p_{k+1}^n(\chi^{n-1-k}(Q, x), Y_{k+1}) - \widehat{p}_{k+1}^n(\chi^{n-1-k}(Q, x), \widehat{Y}_{k+1}) | \widehat{Y}_k \right). \end{aligned}$$

Consequently, still using the elementary inequality $|\sup_{x \in X} a_x - \sup_{x \in X} b_x| \leq \sup_{x \in I} |a_x - b_x|$ for any index set X and, for every $x \in \partial I^{n-1-k}$, that

$$\chi^{n-k-1}(T_{\mathbb{N}}^+(n-k), x) \subset T_{\mathbb{N}}^+(n-k-1)$$

(see the proof of Theorem 2(b)), one has

$$\begin{aligned} \sup_{Q \in T_{\mathbb{N}}^+(n-k)} |p_k^n(Q, Y_k) - \widehat{p}_k^n(Q, \widehat{Y}_k)| &\leq |v_k(Y_k) - v_k(\widehat{Y}_k)| \\ &+ \sup_{Q' \in T_{\mathbb{N}}^+(n-k-1)} |\Theta_k(p_{k+1}^n(Q', \cdot))(Y_k) - \Theta_k(p_{k+1}^n(Q', \cdot))(\widehat{Y}_k)| \\ &+ \mathbb{E} \left(\sup_{Q' \in T_{\mathbb{N}}^+(n-k-1)} |\Theta_k(p_{k+1}^n(Q', \cdot))(\widehat{Y}_k) - \Theta_k(p_{k+1}^n(Q', \cdot))(Y_k)| | \widehat{Y}_k \right) \\ &+ \mathbb{E} \left(\sup_{Q' \in T_{\mathbb{N}}^+(n-k-1)} |p_{k+1}^n(Q', Y_{k+1}) - \widehat{p}_{k+1}^n(Q', \widehat{Y}_{k+1})| | \widehat{Y}_k \right). \end{aligned}$$

Temporarily set for convenience, $\Delta_k^{n,p} := \|\sup_{Q \in T_{\mathbb{N}}^+(n-k)} |p_k^n(Q, Y_k) - \widehat{p}_k^n(Q, \widehat{Y}_k)|\|_p$. One derives that for every $k = 0, \dots, n-1$,

$$\Delta_k^{n,p} \leq ([v_k]_{\text{Lip}} + 2[\Theta_k]_{\text{Lip}}[p_{k+1}^n]_{\text{Lip},y}) \|Y_k - \widehat{Y}_k\|_p + \Delta_{k+1}^{n,p}.$$

Furthermore, $\Delta_{n-1}^{n,p} \leq [v_{n-1}]_{\text{Lip}} \|Y_{n-1} - \widehat{Y}_{n-1}\|_p$. The result follows by induction. \diamond

6.2 Optimal quantization

Theoretical background In this section, we provide a few basic elements about optimal quantization in order to give some error bounds for the premium of the swing option. We refer to [15] for more details about theoretical aspects and to [26] for the algorithmic aspects numerical applications.

Let $p \in [1, +\infty)$. let $Y \in L^p(\Omega, \mathcal{A}, \mathbb{P})$ be an \mathbb{R}^d -valued random vector and let $N \geq 1$ be a given grid size. The best L^p -approximation of Y by a random vector taking its values in a given grid Γ

of size (at most) N is given by a Voronoi quantizer \widehat{Y}^Γ which induces an $L^p(\mathbb{P})$ - mean quantization error

$$e_{N,p}(Y, \Gamma) = \|Y - \widehat{Y}^\Gamma\|_p = \left(\mathbb{E} \min_{y \in \Gamma} |Y - y|^p \right)^{\frac{1}{p}}.$$

It has been shown independently by several authors (in various finite and infinite dimensional frameworks) that when the grids Γ runs over all the subsets of \mathbb{R}^d of size at most N , that $e_{N,p}(Y, \Gamma)$ reaches a minimum denoted $e_{N,p}(Y)$ (see e.g. [15] or [23]) i.e. the minimization problem

$$e_{N,p}(Y) = \min \left\{ e_{N,p}(Y, \Gamma), \Gamma \subset \mathbb{R}^d, \text{card}(\Gamma) \leq N \right\}$$

has at least a solution temporarily denoted $\Gamma^{(N,*)}$. Several algorithms have been designed to compute some optimal or close to optimality quantizers, especially in the quadratic case $p = 2$. They all rely on the stationarity property satisfied by optimal quantizers. In the quadratic case, a grid Γ is stationary

$$\widehat{Y}^\Gamma = \mathbb{E} \left(Y \mid \widehat{Y}^\Gamma \right)$$

This follows from some differentiability property of the L^p -distortion. For a formula in the general case we refer to [16]. In 1-dimension, a regular Newton-Raphson zero search procedure turns out to be quite efficient. In higher dimension (at least when $d \geq 3$ or 4) only stochastic procedures can be implemented like the *CLVQ* (a stochastic gradient descent, see [23] or the Lloyd I procedure (a randomized fixed point procedure, see [13]). For more details and result we refer to [26].

As a result of these methods, some optimized grids of the (centered) normal distribution $\mathcal{N}(0; I_d)$ are available on line at the URL

www.quantize.maths-fi.com

for dimensions $d = 1, \dots, 10$ and sizes from $N = 2$ up to 5 000.

It is clear by considering a sequence of grids $\Gamma^{(N)} := \{r^1, \dots, r^N\}$ where $(r^n)_{n \geq 1}$ is an everywhere dense sequence in \mathbb{R}^d that $e_{N,p}(Y)$ decreases to 0 as $N \rightarrow \infty$.

The rate of convergence of this sequence is ruled by the so-called Zador Theorem (see [30] for a first statement of the result, until the first rigorous proof in [15]).

Theorem 4 (Zador, see [15]) (a) Let $Y \in L^{p+\eta}(\mathbb{P})$, $p \geq 1$, $\eta > 0$, such that $\mathbb{P}_Y(du) = \varphi(du)du + \nu(du)$. Then

$$\lim_N N^{\frac{1}{d}} e_{N,p}(Y) = \widetilde{J}_{p,d} \left(\int_{\mathbb{R}^d} \varphi^{\frac{d}{p+d}}(u) du \right)^{\frac{1}{p} + \frac{1}{d}}.$$

(b) Non asymptotic estimate (see e.g. [21]): Let $p \geq 1$, $\eta > 0$. There exists a real constant $C_{d,p,\eta} > 0$ and an integer $N_{d,p,\eta} \geq 1$ such that for any \mathbb{R}^d -valued random vector Y , for $N \geq N_{d,p,\eta}$,

$$e_{N,p}(Y) \leq C_{d,p,\eta} \|Y\|_{p+\eta} N^{-\frac{1}{d}}.$$

Rate of convergence of the quantization pricing method Now we are in position to apply the above results to provide an error bound for the pricing of swing options by optimal quantization: assume there is a real exponent $p \in [1, +\infty)$ such that the (d -dimensional) Markov structure process $(Y_k)_{0 \leq k \leq n-1}$ satisfies

$$\max_{0 \leq k \leq n-1} |Y_k| \in L^{p+\eta}(\mathbb{P}), \quad \eta > 0$$

At each time $k \in \{0, \dots, n-1\}$, we implement a (quadratic) optimal quantization grid $\Gamma^{\bar{N}}$ of Y_k with constant size \bar{N} . Then the general error bound result (6.6) combined with Theorem 4(b) says that, if $\bar{N} \geq N_{d,p,\eta}$,

$$\left\| \sup_{Q \in T^+(n)} |P_0^n(Q) - \widehat{P}_0^n(Q)| \right\|_p \leq C \frac{n}{\bar{N}^{\frac{1}{d}}}$$

where as the complexity of the procedure is bounded by $n(n+1)\bar{N}^2$ (up to a constant).

In fact this error bound turns out to be conservative and several numerical experiments, as those presented below, suggest that in fact the true rate (for a fixed number n of purchase instants) behaves like $O(\bar{N}^{-\frac{2}{d}})$.

Another approach could be to minimize the complexity of the procedure by considering (optimal) grids with variable sizes N_k satisfying $\sum_{k=0}^{n-1} N_k = n\bar{N}$. We refer to [5] for further results in that direction. However, numerical experiments were carried out with constant size grids for both programming convenience and memory saving.

6.3 A numerical illustration

We considered a two factor continuous model for the price of future contracts which leads to the following dynamics for the spot price

$$S_t = F_{0,t} \exp \left(\sigma_1 \int_0^t e^{-\alpha_1(t-s)} dW_s^1 + \sigma_2 \int_0^t e^{-\alpha_2(t-s)} dW_s^2 - \frac{1}{2} \Lambda_t \right), \quad t \in [0, T],$$

where W^1 and W^2 are two standard Brownian motions with correlation coefficient ρ and

$$\Lambda_t = \frac{\sigma_1^2}{2\alpha_1} (1 - e^{-2\alpha_1 t}) + \frac{\sigma_2^2}{2\alpha_2} (1 - e^{-2\alpha_2 t}) + \frac{2\rho\sigma_1\sigma_2}{\alpha_1 + \alpha_2} (1 - e^{-(\alpha_1 + \alpha_2)t})$$

Then, we consider a (daily) discretization of the Gaussian process $\log(S_t/F_{0,t})$ at times $\frac{kT}{n}$ ($T = 1$, $n = 365$). The sequence $(\log(S_{t_k}/F_{0,t_k}))_{0 \leq k \leq n-1}$ is clearly not Markov. However, adding an appropriate auxiliary processes, one can build a higher dimensional (homogenous) Markov process $(Y_k)_{0 \leq k \leq n-1}$ whom $\log(S_{t_k}/F_{0,t_k})$ is a linear combination. This calls upon classical methods coming from time series analysis. Then a fast quantization method has been developed to make possible a parallel implementation of the quantized probability transitions of $Y = (Y_k)_{0 \leq k \leq n-1}$. For further details about this model and the way it can be quantized, we refer to [5]. In [5], the optimal quantization method described above is extensively tested from a numerical viewpoint (rates of convergence, needed memory, swapping effect, etc). Its performances are compared those of the Longstaff-Schwartz approach introduced in [6]. This comparison emphasizes the accuracy and the velocity of our approach, even if only one contract is to be computed and the computation of the probability transitions is included in the computation time of the quantization method. Furthermore, it seems that it needs significantly less memory capacity when implemented on our tested model.

We simply reproduced below a complete graph of the function $Q := (Q_{\min}, Q_{\max}) \mapsto P_0^n(Q)$ when Q runs over the whole set of admissible global constraints $T^+(n)$. The parameters were settled at the following values

$$n = 30, \alpha_1 = 0.21, \alpha_2 = 5.4, \sigma_1 = 36\%, \sigma_2 = 111\%, \rho = -0.11$$

The graph of the premium function $Q \mapsto P_0^n(Q)$ defined on $T^+(n)$ is depicted in Figure 1.

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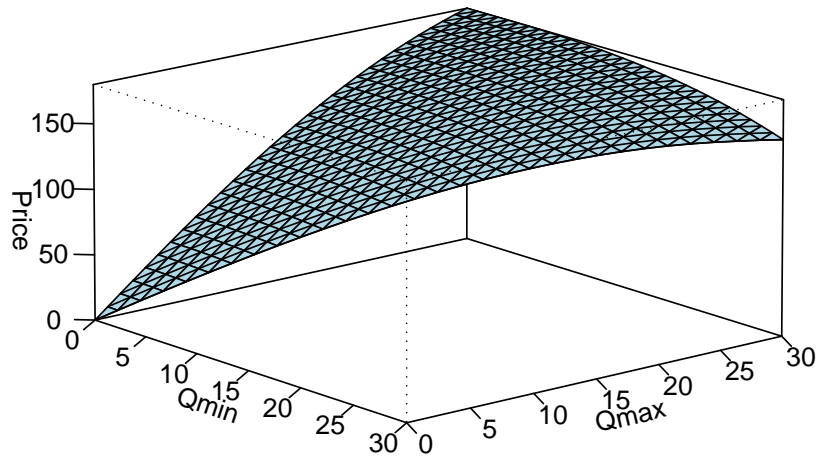


Figure 4: The mapping $Q \mapsto \hat{P}_0^n(Q)$ affinely interpolated from integral-valued global constraints.

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