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Following the traces

An introduction to conjoint measurement without transitivity and additivity¹

D. Bouyssou²
CNRS – LAMSADE

M. Pirlot³
Faculté Polytechnique de Mons

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² CNRS – LAMSADE, Université Paris Dauphine, Place du Maréchal de Lattre de Tassigny, F-75775 Paris Cedex 16, France, tel: +33 1 44 05 48 98, fax: +33 1 44 05 40 91, e-mail: bouyssou@lamsade.dauphine.fr. **Corresponding author.**

³ Faculté Polytechnique de Mons, 9, rue de Houdain, B-7000 Mons, Belgium, tel: +32 65 374682, fax: + 32 65 374689, e-mail: marc.pirlot@fpms.ac.be.

Abstract

This paper presents a self-contained introduction to a general conjoint measurement framework for the analysis of nontransitive and/or incomplete binary relations on product sets. It is based on the use of several kinds of marginal traces on coordinates induced by the binary relation.

This framework leads to defining three general families of models depending on the kind of trace that they use. Contrary to most conjoint measurement models, these models do not involve an addition operation. This allows for a simple axiomatic analysis at the cost of very weak uniqueness results.

Keywords: Multiple criteria analysis, Decision analysis, Measurement theory, Conjoint measurement, Nontransitive preferences.

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1 Introduction

Conjoint measurement (Krantz, Luce, Suppes, & Tversky, 1971; Wakker, 1989) studies binary relations defined on product sets. There are many situations in which such binary relations are of central interest. Among them let us mention:

- *Multiple criteria decision making* using a preference relation comparing alternatives evaluated on several attributes (see Belton & Stewart, 2001; Keeney & Raiffa, 1976; Roy, 1996; Winterfeldt & Edwards, 1986),
- *Decision under uncertainty* using a preference relation comparing alternatives evaluated on several states of nature (see Anscombe & Aumann, 1963; Fishburn, 1970b, 1988; Gul, 1992; Savage, 1954; Shapiro, 1979; Wakker, 1984, 1989),
- *Micro Economics* manipulating preference relations for bundles of several goods (see Debreu, 1959),
- *Dynamic decision making* using a preference relation between alternatives evaluated at several moments in time (see Keeney & Raiffa, 1976; Koopmans, 1960, 1972),
- *Social choice* comparing distribution of wealth across several individuals (see Atkinson, 1970; Ben-Porath & Gilboa, 1994; Ben-Porath, Gilboa, & Schmeidler, 1997; Weymark, 1981)

Given a binary relation \succsim on a set $X = X_1 \times X_2 \times \dots \times X_n$, the theory of conjoint measurement consists in finding conditions under which it is possible to build a convenient numerical representation of \succsim and to study the uniqueness of this representation. Manipulating numbers is clearly much easier than manipulating binary relations. Furthermore, in most cases, the proofs showing the existence of a numerical representation gives very useful hints on how to build this representation and, thus, assess preferences. Finally, the conditions on \succsim guaranteeing the existence of a numerical representation can be subjected to empirical tests (Krantz et al., 1971). This explains why conjoint measurement has attracted much attention in many different fields of research.

In traditional models of conjoint measurement the binary relation studied is most often supposed to be *complete* and *transitive* and the numerical representation is sought to be *additive*. The central model is the *additive*

utility model such that:

$$x \succsim y \Leftrightarrow \sum_{i=1}^n u_i(x_i) \geq \sum_{i=1}^n u_i(y_i) \quad (1)$$

where u_i are real-valued functions on the sets X_i and it is understood that $x = (x_1, x_2, \dots, x_n)$ and $y = (y_1, y_2, \dots, y_n)$.

The axiomatic analysis of this model is now quite firmly established (see Debreu, 1960; Krantz et al., 1971; Wakker, 1989) and additive utilities form the basis of many decision analysis techniques (see French, 1993; Keeney & Raiffa, 1976; Winterfeldt & Edwards, 1986; Wakker, 1989).

This central model mainly raises two types of difficulties. First it excludes all relations that would not be complete or transitive whereas the reasonableness of this hypothesis has been challenged by many authors (see May, 1954; Roy, 1996; Tversky, 1969). Second, its axiomatic analysis raises subtle technical difficulties. When X is finite (but of arbitrary cardinality), it is well-known (see Scott, 1964) that its characterization implies using a system of axioms involving a denumerable number of cancellation conditions guaranteeing the existence of solutions to a system of (finitely many) linear inequalities. Such an axiom system is hardly interpretable and testable. When X is infinite the picture changes provided that conditions are imposed in order to guarantee that the structure of X is “close” to the structure of \mathbb{R} and that \succsim behaves consistently in this continuum. This allows to consider only a limited number of cancellation conditions on \succsim . These necessary conditions however interact with the nonnecessary structural assumptions imposed on X (e.g. solvability Krantz et al. (1971, Chapter 6)), which may obscure their interpretation and test (see Krantz et al., 1971, Chapter 9, Furkhen & Richter, 1991). Furthermore, the analysis of the $n = 2$ case has then to be separated from that of the $n \geq 3$ case.

Several authors have forcefully argued in favor of studying conjoint measurement models that would tolerate intransitive and incomplete binary relations (see Bouyssou & Pirlot, 2002a; Fishburn, 1990a, 1991b, 1991a; May, 1954; Roy, 1996; Roy & Bouyssou, 1993; Tversky, 1969). The ones that we propose here will allow us to capture nontransitive and/or incomplete binary relations while being rather simple to analyze from an axiomatic point of view. In order to get a feeling for these various models, it is useful to envisage the various strategies that are likely to be implemented when comparing objects differing on several dimensions (see Dahlstrand & Montgomery, 1984; Montgomery, 1977; Montgomery & Svenson, 1976; Russo & Doshier, 1983; Svenson, 1979; Tversky, 1969).

Consider two alternatives x and y evaluated on a family of n attributes so that $x = (x_1, x_2, \dots, x_n)$ and $y = (y_1, y_2, \dots, y_n)$. A first strategy that can be

used in order to decide whether or not it can be said that “ x is at least as good as y ” consists in trying to measure the “worth” of each alternative on each attribute and then to combine these evaluations adequately. Abandoning all idea of transitivity and completeness, this suggests a model in which:

$$x \succsim y \Leftrightarrow F(u_1(x_1), u_2(x_2), \dots, u_n(x_n), u_1(y_1), u_2(y_2), \dots, u_n(y_n))) \geq 0 \quad (2)$$

where u_i are real-valued functions on the X_i and F is a real-valued function on $\prod_{i=1}^n u_i(X_i)^2$. Additional properties on F , e.g. its nondecreasingness (resp. nonincreasingness) in its first (resp. last) n arguments, will give rise to a variety of models implementing this first strategy.

A second strategy relies on the idea of measuring “preference differences” separately on each attribute and then combining these (positive or negative) differences in order to know whether the aggregation of these differences leads to an advantage for x over y . More formally, this suggests a model in which:

$$x \succsim y \Leftrightarrow G(p_1(x_1, y_1), p_2(x_2, y_2), \dots, p_n(x_n, y_n)) \geq 0 \quad (3)$$

where p_i are real-valued functions on X_i^2 and G is a real-valued function on $\prod_{i=1}^n p_i(X_i^2)$. Additional properties on G (e.g. its oddness or its nondecreasingness in each of its arguments) or on p_i (e.g. $p_i(x_i, x_i) = 0$ or $p_i(x_i, y_i) = -p_i(y_i, x_i)$) will give rise to a variety of models in line with the above strategy.

Of course these two strategies are nor incompatible and one may well envisage to use the “worth” of each alternative on each attribute to measure “preference differences”. This suggests a model in which:

$$x \succsim y \Leftrightarrow H(\phi_1(u_1(x_1), u_1(y_1)), \phi_2(u_2(x_2), u_2(y_2)), \dots, \phi_n(u_n(x_n), u_n(y_n))) \geq 0 \quad (4)$$

where u_i are real-valued functions on X_i , ϕ_i are real-valued functions on $u_i(X_i)^2$ and H is a real-valued function on $\prod_{i=1}^n \phi_i(u_i(X_i)^2)$.

This paper is devoted to the analysis of several variants of (2), (3) and (4). We shall provide a fairly complete axiomatic analysis of these models and use them as a general framework allowing to understand the characteristics and differences between several aggregation models.

It should be noted that the introduction of intransitivities in conjoint measurement models is not new. A. Tversky (1969) was one of the first to propose such a model generalizing (1), known as the *additive difference model*, in which:

$$x \succsim y \Leftrightarrow \sum_{i=1}^n \Phi_i(u_i(x_i) - u_i(y_i)) \geq 0 \quad (5)$$

where Φ_i are increasing and odd functions. This model was axiomatized in Fishburn (1992a). More recently, additive nontransitive models were proposed in Bouyssou (1986), Fishburn (1990a, 1991b, 1990b, 1992b) and Vind (1991). They are of the following general type:

$$x \succsim y \Leftrightarrow \sum_{i=1}^n p_i(x_i, y_i) \geq 0 \quad (6)$$

where the p_i are real-valued functions on X_i^2 and may have several additional properties (e.g. $p_i(x_i, x_i) = 0$, for all $i \in \{1, 2, \dots, n\}$ and all $x_i \in X_i$).

Comparing (2), (3) and (4) with (5) or (6) shows that the main particularity of our models lies in the use of a general class of functions instead of just addition and subtraction. This is quite reminiscent of the extension of (1) introduced in Krantz et al. (1971, Ch. 7), known as the *decomposable transitive* model, in which:

$$x \succsim y \Leftrightarrow F(u_1(x_1), u_2(x_2), \dots, u_n(x_n)) \geq F(u_1(y_1), u_2(y_2), \dots, u_n(y_n)), \quad (7)$$

where F is increasing in all its arguments.

Clearly the use of very general functional forms, instead of additive ones, will greatly facilitate the axiomatic analysis of our models. It mainly relies on the study of various kinds of *traces* induced by the preference relation on coordinates and does not require a detailed analysis of tradeoffs between attributes.

The price to pay for such an extension of the scope of conjoint measurement is that none of our models is likely to possess any remarkable uniqueness properties. Therefore, although our proofs are constructive, our results will not give hints on how to devise assessment procedures. The general idea is to use numerical representations as guidelines to understand the consequences of a limited number of cancelation conditions, without imposing any transitivity or completeness requirement on the preference relation and any structural assumptions on the set of objects. They are not studied for their own sake and our results are not intended to provide clues on how to build them.

With few exceptions, this paper does not contain new results but synthesizes results scattered in Bouyssou and Pirlot (1999) and Bouyssou and Pirlot (2002a, 2002c, 2002d), Bouyssou et al. (1997). The presentation of the results is however new, emphasizing the underlying strategy and framework.

We shall suppose throughout the paper that binary relations are defined on *countable* (i.e. finite or countably infinite) sets, referring the reader to the original papers for the study of the general case. Contrary to what

happens with conjoint measurement models involving an addition operation, this restrictive hypothesis will allow us to skip many technical details while maintaining the spirit of the results in the general case.

This paper is organized as follows. Section 2 presents some background material: we introduce our vocabulary concerning binary relations and recall some well-known facts on traces. Section 3 studies binary relations defined on product sets and introduces various kinds of *traces on coordinates*. In section 4 we investigate models of type (2), using marginal traces on levels. Section 5 studies models of type (3) using marginal traces on differences. Section 6 combines these two aspects with the study of models of type (4). A final section discusses the results. Examples and technical details are relegated in appendix.

In order to make this paper self-contained, we only omit proofs that are elementary. Therefore, given the variety of models studied, this paper is somewhat long. The reader willing to get a feeling of the results may skip all proofs and remarks without loss of continuity.

2 Background

2.1 Binary relations

A *binary relation* \succsim on a set A is a subset of $A \times A$. We write $a \succsim b$ instead of $(a, b) \in \succsim$. A binary relation \succsim on A is said to be:

- *reflexive* if $[a \succsim a]$,
- *complete* if $[a \succsim b \text{ or } b \succsim a]$,
- *symmetric* if $[a \succsim b] \Rightarrow [b \succsim a]$,
- *asymmetric* if $[a \succsim b] \Rightarrow [\text{Not}[b \succsim a]]$,
- *transitive* if $[a \succsim b \text{ and } b \succsim c] \Rightarrow [a \succsim c]$,
- *Ferrers* if

$$\left. \begin{array}{l} a \succsim b \\ \text{and} \\ c \succsim d \end{array} \right\} \Rightarrow \left\{ \begin{array}{l} a \succsim d \\ \text{or} \\ c \succsim b, \end{array} \right.$$

- *semi-transitive* if

$$\left. \begin{array}{l} a \succsim b \\ \text{and} \\ b \succsim c \end{array} \right\} \Rightarrow \left\{ \begin{array}{l} a \succsim d \\ \text{or} \\ d \succsim c, \end{array} \right.$$

for all $a, b, c, d \in A$.

The *asymmetric* (resp. *symmetric*) part of \succsim is the binary relation \succ (resp. \sim) on A defined letting, for all $a, b \in A$, $a \succ b \Leftrightarrow [a \succsim b \text{ and } \text{Not}(b \succsim a)]$ (resp. $a \sim b \Leftrightarrow [a \succsim b \text{ and } b \succsim a]$). A similar convention will hold when \succsim is subscripted and/or superscripted.

A *weak order* (resp. an *equivalence relation*) is a complete and transitive (resp. reflexive, symmetric and transitive) binary relation. An *interval order* is a complete and Ferrers binary relation; a *semiorder* is a semi-transitive interval order. If \succsim is an equivalence on A , A/\succsim will denote the set of equivalence classes of \succsim on A .

2.2 Traces of binary relations

The idea that any binary relation generates various reflexive and transitive binary relations called *traces* dates back at least to the pioneering work of R.D. Luce (1956). The use of traces have proved especially useful in the study of preference structures tolerating intransitive indifference such as semiorders or interval orders (see Aleskerov & Monjardet, 2002; Fishburn, 1985; Pirlot & Vincke, 1997).

Definition 2.1 (Traces of binary relations)

Let \succsim be a binary relation on a set A . We associate to \succsim three binary relations, called traces, letting, for all $a, b \in A$:

$$\text{Left Trace } a \succsim^+ b \Leftrightarrow [b \succsim c \Rightarrow a \succsim c],$$

$$\text{Right Trace } a \succsim^- b \Leftrightarrow [c \succsim a \Rightarrow c \succsim b],$$

$$\text{Trace } a \succsim^\pm b \Leftrightarrow [a \succsim^+ b \text{ and } a \succsim^- b].$$

Following our conventions, \sim^+ and \succ^+ will denote the symmetric and asymmetric parts of \succsim^+ , the same holding for \succsim^- and \succ^+ . Useful connections between \succsim and its traces are summarized below for the ease of future reference. All of them are straightforward consequences of the preceding definition.

Proposition 2.1 (Properties of traces)

1. \sim^+ , \sim^- and \sim^\pm are equivalence relations (reflexive, symmetric and transitive).
2. \succ^+ , \succ^- and \succ^\pm are reflexive and transitive binary relations.

3. For all $a, b, c, d \in A$:

$$[a \succ b, b \succ^- c] \Rightarrow a \succ c, \quad (8)$$

$$[a \succ b, c \succ^+ a] \Rightarrow c \succ b, \quad (9)$$

$$[d \succ^\pm a, b \succ^\pm c] \Rightarrow \begin{cases} a \succ b \Rightarrow d \succ c, \\ \text{and} \\ a \succ b \Rightarrow d \succ c, \end{cases} \quad (10)$$

$$[d \sim^\pm a, b \sim^\pm c] \Rightarrow \begin{cases} a \succ b \Leftrightarrow d \succ c, \\ \text{and} \\ a \succ b \Leftrightarrow d \succ c. \end{cases} \quad (11)$$

The following proposition summarizes a number of well-known facts about traces (see Fishburn, 1985; Monjardet, 1978; Pirlot & Vincke, 1997; Roubens & Vincke, 1985).

Proposition 2.2 (Completeness of global traces)

1. \succ^+ is complete $\Leftrightarrow \succ^-$ is complete $\Leftrightarrow \succ$ is Ferrers.
2. \succ^\pm is complete $\Leftrightarrow \succ$ is Ferrers and semi-transitive.

For a detailed analysis of the role of traces in various domains of preference modelling we refer to Aleskerov and Monjardet (2002), Doignon, Monjardet, Roubens, and Vincke (1988), Laslier (1997), Monjardet (1978), Pirlot and Vincke (1997), Roubens and Vincke (1985).

2.3 General numerical representations

We present here very general numerical representations of binary relations that will serve later as guidelines. Although the “results” in this section are part of the folklore of binary relations (see Ebert, 1985), we outline their proof, the logic of which being useful in the sequel.

Let \succ be a binary relation on a *countable* (i.e. finite or countably infinite) set A . It is clearly always possible to build a, trivial, numerical representation of \succ such that:

$$a \succ b \Leftrightarrow \mathcal{T}(a, b) \geq 0, \quad (\mathcal{T})$$

where \mathcal{T} is a real-valued function on A^2 defined letting, for all $a, b \in A$:

$$\mathcal{T}(a, b) = \begin{cases} +1 & \text{if } a \succ b, \\ -1 & \text{otherwise.} \end{cases}$$

It is possible to even further specify the trivial numerical representation (\mathcal{T}) as shown below.

Proposition 2.3 (Trivial numerical representation)

Let \succsim be a binary relation on a countable set A . There is a real-valued function u on A and a real-valued function \mathcal{F} on $u(A)^2$ such that, for all $a, b \in A$:

$$a \succsim b \Leftrightarrow \mathcal{F}(u(a), u(b)) \geq 0, \quad (\mathcal{F}0)$$

PROOF

The relation \sim^\pm is an equivalence. Since A is countable (in fact, as soon as the cardinality of A/\sim^\pm is not “too large”), there is a real-valued function u on A such that, for all $a, b \in A$:

$$a \sim^\pm b \Leftrightarrow u(a) = u(b). \quad (12)$$

Take any function u satisfying (12) and define \mathcal{F} letting, for all $a, b \in A$:

$$\mathcal{F}(u(a), u(b)) = \begin{cases} +1 & \text{if } a \succsim b, \\ -1 & \text{otherwise.} \end{cases} \quad (13)$$

We have to show that \mathcal{F} is well-defined, i.e. that $[u(a) = u(c) \text{ and } u(b) = u(d)]$ implies $[a \succsim b \Leftrightarrow c \succsim d]$. This follows from (11). \square

Remark 2.1 (Uniqueness)

The reader will have noticed that the above proof shows a little more than the mere statement of proposition 2.3. We have in fact shown that there always exists a representation in model $(\mathcal{F}0)$ in which u represents \sim^\pm . This is by no means necessary however. In fact, since A is countable there is a one-to-one function v between A and some countable subset of \mathbb{R} . Any such one-to-one function v may clearly be used instead of a function u representing \sim^\pm . Since in that case $[v(a) = v(c) \text{ and } v(b) = v(d)]$ implies $[a = c \text{ and } b = d]$, using (13) always leads to a well-defined function \mathcal{F} . This trivial representation is therefore defiantly irregular Roberts (1979, ch. 2). Observe finally that any function \mathcal{G} assigning arbitrary nonnegative values to ordered pairs $(u(a), u(b))$ such that $a \succsim b$ and arbitrary negative values otherwise may be used instead of \mathcal{F} .

It is not difficult to extend the above trivial representation to the general uncountable case. All that is required is that the cardinal of A should not be “too large”, i.e., there should be enough real numbers to distinguish elements of A that have to be distinguished. This will be the case as soon as there is a one-to-one correspondence between A/\sim^\pm and some subset of \mathbb{R} . \bullet

We now consider nontrivial numerical representations obtained imposing additional constraints on the trivial model $(\mathcal{F}0)$.

Proposition 2.4 (General models)

The function \mathcal{F} in model $(\mathcal{F}0)$ can be chosen so that, for all $x, y \in X$,

1. $\mathcal{F}(u(x), u(x)) \geq 0$ iff \succsim is reflexive, (model $(\mathcal{F}1)$).
2. $\mathcal{F}(u(x), u(y)) = -\mathcal{F}(u(y), u(x))$ iff \succsim is complete (model $(\mathcal{F}2)$).
3. \mathcal{F} is increasing in its first argument and decreasing in its second argument iff \succsim is Ferrers and semi-transitive (model $(\mathcal{F}3)$).
4. $\mathcal{F}(u(x), u(y)) = -\mathcal{F}(u(y), u(x))$ and \mathcal{F} is nondecreasing in its first argument and nonincreasing in its second argument iff \succsim is a semiorder (model $(\mathcal{F}4)$).
5. $\mathcal{F}(u(x), u(y)) = -\mathcal{F}(u(y), u(x))$ and \mathcal{F} is increasing in its first argument and decreasing in its second argument iff \succsim is a weak order (model $(\mathcal{F}5)$). In that case, it is always possible to take $\mathcal{F}(\alpha, \beta) = \alpha - \beta$ (model $(\mathcal{U}0)$).

PROOF

Part 1 is obvious in view of proposition 2.3.

Part 2. The necessity of completeness is clear. In order to show sufficiency, take any function u satisfying (12) and define \mathcal{F} letting, for all $a, b \in A$:

$$\mathcal{F}(u(a), u(b)) = \begin{cases} +1 & \text{if } a \succ b, \\ 0 & \text{if } a \sim b, \\ -1 & \text{otherwise.} \end{cases} \quad (14)$$

Using the completeness of \succsim and (11), it is easy to see that \mathcal{F} is well-defined and skew symmetric.

Part 3. The necessity of Ferrers and semi-transitivity is easily established using the properties of \mathcal{F} . Let us, for instance, show that \succsim is semi-transitive. Suppose that $a \succsim b$ and $b \succsim c$. Hence $\mathcal{F}(u(a), u(b)) \geq 0$ and $\mathcal{F}(u(b), u(c)) \geq 0$. If $u(b) \geq u(d)$ then $\mathcal{F}(u(a), u(d)) \geq \mathcal{F}(u(a), u(b)) \geq 0$ so that $a \succsim d$. Otherwise we have $u(d) > u(b)$, which implies $\mathcal{F}(u(d), u(c)) > \mathcal{F}(u(b), u(c)) \geq 0$ so that $d \succsim c$.

In order to show sufficiency, remember from part 2 of proposition 2.2 that, when \succsim is Ferrers and semi-transitive, \succsim^\pm is a weak order. Since A is countable, there is a real-valued function u such that, for all $a, b \in A$:

$$a \succsim^\pm b \Leftrightarrow u(a) \geq u(b). \quad (15)$$

Using any function u satisfying (15), define \mathcal{F} letting, for all $a, b \in A$,

$$\mathcal{F}(u(a), u(b)) = \begin{cases} +\exp(u(a) - u(b)) & \text{if } a \succsim b, \\ -\exp(u(b) - u(a)) & \text{otherwise.} \end{cases} \quad (16)$$

That \mathcal{F} is well-defined follows from (11). Its monotonicity properties follow from (10) and its definition.

Part 4. The necessity of completeness, Ferrers and semi-transitivity is easily established. We show sufficiency. Since \succsim is Ferrers and semi-transitive and A is countable, there is a function u satisfying (15). Using any such function u , define \mathcal{F} as in (14). That \mathcal{F} is well-defined and skew symmetric follows from (11), since \sim^\pm is the symmetric part of \succsim^\pm . The monotonicity properties of \mathcal{F} follow from (10).

Part 5. Necessity. Suppose that $a \succsim b$ and $b \succsim c$. Hence $\mathcal{F}(u(a), u(b)) \geq 0$ and $\mathcal{F}(u(b), u(c)) \geq 0$. Since \mathcal{F} is skew symmetric, we know that $\mathcal{F}(u(c), u(b)) \leq 0$. Using the increasingness of \mathcal{F} , $\mathcal{F}(u(a), u(b)) \geq 0$ and $\mathcal{F}(u(c), u(b)) \leq 0$ imply $u(a) \geq u(c)$. Since $\mathcal{F}(u(a), u(a)) = 0$, because \mathcal{F} is skew symmetric, we have $\mathcal{F}(u(a), u(c)) \geq 0$ so that $a \succsim c$. Hence, \succsim is transitive.

Sufficiency. Since \succsim is a weak order and A is countable, there is a function u such that, for all $a, b \in A$:

$$a \succsim b \Leftrightarrow u(a) \geq u(b).$$

Using any such function u , define \mathcal{F} letting, for all $a, b \in A$, $\mathcal{F}(u(a), u(b)) = u(a) - u(b)$. The well-definedness of \mathcal{F} easily follows from the fact that \succsim is a weak order. \square

Remark 2.2 (Uniqueness)

A word on the uniqueness of the representations in proposition 2.4 is in order. Observe that in parts 1 and 2, any one-to-one function v between A and some countable subset of \mathbb{R} may be used instead of u representing \sim^\pm . The only difference between the two parts being in the flexibility for the choice of \mathcal{F} .

For parts 3 and 4, we have shown that u may always be chosen so as to represent \succsim^\pm . Again, this is by no means necessary. Any real-valued function v on A such that, for all, $a, b \in A$,

$$a \succ^\pm b \Rightarrow v(a) > v(b), \tag{17}$$

may also clearly be used with the same function \mathcal{F} as in the above proof.

For part 5, the uniqueness depends on whether one wishes a representation in model ($\mathcal{F}5$) or in model ($\mathcal{U}0$). The uniqueness properties of u in model ($\mathcal{U}0$) are well-known. Clearly, model ($\mathcal{F}5$) offers much more flexibility. Observe that (17) is a necessary condition for u to be used in such a representation. It is no more sufficient however since it might happen that $a \sim^\pm b$ and $a \sim c$. In such a situation, we must guarantee that $u(a) = u(b)$ otherwise the increasingness of \mathcal{F} would lead to $c \succ b$ or $b \succ c$ (depending on whether $u(b) < u(a)$ or $u(b) > u(a)$) whereas $a \sim^\pm b$ and $a \sim c$ imply, using (11), $b \sim c$.

In all parts, the possible degrees of freedom in the choice of \mathcal{F} are obvious.

Again, the extension to the general uncountable case should be clear. For parts 1 and 2, it suffices to require that there is a one-to-one correspondence between A/\sim^\pm and some subset of \mathbb{R} . For all other parts, it is necessary and sufficient to require that the weak order \succsim^\pm has a numerical representation (i.e., A must have a countable order dense subset w.r.t. \succsim^\pm) since the weak order induced by u refines \succsim^\pm . •

The numerical representations envisaged so far are summarized in table 1

Table 1: General numerical representations

Models	Definition	Conditions
(\mathcal{T})	$a \succsim b \Leftrightarrow \mathcal{T}(a, b) \geq 0$	
\Updownarrow		\emptyset
($\mathcal{F}0$)	$a \succsim b \Leftrightarrow \mathcal{F}(u(a), u(b)) \geq 0$	
.....
($\mathcal{F}1$)	($\mathcal{F}0$) with $\mathcal{F}(\alpha, \alpha) \geq 0$	refl.
.....
($\mathcal{F}2$)	($\mathcal{F}0$) with $\mathcal{F}(\alpha, \beta) = -\mathcal{F}(\beta, \alpha)$	cpl.
.....
($\mathcal{F}3$)	($\mathcal{F}0$) with $\mathcal{F}(\nearrow, \searrow)$	Ferrers and semi-transitivity
.....
($\mathcal{F}4$)	($\mathcal{F}2$) with $\mathcal{F}(\nearrow, \searrow)$	semiorde
.....
($\mathcal{F}5$)	($\mathcal{F}2$) with $\mathcal{F}(\nearrow, \searrow)$	
\Updownarrow		weak order
($\mathcal{U}0$)	$a \succsim b \Leftrightarrow u(a) \geq u(b)$	

\nearrow means increasing, \searrow means decreasing
 refl. means reflexive, cpl. means complete

3 Binary relations on product sets

In the rest of this paper, we consider a *countable* set $X = \prod_{i=1}^n X_i$ with $n \geq 2$. Elements x, y, z, \dots of X will be interpreted as alternatives evaluated on a set $N = \{1, 2, \dots, n\}$ of attributes. A typical binary relation of X is still denoted as \succsim with \sim and \succ as symmetric and asymmetric parts. Again, a similar convention holds when \succsim is superscripted and/or subscripted. It is useful to interpret \succsim as an “at least as good as” preference relation between multi-attributed alternatives.

For any non empty subset J of the set of attributes N , we denote by X_J (resp. X_{-J}) the set $\prod_{i \in J} X_i$ (resp. $\prod_{i \notin J} X_i$). With customary abuse of notation, (x_J, y_{-J}) will denote the element $w \in X$ such that $w_i = x_i$ if $i \in J$ and $w_i = y_i$ otherwise. When $J = \{i\}$ we shall simply write X_{-i} and (x_i, y_{-i}) . We say that \succsim is *marginally complete* for $i \in N$ if $(x_i, a_{-i}) \succsim (y_i, a_{-i})$ or $(y_i, a_{-i}) \succsim (x_i, a_{-i})$, for all $x_i, y_i \in X_i$ and all $a_{-i} \in X_{-i}$, i.e. if no incomparability occurs when comparing alternatives differing only on attribute $i \in N$.

3.1 Independence and marginal preferences

In conjoint measurement, one starts with a preference relation \succsim on X . It is then important to investigate how this information enables to define preference relations on attributes or subsets of attributes. Let $J \subseteq N$ be a nonempty set of attributes. We define the *marginal relation* \succsim_J induced on X_J by \succsim letting, for all $x_J, y_J \in X_J$:

$$x_J \succsim_J y_J \Leftrightarrow (x_J, z_{-J}) \succsim (y_J, z_{-J}), \text{ for all } z_{-J} \in X_{-J},$$

with asymmetric (resp. symmetric) part \succ_J (resp. \sim_J). Note that if \succsim is reflexive (resp. transitive), the same will be true for \succsim_J . This is clearly not true for completeness however.

We define two other binary relations induced by \succsim on X_J , letting for all $x_J, y_J \in X_J$,

$$x_J \mathcal{S}_J y_J \Leftrightarrow (x_J, z_{-J}) \succsim (y_J, z_{-J}), \text{ for some } z_{-J} \in X_{-J},$$

and

$$x_J \mathcal{P}_J y_J \Leftrightarrow (x_J, z_{-J}) \succ (y_J, z_{-J}), \text{ for some } z_{-J} \in X_{-J}.$$

Definition 3.1 (Independence and separability)

Consider a binary relation \succsim on a set $X = \prod_{i=1}^n X_i$ and let $J \subseteq N$ be a nonempty subset of attributes. We say that \succsim is:

1. independent for J if $\mathcal{S}_J \subseteq \succsim_J$,
2. separable for J if \mathcal{P}_J is asymmetric.

If \succsim is independent (resp. separable) for all non empty subsets of N , we say that \succsim is independent (resp. separable). If \succsim is independent (resp. separable) for all subsets containing a single attribute, we say that \succsim is weakly independent (resp. weakly separable).

Independence is a classical notion in conjoint measurement. It states that common evaluations on some attributes do not influence preference. Whereas independence implies weak independence, it is well-known that the converse is not true (see Wakker, 1989).

Independence implies separability but not *vice versa*. Separability is a weakening of independence that can be motivated considering aggregation models based on “Max” or “Min”. It forbids strict reversals of preference when varying common evaluations on some attribute. It is easy to see that weak separability does not entail separability.

Remark 3.1 (Completeness of \succsim_i)

Let us observe that when \succsim is complete and independent for $i \in N$ then \succsim_i is clearly complete. It is not difficult to see that \succsim_i is complete if and only if \succsim is marginally complete and weakly separable for $i \in N$. The marginal preference relation \succsim_i reflects the comparison of x_i and y_i when they are adjoined with the same evaluations on X_{-i} . This does not exploit all the information contained in \succsim . The marginal traces on levels will do so. •

3.2 Marginal traces on levels

The definitions and results from section 2.2 clearly apply when \succsim is a binary relation on a product set $X = \prod_{i=1}^n X_i$. Hence, the binary relation \succsim has a *left trace* (resp. *right trace* and *trace*) \succsim^+ (resp. \succsim^- and \succsim^\pm) that is reflexive and transitive.

Consider an attribute $i \in N$. Sticking to the notation introduced above, \succsim_i^+ (resp. \succsim_i^- and \succsim_i^\pm) will denote the marginal preference relation induced on X_i by \succsim^+ (resp. \succsim^- and \succsim^\pm), i.e.

$$\begin{aligned} x_i \succsim_i^+ y_i &\Leftrightarrow [(x_i, z_{-i}) \succsim^+ (y_i, z_{-i}), \text{ for all } z_{-i} \in X_{-i}], \\ x_i \succsim_i^- y_i &\Leftrightarrow [(x_i, z_{-i}) \succsim^- (y_i, z_{-i}), \text{ for all } z_{-i} \in X_{-i}], \\ x_i \succsim_i^\pm y_i &\Leftrightarrow [(x_i, z_{-i}) \succsim^\pm (y_i, z_{-i}), \text{ for all } z_{-i} \in X_{-i}]. \end{aligned}$$

Since, by construction, \succsim^+ , \succsim^- and \succsim^\pm are reflexive and transitive, the same is true for \succsim_i^+ , \succsim_i^- and \succsim_i^\pm . As shown in the following lemma, \succsim_i^+ (resp. \succsim_i^- and \succsim_i^\pm), the marginal relation induced on $X_i, i \in N$, by the global trace \succsim^+ (resp. \succsim^- and \succsim^\pm) can also be usefully interpreted as a *marginal trace on levels* on attribute $i \in N$.

Lemma 3.1 (Marginal traces on levels induced by traces)

For all $i \in N$, all $x_i, y_i \in X_i$, all $a_{-i} \in X_{-i}$ and all $z \in X$:

1. $x_i \succsim_i^+ y_i \Leftrightarrow [(y_i, a_{-i}) \succsim z \Rightarrow (x_i, a_{-i}) \succsim z]$,

$$\begin{aligned}
2. \quad x_i \succsim_i^- y_i &\Leftrightarrow [z \succsim (x_i, a_{-i}) \Rightarrow z \succsim (y_i, a_{-i})], \\
3. \quad x_i \succsim_i^\pm y_i &\Leftrightarrow \begin{cases} (y_i, a_{-i}) \succsim z \Rightarrow (x_i, a_{-i}) \succsim z, \\ \text{and} \\ z \succsim (x_i, a_{-i}) \Rightarrow z \succsim (y_i, a_{-i}). \end{cases}
\end{aligned}$$

PROOF

We give the proof of part 1, the proof of the other parts being similar. By definition we have: $x_i \succsim_i^+ y_i \Leftrightarrow [(x_i, a_{-i}) \succsim^+ (y_i, a_{-i}), \text{ for all } a_{-i} \in X_{-i}] \Leftrightarrow [(y_i, a_{-i}) \succsim z \Rightarrow (x_i, a_{-i}) \succsim z, \text{ for all } a_{-i} \in X_{-i} \text{ and all } z \in X].$ \square

Remark 3.2 (Marginal traces on levels and marginal preferences)

The above result should clarify the similarities and differences between the marginal traces on levels \succsim_i^\pm and the marginal preferences \succsim_i . They are both defined on X_i . However marginal traces use all information contained in \succsim in order to compare x_i with y_i including how they compare, when adjoined with the same evaluations on X_{-i} , with every other alternatives. On the contrary the definition of \succsim_i only relies on the comparison between x_i with y_i when adjoined with the same evaluations on X_{-i} . No other alternative is involved here. \bullet

As before, the symmetric and asymmetric parts of \succsim_i^+ are respectively denoted \sim_i^+ and \succ_i^+ , the same convention applying to \succsim_i^- and \succsim_i^\pm (although it is clearly possible to define marginal traces on a subset of attributes other than singletons, we do not envisage this possibility in this paper).

As in proposition 2.1, there are many interesting connections between marginal traces and \succsim . We list some of them in the following lemma, leaving its easy proof as an exercise for the reader.

Lemma 3.2 (Properties of marginal traces on levels)

For all $i \in N$ and $x, y, z, w \in X$:

$$[x \succsim y, z_i \succsim_i^+ x_i] \Rightarrow (z_i, x_{-i}) \succsim y, \quad (18)$$

$$[x \succsim y, y_i \succsim_i^- w_i] \Rightarrow x \succsim (w_i, y_{-i}), \quad (19)$$

$$[z_i \succsim_i^\pm x_i, y_i \succsim_i^\pm w_i] \Rightarrow \begin{cases} x \succsim y \Rightarrow (z_i, x_{-i}) \succsim (w_i, y_{-i}), \\ \text{and} \\ x \succ y \Rightarrow (z_i, x_{-i}) \succ (w_i, y_{-i}), \end{cases} \quad (20)$$

$$[z_i \sim_i^\pm x_i, y_i \sim_i^\pm w_i, \text{ for all } i \in N] \Rightarrow \begin{cases} x \succsim y \Leftrightarrow z \succsim w, \\ \text{and} \\ x \succ y \Leftrightarrow z \succ w. \end{cases} \quad (21)$$

It is clear that the marginal traces \succsim_i^+ , \succsim_i^- and \succsim_i^\pm need not be complete. Interesting consequences will arise when this is the case. This is explored in section 4.

3.3 Marginal traces on differences

P. P. Wakker (1988, 1989) has powerfully shown how the consideration of induced relations comparing “preference differences” on each attribute may illuminate the analysis of conjoint measurement models. We follow the same path in this section.

Given a binary relation \succsim on X , we define the binary relation \succsim_i^* on X_i^2 letting, for all $x_i, y_i, z_i, w_i \in X_i$,

$$(x_i, y_i) \succsim_i^* (z_i, w_i) \text{ iff} \\ [\text{for all } a_{-i}, b_{-i} \in X_{-i}, (z_i, a_{-i}) \succsim (w_i, b_{-i}) \Rightarrow (x_i, a_{-i}) \succsim (y_i, b_{-i})].$$

Intuitively, if $(x_i, y_i) \succsim_i^* (z_i, w_i)$, it seems reasonable to conclude that the preference difference between x_i and y_i is not smaller than the preference difference between z_i and w_i . Notice that, by construction, \succsim_i^* is reflexive and transitive.

Contrary to our intuition concerning preference differences, the definition of \succsim_i^* does not imply that the two “opposite” differences (x_i, y_i) and (y_i, x_i) are linked. Henceforth we introduce the binary relation \succsim_i^{**} on X_i^2 letting, for all $x_i, y_i, z_i, w_i \in X_i$,

$$(x_i, y_i) \succsim_i^{**} (z_i, w_i) \text{ iff } [(x_i, y_i) \succsim_i^* (z_i, w_i) \text{ and } (w_i, z_i) \succsim_i^* (y_i, x_i)].$$

It is easy to see that \succsim_i^{**} is reflexive, transitive and *reversible*, i.e. $(x_i, y_i) \succsim_i^{**} (z_i, w_i) \Leftrightarrow (w_i, z_i) \succsim_i^{**} (y_i, x_i)$.

The relations \succsim_i^* and \succsim_i^{**} both appear to capture the idea of comparison of preference differences between elements of X_i induced by the relation \succsim . For the sake of easy reference, we note a few useful connections between \succsim_i^* , \succsim_i^{**} and \succsim in the following lemma.

Lemma 3.3 (Properties of marginal traces on differences)

For all $x, y \in X$ and all $z_i, w_i \in X_i$,

$$\succsim \text{ is independent iff } (x_i, x_i) \sim_i^* (y_i, y_i) \text{ for all } i \in N, \quad (22)$$

$$[x \succsim y \text{ and } (z_i, w_i) \succsim_i^* (x_i, y_i)] \Rightarrow (z_i, x_{-i}) \succsim (w_i, y_{-i}), \quad (23)$$

$$[(z_i, w_i) \sim_i^* (x_i, y_i) \text{ for all } i \in N] \Rightarrow [x \succsim y \Leftrightarrow z \succsim w], \quad (24)$$

$$[x \succ y \text{ and } (z_i, w_i) \succsim_i^{**} (x_i, y_i)] \Rightarrow (z_i, x_{-i}) \succ (w_i, y_{-i}), \quad (25)$$

$$[(z_i, w_i) \sim_i^{**} (x_i, y_i) \text{ for all } i \in N] \Rightarrow \begin{cases} [x \succsim y \Leftrightarrow z \succsim w] \\ \text{and} \\ [x \succ y \Leftrightarrow z \succ w], \end{cases} \quad (26)$$

PROOF

(22). It is clear that $[\succsim \text{ is independent}] \Leftrightarrow [\succsim \text{ is independent for } N \setminus \{i\}, \text{ for all } i \in N]$. Observe that $[\succsim \text{ is independent for } N \setminus \{i\}, \text{ for all } i \in N] \Leftrightarrow [(x_i, a_{-i}) \succsim (x_i, b_{-i}) \Leftrightarrow (y_i, a_{-i}) \succsim (y_i, b_{-i}), \text{ for all } x_i, y_i \in X_i \text{ and all } a_{-i}, b_{-i} \in X_{-i}] \Leftrightarrow [(x_i, x_i) \sim_i^* (y_i, y_i) \text{ for all } x_i, y_i \in X_i]$.

(23) is obvious from the definition of \succsim_i^* . By induction, (24) is immediate.

(25). Using (23), we know that $(z_i, x_{-i}) \succsim (w_i, y_{-i})$. Suppose that $(w_i, y_{-i}) \succsim (z_i, x_{-i})$. Since $(z_i, w_i) \succsim_i^{**} (x_i, y_i)$ implies $(y_i, x_i) \succsim_i^* (w_i, z_i)$, (23) implies $y \succsim x$, a contradiction. (26) is immediate from (25). \square

It is clear that the marginal traces on differences \succsim_i^* and \succsim_i^{**} need not be complete. Interesting consequences will arise when this is the case. This is explored in section 5.

3.4 Relations between traces on levels and traces on differences

All the marginal traces envisaged up to this point (\succsim_i^+ , \succsim_i^- , \succsim_i^\pm , \succsim_i^* and \succsim_i^{**}) are induced by the same relation \succsim . Clearly they are not unrelated. The following lemma shows that they relate exactly as could be expected.

Lemma 3.4 (Relations between marginal traces)

For all $i \in N$ and all $x_i, y_i \in X_i$,

1. $x_i \succsim_i^+ y_i \Leftrightarrow [(x_i, w_i) \succsim_i^* (y_i, w_i), \text{ for all } w_i \in X_i]$,
2. $x_i \succsim_i^- y_i \Leftrightarrow [(w_i, y_i) \succsim_i^* (w_i, x_i), \text{ for all } w_i \in X_i]$,
3. $x_i \succsim_i^\pm y_i \Leftrightarrow [(x_i, w_i) \succsim_i^{**} (y_i, w_i), \text{ for all } w_i \in X_i]$,
4. $[\ell_i \succsim_i^+ x_i \text{ and } (x_i, y_i) \succsim_i^* (z_i, w_i)] \Rightarrow (\ell_i, y_i) \succsim_i^* (z_i, w_i)$,
5. $[y_i \succsim_i^- \ell_i \text{ and } (x_i, y_i) \succsim_i^* (z_i, w_i)] \Rightarrow (x_i, \ell_i) \succsim_i^* (z_i, w_i)$,
6. $[z_i \succsim_i^+ \ell_i \text{ and } (x_i, y_i) \succsim_i^* (z_i, w_i)] \Rightarrow (x_i, y_i) \succsim_i^* (\ell_i, w_i)$,
7. $[\ell_i \succsim_i^- w_i \text{ and } (x_i, y_i) \succsim_i^* (z_i, w_i)] \Rightarrow (x_i, y_i) \succsim_i^* (z_i, \ell_i)$,
8. $[x_i \sim_i^+ z_i \text{ and } y_i \sim_i^- w_i] \Rightarrow (x_i, y_i) \sim_i^* (z_i, w_i)$,
9. $[x_i \sim_i^\pm z_i \text{ and } y_i \sim_i^\pm w_i] \Rightarrow (x_i, y_i) \sim_i^{**} (z_i, w_i)$,

PROOF

Part 1. By definition, $[x_i \succsim_i^+ y_i] \Leftrightarrow [(y_i, a_{-i}) \succsim z \Rightarrow (x_i, a_{-i}) \succsim z] \Leftrightarrow [(x_i, z_i) \succsim_i^* (y_i, z_i), \text{ for all } z_i \in X_i]$. The proof of part 2 is similar. Part 3 easily follows from parts 1 and 2.

Part 4. By definition, $(x_i, y_i) \succsim_i^* (z_i, w_i)$ iff $[(z_i, a_{-i}) \succsim (w_i, b_{-i}) \Rightarrow (x_i, a_{-i}) \succsim (y_i, b_{-i})]$. Since $\ell_i \succsim_i^+ x_i$, $(x_i, a_{-i}) \succsim (y_i, b_{-i}) \Rightarrow (\ell_i, a_{-i}) \succsim (y_i, b_{-i})$. Therefore $(\ell_i, y_i) \succsim_i^* (z_i, w_i)$. The proof of parts 5, 6 and 7 is similar. Part 8 easily follows from parts 4, 5, 6, 7 since \succsim_i^* is reflexive. Part 9 is immediate from part 8. \square

3.5 Trivial numerical representations on product sets

Arbitrary binary relations on product sets have trivial numerical representations of many different kinds (see Bouyssou & Pirlot, 2002d, 2002a, 2002c). We present three of them below that will be easily compared with the general representations introduced above. We abuse notations in the sequel, writing $[u_i(x_i)]$ instead of $(u_1(x_1), u_2(x_2), \dots, u_n(x_n))$. This should not cause confusion.

Proposition 3.1 (Trivial representations on product sets)

Let \succsim be a binary relation on a countable set $X = \prod_{i=1}^n X_i$.

1. There are real-valued functions u_i on X_i and a real valued function F on $[\prod_{i=1}^n u_i(X_i)]^2$ such that, for all $x, y \in X$,

$$x \succsim y \Leftrightarrow F([u_i(x_i)]; [u_i(y_i)]) \geq 0. \quad (L0)$$

2. There are real-valued functions p_i on X_i^2 and a real valued function G on $\prod_{i=1}^n p_i(X_i^2)$ such that, for all $x, y \in X$,

$$x \succsim y \Leftrightarrow G([p_i(x_i, y_i)]) \geq 0 \quad (D0)$$

3. There is a real-valued function u_i on X_i , a real-valued function ϕ_i on $u_i(X_i)^2$ and a real-valued function H on $\prod_{i=1}^n \phi_i(u_i(X_i)^2)$ such that, for all $x, y \in X$,

$$x \succsim y \Leftrightarrow H([\phi_i(u_i(x_i), u_i(y_i))]) \geq 0, \quad (L0D0)$$

PROOF

Part 1. Let $i \in N$. By construction, \sim_i^\pm is an equivalence being a reflexive, symmetric and transitive binary relation. Since X_i is countable, we know that there is a real-valued function u_i on X_i such that, for all $x_i, y_i \in X_i$:

$$x_i \sim_i^\pm y_i \Leftrightarrow u_i(x_i) = u_i(y_i). \quad (27)$$

On each $i \in N$, take any function u_i satisfying (27). Define F on $[\prod_{i=1}^n u_i(X_i)]^2$ letting:

$$F([u_i(x_i)]; [u_i(y_i)]) = \begin{cases} +1 & \text{if } x \succsim y, \\ -1 & \text{otherwise.} \end{cases} \quad (28)$$

The well-definedness of F follows from (21).

Part 2. Let $i \in N$. By construction, \sim_i^{**} is an equivalence being a reflexive, symmetric and transitive binary relation. Since X_i is countable, we know that there is a real-valued function p_i on X_i^2 such that, for all $x_i, y_i, z_i, w_i \in X_i$:

$$(x_i, y_i) \sim_i^{**} (z_i, w_i) \Leftrightarrow p_i(x_i, y_i) = p_i(z_i, w_i). \quad (29)$$

On each $i \in N$, take any function p_i satisfying (29). Define G on $\prod_{i=1}^n p_i(X_i^2)$ letting:

$$G([p_i(x_i, y_i)]) = \begin{cases} +1 & \text{if } x \succsim y, \\ -1 & \text{otherwise.} \end{cases} \quad (30)$$

Using (26), it is easy to show that G is well-defined.

Part 3. Take any function u_i satisfying (27) and any function p_i satisfying (29). Define ϕ_i on $u_i(X_i)^2$ letting, for all $x_i, y_i \in X_i$,

$$\phi_i(u_i(x_i), u_i(y_i)) = p_i(x_i, y_i). \quad (31)$$

Let us show that ϕ_i is well-defined i.e. that $u_i(x_i) = u_i(z_i)$ and $u_i(y_i) = u_i(w_i)$ imply $p_i(x_i, y_i) = p_i(z_i, w_i)$. By construction, we have $x_i \sim_i^{\pm} z_i$ and $y_i \sim_i^{\pm} w_i$ so that, by part 9 of lemma 3.4, $(x_i, y_i) \sim_i^{**} (z_i, w_i)$. Therefore we have $p_i(x_i, y_i) = p_i(z_i, w_i)$, as required.

Define H on $\prod_{i=1}^n \phi_i(u_i(X_i), u_i(X_i))$ letting:

$$H([\phi_i(u_i(x_i), u_i(y_i))]) = \begin{cases} +1 & \text{if } x \succsim y, \\ -1 & \text{otherwise.} \end{cases} \quad (32)$$

Using (24), it is easy to show that H is well-defined. \square

Remark 3.3 (Uniqueness)

As was the case in section 2.3, these various trivial representations are definitely irregular. A word on uniqueness may clarify things and avoid misunderstandings.

Consider first model (L0). Since X_i is countable, consider any one-to-one correspondence v_i between X_i and some countable subset of \mathbb{R} . Any such correspondence on each attribute may be used instead of the functions u_i representing \sim_i^{\pm} used in the above proof. Clearly any F assigning arbitrary nonnegative values to $2n$ -tuples $([u_i(x_i)]; [u_i(y_i)])$ such that $x \succsim y$ and arbitrary negative values otherwise is an acceptable choice.

In model (*D0*) any one-to-one function q_i between X_i^2 and some countable subset of \mathbb{R} may clearly be used instead of the functions p_i representing \sim_i^{**} used in the above proof. Notice that, in the particular case in which, for all $i \in N$ and all $x_i, y_i \in X_i$, $(x_i, x_i) \sim_i^{**} (y_i, y_i)$, it is always possible to use a skew symmetric function q_i (i.e. $q_i(x_i, y_i) = -q_i(y_i, x_i)$, for all $x_i, y_i \in X_i$) on each attribute. Clearly any function G assigning arbitrary nonnegative values to n -tuples $[p_i(x_i, y_i)]$ such that $x \succsim y$ and arbitrary negative values otherwise, is an acceptable choice.

If possible, the representation is even more irregular in model (*L0D0*). Any one-to-one correspondence v_i between X_i and some countable subset of \mathbb{R} may be used instead of u_i representing \sim_i^\pm used in the above proof. Furthermore any one-to-one correspondence φ_i between $v_i(X_i)^2$ and some countable subset of \mathbb{R} may be used instead of ϕ_i representing \sim_i^{**} as done in the above proof. Note, in particular, that it is always possible to choose φ_i in such a way that it is increasing in its first argument and decreasing in its second argument. Furthermore, as above, if, for all $i \in N$ and all $x_i, y_i \in X_i$, $(x_i, x_i) \sim_i^{**} (y_i, y_i)$ it is always possible to choose a skew symmetric representation φ_i (i.e. $\varphi_i(u_i(x_i), u_i(y_i)) = -\varphi_i(u_i(y_i), u_i(x_i))$). Again, any function H assigning arbitrary nonnegative values to n -tuples $[\phi_i(u_i(x_i), u_i(y_i))]$ such that $x \succsim y$ and arbitrary negative values otherwise, is an acceptable choice. •

Remark 3.4 (Extension to the general case)

The above result extends easily to the general uncountable case. In each case, it suffices to require that there are enough real numbers to distinguish the elements that have to be distinguished. This amounts to requiring for the second case that there is a one-to-one correspondence between X_i^2/\sim_i^{**} and some subset of \mathbb{R} . For the other two cases, the, stronger, requirement that there is a one-to-one correspondence between X_i/\succsim_i^\pm is clearly necessary and sufficient. •

We summarize the various trivial models envisaged to this point in table 2.

Table 2: Trivial numerical representations on product sets

Models	Definition
<i>(L0)</i>	$x \succsim y \Leftrightarrow F([u_i(x_i)]; [u_i(y_i)]) \geq 0$
\Updownarrow	
<i>(D0)</i>	$x \succsim y \Leftrightarrow G([p_i(x_i, y_i)]) \geq 0$
\Updownarrow	
<i>(L0D0)</i>	$x \succsim y \Leftrightarrow H([\phi_i(u_i(x_i), u_i(y_i))]) \geq 0$

Our strategy in the sequel is quite similar to the one used in section 2.3. We shall start with each one of the three trivial models ($L0$), ($D0$) or ($L0D0$) and study the consequences of imposing additional constraints on these models. We start with model ($L0$).

4 Models using marginal traces on levels

4.1 Complete marginal traces on levels

As was the case with the Ferrers and semi-transitivity conditions when studying traces, we envisage here conditions that will guarantee that marginal traces on levels are complete and, hence, weak orders. As with interval orders and semiorders, these conditions will prove useful to analyze the underlying structures and to build numerical representations.

Definition 4.1 (Conditions $AC1$, $AC2$ and $AC3$)

Let \succsim be a binary relation on a set $X = \prod_{i=1}^n X_i$. For $i \in N$, this relation is said to satisfy:

$AC1_i$ if

$$\left. \begin{array}{l} x \succsim y \\ \text{and} \\ z \succsim w \end{array} \right\} \Rightarrow \left\{ \begin{array}{l} (z_i, x_{-i}) \succsim y \\ \text{or} \\ (x_i, z_{-i}) \succsim w, \end{array} \right.$$

$AC2_i$ if

$$\left. \begin{array}{l} x \succsim y \\ \text{and} \\ z \succsim w \end{array} \right\} \Rightarrow \left\{ \begin{array}{l} x \succsim (w_i, y_{-i}) \\ \text{or} \\ z \succsim (y_i, w_{-i}), \end{array} \right.$$

$AC3_i$ if

$$\left. \begin{array}{l} z \succsim (x_i, a_{-i}) \\ \text{and} \\ (x_i, b_{-i}) \succsim y \end{array} \right\} \Rightarrow \left\{ \begin{array}{l} z \succsim (w_i, a_{-i}) \\ \text{or} \\ (w_i, b_{-i}) \succsim y, \end{array} \right.$$

for all $x, y, z, w \in X$, all $a_{-i}, b_{-i} \in X_{-i}$ and all $x_i, w_i \in X_i$.

We say that \succsim satisfies $AC1$ (resp. $AC2$, $AC3$) if it satisfies $AC1_i$ (resp. $AC2_i$, $AC3_i$) for all $i \in N$. We sometimes write $AC123$ as shorthand for $AC1$, $AC2$ and $AC3$.

These three conditions are transparent variations on the theme of the Ferrers ($AC1$ and $AC2$) and semi-transitivity ($AC3$) conditions that are made possible by the product structure of X . The rationale for the name “ AC ” is that these conditions are “intra-attribute Cancellation” conditions.

Condition $AC1_i$ suggests that the elements of X_i (instead of the elements of X had the original Ferrers condition been invoked) can be linearly ordered considering “upward dominance”: if x_i “upward dominates” z_i then $(z_i, c_{-i}) \succsim w$ entails $(x_i, c_{-i}) \succsim w$. Condition $AC2_i$ has a similar interpretation considering now “downward dominance”. Condition $AC3_i$ ensures that the linear arrangements of the elements of X_i obtained considering upward and downward dominance are not incompatible.

Conditions $AC1$, $AC2$ and $AC3$ were first introduced in Bouyssou and Pirlot (1999), Bouyssou et al. (1997) and later used in Greco, Matarazzo, and Słowiński (2002). The strong links between $AC1$, $AC2$, $AC3$ and marginal traces are noted in the following:

Lemma 4.1 (Completeness of marginal traces on levels)

We have:

1. \succsim_i^+ is complete iff $AC1_i$ holds,
2. \succsim_i^- is complete iff $AC2_i$ holds,
3. $[Not[x_i \succsim_i^+ y_i] \Rightarrow y_i \succsim_i^- x_i]$ iff $AC3_i$ holds,
4. \succsim_i^\pm is complete iff $AC1_i, AC2_i$ and $AC3_i$ hold,
5. For a binary relation \succsim on $X = \prod_{i=1}^n X_i$, completeness, $AC1_i, AC2_i$ and $AC3_i$ are independent conditions.

PROOF

Part 1 is proved observing that the negation of $AC1_i$ is equivalent to the negation of the completeness of \succsim_i^+ . The proof of part 2 is similar.

Part 3. Suppose that $Not[x_i \succsim_i^+ y_i]$ so that $z \succ (x_i, a_{-i})$ and $Not[z \succ (y_i, a_{-i})]$, for some $z \in X$ and some $a_{-i} \in X_{-i}$. If $(x_i, b_{-i}) \succ w$ then $AC3_i$ implies $(y_i, b_{-i}) \succ w$ or $z \succ (y_i, a_{-i})$. Since, by hypothesis, $Not[z \succ (y_i, a_{-i})]$, we must have $(y_i, b_{-i}) \succ w$, so that $y_i \succsim_i^- x_i$. The reverse implication is proved observing that the negation of $AC3_i$ is equivalent to $Not[y_i \succsim_i^+ x_i]$ and $Not[x_i \succsim_i^- y_i]$ for some $x_i, y_i \in X_i$. Part 4 immediately results from parts 1, 2 and 3.

Part 5: see examples A.1 to A.3 in appendix A. □

Comparing lemma 4.1 with proposition 2.2 shows an important difference between traces and marginal traces on levels: in the latter case, the right trace may be complete without implying the completeness of the left trace. Hence, the use of three conditions ($AC1$, $AC2$ and $AC3$) when studying marginal traces on levels instead of the two classical conditions (Ferrers and semi-transitivity) used when studying traces.

The combination of our three conditions ($AC1$, $AC2$ and $AC3$) implies that the marginal traces on levels induced by \succsim are weak orders. Unsurprisingly, this implies that marginal relations \succsim_i do have special properties even when they differ from marginal traces (which is the general case). We summarize them in the following:

Proposition 4.1 (Properties of marginal preferences)

1. If \succsim is reflexive and either $AC1_i$ or $AC2_i$ holds then \succsim is marginally complete and weakly separable for $i \in N$.
2. If \succsim is reflexive and either $AC1_i$ or $AC2_i$ holds then \succsim_i is an interval order.
3. If, in addition, \succsim satisfies $AC3_i$ then \succsim_i is a semiorder.

PROOF

Part 1. We give the proof using $AC1_i$, the proof using $AC2_i$ being similar. Using the reflexivity of \succsim , we know that $(x_i, a_{-i}) \succsim (x_i, a_{-i})$ and $(y_i, a_{-i}) \succsim (y_i, a_{-i})$. Since $AC1_i$ holds, \succsim_i^+ is complete so that $x_i \succsim_i^+ y_i$ or $y_i \succsim_i^+ x_i$. If $x_i \succsim_i^+ y_i$ then, using (18), we have $(x_i, a_{-i}) \succsim (y_i, a_{-i})$. Similarly if $y_i \succsim_i^+ x_i$ then $(y_i, a_{-i}) \succsim (x_i, a_{-i})$. Hence, \succsim is marginally complete for $i \in N$.

Suppose now that \succsim is not weakly separable for $i \in N$. Then we have $(x_i, a_{-i}) \succ (y_i, a_{-i})$ and $(y_i, b_{-i}) \succ (x_i, b_{-i})$, for some $x_i, y_i \in X_i$ and some $a_{-i}, b_{-i} \in X_{-i}$. Since \succsim is reflexive, we have $(y_i, a_{-i}) \succsim (y_i, a_{-i})$ and $(x_i, b_{-i}) \succsim (x_i, b_{-i})$. This would imply $Not[x_i \succsim_i^+ y_i]$ and $Not[y_i \succsim_i^+ x_i]$, violating $AC1_i$. Hence, \succsim is weakly separable for $i \in N$.

Part 2. We know from part 1 that \succsim is marginally complete and weakly separable for $i \in N$. Hence, \succsim_i is complete. It remains to prove that \succsim_i is Ferrers. It is well-known (see Pirlot & Vincke, 1997, theorem 3.11) that, since \succsim_i is complete, the Ferrers condition is equivalent to: $[x_i \succ_i y_i, y_i \sim_i z_i \text{ and } z_i \succ_i w_i] \Rightarrow x_i \succ_i w_i$, for all $x_i, y_i, z_i, w_i \in X_i$. We prove that this condition holds.

We claim that if $AC1_i$ holds then $[x_i \sim_i z_i \text{ and } x_i \succ_i y_i] \Rightarrow z_i \succ_i^+ y_i$. In contradiction with the thesis, since \succsim_i^+ is complete, suppose that $y_i \succsim_i^+ z_i$. Since $x_i \sim_i z_i$, we have $z_i \succsim_i x_i$ so that $(z_i, a_{-i}) \succsim (x_i, a_{-i})$, for all $a_{-i} \in X_{-i}$. But $y_i \succsim_i^+ z_i$ implies $(y_i, a_{-i}) \succsim (x_i, a_{-i})$, for all $a_{-i} \in X_{-i}$, contradicting $x_i \succ_i y_i$. This proves the claim.

Suppose that $x_i \succ_i y_i$, $y_i \sim_i z_i$ and $z_i \succ_i w_i$. Using the above claim, we have $y_i \succ_i^+ w_i$. If $w_i \succsim_i x_i$, so that $(w_i, a_{-i}) \succsim (x_i, a_{-i})$, for all $a_{-i} \in X_{-i}$, $y_i \succ_i^+ w_i$ implies $(y_i, a_{-i}) \succsim (x_i, a_{-i})$, for all $a_{-i} \in X_{-i}$, contradicting $x_i \succ_i y_i$. Hence $x_i \succ_i w_i$. The proof using $AC2_i$ is similar.

Part 3. In view of part 2 above, all we have to show is that \succsim_i is semi-transitive. It is well-known (see Pirlot & Vincke, 1997, theorem 3.1) that, for a complete relation, semi-transitivity is equivalent to saying that no pair of elements of X_i can be linked both by a path of two consecutive \succ_i -arcs and by a path of two consecutive \sim_i -arcs, i.e., for all $x_i, y_i, z_i, w_i \in X_i$, we may not have $x_i \succ_i y_i, y_i \succ_i z_i, x_i \sim_i w_i$ and $w_i \sim_i z_i$. We prove that this condition holds.

Suppose that $x_i \succ_i y_i, y_i \succ_i z_i, x_i \sim_i w_i$ and $w_i \sim_i z_i$. Using $AC1_i$ and the claim proved in part 2 above, we know that this implies $w_i \succ_i^+ y_i$. In view of part 3 of lemma 4.1, $AC3_i$ then implies $w_i \succsim_i^- y_i$. Since $w_i \sim_i z_i$, we have $(z_i, a_{-i}) \succsim (w_i, a_{-i})$, for all $a_{-i} \in X_{-i}$ and $w_i \succsim_i^- y_i$ leads to $(z_i, a_{-i}) \succsim (y_i, a_{-i})$, for all $a_{-i} \in X_{-i}$, contradicting $y_i \succ_i z_i$. The proof using $AC2_i$ is similar and is omitted. \square

4.2 Strict responsiveness to traces on levels

Keeping in mind the classical constant threshold numerical representation for finite semiorders (see Pirlot & Vincke, 1997; Scott & Suppes, 1958), it is clear that, in general, in a semiorder we may have $x \succsim y, y \succ^\pm z$ and $x \sim z$. Hence, \succsim may not be strictly responsive to \succ^\pm even when \succsim and \succ^\pm are complete. Indeed, it is easy to see that a semiorder for which

$$[x \succsim y \text{ and } y \succ^\pm z] \Rightarrow x \succ z, \quad (33)$$

must be a weak order.

Considering marginal traces on levels, it is now possible to envisage binary relations that are strictly responsive to each of their marginal traces without implying that they are (semi-)transitive or Ferrers.

Definition 4.2 (Conditions *TAC1*, *TAC2*)

We say that \succsim satisfies:

TAC1_i if

$$\left. \begin{array}{l} (x_i, a_{-i}) \succsim y \\ \text{and} \\ y \succsim (z_i, a_{-i}) \\ \text{and} \\ (z_i, b_{-i}) \succsim w \end{array} \right\} \Rightarrow (x_i, b_{-i}) \succsim w,$$

TAC2_i if

$$\left. \begin{array}{l} (x_i, a_{-i}) \succsim y \\ \text{and} \\ y \succsim (z_{-i}, a_{-i}) \\ \text{and} \\ w \succsim (x_i, b_{-i}) \end{array} \right\} \Rightarrow w \succsim (z_i, b_{-i}),$$

for all $x_i, z_i \in X_i$, all $a_{-i}, b_{-i} \in X_{-i}$ and all $y, w \in X$.

We say that \succsim satisfies *TAC1* (resp. *TAC2*) if it satisfies *TAC1_i* (resp. *TAC2_i*) for all $i \in N$. We sometimes write *TAC12* instead of *TAC1* and *TAC2*.

The first two conditions in the premise of *TAC1_i* and *TAC2_i* (the rationale for the names being that *TAC1* and *TAC2* are “intrA-attribute Cancellation” involving Three conditions in their premise) suggest that the level x_i is not worse than the level z_i . *TAC1_i* (resp. *TAC2_i*) then implies that x_i should upward dominate (resp. downward dominate) z_i .

Lemma 4.2 (Strict responsiveness to marginal traces on levels)

If \succsim is a complete binary relation on $X = \prod_{i=1}^n X_i$ then:

1. *TAC1_i* \Rightarrow [*AC1_i* and *AC3_i*]
2. *TAC2_i* \Rightarrow [*AC2_i* and *AC3_i*]
3. *TAC1_i* is equivalent to the completeness of \succsim_i^+ and the following condition:

$$[x \succsim y \text{ and } z_i \succ_i^+ x_i] \Rightarrow (z_i, x_{-i}) \succ y. \quad (34)$$

4. *TAC2_i* is equivalent to the completeness of \succsim_i^- and the following condition:

$$[x \succsim y \text{ and } y_i \succ_i^- w_i] \Rightarrow x \succ (w_i, y_{-i}). \quad (35)$$

5. If *TAC1_i* or *TAC2_i* hold then \succsim is independent for $\{i\}$ and \succsim_i is a weak-order. Furthermore when *TAC12* hold then $\succsim_i = \succsim_i^\pm$.
6. In the class of complete relations, *TAC1* and *TAC2* are independent conditions.
7. There are weakly independent semiorders verifying *TAC1* and *TAC2* that are not weak orders.

PROOF

Part 1. Let us first show that when \succsim is complete, *TAC1_i* \Rightarrow *AC1_i*. Suppose that *AC1_i* is violated so that $(x_i, a_{-i}) \succsim y$, $(z_i, b_{-i}) \succsim w$, *Not* [$(z_i, a_{-i}) \succsim y$] and *Not* [$(x_i, b_{-i}) \succsim w$]. Since \succsim is complete, we know that $y \succsim (z_i, a_{-i})$. Using *TAC1_i*, $(x_i, a_{-i}) \succsim y$, $y \succsim (z_i, a_{-i})$ and $(z_i, b_{-i}) \succsim w$ imply $(x_i, b_{-i}) \succsim w$, a contradiction. Hence *AC1_i* holds and \succsim_i^+ is complete.

Similarly suppose that *AC3_i* is violated so that $z \succsim (x_i, a_{-i})$, $(x_i, b_{-i}) \succsim y$, *Not* [$z \succsim (w_i, a_{-i})$] and *Not* [$(w_i, b_{-i}) \succsim y$]. Since \succsim is complete, we know

that $(w_i, a_{-i}) \succsim z$. Using $TAC1_i$, $(w_i, a_{-i}) \succsim z$, $z \succsim (x_i, a_{-i})$ and $(x_i, b_{-i}) \succsim y$ imply $(w_i, b_{-i}) \succsim y$, a contradiction. The proof of part 2 is similar.

Part 3. $[\Rightarrow]$. We know from part 1 that $AC1_i$ holds so that \succsim_i^+ is complete. Suppose now, in contradiction with (34) that $x \succsim y$, $z_i \succ_i^+ x_i$ and $y \succsim (z_i, x_{-i})$. We know that $Not[x_i \succsim_i^+ z_i]$, so that $(z_i, a_{-i}) \succsim w$ and $w \succ (x_i, a_{-i})$, for some $w \in X$ and some $a_{-i} \in X_{-i}$. Using $TAC1_i$, $x \succsim y$, $y \succsim (z_i, x_{-i})$ and $(z_i, a_{-i}) \succsim w$ imply $(x_i, a_{-i}) \succsim w$, a contradiction.

$[\Leftarrow]$. Suppose that $TAC1_i$ is violated so that $(x_i, a_{-i}) \succsim y$, $y \succsim (z_i, a_{-i})$ $(z_i, b_{-i}) \succsim w$ and $w \succ (x_i, b_{-i})$. This implies $Not[x_i \succsim_i^+ z_i]$. Since \succsim_i^+ is complete, we have $z_i \succ_i^+ x_i$. Using (34), $(x_i, a_{-i}) \succsim y$ and $z_i \succ_i^+ x_i$ would imply $(z_i, a_{-i}) \succ y$, a contradiction.

The proof of part 4 is similar.

Part 5. We give the proof using $TAC1_i$, the proof using $TAC2_i$ being similar. Suppose that $(x_i, a_{-i}) \succsim (y_i, a_{-i})$ and $Not[(x_i, b_{-i}) \succsim (y_i, b_{-i})]$, for some $x_i, y_i \in X_i$ and some $a_{-i}, b_{-i} \in X_{-i}$. Since \succsim is complete, we know that $(y_i, b_{-i}) \succ (y_i, a_{-i})$. Thus, since we know from part 1 that \succsim_i^+ is complete, we have $y_i \succ_i^+ x_i$. Using (34), $y_i \succ_i^+ x_i$ and $(x_i, a_{-i}) \succsim (y_i, a_{-i})$ would imply $(y_i, a_{-i}) \succ (y_i, a_{-i})$, a contradiction. Hence, \succsim is independent for $\{i\}$.

Using the reflexivity of \succsim , and (18), we have: $x_i \succsim_i^+ y_i \Rightarrow x_i \succsim_i y_i$. Let us show that $x_i \succ_i^+ y_i \Rightarrow x_i \succ_i y_i$, which will complete the proof since \succsim_i^+ is complete. Suppose that $x_i \succ_i^+ y_i$. Since \succsim is reflexive, we have $(y_i, a_{-i}) \succsim (y_i, a_{-i})$, for all $a_{-i} \in X_{-i}$. Using (34), we obtain $(x_i, a_{-i}) \succ (y_i, a_{-i})$, for all $a_{-i} \in X_{-i}$. We thus have $x_i \succ_i y_i$. Note that when both $TAC1_i$ and $TAC2_i$ hold, the above reasoning shows that $\succsim_i^\pm = \succsim_i$.

Parts 6 and 7: see examples A.4 and A.5 in appendix A. \square

As soon as \succsim is complete, the conjunction of $TAC1_i$ and $TAC2_i$ is therefore exactly what is needed to ensure the strict responsiveness of \succsim with respect to \succ_i^\pm . It also implies that \succsim is independent for $\{i\}$ and that \succsim_i is a weak order that is identical to \succsim_i^\pm . It does not imply however that \succsim is (semi-)transitive or Ferrers. Using (34) and (35) will facilitate the test of $TAC1_i$ and $TAC2_i$.

4.3 Marginal traces on levels and numerical representations

4.3.1 Marginal traces and numerical representations

In model $(L0)$, the role of u_i is merely to attach a number to each equivalence class of X_i / \sim_i^\pm while F passively recodes as +1's and -1's the presence or absence of \succsim for every possible combination of elements of X_i / \sim_i^\pm . Clearly,

as was the case in section 2.3, the situation changes as soon as additional properties are imposed on F .

Starting from the trivial model ($L0$) in which:

$$x \succsim y \Leftrightarrow F([u_i(x_i)]; [u_i(y_i)]) \geq 0,$$

we envisage:

- ($L1$) adding to ($L0$) the fact that $F([u_i(x_i)]; [u_i(x_i)]) \geq 0$,
- ($L2$) adding to ($L1$) the fact that $F([u_i(x_i)]; [u_i(y_i)]) = -F([u_i(y_i)]; [u_i(x_i)])$.

Furthermore, in each of ($L0$), ($L1$) and ($L2$), we envisage here the consequences of supposing that F is nondecreasing (resp. increasing) in its first n arguments and nonincreasing (resp. decreasing) in its last n arguments. These eight models are defined in table 3.

Table 3: Models involving traces on levels

$(L0)$	$x \succsim y \Leftrightarrow F([u_i(x_i)]; [u_i(y_i)]) \geq 0$
$(L1)$	$(L0)$ with $F([u_i(x_i)]; [u_i(x_i)]) \geq 0$
$(L2)$	$(L0)$ with $F([u_i(x_i)]; [u_i(y_i)]) = -F([u_i(y_i)]; [u_i(x_i)])$
.....	
$(L3)$	$(L0)$ with F nondecreasing, nonincreasing,
$(L4)$	$(L0)$ with F increasing, decreasing,
.....	
$(L5)$	$(L1)$ with F nondecreasing, nonincreasing,
$(L6)$	$(L1)$ with F increasing, decreasing,
.....	
$(L7)$	$(L2)$ with F nondecreasing, nonincreasing,
$(L8)$	$(L2)$ with F increasing, decreasing,

The implications between these various models are clear from their definitions. They are summarized in figure 1.

Lemma 4.3 (Necessary conditions)

1. Model ($L1$) implies that \succsim is reflexive,
2. Model ($L2$) implies that \succsim is complete,
3. Model ($L3$) implies AC1, AC2, AC3,
4. Model ($L8$) implies TAC1 and TAC2.

Figure 1: Implication between models involving traces on levels

$$\begin{array}{ccccccc}
(L8) & \Rightarrow & (L6) & \Rightarrow & (L4) & (\nearrow, \searrow) \\
\downarrow & & \downarrow & & \downarrow & \\
(L7) & \Rightarrow & (L5) & \Rightarrow & (L3) & (\nearrow, \searrow) \\
\downarrow & & \downarrow & & \downarrow & \\
(L2) & \Rightarrow & (L1) & \Rightarrow & (L0) & \emptyset
\end{array}$$

$$F(\alpha, \beta) = -F(\beta, \alpha) \quad F(\alpha, \alpha) \geq 0 \quad \emptyset$$

(\nearrow, \searrow) : increasing (resp. decreasing) in first (resp. last) n arguments

(\nearrow, \searrow) : nondecreasing (resp. nonincreasing) in first (resp. last) n arguments

PROOF

Parts 1 and 2 are obvious.

Part 3. The fact that $AC1$, $AC2$ and $AC3$ hold is easily shown using the properties of F . We take the case of $AC3$. Suppose that $(x_i, a_{-i}) \succsim y$ and $w \succsim (x_i, b_{-i})$ so that, abusing notations, $F([u_i(x_i), (u_j(a_j))_{j \neq i}]; [u_i(y_i)]) \geq 0$ and $F([u_i(w_i)]; [u_i(x_i), (u_j(b_j))_{j \neq i}]) \geq 0$.

If $u_i(z_i) > u_i(x_i)$ then $F([u_i(z_i), (u_j(a_j))_{j \neq i}]; [u_i(y_i)]) \geq 0$ so that $(z_i, a_{-i}) \succsim y$. Otherwise $u_i(x_i) \geq u_i(z_i)$ leads to $F([u_i(w_i)]; [u_i(z_i), (u_j(b_j))_{j \neq i}]) \geq 0$ so that $w \succsim (z_i, b_{-i})$.

Part 4. Suppose that $(x_i, a_{-i}) \succsim y$, $y \succsim (z_i, a_{-i})$, $(z_i, b_{-i}) \succsim w$ and $\text{Not}[(x_i, b_{-i}) \succsim w]$. Using the increasingness of F in its first n arguments, the last two conditions imply that $u_i(z_i) > u_i(x_i)$. But $(x_i, a_{-i}) \succsim y$ and $u_i(z_i) > u_i(x_i)$ imply $(z_i, a_{-i}) \succ y$, a contradiction. Hence the necessity of $TAC1$. The necessity of $TAC2$ is proved similarly. \square

We are now in position to characterize all our models involving traces on levels.

Theorem 4.1 (Models using traces on levels)

Let \succsim be a binary relation on a countable set $X = \prod_{i=1}^n X_i$.

1. Model (L1) holds iff \succsim is reflexive.
2. Model (L2) holds iff \succsim is complete.
3. Model (L4) holds iff \succsim satisfies $AC1$, $AC2$ and $AC3$.
4. Model (L6) holds iff \succsim is reflexive and satisfies $AC1$, $AC2$ and $AC3$.
5. Model (L7) holds iff \succsim is complete and satisfies $AC1$, $AC2$ and $AC3$.
6. Model (L8) holds iff \succsim is complete and satisfies $TAC1$ and $TAC2$.

PROOF

For all parts, necessity results from lemma 4.3 and the implications between our models. We show sufficiency.

Part 1 is immediate from part 1 of proposition 3.1.

Part 2. Take, on each $i \in N$, a function u_i satisfying (27) and define F on $[\prod_{i=1}^n u_i(X_i)]^2$ letting:

$$F([u_i(x_i)]; [u_i(y_i)]) = \begin{cases} +1 & \text{if } x \succ y, \\ 0 & \text{if } x \sim y, \\ -1 & \text{otherwise.} \end{cases} \quad (36)$$

The well-definedness of F follows from (21).

Part 3. Since $AC1$, $AC2$ and $AC3$ hold, we know from part 4 of lemma 4.1 that \succsim_i^\pm is a weak order. Since X_i is countable, there is a real-valued function u_i on X_i such that, for all $x_i, y_i \in X_i$:

$$x_i \succsim_i^\pm y_i \Leftrightarrow u_i(x_i) \geq u_i(y_i). \quad (37)$$

Take, on each $i \in N$, a function u_i satisfying (37) and define F on $[\prod_{i=1}^n u_i(X_i)]^2$ letting:

$$F([u_i(x_i)]; [u_i(y_i)]) = \begin{cases} + \exp(\sum_{i=1}^n (u_i(x_i) - u_i(y_i))) & \text{if } x \succsim y, \\ - \exp(\sum_{i=1}^n (u_i(y_i) - u_i(x_i))) & \text{otherwise.} \end{cases} \quad (38)$$

The well-definedness of F results from (21). The monotonicity properties of F follow from (20) and its definition. The proof of part 4 is immediate from part 3.

Part 5. Take, on each $i \in N$, a function u_i satisfying (37) and define F on $[\prod_{i=1}^n u_i(X_i)]^2$ letting:

$$F([u_i(x_i)]; [u_i(y_i)]) = \begin{cases} + \exp(\sum_{i=1}^n (u_i(x_i) - u_i(y_i))) & \text{if } x \succ y, \\ 0 & \text{if } x \sim y, \\ - \exp(\sum_{i=1}^n (u_i(y_i) - u_i(x_i))) & \text{otherwise.} \end{cases} \quad (39)$$

The well-definedness of F results from (21). It is skew symmetric by construction since \succsim is complete. Let us show that F is nondecreasing in its first n arguments. Suppose that $u_i(z_i) > u_i(x_i)$ so that $z_i \succ_i^\pm x_i$. If $x \succ y$, we know, using (20), that $(z_i, x_{-i}) \succ y$ and the conclusion follows from the definition of F . If $x \sim y$, we have, using (20), $(z_i, x_{-i}) \succsim y$ and the conclusion follows from the definition of F . If $\text{Not}[x \succsim y]$ we have either $(z_i, x_{-i}) \succ y$, $(z_i, x_{-i}) \sim y$, or $\text{Not}[(z_i, x_{-i}) \succsim y]$. In either case, the conclusion follows from the definition of F . The proof that F is nonincreasing in its last n argument is similar and is omitted.

Part 6. In view of the proof of part 5 above, since $TAC1$ and $TAC2$ imply $AC1$, $AC2$ and $AC3$ when \succsim is complete, sufficiency follows from lemma 4.2. \square

Corollary 4.1 (Links between models)

1. Models (L4) and (L3) are equivalent.
2. Models (L6) and (L5) are equivalent.
3. Model (L7) \Rightarrow Model (L6).

PROOF

All parts directly results from lemma 4.3 and theorem 4.1 \square

4.3.2 Remarks

Remark 4.1

A somewhat weaker form (using nondecreasingness and nonincreasingness) of part 4 of theorem 4.1 was noted in Greco et al. (2002) using our conditions $AC1$, $AC2$ and $AC3$. \bullet

Remark 4.2 (Skew symmetry and the role of 0)

It should be observed that increasingness and nondecreasingness are equivalent in our models except in the case of a “skew symmetric” function F (i.e., such that $F([u_i(x_i)]; [u_i(y_i)]) = -F([u_i(y_i)]; [u_i(x_i)])$). When F is skew symmetric, the value “0” plays a special role. This leads to distinguish the increasing case from the nondecreasing one, as in proposition 2.4 with semiorders and weak orders. This will often be the case below. \bullet

Remark 4.3 (Uniqueness)

It should be clear that the numerical representations envisaged in this section (see theorem 4.1) do not possess any remarkable uniqueness properties. It is not difficult however to analyze them. We take the case of model (L4).

Our proof shows that it is always possible to use functions u_i such that:

$$x_i \succsim_i^\pm y_i \Leftrightarrow u_i(x_i) \geq u_i(y_i). \tag{40}$$

This defines what could be called a regular representation of (L4). From the proof of theorem 4.1, it is clear that *any* u_i satisfying (40) may be used, i.e. we may apply, independently on each attribute, any increasing transformation to the functions u_i without altering the representation.

Other choices for u_i are possible however. Let us show that any function u_i such that:

$$x_i \succ_i^\pm y_i \Rightarrow u_i(x_i) > u_i(y_i), \tag{41}$$

can be used in a representation of model (L4).

The necessity of (41) is clear since $x_i \succ_i^\pm y_i$ implies either $x_i \succ_i^+ y_i$ or $x_i \succ_i^- y_i$. In the first case, we know that $(x_i, a_{-i}) \succ z$ and $\text{Not}[(y_i, a_{-i}) \succ z]$, for some $z \in X$ and some $a_{-i} \in X_{-i}$. In the second case, we obtain $w \succ (y_i, b_{-i})$ and $\text{Not}[w \succ (x_i, b_{-i})]$, for some $w \in X$ and some $b_{-i} \in X_{-i}$. Using the increasingness of F , either case implies $u_i(x_i) > u_i(y_i)$.

Conversely, it is clear that if u_i satisfies (41) then

$$u_i(x_i) = u_i(y_i) \Rightarrow x_i \sim_i^\pm y_i, \quad (42)$$

so that defining F using (38) as in the proof of theorem 4.1 leads to a well-defined function being increasing in its first n arguments and decreasing in its last n arguments.

It should be noted that any positive real-valued functions on \mathbb{R}^{2n} that is increasing in its first n arguments and decreasing in its last n arguments when restricted to $[\prod_{i=1}^n u_i(X_i)]^2$ may be used instead of the exponential chosen in (38) in order to define F . It is furthermore clear that only such functions may be used.

We have therefore described the set of all possible numerical representations in model (L4). It is not difficult to adapt the above reasoning to cover all the models envisaged in this section (see Bouyssou & Pirlot, 2002c, for details). •

Remark 4.4 (Weak orders)

When \succsim is a weak order, marginal traces on levels are confounded with marginal preferences. Hence, the above results can be greatly simplified. It is easy to see that, in this case,

- $[\succsim \text{ is weakly separable}] \Leftrightarrow [\succsim \text{ satisfies AC1}] \Leftrightarrow [\succsim \text{ satisfies AC2}] \Leftrightarrow [\succsim \text{ satisfies AC3}]$,
- $[\succsim \text{ is weakly independent}] \Leftrightarrow [\succsim \text{ satisfies TAC12}]$.

Using these observations, it is easy to show that, when \succsim is a weak order on a countable set $X = \prod_{i=1}^n X_i$, it always has a representation such that:

$$x \succsim y \Leftrightarrow U(u_1(x_1), \dots, u_n(x_n)) \geq U(u_1(y_1), \dots, u_n(y_n)) \geq 0. \quad (U1)$$

Furthermore, U can be chosen so that it is nondecreasing in all its arguments (resp. increasing) if and only if \succsim is weakly separable (resp. weakly independent).

Whereas the increasing case is well-known Krantz et al. (1971, theorem 7.1), the result in the nondecreasing case generalizes a result obtained in Blackorby, Primont, and Russell (1978) in the case in which $X \subseteq \mathbb{R}^n$. •

Remark 4.5 (Extension to the general case)

Most of our results are technically simple. Their extension to the general case (i.e. removing the hypothesis that X is countable) do not raise any serious difficulty beyond the well-known one of guaranteeing that equivalences and/or weak orders have a numerical representation (see Bouyssou & Pirlot, 2002c). •

Remark 4.6 (Left and right traces on levels)

For the sake of conciseness, we do not investigate here all possible models involving marginal traces on levels. Let us simply mention that the case in which $AC1$ and $AC2$ hold but not $AC3$ is of particular interest. The similarity of this case with that of interval orders (see Fishburn, 1970a, 1973, 1985) should be clear at this point. Most of our results can easily be modified to cover this case using a representation of the type:

$$x \succsim y \Leftrightarrow F([v_i(x_i)]; [w_i(y_i)]) \geq 0,$$

leaving room for relations in which the marginal traces \succsim_i^+ and \succsim_i^- may not be compatible. •

Remark 4.7 (Other extensions)

We restricted our attention in this paper to the analysis of conditions $AC1_i$, $AC2_i$, $AC3_i$, $TAC1_i$ and $TAC2_i$ when imposed on all $i \in N$. As observed in Greco et al. (2002), this might be overly restrictive. It is not difficult however to study the, rather awkward, models that obtain when these conditions are only imposed on some, but not all attributes. They amount to supposing that F is increasing (resp. decreasing) in some, but not necessarily all, of its n first (resp. last arguments).

Similarly, it is easy to generalize our conditions to subsets of attributes more general than a singleton. The study of the resulting models certainly deserves attention. In fact, when aggregating attributes, it might well happen that attributes interact in such a way that weak separability is violated (remember the famous example of the choice a dinner involving two attributes: the main course and the wine). This would forbid the use of $AC1$ or $AC2$ as done here. Imposing these conditions on the groups of “strongly interacting” attributes might however lead to useful models. Such models would be in the spirit of the process of “building criteria” by sub-aggregation as described in Bouyssou (1990), Roy (1996). •

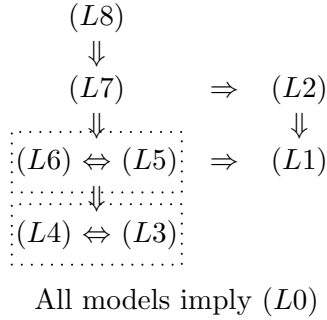
The main results in this section are summarized in table 4 and figure 2.

Table 4: Main results using traces on levels

Models	Definition	Conditions
(L0)	$x \succsim y \Leftrightarrow F([u_i(x_i)], [u_i(y_i)]) \geq 0$	\emptyset
(L1)	(L0) with $F([u_i(x_i)], [u_i(x_i)]) = 0$	refl.
(L2)	(L1) with $F([u_i(x_i)]; [u_i(y_i)]) = -F([u_i(y_i)]; [u_i(x_i)])$	cpl.
(L3) \Leftrightarrow (L4)	(L0) with $F(\nearrow, \searrow)$	AC123
(L5) \Leftrightarrow (L6)	(L1) with $F(\nearrow, \searrow)$	refl., AC123
(L7)	(L2) with $F(\nearrow, \searrow)$	cpl., AC123
(L8)	(L2) with $F(\nearrow, \searrow)$	cpl., TAC12

\nearrow means nondecreasing, \searrow means nonincreasing
 \nearrow means increasing, \searrow means decreasing
 refl. means reflexive, cpl. means complete

Figure 2: Implication between models resulting from theorem 4.1



5 Models using marginal traces on differences

This section follows the same path as in section 4 replacing marginal traces on levels by marginal traces on differences.

5.1 Complete marginal traces on differences

Definition 5.1 (Conditions RC1 and RC2)

Let \succsim be a binary relation on a set $X = \prod_{i=1}^n X_i$. For $i \in N$, this relation is said to satisfy:

$RC1_i$ if

$$\left. \begin{array}{l} (x_i, a_{-i}) \succsim (y_i, b_{-i}) \\ \text{and} \\ (z_i, c_{-i}) \succsim (w_i, d_{-i}) \end{array} \right\} \Rightarrow \left\{ \begin{array}{l} (x_i, c_{-i}) \succsim (y_i, d_{-i}) \\ \text{or} \\ (z_i, a_{-i}) \succsim (w_i, b_{-i}), \end{array} \right.$$

$RC2_i$ if

$$\left. \begin{array}{l} (x_i, a_{-i}) \succsim (y_i, b_{-i}) \\ \text{and} \\ (y_i, c_{-i}) \succsim (x_i, d_{-i}) \end{array} \right\} \Rightarrow \left\{ \begin{array}{l} (z_i, a_{-i}) \succsim (w_i, b_{-i}) \\ \text{or} \\ (w_i, c_{-i}) \succsim (z_i, d_{-i}), \end{array} \right.$$

for all $x_i, y_i, z_i, w_i \in X_i$ and all $a_{-i}, b_{-i}, c_{-i}, d_{-i} \in X_{-i}$. We say that \succsim satisfies $RC1$ (resp. $RC2$) if it satisfies $RC1_i$ (resp. $RC2_i$) for all $i \in N$. We sometimes write $RC12$ as shorthand for $RC1$ and $RC2$.

Condition $RC1_i$ suggests that either (x_i, y_i) is a larger preference difference than (z_i, w_i) or vice versa. Indeed, it is easy to see that supposing $\text{Not}[(x_i, y_i) \succsim_i^* (z_i, w_i)]$ and $\text{Not}[(z_i, w_i) \succsim_i^* (x_i, y_i)]$ leads to a violation of $RC1_i$. Hence, $RC1_i$ is equivalent to supposing the completeness of \succsim_i^* . Similarly, $RC2_i$ suggests that the two opposite differences (x_i, y_i) and (y_i, x_i) are linked. In terms of the relation \succsim_i^* , it says that if the preference difference between x_i and y_i is not at least as large as the preference difference between z_i and w_i then the preference difference between y_i and x_i should be at least as large as the preference difference between w_i and z_i . We summarize these observations in the following lemma; we omit its straightforward proof.

Lemma 5.1 (Complete Traces on differences)

We have:

1. $[\succsim_i^* \text{ is complete}] \Leftrightarrow RC1_i$,
2. $RC2_i \Leftrightarrow$
 $[\text{for all } x_i, y_i, z_i, w_i \in X_i, \text{Not}[(x_i, y_i) \succsim_i^* (z_i, w_i)] \Rightarrow (y_i, x_i) \succsim_i^* (w_i, z_i)],$
3. $[\succsim_i^{**} \text{ is complete}] \Leftrightarrow [RC1_i \text{ and } RC2_i].$

Condition $RC1$ was introduced in Bouyssou (1986) under the name “weak cancellation”. Technically $RC1_i$ amounts to defining a biorder, in the sense of Ducamp and Falmagne (1969) and Doignon, Ducamp, and Falmagne (1984), between the sets X_i^2 and X_{-i}^2 . The extension of condition $RC1$ to subsets of attributes is central in Vind (1991) where this condition is called “independence”. Condition $RC2$ was first proposed in Bouyssou and Pirlot (2002a, 1999), Bouyssou et al. (1997). We note below a number of other important facts about $RC1$ and $RC2$.

Lemma 5.2 (Consequences of RC1 and RC2)

1. If \succsim satisfies RC1_{*i*} then it is weakly separable for $i \in N$,
2. If \succsim satisfies RC2 then it is independent and either reflexive or irreflexive,
3. Reflexivity, independence and RC1 are independent conditions,
4. In the class of complete relations, RC1 and RC2 are independent conditions.

PROOF

Part 1. Suppose that $(x_i, a_{-i}) \succ (x_i, b_{-i})$ and $(y_i, b_{-i}) \succ (y_i, a_{-i})$. This implies $(x_i, a_{-i}) \succsim (x_i, b_{-i})$ and $(y_i, b_{-i}) \succsim (y_i, a_{-i})$ so that RC1_{*i*} implies either $(y_i, a_{-i}) \succsim (y_i, b_{-i})$ or $(x_i, b_{-i}) \succsim (x_i, a_{-i})$, a contradiction.

Part 2. If $(x_i, a_{-i}) \succsim (x_i, b_{-i})$, RC2_{*i*} implies $(y_i, a_{-i}) \succsim (y_i, b_{-i})$ for all $y_i \in X_i$ so that \succsim is independent. It is clear that an independent relation is either reflexive or irreflexive.

Part 3. In order to show that these three properties are completely independent, we need $2^3 = 8$ examples. It is easy to build a relation \succsim that does not satisfy RC1 and is neither reflexive nor independent (e.g. take $X = \{a, b\} \times \{z, w\}$ and let \succsim be an empty relation on X except that $(a, z) \succsim (b, z)$ and $(b, w) \succsim (a, w)$). Any relation \succsim satisfying the additive utility model satisfies the three properties. The other six examples are provided in appendix B.

Part 4. Any relation \succsim satisfying the additive utility model is complete and satisfies both RC1 and RC2. We provide in appendix B the three remaining examples. \square

5.2 Strict responsiveness to traces on differences

Definition 5.2 (Condition TC)

Let \succsim be a binary relation on a set $X = \prod_{i=1}^n X_i$. For $i \in N$, this relation is said to satisfy:

TC_{*i*} if

$$\left. \begin{array}{l} (x_i, a_{-i}) \succsim (y_i, b_{-i}) \\ \text{and} \\ (z_i, b_{-i}) \succsim (w_i, a_{-i}) \\ \text{and} \\ (w_i, c_{-i}) \succsim (z_i, d_{-i}) \end{array} \right\} \Rightarrow (x_i, c_{-i}) \succsim (y_i, d_{-i}),$$

for all $x_i, y_i, z_i, w_i \in X_i$ and all $a_{-i}, b_{-i}, c_{-i}, d_{-i} \in X_{-i}$. We say that \succsim satisfies TC if it satisfies TC_{*i*} for all $i \in N$.

Condition TC_i (Triple Cancellation) is a classical cancellation condition that has been often used (see Krantz et al., 1971; Wakker, 1989) in the analysis of the additive utility model (1). As shown below, it implies both $RC1$ and $RC2$ when \succsim is complete. We refer to Wakker (1988, 1989) for a detailed analysis of TC including its interpretation in terms of difference of preference.

As soon as \succsim is complete, the following lemma shows that TC_i is exactly what is needed to ensure the strict responsiveness of \succsim w.r.t. \succsim_i^{**} .

Lemma 5.3 (Strict responsiveness to marginal traces on differences)

1. If \succsim is complete, TC_i implies $RC1_i$ and $RC2_i$,
2. If TC_i holds and \succsim is complete, $[x \succsim y \text{ and } (z_i, w_i) \succ_i^{**} (x_i, y_i)] \Rightarrow (z_i, x_{-i}) \succ (w_i, y_{-i})$.

PROOF

Part 1. In contradiction with $RC1_i$, suppose that $(x_i, a_{-i}) \succsim (y_i, b_{-i})$, $(z_i, c_{-i}) \succsim (w_i, d_{-i})$, $Not[(z_i, a_{-i}) \succsim (w_i, b_{-i})]$ and $Not[(x_i, c_{-i}) \succsim (y_i, d_{-i})]$. Since \succsim is complete, we have $(w_i, b_{-i}) \succ (z_i, a_{-i})$. Using TC_i , $(x_i, a_{-i}) \succsim (y_i, b_{-i})$, $(w_i, b_{-i}) \succ (z_i, a_{-i})$ and $(z_i, c_{-i}) \succsim (w_i, d_{-i})$ imply $(x_i, c_{-i}) \succsim (y_i, d_{-i})$, a contradiction.

Similarly suppose, in contradiction with $RC2_i$ that $(x_i, a_{-i}) \succsim (y_i, b_{-i})$, $(y_i, c_{-i}) \succsim (x_i, d_{-i})$, $Not[(z_i, a_{-i}) \succsim (w_i, b_{-i})]$ and $Not[(w_i, c_{-i}) \succsim (z_i, d_{-i})]$. Since \succsim is complete, we know that $(w_i, b_{-i}) \succ (z_i, a_{-i})$. Using TC_i , $(w_i, b_{-i}) \succ (z_i, a_{-i})$, $(x_i, a_{-i}) \succsim (y_i, b_{-i})$ and $(y_i, c_{-i}) \succsim (x_i, d_{-i})$ imply $(w_i, c_{-i}) \succsim (z_i, d_{-i})$, a contradiction.

Part 2. Using (23), we get $(z_i, x_{-i}) \succsim (w_i, y_{-i})$. Suppose that $(w_i, y_{-i}) \succsim (z_i, x_{-i})$. From part 1 and lemma 5.1, we know that \succsim_i^{**} is complete. We thus have $(z_i, w_i) \succ_i^{**} (x_i, y_i) \Leftrightarrow Not[(x_i, y_i) \succsim_i^{**} (z_i, w_i)] \Leftrightarrow [Not[(x_i, y_i) \succsim_i^* (z_i, w_i)] \text{ or } Not[(w_i, z_i) \succsim_i^* (y_i, x_i)]]$. In the first case we have $Not[(x_i, c_{-i}) \succsim (y_i, d_{-i})]$ and $(z_i, c_{-i}) \succsim (w_i, d_{-i})$, for some $c_{-i}, d_{-i} \in X_{-i}$. Using TC_i , $x \succsim y$, $(w_i, y_{-i}) \succsim (z_i, x_{-i})$ and $(z_i, c_{-i}) \succsim (w_i, d_{-i})$ imply $(x_i, c_{-i}) \succsim (y_i, d_{-i})$, a contradiction. The other case is similar. \square

5.3 Marginal traces on differences and numerical representations

Starting from the trivial model ($D0$), introduced in section 3.5, in which:

$$x \succsim y \Leftrightarrow G([p_i(x_i, y_i)]) \geq 0,$$

we envisage:

- model (D1) adding to (D0) the fact that $p_i(x_i, x_i) = 0$, for all $i \in N$ and all $x_i \in X_i$,
- model (D2) adding to (D1) the skew symmetry of each p_i , i.e. $p_i(x_i, y_i) = -p_i(y_i, x_i)$, for all $i \in N$ and all $x_i, y_i \in X_i$,
- model (D3) adding to (D2) the oddness of G , i.e. $G(\mathbf{x}) = -G(-\mathbf{x})$, abusing notation in an obvious way.

As before, we furthermore envisage the consequences of supposing that in each of the above four models ((D0), (D1), (D2) and (D3)), G is nondecreasing or increasing in each of its arguments. This leads to a total of 12 different models, the definition of which being summarized in table 5

Table 5: Models involving traces on differences

(D0)	$x \succsim y \Leftrightarrow G([p_i(x_i, y_i)]) \geq 0$
(D1)	(D0) with $p_i(x_i, x_i) = 0$
(D2)	(D1) with $p_i(x_i, y_i) = -p_i(y_i, x_i)$
(D3)	(D2) with G odd
.....	
(D4)	(D0) with G nondecreasing
(D8)	(D0) with G increasing
.....	
(D5)	(D1) with G nondecreasing
(D9)	(D1) with G increasing
.....	
(D6)	(D2) with G nondecreasing
(D10)	(D2) with G increasing
.....	
(D7)	(D3) with G nondecreasing
(D11)	(D3) with G increasing

The implications between these various models are clear from their definitions. They are summarized in figure 3. The following lemma takes note of the relations between these models and our conditions involving traces on differences.

Lemma 5.4 (Necessary conditions)

1. Model (D1) implies that \succsim is independent,
2. Model (D3) implies that \succsim is complete,
3. Model (D4) implies RC1,

Figure 3: Implication between models involving traces on differences

$$\begin{array}{cccccc}
(D11) & \Rightarrow & (D10) & \Rightarrow & (D9) & \Rightarrow & (D8) & \text{inc.} \\
\downarrow & & \downarrow & & \downarrow & & \downarrow & \\
(D7) & \Rightarrow & (D6) & \Rightarrow & (D5) & \Rightarrow & (D4) & \text{nondec.} \\
\downarrow & & \downarrow & & \downarrow & & \downarrow & \\
(D3) & \Rightarrow & (D2) & \Rightarrow & (D1) & \Rightarrow & (D0) & \emptyset
\end{array}$$

$$G \text{ odd} \quad p_i \text{ sk-sym.} \quad p_i(x_i, x_i) = 0 \quad \emptyset$$

inc. means increasing in all arguments
nondec. means nondecreasing in all arguments
sk-sym. means skew symmetric

4. Model (D6) implies RC2,

5. Model (D11) implies TC,

PROOF

Part 1. In model (D1), we have $p_i(x_i, x_i) = 0$, so that $(x_i, a_{-i}) \succsim (x_i, b_{-i}) \Leftrightarrow G(0, (p_j(a_j, b_j))_{j \neq i}) \geq 0 \Leftrightarrow (y_i, a_{-i}) \succsim (y_i, b_{-i})$. Hence \succsim is independent.

Part 2. By definition of model (D3), $\text{Not}[x \succsim y] \Leftrightarrow G([p_i(x_i, y_i)]) < 0$ so that, using the skew symmetry of the p_i 's and the oddness of G , $G([p_i(y_i, x_i)]) > 0$ which implies $y \succsim x$. Hence \succsim is complete.

Part 3. Suppose that $(x_i, a_{-i}) \succsim (y_i, b_{-i})$ and $(z_i, c_{-i}) \succsim (w_i, d_{-i})$. Using model (D4) we have:

$$\begin{aligned}
G(p_i(x_i, y_i), (p_j(a_j, b_j))_{j \neq i}) &\geq 0 \text{ and} \\
G(p_i(z_i, w_i), (p_j(c_j, d_j))_{j \neq i}) &\geq 0.
\end{aligned}$$

If $p_i(x_i, y_i) \geq p_i(z_i, w_i)$, then using the nondecreasingness of G , we have $G(p_i(x_i, y_i), (p_j(c_j, d_j))_{j \neq i}) \geq 0$ so that $(x_i, c_{-i}) \succsim (y_i, d_{-i})$. If $p_i(z_i, w_i) > p_i(x_i, y_i)$, we have $G(p_i(z_i, w_i), (p_j(a_j, b_j))_{j \neq i}) \geq 0$ so that $(z_i, a_{-i}) \succsim (w_i, b_{-i})$. Therefore RC1 holds.

Part 4. Suppose that $(x_i, a_{-i}) \succsim (y_i, b_{-i})$ and $(y_i, c_{-i}) \succsim (x_i, d_{-i})$. Using model (D6) we have:

$$\begin{aligned}
G(p_i(x_i, y_i), (p_j(a_j, b_j))_{j \neq i}) &\geq 0 \text{ and} \\
G(p_i(y_i, x_i), (p_j(c_j, d_j))_{j \neq i}) &\geq 0.
\end{aligned}$$

If $p_i(x_i, y_i) \geq p_i(z_i, w_i)$, the skew symmetry of p_i implies $p_i(w_i, z_i) \geq p_i(y_i, x_i)$ so that $(w_i, c_{-i}) \succsim (z_i, d_{-i})$ using the nondecreasingness of G . Similarly, if $p_i(z_i, w_i) > p_i(x_i, y_i)$ we have, using the nondecreasingness of G , $(z_i, a_{-i}) \succsim (w_i, b_{-i})$. Therefore, RC2 holds.

Part 5. Suppose that $(x_i, a_{-i}) \succsim (y_i, b_{-i})$, $(z_i, b_{-i}) \succsim (w_i, a_{-i})$, $(w_i, c_{-i}) \succsim (z_i, d_{-i})$ and $\text{Not}[(x_i, c_{-i}) \succsim (y_i, d_{-i})]$. Using model (D11) we know that:

$$\begin{aligned} G(p_i(x_i, y_i), (p_j(a_j, b_j))_{j \neq i}) &\geq 0, \\ G(p_i(z_i, w_i), (p_j(b_j, a_j))_{j \neq i}) &\geq 0, \\ G(p_i(w_i, z_i), (p_j(c_j, d_j))_{j \neq i}) &\geq 0 \text{ and} \\ G(p_i(x_i, y_i), (p_j(c_j, d_j))_{j \neq i}) &< 0. \end{aligned}$$

Using the oddness of G , its increasingness and the skew symmetry of the p_i 's, the first two inequalities imply $p_i(x_i, y_i) \geq p_i(w_i, z_i)$ whereas the last two imply that $p_i(x_i, y_i) < p_i(w_i, z_i)$, a contradiction. \square

We are now in position to characterize our models involving traces on differences.

Theorem 5.1 (Models using traces on differences)

Let \succsim be a binary relation on a countable set $X = \prod_{i=1}^n X_i$.

1. Model (D2) holds iff \succsim is independent,
2. Model (D3) holds iff \succsim is complete and independent,
3. Model (D8) holds iff \succsim satisfies RC1,
4. Model (D9) holds iff \succsim is independent and satisfies RC1,
5. Model (D10) holds iff \succsim satisfies RC1 and RC2,
6. Model (D7) holds iff \succsim is complete and satisfies RC1 and RC2,
7. Model (D11) holds iff \succsim is complete and satisfies TC.

PROOF

Necessity of all parts results from lemma 5.4 and the implications between our models. We show sufficiency.

Part 1. Since \succsim is independent, we know that, for all $i \in N$ and all $x_i, y_i \in X_i$, $(x_i, x_i) \sim_i^* (y_i, y_i)$ and, thus, $(x_i, x_i) \sim_i^{**} (y_i, y_i)$. Since, X_i is countable, there is a one-to-one correspondence p_i between $X_i^2 \setminus \{(x_i, x_i) : x_i \in X_i\}$ and some countable subset of \mathbb{R} . It can always be chosen so that, for all $x_i, y_i \in X_i$, $p_i(x_i, y_i) = -p_i(y_i, x_i)$. Let $p_i(x_i, x_i) = 0$, for all $x_i \in X_i$. By construction, p_i is skew symmetric. Furthermore, we have:

$$p_i(x_i, y_i) = p_i(z_i, w_i) \Leftrightarrow \begin{cases} x_i = y_i \text{ and } z_i = w_i \\ \text{or} \\ x_i = z_i \text{ and } y_i = w_i. \end{cases} \quad (43)$$

Consider, on all $i \in N$, a function p_i satisfying (43) and define G on $\prod_{i=1}^n p_i(X_i^2)$ letting:

$$G([p_i(x_i, y_i)]) = \begin{cases} +1 & \text{if } x \succsim y, \\ -1 & \text{otherwise.} \end{cases} \quad (44)$$

Since, by hypothesis, $(x_i, x_i) \sim_i^* (y_i, y_i)$, it is easy to see that that G is well-defined.

Part 2. Consider functions p_i defined as in part 2 and define G on $\prod_{i=1}^n p_i(X_i^2)$ letting:

$$G([p_i(x_i, y_i)]) = \begin{cases} +1 & \text{if } x \succ y, \\ 0 & \text{if } x \sim y, \\ -1 & \text{otherwise.} \end{cases} \quad (45)$$

Since, by hypothesis, $(x_i, x_i) \sim_i^* (y_i, y_i)$, it is easy to see that G is well-defined. It is odd since \succsim is complete.

Part 3. Since $RC1$ holds, we know that \succsim_i^* is a weak order. Since X_i is countable, there is a real-valued function p_i on X_i^2 such that:

$$(x_i, y_i) \succsim_i^* (z_i, w_i) \Leftrightarrow p_i(x_i, y_i) \geq p_i(z_i, w_i). \quad (46)$$

Given such a particular numerical representation p_i of \succsim_i^* for $i = 1, 2, \dots, n$, define G on $p_i(X_i^2)$ as follows:

$$G([p_i(x_i, y_i)]) = \begin{cases} +\exp(\sum_{i=1}^n p_i(x_i, y_i)) & \text{if } x \succsim y, \\ -\exp(-\sum_{i=1}^n p_i(x_i, y_i)) & \text{otherwise.} \end{cases} \quad (47)$$

The well-definedness of G follows from (24) and the definition of the p_i 's. To show that G is increasing, suppose that $p_i(z_i, w_i) > p_i(x_i, y_i)$, i.e. that $(z_i, w_i) \succ_i^* (x_i, y_i)$. If $x \succsim y$, we know from (23) that $(z_i, x_{-i}) \succsim (w_i, y_{-i})$ and the conclusion follows from the definition of G . If $\text{Not}[x \succsim y]$ we have either $\text{Not}[(z_i, x_{-i}) \succsim (w_i, y_{-i})]$ or $(z_i, x_{-i}) \succsim (w_i, y_{-i})$. In either case, the conclusion follows from the definition of G .

Part 4 immediately follows from part 3. Indeed, when \succsim is independent, $(x_i, x_i) \sim_i^* (y_i, y_i)$, for all $i \in N$ and all $x_i, y_i \in X_i$. It is therefore always possible to choose a function p_i satisfying (46) so that $p_i(x_i, x_i) = 0$.

Part 5. Since $RC1_i$ and $RC2_i$ hold, we know from lemma 5.2 that \succsim_i^{**} is complete so that it is a weak order. This implies that \succsim_i^* is a weak order and, since X_i is countable, there is a real-valued function q_i on X_i^2 satisfying (46). Given a particular numerical representation q_i of \succsim_i^* , let $p_i(x_i, y_i) = q_i(x_i, y_i) - q_i(y_i, x_i)$. It is obvious that p_i is skew symmetric and represents \succsim_i^{**} . Define G as in (47). Its well-definedness results from (26). To show

that G is increasing, suppose that $p_i(z_i, w_i) > p_i(x_i, y_i)$, i.e. that $(z_i, w_i) \succ_i^{**}(x_i, y_i)$. By construction, this implies $(z_i, w_i) \succ_i^*(x_i, y_i)$. The increasingness of G is proved as in part 3 using (23).

Part 6. Define the functions p_i as in part 5 and define G letting

$$G([p_i(x_i, y_i)]) = \begin{cases} + \exp(\sum_{i=1}^n p_i(x_i, y_i)) & \text{if } x \succ y, \\ 0 & \text{if } x \sim y, \\ - \exp(-\sum_{i=1}^n p_i(x_i, y_i)) & \text{otherwise,} \end{cases} \quad (48)$$

Since \succsim is complete, G is odd. It is well-defined in view of (26) and the definition of the p_i 's. Its nondecreasingness follows from (23) and (25).

Part 7. Since TC holds and \succsim is complete, we know from part 1 of lemma 5.3 that $RC1$ and $RC2$ hold. Define p_i and G as in part 6. The increasingness of F follows from part 2 of lemma 5.3. \square

Corollary 5.1 (Links between models)

1. Models (D1) and (D2) are equivalent.
2. Models (D4) and (D8) are equivalent.
3. Models (D5) and (D9) are equivalent.
4. Models (D6) and (D10) are equivalent.
5. Model (D7) \Rightarrow Model (D10).

PROOF

All parts result directly from lemma 5.4 and theorem 5.1. \square

5.4 Remarks

Remark 5.1 (Goldstein (1991))

Models (D8) and (D4) were introduced by Goldstein (1991) as particular cases of his “decomposable thresholds models”. He already noted their equivalence. \bullet

Remark 5.2 (Equivalence between models)

Some care must be exercised when adding additional conditions to equivalent models. These additions may turn equivalent models into distinct ones. An example of this situation occurred in theorem 5.1. While models (D1) and (D2) are equivalent this is no more the case when it is required that G is nondecreasing in all its arguments. Models (D5) and (D6) are indeed distinct. \bullet

Remark 5.3 (Marginal preferences)

The nontransitivity and/or noncompleteness of \succsim combined with that of \succsim_i may obscure some features of models involving traces on differences. For the sake of completeness, we sum up a few useful observations in the following proposition.

Proposition 5.1 (Properties of models using differences)

Let \succsim be a binary relation on $X = \prod_{i=1}^n X_i$ and $J \subseteq N$.

1. If \succsim satisfies model (D5) then: $[x_i \succ_i y_i \text{ for all } i \in J \subseteq N] \Rightarrow \text{Not}[y_J \succsim_J x_J]$.
2. If \succsim satisfies model (D6) then:
 - \succsim_i is complete,
 - $[x_i \succ_i y_i \text{ for all } i \in J \subseteq N] \Rightarrow [x_J \succ_J y_J]$.
3. If \succsim satisfies model (D11) then:
 - $[x_i \succsim_i y_i \text{ for all } i \in J \subseteq N] \Rightarrow [x_J \succsim_J y_J]$,
 - $[x_i \succsim_i y_i \text{ for all } i \in J \subseteq N, x_j \succ_j y_j, \text{ for some } j \in J] \Rightarrow [x_J \succ_J y_J]$.

PROOF

Part 1. Using obvious notation, $x_i \succ_i y_i$ implies $\text{Not}[y_i \succsim_i x_i]$ so that $G(p_i(y_i, x_i), \mathbf{0}) < 0$. Since $G(\mathbf{0}) \geq 0$, we know that $p_i(y_i, x_i) < 0$, using the nondecreasingness of G . Select any $j \in J$. Starting from $G(p_j(y_j, x_j), \mathbf{0}) < 0$, using the nondecreasingness of G and the fact that $p_i(y_i, x_i) < 0$, for all $i \in J$, we obtain $G((p_j(y_j, x_j))_{i \in J}, \mathbf{0}) < 0$. This implies $\text{Not}[y_J \succsim_J x_J]$.

Part 2. $\text{Not}[x_i \succsim_i y_i]$ and $\text{Not}[y_i \succsim_i x_i]$ imply $G(p_i(x_i, y_i), \mathbf{0}) < 0$ and $G(p_i(y_i, x_i), \mathbf{0}) < 0$. Since $G(\mathbf{0}) \geq 0$ and G is nondecreasing, this implies $p_i(x_i, y_i) < 0$ and $p_i(y_i, x_i) < 0$, which contradicts the skew symmetry of p_i . Hence \succsim_i is complete.

Observe that $x_i \succ_i y_i$ is equivalent to $G(p_i(x_i, y_i), \mathbf{0}) \geq 0$ and $G(p_i(y_i, x_i), \mathbf{0}) < 0$. Since $G(\mathbf{0}) \geq 0$, we know that $p_i(y_i, x_i) < 0$, using the nondecreasingness of G . The skew symmetry of p_i implies $p_i(x_i, y_i) > 0 > p_i(y_i, x_i)$ and the desired property easily follows using the nondecreasingness of G .

Part 3. Since G is increasing and odd, we have $x_i \succsim_i y_i \Leftrightarrow p_i(x_i, y_i) \geq 0$. The desired properties easily follow from the increasingness of G and $G(\mathbf{0}) = 0$. \square

Except for model (D11), the monotonicity properties of our models linking \succsim and \succsim_i may seem disappointing. Such properties should however be

analyzed keeping in mind that we are dealing with possibly nontransitive and/or incomplete preferences. In such a framework, some “obvious properties” may not always be desirable. For example, when the relations \sim_i are not transitive, it may not be reasonable to impose that:

$$[x_i \sim_i y_i \text{ for all } i \in J] \Rightarrow [x_J \sim_J y_J],$$

which would forbid any interaction between separately non-noticeable differences on each attribute (on this point see Gilboa & Lapsan, 1995 or Pirlot & Vincke, 1997). Clearly, if nice monotonicity properties are looked for, one should use traces (see lemmas 3.3 and 5.3). •

Remark 5.4 (Interpretation of TC)

It is not difficult to show that, when \succsim is complete, $[RC1, RC2 \text{ and } (x \sim y \text{ and } (z_i, w_i) \succ_i^{**} (x_i, y_i) \Rightarrow (z_i, x_{-i}) \succ (w_i, y_{-i}))] \Leftrightarrow TC$. This offers an additional interpretation of TC and shows that the only difference between $(D11)$ and $(D7)$ is the possible failure in $(D7)$ of “strict monotonicity” with respect to \succ_i^{**} for pairs such that $x \sim y$. •

Remark 5.5 (Uniqueness)

As in section 4.3, the uniqueness properties of the representations exhibited in theorem 5.1 are very weak. Again, numerical representations are only used here as guidelines to investigate the consequences of some particular conditions on \succsim and not as a basis to derive assessment procedures.

Let us analyze the uniqueness properties of model $(D8)$. Our proof shows that it is always possible to use functions p_i such that:

$$(x_i, y_i) \succsim_i^* (z_i, w_i) \Leftrightarrow p_i(x_i, y_i) \geq p_i(z_i, w_i). \quad (49)$$

This could be called a regular representation of model $(D8)$. From the proof of theorem 5.1, it is clear that *any* p_i satisfying (49) may be used, i.e. we may apply, independently on each attribute, any increasing transformation to the functions p_i without altering the representation.

Other choices for p_i are possible however. Let us show that any function p_i such that:

$$(x_i, y_i) \succ_i^* (z_i, w_i) \Rightarrow p_i(x_i, y_i) > p_i(z_i, w_i), \quad (50)$$

can be used in a representation of model $(D8)$.

The necessity of (50) is clear since $(x_i, y_i) \succ_i^* (z_i, w_i)$ implies $(x_i, a_{-i}) \succ (y_i, b_{-i})$ and $\text{Not}[(z_i, a_{-i}) \succ (w_i, b_{-i})]$, for some $a_{-i}, b_{-i} \in X_{-i}$. Using the increasingness of G in model $(D8)$ this implies $p_i(x_i, y_i) > p_i(z_i, w_i)$.

Conversely, it is clear that if p_i satisfies (50) then

$$p_i(x_i, y_i) = p_i(z_i, w_i) \Rightarrow (x_i, y_i) \sim_i^* (z_i, w_i), \quad (51)$$

so that defining G using (47) as in the proof of theorem 5.1 leads to a well-defined function being increasing in its arguments.

It should be noted that any function h from \mathbb{R}^n into $(0, +\infty)$ that is increasing in all arguments when restricted to $\prod_{i=1}^n p_i(X_i^2)$ may be used in order to define G letting:

$$G([p_i(x_i, y_i)]) = \begin{cases} +h([p_i(x_i, y_i)]) & \text{if } x \succsim y, \\ -h([-p_i(x_i, y_i)]) & \text{otherwise.} \end{cases} \quad (52)$$

It is furthermore clear that only such functions may be used.

We have therefore described the set of all possible numerical representations in model model (D8). We shall need below the exact statement of the degrees of freedom we have for choosing the functions p_i in our models. This is summarized in the next lemma.

Lemma 5.5 (Uniqueness of p_i)

1. Let \succsim satisfy model (D8). A real-valued function p_i on X_i^2 may be used in the representation of \succsim in model (D8) iff

$$(z_i, w_i) \succ_i^* (x_i, y_i) \Rightarrow p_i(z_i, w_i) > p_i(x_i, y_i). \quad (53)$$

2. Let \succsim satisfy model (D9). A real-valued function p_i on X_i^2 may be used in the representation of \succsim in model (D9) iff

$$p_i(x_i, x_i) = 0 \text{ and} \quad (54)$$

$$(z_i, w_i) \succ_i^{**} (x_i, y_i) \Rightarrow p_i(z_i, w_i) > p_i(x_i, y_i). \quad (55)$$

3. Let \succsim satisfy model (D7) or (D10). A real-valued function p_i on X_i^2 may be used in the representation of \succsim in model (D7) or (D10) iff

$$p_i(x_i, y_i) = -p_i(y_i, x_i) \text{ and} \quad (56)$$

$$(z_i, w_i) \succ_i^{**} (x_i, y_i) \Rightarrow p_i(z_i, w_i) > p_i(x_i, y_i). \quad (57)$$

4. Let \succsim satisfy model (D11). A real-valued function p_i on X_i^2 may be used in the representation of \succsim in model (D11) iff

$$p_i(x_i, y_i) = -p_i(y_i, x_i), \quad (58)$$

$$(z_i, w_i) \succ_i^{**} (x_i, y_i) \Rightarrow p_i(z_i, w_i) > p_i(x_i, y_i) \text{ and} \quad (59)$$

$$\left. \begin{array}{l} (z_i, w_i) \sim_i^{**} (x_i, y_i) \text{ and} \\ \exists a_{-i}, b_{-i} \in X_{-i} \text{ s.t. } (x_i, a_{-i}) \sim (y_i, b_{-i}) \end{array} \right\} \Rightarrow p_i(z_i, w_i) = p_i(x_i, y_i). \quad (60)$$

PROOF

The observations preceding the statement of the lemma prove part 1. The proof of parts 2 and 3 follows from that of part 1.

Part 4. The necessity of (58) and (59) is clear. Suppose that (60) is violated. One would then have $(z_i, w_i) \sim_i^{**} (x_i, y_i)$, $(x_i, a_{-i}) \sim (y_i, b_{-i})$ for some $a_{-i}, b_{-i} \in X_{-i}$ and $p_i(z_i, w_i) \neq p_i(x_i, y_i)$. Since G is strictly increasing, $G(p_i(z_i, w_i), (p_j(a_j, b_j))_{j \neq i}) \neq 0$ while $(z_i, a_{-i}) \sim (w_i, b_{-i})$, a contradiction.

Sufficiency. Consider, on each $i \in N$, any function p_i satisfying (58), (59) and (60) and define G as in (48). The well-definedness of G follows from (58) and (26) since $p_i(z_i, w_i) = p_i(x_i, y_i) \Rightarrow (z_i, w_i) \sim_i^{**} (x_i, y_i)$. For proving increasingness, suppose $p_i(z_i, w_i) > p_i(x_i, y_i)$. This implies that $(z_i, w_i) \succ_i^{**} (x_i, y_i)$. If $x \succ y$, (25) implies $(z_i, x_{-i}) \succ (w_i, y_{-i})$ and the conclusion follows from the definition of G . If $x \sim y$ we have $G([p_i(x_i, y_i)]) = 0$. Consider two cases. If $(z_i, w_i) \succ_i^{**} (x_i, y_i)$, then part 2 of lemma 5.3 implies $(z_i, x_{-i}) \succ (w_i, y_{-i})$ and the conclusion follows from the definition of G . If $(z_i, w_i) \sim_i^{**} (x_i, y_i)$ then, by (26) we obtain $(z_i, x_{-i}) \sim (w_i, y_{-i})$, violating (60). Finally, the case $Not[x \succsim y]$ is dealt with like as in the proof of theorem 5.1. This completes the proof. \square

Remark 5.6 ($n = 2$ case)

It is easy to see that $RC1_i$ amounts to defining a biorder (see Doignon et al., 1984; Ducamp & Falmagne, 1969) between the sets X_i^2 and X_{-i}^2 . Therefore $RC1_i$ on its own implies, when X is finite or countably infinite, the existence of two real-valued functions p_i and P_{-i} respectively on X_i^2 and X_{-i}^2 such that, for all $x, y \in X$, $x \succsim y$ iff $p_i(x_i, y_i) + P_{-i}(x_{-i}, y_{-i}) \geq 0$ (see Ducamp & Falmagne, 1969, Proposition 3).

Therefore models using traces on differences closely relate to ordinal measurement when $n = 2$.

In a similar vein, Bouyssou (1986, Theorem 1) noted an interesting implication of TC_i on its own. When X is countable, TC_i implies the existence of two real-valued skew symmetric functions p_i and P_{-i} respectively on X_i^2 and X_{-i}^2 such that, for all $x, y \in X$, $x \succsim y \Leftrightarrow p_i(x_i, y_i) + P_{-i}(x_{-i}, y_{-i}) \geq 0$. This result can easily be extended to sets of arbitrary cardinality (see remark 5.8). When $n = 2$ this offers an alternative to Fishburn (1991b, theorem B). \bullet

Remark 5.7 (Extensions to subsets)

The obvious extension of $RC1$ to subsets of attributes is the main necessary condition used by Vind (1991) together with topological assumptions on X to axiomatize a model such that:

$$x \succsim y \Leftrightarrow \sum_{i=1}^n p_i(x_i, y_i) \geq 0 \tag{61}$$

with $p_i(x_i, x_i) = 0$.

Similarly, it is easy to see that the extension of TC to subsets of attributes is necessary for a model of type (61) with all p_i 's being skew symmetric. A complete axiomatic treatment of this model may be found in Fishburn (1990a, 1990b, 1991b).

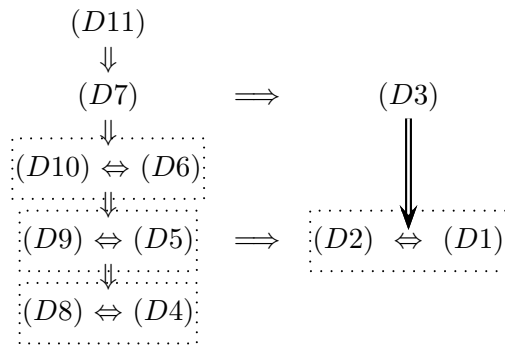
It should be noted that our use here of a function G instead of a sum allows to considerably simplify our analysis. Clearly as soon as additivity is required, axioms involve either a denumerable scheme of conditions guaranteeing the existence of solutions to a set of linear equations or a limited number of conditions together with unnecessary structural assumptions on the set of alternatives (e.g. solvability). We refer to Bouyssou and Pirlot (2002a) for a thorough comparison between models involving traces on differences and their additive specializations as well as an evaluation of the contribution of theorem 5.1 to the general theory of conjoint measurement. •

Remark 5.8 (The general case)

The results in this section can easily be extended to cover the general case. This requires the addition of, necessary, conditions guaranteeing that the equivalences and weak orders encountered in this section have a numerical representation. •

The main results in this section are summarized in table 6 and figure 4.

Figure 4: Implication between models resulting from theorem 5.1



All models imply (D0)

Table 6: Main results using traces on differences

Models	Definition	Conditions
(D0)	$x \succsim y \Leftrightarrow G([p_i(x_i, y_i)]) \geq 0$	\emptyset
	
(D1)	(D0) with $p_i(x_i, x_i) = 0$	
\Updownarrow		ind.
(D2)	(D0) with p_i skew symmetric	
	
(D3)	(D0) with p_i skew symmetric and G odd	cpl., ind.
	
(D8) \Leftrightarrow (D4)	(D0) with $G(\nearrow \nearrow)$	RC1
	
(D9) \Leftrightarrow (D5)	(D1) with $G(\nearrow \nearrow)$	RC1, ind.
	
(D10) \Leftrightarrow (D6)	(D2) with $G(\nearrow \nearrow)$	RC12
	
(D7)	(D3) with $G(\nearrow)$	cpl., RC12
	
(D11)	(D3) with $G(\nearrow \nearrow)$	cpl., TC
	

\nearrow means nondecreasing, $\nearrow \nearrow$ means increasing
cpl. means completeness, ind. means independence

6 Models using marginal traces on levels and on differences

This section studies models based on traces on differences in which the magnitude of differences may be modelled through traces on levels. We have already encountered the trivial model (L0D0) in which:

$$x \succsim y \Leftrightarrow H([\phi_i(u_i(x_i), u_i(y_i))]) \geq 0, \quad (L0D0)$$

This model may be seen as a particular case of model (D0) in which the term $p(x_i, y_i)$ is substituted with the term $\phi_i(u_i(x_i), u_i(y_i))$. Alternatively it may be seen as a generalization of the additive difference model (5) replacing addition and subtraction by more general functions.

To each of the 12 models (D0) to (D11) studied in section 5 corresponds a model in which the term $p_i(x_i, y_i)$ is substituted with $\phi_i(u_i(x_i), u_i(y_i))$. This defines models (L0D0) to (L0D11).

In order to bring the function ϕ “closer” to a subtraction, we envisage two variants of each of these models. In the first one, we impose that ϕ is

nondecreasing in its first argument and nonincreasing in its second argument. This defines models (L1D0) to (L1D11). In the other variant we impose that ϕ is increasing in its first argument and decreasing in its second argument. This defines models (L2D0) to (L2D11). The definition of all these models is summarized in table 7. We have thus defined a total of $3 \times 12 = 36$ models involving differences and levels. We study them in this section.

Table 7: Models involving traces on levels and on differences

(L0D0)	$x \succsim y \Leftrightarrow H([\phi_i(u_i(x_i), u_i(y_i))]) \geq 0$
(L0D1)	(L0D0) with $\phi_i(u_i(x_i), u_i(x_i)) = 0$
(L0D2)	(L0D1) with ϕ_i skew symmetric
(L0D3)	(L0D2) with H odd
.....	
(L0D4)	(L0D0) with H nondecreasing
(L0D5)	(L0D0) with H increasing
.....	
(L0D6)	(L0D1) with H nondecreasing
(L0D7)	(L0D1) with H increasing
.....	
(L0D8)	(L0D2) with H nondecreasing
(L0D9)	(L0D2) with H increasing
.....	
(L0D10)	(L0D3) with H nondecreasing
(L0D11)	(L0D3) with H increasing

Models (L1Dx) correspond to models (L0Dx) with $\phi_i(\nearrow, \searrow)$
 Models (L2Dx) correspond to models (L0Dx) with $\phi_i(\nearrow\nearrow, \searrow\nwarrow)$

6.1 Models (L0D1) to (L0D11)

These 12 models (L0D0) to (L0D11) correspond to models (D0) to (D11) involving differences in which the term $p_i(x_i, y_i)$ is substituted with the term $\phi_i(u_i(x_i), u_i(y_i))$ with no monotonicity property required for the functions ϕ_i . These models are easily analyzed using the following elementary observation.

Lemma 6.1 (Numerical representation of a weak order on pairs)

Let \succsim^A be a weak order on a countable set A^2 . Consider any real valued function f on A^2 representing \succsim^A , i.e. such that, for all $a, b, c, d \in A$,

$$(a, b) \succsim^A (c, d) \Leftrightarrow f(a, b) \geq f(c, d). \quad (62)$$

There is a real-valued function u on A and a real-valued function g on $u(A) \times u(A)$ such that, for all $a, b \in A$,

$$f(a, b) = g(u(a), u(b)). \quad (63)$$

PROOF

Define the binary relation E on A letting, for all $a, b \in A$,

$$a E b \Leftrightarrow (a, c) \sim^A (b, c), \text{ and } (c, b) \sim^A (c, a), \text{ for all } c \in A, \quad (64)$$

where \sim^A denotes the symmetric part of \succsim^A . Since \sim^A is an equivalence, it is easy to show that E is an equivalence. Therefore, since A is countable, there is a real-valued function u on A such that, for all $a, b \in A$,

$$a E b \Leftrightarrow u(a) = u(b). \quad (65)$$

Take any function f such that (62) holds and define g on $u(A) \times u(A)$ letting, for all $a, b \in A$, $g(u(a), u(b)) = f(a, b)$. We have to show that g is well-defined, i.e. that $u(a) = u(c)$ and $u(b) = u(d)$ imply $f(a, b) = f(c, d)$. By construction, we have $a E c$ and $b E d$. This implies $(a, \ell) \sim^A (c, \ell)$ and $(\ell', b) \sim^A (\ell', d)$, for all $\ell, \ell' \in A$. Taking $\ell = b$ and $\ell' = c$ implies $(a, b) \sim^A (c, b)$ and $(c, b) \sim^A (c, d)$. Using the transitivity of \sim^A we obtain $(a, b) \sim^A (c, d)$ so that $f(a, b) = f(c, d)$, as required. \square

Theorem 6.1 (Models (L0D0) to (L0D11))

Models (L0D0) to (L0D11) are respectively equivalent to models (D0) to (D11).

PROOF

Consider any representation of models (D0) to (D11) and apply lemma 6.1 to the weak order on X_i^2 induced by p_i . \square

Models (L0D0) to (L0D11) are therefore nothing more than a different presentation of models (D0) to (D11). Clearly, the equivalences between models (D0) to (D11) noted in corollary 5.1 carry over to models (L0D0) to (L0D11). The relations between these models is therefore given by figure 4 replacing Dk by $L0Dk$, for $k = 0, 1, \dots, 11$.

6.2 Models (L1D0)-(L1D3), (L2D0)-(L2D3)

We first take up the case of models (L1D0) to (L1D3) and (L2D0) to (L2D3) in which H has no particular monotonicity properties. Substituting $p_i(x_i, y_i)$

with $\phi_i(u_i(x_i), u_i(y_i))$, ϕ_i being nondecreasing in its first argument and non-increasing in its second argument is likely to have little impact since the monotonicity properties of ϕ_i can be “absorbed” by H . A similar reasoning applies if ϕ_i is supposed to be increasing in its first argument and decreasing in its second argument. As already mentioned in remark 3.3, this is indeed the case.

Theorem 6.2 (Models (L1D0)–(L1D3) and (L2D0)–(L2D3))

1. $(D0) \Leftrightarrow (L1D0) \Leftrightarrow (L2D0)$,
2. $(D1) \Leftrightarrow (L1D1) \Leftrightarrow (L2D1) \Leftrightarrow (D2) \Leftrightarrow (L1D2) \Leftrightarrow (L2D2)$,
3. $(D3) \Leftrightarrow (L1D3) \Leftrightarrow (L2D3)$.

PROOF

Part 1. By construction, $(L2D0) \Rightarrow (L1D0) \Rightarrow (D0)$. We show that $(D0) \Rightarrow (L2D0)$.

Consider any one-to-one correspondence u_i between X_i and a subset of the set of integers ≥ 2 . Define ϕ_i letting, for all $x_i, y_i \in X_i$,

$$\phi_i(u_i(x_i), u_i(y_i)) = u_i(x_i) + \frac{1}{u_i(y_i)}. \quad (66)$$

By construction, ϕ_i is increasing in its first argument and decreasing in its second argument. Observe that

$$\phi_i(u_i(x_i), u_i(y_i)) = \phi_i(u_i(z_i), u_i(w_i)) \Rightarrow x_i = z_i \text{ and } y_i = w_i \quad (67)$$

Define H on $\prod_{i=1}^n \phi_i(u_i(X_i), u_i(X_i))$ letting:

$$H([\phi_i(u_i(x_i), u_i(y_i))]) = \begin{cases} +1 & \text{if } x \succ y, \\ -1 & \text{otherwise.} \end{cases} \quad (68)$$

We have to show that H is well-defined. This is obvious in view of (67).

Part 2. We know from corollary 5.1 that $(D2) \Leftrightarrow (D1)$. By construction $(L2D2) \Rightarrow (L1D2) \Rightarrow (D2)$ and $(L2D1) \Rightarrow (L1D1) \Rightarrow (D1)$. We show that $(D1) \Rightarrow (L2D2)$.

Consider on each X_i any function ϕ_i satisfying (66) and define φ_i letting, for all $x_i, y_i \in X_i$,

$$\varphi_i(u_i(x_i), u_i(y_i)) = \phi_i(u_i(x_i), u_i(y_i)) - \phi_i(u_i(y_i), u_i(x_i)). \quad (69)$$

By construction, φ_i is skew symmetric, increasing in its first argument and decreasing in its second argument. We clearly have:

$$\varphi_i(u_i(x_i), u_i(y_i)) = \varphi_i(u_i(z_i), u_i(w_i)) \Rightarrow \begin{cases} x_i = z_i \text{ and } y_i = w_i \\ \text{or} \\ x_i = y_i \text{ and } z_i = w_i. \end{cases} \quad (70)$$

Define H on $\prod_{i=1}^n \varphi_i(u_i(X_i), u_i(X_i))$ letting:

$$H([\varphi_i(u_i(x_i), u_i(y_i))]) = \begin{cases} +1 & \text{if } x \succsim y, \\ -1 & \text{otherwise.} \end{cases} \quad (71)$$

Since model $(D1)$ holds, we know that \succsim is independent so that $(x_i, x_i) \sim_i^* (y_i, y_i)$. In view of (70), it is then clear that H is well-defined.

Part 3. By construction, $(L2D3) \Rightarrow (L1D3) \Rightarrow (D3) \Rightarrow (L2D3)$. Define, on each $i \in N$, a function φ_i using (69). Define H on $\prod_{i=1}^n \varphi_i(u_i(X_i), u_i(X_i))$ letting:

$$H([\varphi_i(u_i(x_i), u_i(y_i))]) = \begin{cases} +1 & \text{if } x \succ y, \\ 0 & \text{if } x \sim y, \\ -1 & \text{otherwise.} \end{cases} \quad (72)$$

Since \succsim is complete, H is odd. Its well-definedness is proved as in part 2. \square

6.3 Models (L1D4) to (L1D11)

These eight models are the counterparts of models $(D4)$ to $(D11)$ in which the term $p_i(x_i, y_i)$ is substituted with the term $\phi_i(u_i(x_i), u_i(y_i))$, ϕ_i being nondecreasing in its first argument and nonincreasing in its second argument. In all models $(D4)$ to $(D11)$, the function G is nondecreasing in all its arguments. Therefore, it is clear that imposing that ϕ_i is nondecreasing in its first argument and nonincreasing in its second argument is no more innocuous. The conditions that were used to analyze models using traces on levels reappear here. The least constrained of the models considered here, $(L1D4)$, already implies all of these conditions.

Lemma 6.2 (Necessity of AC1, AC2 and AC3)

If \succsim has a representation in model $(L1D4)$ then it satisfies AC1, AC2 and AC3.

PROOF

In model $(L1D4)$, we have:

$$x \succsim y \Leftrightarrow H([\phi_i(u_i(x_i), u_i(y_i))]) \geq 0,$$

with H being nondecreasing in all its arguments and all ϕ_i being nondecreasing in their first argument and nonincreasing in their second argument. The proof that AC1, AC2 and AC3 hold easily follows from the numerical representation. We show that AC1 holds. The premise of AC1 _{i} yields in terms

of model (L1D4):

$$\begin{aligned} H(\phi_i(u_i(x_i), u_i(y_i)), (\phi_j(u_j(x_j), u_j(y_j)))_{j \neq i}) &\geq 0 \text{ and} \\ H(\phi_i(u_i(z_i), u_i(w_i)), (\phi_j(u_j(z_j), u_j(w_j)))_{j \neq i}) &\geq 0. \end{aligned}$$

We have either $u_i(z_i) \geq u_i(x_i)$ or $u_i(x_i) \geq u_i(z_i)$. In the first case, the monotonicity of H and ϕ_i , implies:

$$H(\phi_i(u_i(z_i), u_i(y_i)), (\phi_j(u_j(x_j), u_j(y_j)))_{j \neq i}) \geq 0.$$

In the second case, we have:

$$H(\phi_i(u_i(x_i), u_i(w_i)), (\phi_j(u_j(z_j), u_j(w_j)))_{j \neq i}) \geq 0.$$

Hence, $AC1_i$ holds. The proof for $AC2_i$ and $AC3_i$ is similar. \square

As was the case in section 6.1, a simple lemma on the numerical representation of a weak order on ordered pairs will allow us to deal with all our models in a simple way. This will require some new vocabulary however.

Definition 6.1 (Strong linearity)

Let \succsim^A be a binary relation on a set A^2 . We say that:

1. \succsim^A is right-linear iff $[Not[(b, c) \succsim^A (a, c)]] \Rightarrow (a, d) \succsim^A (b, d)$,
2. \succsim^A is left-linear iff $[Not[(c, a) \succsim^A (c, b)]] \Rightarrow (d, b) \succsim^A (d, a)$,
3. \succsim^A is strongly linear iff $[Not[(b, c) \succsim^A (a, c)]]$ or $Not[(c, a) \succsim^A (c, b)] \Rightarrow [(a, d) \succsim^A (b, d) \text{ and } (d, b) \succsim^A (d, a)]$,

for all $a, b, c, d \in A$

Lemma 6.3 (Numerical representation of a weak order on pairs)

Let \succsim^A be a weak order on a countable set A^2 . Let f be any real-valued function on A^2 such that, for all $a, b, c, d \in A$,

$$(a, b) \succsim^A (c, d) \Leftrightarrow f(a, b) \geq f(c, d). \quad (73)$$

There is a real-valued function u on A and a real-valued function g on $u(A) \times u(A)$ nondecreasing in its first argument, nonincreasing in its second argument, such that for all $a, b, c, d \in A$,

$$f(a, b) = g(u(a), u(b)) \quad (74)$$

iff \succsim^A is strongly linear.

PROOF

Necessity of strong linearity is easily shown. Suppose that $\text{Not}[(b, c) \succsim^A (a, c)]$ or $\text{Not}[(c, a) \succsim^A (c, b)]$. This implies $g(u(b), u(c)) < g(u(a), u(c))$ or $g(u(c), u(a)) < g(u(c), u(b))$. In either case, the monotonicity properties of g imply $u(a) > u(b)$. Starting now from $g(u(b), u(d))$ and using the nondecreasingness of g in its first argument, we obtain $g(u(a), u(d)) \geq g(u(b), u(d))$ and, hence, $(a, d) \succsim^A (b, d)$. The proof that $(d, b) \succsim^A (d, a)$ holds is similar.

Sufficiency. Define the binary relation T on A letting:

$$a T b \Leftrightarrow [(a, c) \succ (b, c) \text{ and } (c, b) \succ (c, a), \text{ for all } c \in A].$$

It is clear that T is reflexive and transitive. An easy proof shows that it is complete if and only if \succ is strongly linear.

Since A is countable and T is a weak order, there is a real-valued function u on A such that, for all $a, b \in A$,

$$a T b \Leftrightarrow u(a) \geq u(b). \quad (75)$$

Let f be any real-valued function on A^2 such that (73) holds. Define the real-valued function g on $u(A)^2$ letting, for all $a, b, c, d \in A$,

$$g(u(a), u(b)) = f(a, b).$$

Using the definition of T , it is routine to show that g is well-defined, nondecreasing in its first argument and nonincreasing in its second argument. \square

The following lemma interprets conditions $AC1$, $AC2$ and $AC3$ in terms of linearity properties of \succsim_i^* .

Lemma 6.4 ($AC1_i$, $AC2_i$, $AC3_i$ and strong linearity)

Let \succ be a binary relation on $X = \prod_{i=1}^n X_i$.

1. $AC1_i$ holds iff \succsim_i^* is right linear,
2. $AC2_i$ holds iff \succsim_i^* is left linear,
3. $AC3_i$ holds iff $[\text{Not}[(x_i, z_i) \succsim_i^* (y_i, z_i)], \text{ for some } z_i \in X_i] \Rightarrow (w_i, x_i) \succsim_i^* (w_i, y_i), \text{ for all } w_i \in X_i,$
4. $AC1_i$, $AC2_i$ and $AC3_i$ hold iff \succsim_i^* is strongly linear iff \succsim_i^{**} is strongly linear.

PROOF

Part 1. We show equivalently that $Not[AC1_i]$ iff for some $x_i, y_i, z_i, w_i \in X_i$, $Not[(z_i, y_i) \succsim_i^* (x_i, y_i)]$ and $Not[(x_i, w_i) \succsim_i^* (z_i, w_i)]$. This last statement means, by definition of \succsim_i^* , that for some $a_{-i}, b_{-i}, c_{-i}, d_{-i} \in X_{-i}$, we have

$$[(x_i, a_{-i}) \succ (y_i, b_{-i})] \text{ and } Not[(z_i, a_{-i}) \succ (y_i, b_{-i})], \text{ and} \\ [(z_i, c_{-i}) \succ (w_i, d_{-i})] \text{ and } Not[(x_i, c_{-i}) \succ (w_i, d_{-i})],$$

which is exactly $Not[AC1_i]$. The proof of part 2 is similar.

Part 3. We show that $Not[AC3_i] \Leftrightarrow [Not[(x_i, z_i) \succsim_i^* (y_i, z_i)] \text{ and } Not[(w_i, x_i) \succsim_i^* (w_i, y_i)]]$, for some $z_i, w_i \in X_i$. The last expression means that, for some $a_{-i}, b_{-i}, c_{-i}, d_{-i} \in X_{-i}$, we have

$$[(y_i, a_{-i}) \succ (z_i, b_{-i})] \text{ and } Not[(x_i, a_{-i}) \succ (z_i, b_{-i})], \\ [(w_i, c_{-i}) \succ (y_i, d_{-i})] \text{ and } Not[(w_i, c_{-i}) \succ (x_i, d_{-i})],$$

which is exactly $Not[AC3_i]$.

Part 4. Combining the equivalences in parts 1, 2 and 3 leads to the equivalence between $AC1_i, AC2_i$ and $AC3_i$ with the strong linearity of \succsim_i^* . By construction this also shows that \succsim_i^{**} is strongly linear. \square

We are now in position to characterize models (L1D4) to (L1D11)

Theorem 6.3

Let \succsim be a binary relation on a countable set $X = \prod_{i=1}^n X_i$.

1. Model (L1D8) holds iff \succsim satisfies RC1, AC1, AC2 and AC3.
2. Model (L1D9) holds iff \succsim is independent and satisfies RC1, AC1, AC2 and AC3.
3. Model (L1D10) holds iff \succsim satisfies RC1, RC2, AC1, AC2 and AC3.
4. Model (L1D7) holds iff \succsim is complete and satisfies RC1, RC2, AC1, AC2 and AC3.
5. Model (L1D11) holds iff \succsim is complete and satisfies TC, AC1, AC2 and AC3.

PROOF

The necessity of all parts follows from theorem 5.1 and lemma 6.2.

Observe that in the proof of parts 3 to 7 of theorem 5.1, we have exhibited a regular representation, i.e. a representation in which p_i represents \succsim_i^* in models (D8) and (D9) and \succsim_i^{**} in models (D10), (D7) and (D11). The sufficiency of each part therefore follows applying lemma 6.3 to the weak order induced on X_i^2 by the function p_i used in the proof of theorem 5.1. \square

Corollary 6.1 (Links between models)

1. Model (L1D4) and (L1D8) are equivalent.
2. Model (L1D5) and (L1D9) are equivalent.
3. Models (L1D6) and (L1D10) are equivalent.
4. Model (L1D7) \Rightarrow Model (L1D10).

Remark 6.1 (Independence of axioms)

We refer to Bouyssou and Pirlot (2002d) for a complete analysis of the independence of the axioms used in theorem 6.3. It turns out that conditions implying the completeness of marginal traces on levels ($AC1$, $AC2$ and $AC3$) are completely independent from the conditions implying that marginal traces on differences are complete ($RC1$ and $RC2$). In order not to multiply examples, we simply show in appendix C that, in the class of complete relations, conditions TC , $AC1$, $AC2$ and $AC3$ are independent. \bullet

6.4 Models (L2D4) to (L2D11)

Except for the most constrained model (L2D11), it turns out that no additional constraint is brought into the picture supposing that ϕ_i increases in its first arguments and decreases in its second argument. We summarize our results below.

Theorem 6.4

Let \succsim be a binary relation on a countable set $X = \prod_{i=1}^n X_i$.

1. Models (L2D8) and (L1D8) are equivalent.
2. Models (L2D9) and (L1D9) are equivalent.
3. Models (L2D10) and (L1D10) are equivalent.
4. Models (L2D7) and (L1D7) are equivalent.
5. Model (L2D11) holds iff \succsim is complete and satisfies TC , $TAC1$ and $TAC2$.

In the proof of theorem 6.3 the strategy was to consider the underlying models involving differences and to factorize the function p_i using lemma 6.3 using the fact that it is always possible to build a regular representation. Hence in all the representations used in theorem 6.3 the functions ϕ_i represent \succsim_i^* or \succsim_i^{**} depending on the model. Our strategy here is rather different and amounts to exploiting the fact, already stressed, that there is no need for the

functions p_i in models involving differences to represent \succsim_i^* or \succsim_i^{**} . We use this degree of freedom to modify nondecreasing/nonincreasing functions ϕ_i into increasing and decreasing functions φ_i .

The following lemma states the conditions under which a function f that is nondecreasing in its first argument and nonincreasing in its second argument, can be appropriately transformed into a strictly monotonic function g .

Consider a function $f : U \times U \rightarrow \mathbb{R}$, with U a countable subset of \mathbb{R} and suppose that f is nondecreasing in its first argument and nonincreasing in its second argument. There are two types of situations that may cause the lack of strict monotonicity of f . We denote by S , the set of values r of f for which either there are $a, b, c \in U$ such that:

$$f(a, c) = f(b, c) = r \text{ with } a > b, \quad (76)$$

or there are $a, c, d \in U$ such that:

$$f(a, c) = f(a, d) = r \text{ with } c > d. \quad (77)$$

Clearly, f is strictly monotonic iff S is empty.

Lemma 6.5 (Modifying a nondecreasing function)

Let U be a countable subset of the $]0, 1[$ interval and $f : U \times U \rightarrow \mathbb{R}$ a function that is nondecreasing in its first argument and nonincreasing in its second argument.

There exists a function $g : U \times U \rightarrow \mathbb{R}$ that is increasing in its first argument and decreasing in its second argument and such that, for all $u, v, u', v' \in U$,

$$[f(u, v) > f(u', v')] \Rightarrow [g(u, v) > g(u', v')], \text{ and} \quad (78)$$

$$[f(u, v) = f(u', v')] \Rightarrow [g(u, v) = g(u', v')] \text{ iff } f(u, v) \notin S. \quad (79)$$

If, in addition, f vanishes on the diagonal ($f(u, u) = 0$, for all $u \in U$) (resp. is skew symmetric) there exists a function g satisfying (78) and (79) that vanishes on the diagonal (resp. is skew symmetric).

The proof of this lemma, being rather uninformative, is relegated in appendix D.

PROOF (THEOREM 6.4)

Parts 1 to 4. We show that model (L2D8) holds if and only if RC1 and AC123 hold.

The necessity of *RC1* and *AC123* for model (*L2D8*) is clear, using theorem 6.3. We show sufficiency. Consider a representation of \succsim in model (*L1D8*). From the proof of theorem 6.3, we have:

$$(z_i, w_i) \succsim_i^* (x_i, y_i) \Leftrightarrow \phi_i(u_i(z_i), u_i(w_i)) \geq \phi_i(u_i(x_i), u_i(y_i)). \quad (80)$$

Without loss of generality, we may suppose that $U = u_i(X_j)$ is included in $]0; 1[$. We may then apply lemma 6.5 to obtain a function φ_i that is increasing in its first argument and decreasing in its second argument. According to (78), we have:

$$(z_i, w_i) \succsim_i^* (x_i, y_i) \Rightarrow \varphi_i(u_i(z_i), u_i(w_i)) > \varphi_i(u_i(x_i), u_i(y_i)). \quad (81)$$

Hence, this function can be used as a basis of the definition of H in view of lemma 5.5. This shows sufficiency.

Combining theorem 6.3 and lemmas 5.5 and 6.5, the proof for models (*L2D9*), (*L2D10*) and (*L2D7*) is similar.

Part 5. Necessity. In view of theorem 6.3 we only have to show that *TAC1* and *TAC2* are necessary. Suppose that \succsim has a representation in model (*L2D11*). The premise of *TAC1_i*, interpreted in terms of the model, yields three inequalities:

$$H(\phi_i(u_i(x_i), u_i(y_i)), (\phi_j(u_j(a_j), u_j(y_j))_{j \neq i}) \geq 0 \quad (82)$$

$$H(\phi_i(u_i(y_i), u_i(z_i)), (\phi_j(u_j(y_j), u_j(a_j))_{j \neq i}) \geq 0 \quad (83)$$

$$H(\phi_i(u_i(z_i), u_i(w_i)), (\phi_j(u_j(b_j), u_j(w_j))_{j \neq i}) \geq 0 \quad (84)$$

Due to skew symmetry of ϕ_i and oddness of H , (83) may be rewritten as:

$$H(\phi_i(u_i(z_i), u_i(y_i)), (\phi_j(u_j(a_j), u_j(y_j))_{j \neq i}) \leq 0 \quad (85)$$

Using the increasingness of H (resp. ϕ_i) in its i th (resp. first) argument, (82) and (85) imply $u_i(x_i) \geq u_i(z_i)$. Substituting $u_i(z_i)$ with $u_i(x_i)$ in equation (84) yields:

$$H(\phi_i(u_i(x_i), u_i(w_i)), (\phi_j(u_j(b_j), u_j(w_j))_{j \neq i}) \geq 0$$

which establishes *TAC1_i*. The proof for *TAC2_i* is similar.

Sufficiency. From the proof of theorem 6.3, we know that \succsim has a representation in model (*L1D11*) such that

$$(x_i, y_i) \succsim_i^{**} (z_i, w_i) \Leftrightarrow \phi_i(u_i(x_i), u_i(y_i)) \geq \phi_i(u_i(z_i), u_i(w_i)) \text{ and} \quad (86)$$

$$x_i \succsim_i^\pm y_i \Leftrightarrow u_i(x_i) \geq u_i(y_i) \quad (87)$$

In view of lemmas 5.5 and 6.5, the proof will be complete if we show that after transforming the functions ϕ_i into functions φ_i according to lemma 6.5, it is still true that (60) holds, i.e. that:

$$\left. \begin{array}{l} (z_i, w_i) \sim_i^{**} (x_i, y_i) \text{ and} \\ \exists a_{-i}, b_{-i} \in X_{-i} \text{ s.t. } (x_i, a_{-i}) \sim (y_i, b_{-i}) \end{array} \right\} \Rightarrow \varphi_i(u_i(z_i), u_i(w_i)) = \varphi_i(u_i(x_i), u_i(y_i)).$$

In view of (79), this will be true if, as soon as $(z_i, w_i) \sim_i^{**} (x_i, y_i)$ and $\exists a_{-i}, b_{-i} \in X_{-i}$ s.t. $(x_i, a_{-i}) \sim (y_i, b_{-i})$, it is not true that $\phi_i(u_i(x_i), u_i(y_i)) \in S$. This results from (34) and (35). \square

Remark 6.2 (Regular representations)

For models (L1D8), (L1D7), (L1D9), (L1D10), (L1D7) and (L1D11), theorem 6.3 shows that it is always possible to build a representation of these models in which:

- u_i is a numerical representation of the weak order \succsim_i^\pm and
- ϕ_i is a numerical representation of the weak order \succsim_i^* , in model (L1D8) and of the weak order \succsim_i^{**} in the more constrained models.

We call *regular* a representation in which this is the case (see Roberts (1979, ch. 2) about regularization of a scale of measurement).

In theorem 6.4, the representations that we build start from the regular representations provided by theorem 6.3 and we modify them, in accordance with the restrictions of lemma 5.5, so as to make ϕ_i increasing in its first argument and decreasing in its second argument. This modification is done, on each $i \in N$, on a set S_i containing all the values s such that:

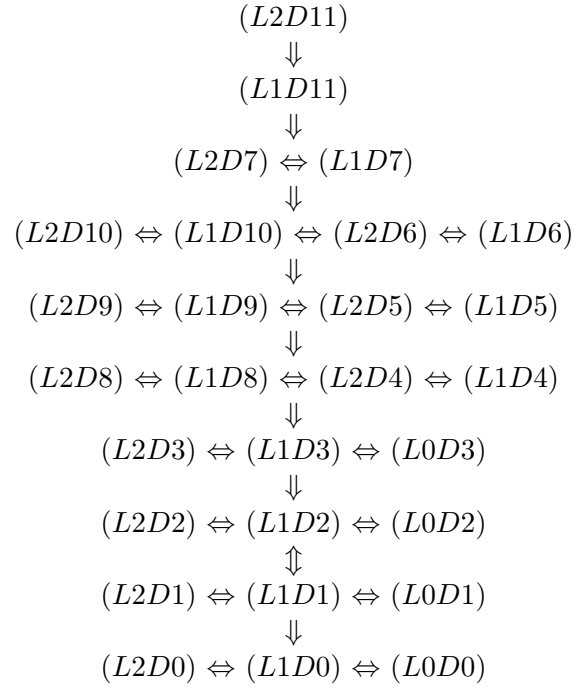
$$\begin{aligned} s &= \phi_i(u_i(x_i), u_i(z_i)) = \phi_i(u_i(y_i), u_i(z_i)), \text{ or} \\ s &= \phi_i(u_i(z_i), u_i(x_i)) = \phi_i(u_i(z_i), u_i(y_i)), \end{aligned}$$

for some $x_i, y_i, z_i \in X_i$ such that $x_i \succ_i^\pm y_i$. It is not difficult to see that the emptiness of the sets S_i is a necessary and sufficient condition to obtain a regular representation in the models envisaged in theorem 6.4 (see Bouyssou & Pirlot, 2002d, for details). \bullet

Remark 6.3 (Extension to the general case)

Because we have deliberately used above representations that are not regular (i.e. ϕ_i do not necessarily represent \succsim_i^* or \succsim_i^{**}), the extension of the results in this section to the general uncountable case is slightly more involved than for other models (see Bouyssou & Pirlot, 2002d). This does not raise major difficulties however. \bullet

Figure 5: Implication between models resulting from theorems 6.3 and 6.4



The main results in this section are summarized in tables 8 and 9 and figure 5.

Table 8: Models (*L1D0*) to (*L1D11*)

Models	Definition	Conditions
(<i>L1D0</i>)	$x \succsim y \Leftrightarrow H([\phi_i(u_i(x_i), u_i(y_i))]) \geq 0$ $\phi_i(\nearrow, \searrow)$	\emptyset
(<i>L1D1</i>)	(<i>L1D0</i>) with $\phi_i(u(x_i), u_i(x_i)) = 0$	ind.
(<i>L1D2</i>)	(<i>L1D1</i>) with ϕ_i skew symmetric	
(<i>L1D3</i>)	(<i>L1D2</i>) with H odd	cpl., ind.
(<i>L1D4</i>)	(<i>L1D0</i>) with $H(\nearrow)$	<i>RC1, AC123</i>
(<i>L1D8</i>)	(<i>L1D0</i>) with $H(\nearrow \nearrow)$	
(<i>L1D5</i>)	(<i>L1D1</i>) with $H(\nearrow)$	<i>RC1, ind., AC123</i>
(<i>L1D9</i>)	(<i>L1D1</i>) with $H(\nearrow \nearrow)$	
(<i>L1D6</i>)	(<i>L1D2</i>) with $H(\nearrow)$	<i>RC12, AC123</i>
(<i>L1D10</i>)	(<i>L1D2</i>) with $H(\nearrow \nearrow)$	
(<i>L1D7</i>)	(<i>L1D3</i>) with $H(\nearrow)$	cpl., <i>RC12, AC123</i>
(<i>L1D11</i>)	(<i>L1D3</i>) with $H(\nearrow \nearrow)$	cpl., <i>TC, AC123</i>

$\nearrow \nearrow$ means increasing, \nearrow means nondecreasing, \searrow means nonincreasing
cpl. means completeness, ind. means independence

Conditions for the first four rows are identical to those of table 6

Table 9: Models (*L2D0*) to (*L2D11*)

Models	Definition	Conditions
(<i>L2D0</i>)	$x \succsim y \Leftrightarrow H([\phi_i(u_i(x_i), u_i(y_i))]) \geq 0$ $\phi_i(\nearrow, \searrow)$	\emptyset
(<i>L2D1</i>)	(<i>L1D0</i>) with $\phi_i(u(x_i), u_i(x_i)) = 0$	ind.
(<i>L2D2</i>)	(<i>L2D1</i>) with ϕ_i skew symmetric	
(<i>L2D3</i>)	(<i>L2D2</i>) with H odd	cpl., ind.
(<i>L2D4</i>)	(<i>L2D0</i>) with $H(\nearrow)$	<i>RC1, AC123</i>
(<i>L2D8</i>)	(<i>L2D0</i>) with $H(\nearrow)$	
(<i>L2D5</i>)	(<i>L2D1</i>) with $H(\nearrow)$	<i>RC1, ind., AC123</i>
(<i>L2D9</i>)	(<i>L2D1</i>) with $H(\nearrow)$	
(<i>L2D6</i>)	(<i>L2D2</i>) with $H(\nearrow)$	<i>RC12, AC123</i>
(<i>L2D10</i>)	(<i>L2D2</i>) with $H(\nearrow)$	
(<i>L2D7</i>)	(<i>L2D3</i>) with $H(\nearrow)$	cpl., <i>RC12, AC123</i>
(<i>L2D11</i>)	(<i>L2D3</i>) with $H(\nearrow)$	cpl., <i>TC, TAC12</i>

\nearrow means increasing, \nearrow means nondecreasing, \searrow means decreasing
cpl. means completeness, ind. means independence

Conditions are identical to those of table 8 except for the last row

7 Discussion

This paper has proposed a general approach of conjoint measurement models tolerating intransitivity and/or incompleteness using simple tools based on several kinds of marginal traces on coordinates induced by the binary relation on the product set. We have provided, when X is supposed to be countable, a fairly complete analysis of a large variety of models. Our project was to investigate how far it was possible to go in terms of numerical representations using a limited number of cancelation conditions without imposing any transitivity requirement on the preference relation and any structural assumptions on the set of objects. Rather surprisingly, as we saw, such a poor framework allows us to go rather far. Furthermore the cancelation conditions that we used ($RC1$, $RC2$, Independence, TC , $AC1$, $AC2$, $AC3$, $TAC1$, $TAC2$) are reasonably simple and have close relations with the conditions used in the analysis of traditional conjoint measurement models.

Our framework and results have many possible applications. Among them let us mention:

- the characterization of all relations compatible with a dominance relation, using our models based on marginal traces on levels (see Bouyssou & Pirlot, 2002c),
- the characterization of preference relations that can be obtained using an “ordinal aggregation model” using our models based on marginal traces on levels (see Bouyssou & Pirlot, 2002e, 2002f, 2002b),
- the characterization of “ordinal” models for decision under uncertainty (see Bouyssou, Perny, & Pirlot, 2000; Bouyssou & Pirlot, 2003a, 2003b),
- the characterization of various functional forms for F , G or H (see Bouyssou, Greco, Matarazzo, Pirlot, & Słowiński, 2002),

It is clearly impossible to develop these points here. The patient reader who has followed us till now should be in position to guess the general spirit of these results.

Summarizing the main messages in a few words, we would say that:

- when confronted to a nontransitive/noncomplete relation, it is always profitable to investigate its traces and/or marginal traces,
- the use of conjoint measurement techniques is not restricted to the study of complete and transitive binary relations,

- if assessment procedures are not looked for, replacing additivity by mere decomposability requirements often allows to grasp in a very simple way the essence of a model,
- replacing additivity by mere decomposability requirements amounts to using models which are intimately related with “rule-based” preference modeling Greco, Matarazzo, and Słowiński (1999, 2001), Greco et al. (2002). This allows to consider the possibility of deriving assessment procedures using the machinery of “rule induction” in Artificial Intelligence.

Our framework and results are also intended to contribute to the general theory of conjoint measurement. They allow to draw the following general picture of conjoint measurement models (see figure 6), where models are classified, when studying the proposition $x \succsim y$, according to whether:

- they use traces on differences, i.e. their functional form can be written so as to be nondecreasing in functions $p_i(x_i, y_i)$,
- they use traces on levels, i.e. their functional form can be written so as to be nondecreasing in functions $u_i(x_i)$ and nonincreasing in functions $u_i(y_i)$,
- they are transitive.

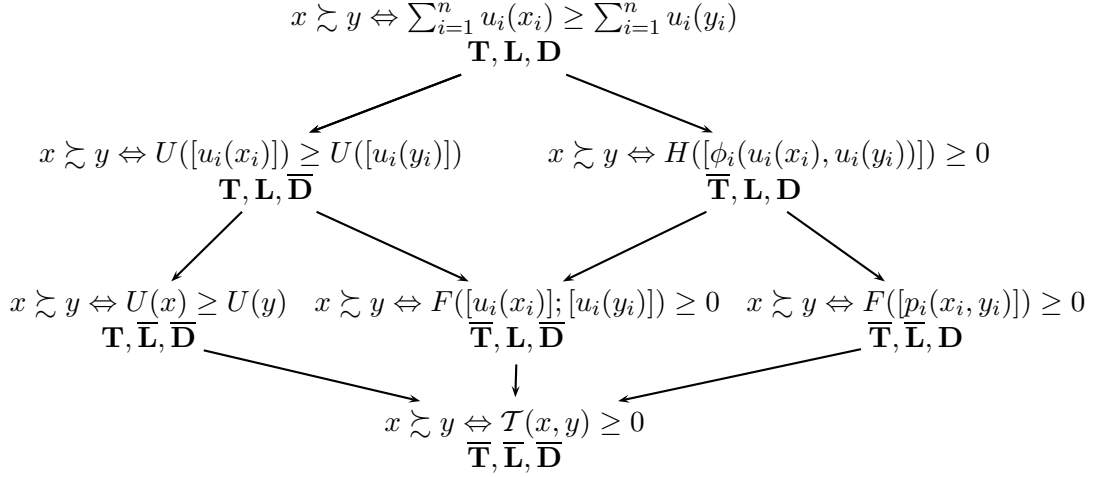
These various models are summarized in figure 6, where **T** denotes a transitive model, **L** a model involving traces on levels and **D** a model involving traces on differences.

In the **L** family all relations are weakly separable but may not be weakly independent (and, even less, independent). On the contrary, the **D** family includes only independent relations, as soon as axiom *RC2* is invoked. The marginal preference relations will be rather well-behaved in the **L** family being complete and most often semiorders (as soon as *AC3* and either *AC1* or *AC2* are invoked). This will not be the case in the **D** family.

It is worth noting that all combinations of **T**, **L** and **D** have been studied in the literature except for the combination **T**, $\bar{\mathbf{L}}$, **D**. This is no surprise since when **D** applies, most models appeal to *RC2* and, hence, are independent. When this is combined with transitivity and completeness of \succsim , \succsim_i is a weak order and is confounded with \succsim_i^\pm . Hence, such models also involve traces on levels.

Our hope is that the proposed framework and results will stimulate research in the area.

Figure 6: Summary of models



T means “transitive”

D means “uses traces on differences”

L means “uses traces on levels”

\bar{x} means “Not x ”

Examples

A AC1, AC2, AC3, TAC1 and TAC2

We first give 3 examples showing that, in the class of complete binary relations on X , $AC1_i$, $AC2_i$ and $AC3_i$ are independent conditions. We leave to the reader the easy task of finding similar examples in the case of incomplete (e.g., irreflexive) binary relations. This will prove part 5 of lemma 4.1. It is tedious, but easy, to check that completeness, $AC1_i$, $AC2_i$ and $AC3_i$ are in fact *completely independent* conditions.

Examples A.1 to A.3 have a common structure. In all of them $X = X_1 \times X_2$ with $X_1 = \{w_1, x_1, y_1, z_1\}$ and $X_2 = \{w_2, x_2, y_2, z_2\}$. We define \succsim on X by:

$$(r_1, r_2) \succsim (s_1, s_2) \Leftrightarrow F[f(r_1, s_1); g(r_2, s_2)] \geq 0 \quad (88)$$

where f (resp. g) is a real-valued function on X_1^2 (resp. on X_2^2) and F is a

real-valued function on \mathbb{R}^2 .

It is clear that, using obvious notation:

1. When F is odd ($F(\mathbf{x}) = -F(-\mathbf{x})$) and f and g are skew symmetric ($f(r_1, s_1) = -f(s_1, r_1)$), \succsim is complete.
2. When F is nondecreasing in its first argument, $[f(r_1, t_1) \geq f(s_1, t_1)$ for all $t_1 \in X_1] \Rightarrow r_1 \succsim_1^+ s_1$ and $[f(t_1, r_1) \leq f(t_1, s_1)$ for all $t_1 \in X_1] \Rightarrow r_1 \succsim_1^- s_1$. Similar conditions hold for the second component when F is nonincreasing in its second argument.

In all the following examples, f and g will be skew symmetric and $F(\alpha, \beta) = \alpha + \beta$, so that F is odd and increasing in its two arguments.

Example A.1 (AC1, AC2, Not AC3)

Define f and g by the following table (all tables are to be read from row to column; since the functions are skew symmetric only the upper diagonal parts of the tables are given):

f	w_1	x_1	y_1	z_1
w_1	0	0	0	0
x_1	0	0	2	2
y_1	0	-2	0	1
z_1	0	-2	-1	0
g	w_2	x_2	y_2	z_2
w_2	0	0	3	4
x_2	0	0	3	4
y_2	-3	-3	0	4
z_2	-4	-4	-4	0

Since the table for g is step-typed, it is clear that \succsim_2^\pm is complete. Hence, $AC1_2$, $AC2_2$ and $AC3_2$ hold.

It is easily checked that $AC1_1$ holds with \succsim_1^+ as (using obvious simplified notations for weak orders): $[x_1; w_1] \succsim_1^+ y_1 \succsim_1^+ z_1$. Similarly, $AC2_1$ holds with: $x_1 \succsim_1^- y_1 \succsim_1^- [w_1; z_1]$.

Because $w_1 \succsim_1^+ y_1$ and $y_1 \succsim_1^- w_1$, $AC3_1$ is violated.

Hence we have an example of a complete binary relation satisfying $AC1$, $AC2$ and $AC3_i$ on all but one attribute. \diamond

Example A.2 (Not AC1, AC2, AC3)

Define f by the following table:

f	w_1	x_1	y_1	z_1
w_1	0	0	0	1
x_1	0	0	2	3
y_1	0	-2	0	3
z_1	-1	-3	-3	0

and use the same table for g as in example A.1. We have: $Not[y_1 \succ_1^+ w_1]$ (because $(w_1, w_2) \succ (x_1, w_2)$ and $Not[(y_1, w_2) \succ (x_1, w_2)]$) and $Not[w_1 \succ_1^+ y_1]$ (because $(y_1, y_2) \succ (z_1, w_2)$ and $Not[(w_1, y_2) \succ (z_1, w_2)]$). Hence $AC1_1$ is violated. In fact it is easy to check that \succ_1^+ is such that $x_1 \succ_1^+ y_1$, $x_1 \succ_1^+ w_1$, $x_1 \succ_1^+ z_1$, $y_1 \succ_1^+ z_1$ and $w_1 \succ_1^+ z_1$.

It is easily checked that $AC2_1$ holds with $x_1 \succ_1^- [y_1; w_1] \succ_1^- z_1$. Using part 3 of lemma 4.1, we can check that $AC3_1$ holds.

Hence we have an example of a complete binary relation satisfying $AC2$, $AC3$ and $AC1_i$ on all but one attribute. \diamond

Example A.3 (AC1, Not AC2, AC3)

Transposing the tables defining f and g in example A.2 gives an example of a complete binary relation satisfying $AC1$, $AC3$ and $AC2_i$ on all but one attribute. \diamond

The next two examples are related to lemma 4.2. We first show that there are weakly independent semiorders satisfying $TAC12$ that are not weak orders.

Example A.4 (Nontransitive semiorder satisfying TAC12)

Let $X = X_1 \times X_2$ with $X_1 = \{x_1, y_1, z_1\}$ and $X_2 = \{x_2, y_2, z_2\}$. Consider the binary relation \succ identical to the weak order: $(x_1, x_2) > (x_1, y_2) > (y_1, x_2) > (x_1, z_2) > (y_1, y_2) > (z_1, x_2) > (y_1, z_2) > (z_1, y_2) > (z_1, z_2)$, except that $(y_1, y_2) \sim (x_1, z_2)$ and $(z_1, x_2) \sim (y_1, y_2)$.

This relation is clearly complete. It is not transitive since $(z_1, x_2) \succ (y_1, y_2)$, $(y_1, y_2) \succ (x_1, z_2)$ but $(x_1, z_2) \succ (z_1, x_2)$.

It is easily checked that this relation is a semiorder having the preceding weak order for trace. This semiorder is independent. Its marginal relations are weak orders identical to its marginal traces. We have $x_1 > y_1 > z_1$ and $x_2 > y_2 > z_2$.

This relation has only a few pairs of alternatives linked by \sim . It is then easy to check that $TAC12$ holds. For instance, starting with $(y_1, y_2) \succ (x_1, z_2)$ we should have $(x_1, y_2) \succ (x_1, z_2)$, $(y_1, x_2) \succ (x_1, z_2)$ and $(y_1, y_2) \succ (y_1, z_2)$, because $x_1 \succ_1^\pm y_1$ and $x_2 \succ_2^\pm y_2$. This is indeed the case.

Hence we have an example of a nontransitive weakly independent semiorder satisfying $TAC12$. \diamond

The final example shows that for complete relations, *TAC2* may hold without *TAC1*. An example of complete relation verifying *TAC1* but not *TAC2* is easily built using a similar principle.

Example A.5 (Not *TAC1*, *TAC2*)

Let $X = X_1 \times X_2$ with $X_1 = \mathbb{R}^2$ and $X_2 = \mathbb{R}$.

Define \succsim letting:

$$((a_1, b_1), x_2) \succsim ((c_1, d_1), y_2) \Leftrightarrow a_1 + x_2 > c_1 + y_2 \text{ or } \begin{cases} a_1 + x_2 = c_1 + y_2 \\ \text{and} \\ a_1 + b_1 \geq c_1. \end{cases}$$

It is clear that \succsim is complete.

We have $(a_1, b_1) \succsim_1^- (c_1, d_1) \Leftrightarrow a_1 \geq c_1$ and $(a_1, b_1) \succsim_1^+ (c_1, d_1) \Leftrightarrow (a_1, b_1) \geq_L (c_1, d_1)$, where \geq_L denotes the lexicographic order on \mathbb{R}^2 . On the second attribute, it is clear that $x_2 \succsim_2^+ y_2 \Leftrightarrow x_2 \succsim_2^- y_2 \Leftrightarrow x_2 > y_2$.

A simple check shows that \succsim is strictly responsive to \succsim_2^+ , \succsim_2^- and \succsim_1^- . This not so for \succsim_1^+ since, for instance, $((10, 0), 10) \sim ((8, 2), 12)$ and $((10, 2), 10) \sim ((8, 2), 12)$, while $(10, 2) \succ_1^+ (10, 0)$ (because $((10, 2), 10) \succ ((11, 0), 9)$ and $\text{Not}[(10, 0), 10] \succ ((11, 0), 9)$]).

Hence we have an example of a complete relation satisfying *TAC2* and *TAC1₂* but violating *TAC1₁*. \diamond

B *RC1* and *RC2*

We first give six examples complete the proof of part 3 of lemma 5.2 showing that independence, reflexivity and *RC1_i* are completely independent conditions.

Example B.1 (reflexive, *RC1*, not independent)

Let $X = \{a, b\} \times \{z, w\}$ and consider \succsim on X defined by: for all $(\alpha, \beta), (\gamma, \delta) \in X$, $(\alpha, \beta) \succsim (\gamma, \delta) \Leftrightarrow f(\alpha, \gamma) + g(\beta, \delta) \geq 0$, where f and g are such that: $f(a, a) = -1$, $f(a, b) = 0.5$, $f(b, a) = -0.5$, $f(b, b) = 1$, $g(z, z) = g(w, w) = g(w, z) = 1$, $g(z, w) = 0$

It is easy to see that \succsim is reflexive and satisfies *RC1*. It is not independent since $(b, z) \succsim (b, w)$ and $\text{Not}[(a, z) \succsim (a, w)]$. \diamond

Example B.2 (*RC1*, not reflexive, not independent)

In example B.1, taking $f(a, a) = -2$ leads to a relation \succsim that verifies *RC1* but is neither independent nor reflexive (since $\text{Not}[(a, z) \succsim (a, z)]$). \diamond

Example B.3 (RC1, not reflexive, independent)

Let $X = \{a, b\} \times \{z, w\}$ and consider \succsim on X defined by: for all $(\alpha, \beta), (\gamma, \delta) \in X$, $(\alpha, \beta) \succsim (\gamma, \delta) \Leftrightarrow f(\alpha, \gamma) + g(\beta, \delta) \geq 0$, where f and g are such that: $f(a, a) = f(b, b) = f(b, a) = -1$, $f(a, b) = 1$, $g(z, z) = g(w, w) = 0$, $g(z, w) = 1$, $g(w, z) = -1$

It is easy to see that \succsim is not reflexive (it is in fact irreflexive). It is easily seen to satisfy *RC1*. Since $f(a, a) = f(b, b)$ and $g(z, z) = g(w, w)$, \succsim is clearly independent. \diamond

Example B.4 (not RC1, reflexive, independent)

Let $X = \{a, b, c\} \times \{z, w\}$ and consider \succsim on X that is a clique (with all loops) except that $\text{Not}[(a, z) \succsim (c, w)]$ and $\text{Not}[(a, w) \succsim (b, z)]$.

It is clear that \succsim is reflexive. It can easily be checked that \succsim is independent. It does not satisfy *RC1* since: $(a, z) \succsim (b, w)$, $(a, w) \succsim (c, z)$, $\text{Not}[(a, z) \succsim (c, w)]$ and $\text{Not}[(a, w) \succsim (b, z)]$. \diamond

Example B.5 (not RC1, not reflexive, independent)

Modifying example B.4 in order to have \succsim irreflexive gives an example of a relation that is independent but violates *RC1* and reflexivity. \diamond

Example B.6 (not RC1, reflexive, not independent)

Modifying example B.4 in order to have $\text{Not}[(b, z) \succsim (b, w)]$ leads to a relation \succsim that is reflexive but violates independence and *RC1*. \diamond

We now give three examples that complete the proof of part 4 of lemma 5.2 showing that *RC1_i* and *RC2_i* are completely independent conditions in the class of complete relations.

Example B.7 (not RC1, not RC2)

Let $X = \{a, b, c\} \times \{z, w, k\}$ and consider \succsim on X that is a clique (with all loops) except that $\text{Not}[(a, z) \succsim (c, w)]$, $\text{Not}[(a, k) \succsim (b, z)]$ and $\text{Not}[(c, z) \succsim (a, w)]$. It is clear that \succsim is complete. Since $(a, z) \succsim (b, w)$, $(c, k) \succsim (a, z)$, $\text{Not}[(a, k) \succsim (b, z)]$ and $\text{Not}[(c, z) \succsim (a, w)]$, \succsim violates *RC1*. Since $(a, z) \succsim (b, w)$, $(b, z) \succsim (a, w)$, $\text{Not}[(a, z) \succsim (c, w)]$ and $\text{Not}[(c, z) \succsim (a, w)]$, \succsim violates *RC2*. \diamond

Example B.8 (not RC1, RC2)

Modify example B.7 adding the relation $(a, z) \succsim (c, w)$. It is clear that \succsim is complete and violates *RC1*. Using lemma 5.1, it is not difficult to see that it satisfies *RC2*. \diamond

Example B.9 (RC1, not RC2)

Let $X = \{a, b\} \times \{z, w\}$ and consider \succsim on X defined by: for all $(\alpha, \beta), (\gamma, \delta) \in X$, $(\alpha, \beta) \succsim (\gamma, \delta) \Leftrightarrow f(\alpha, \gamma) + g(\beta, \delta) \geq 0$, where f and g are such

that: $f(a, a) = -1, f(a, b) = f(b, a) = f(b, b) = 1, g(z, w) = 0, g(z, z) = g(w, w) = g(w, z) = 1$.

It is easy to see that \succsim is complete and satisfies *RC1*. It is not independent since $(b, z) \succsim (b, w)$ and $\text{Not}[(a, z) \succsim (a, w)]$. In view of part 2 of lemma 5.2, this shows that *RC2* is violated. \diamond

C *TC, AC1, AC2, AC3*

We show below that, in the class of complete binary relations conditions *TC_i*, *AC1_i*, *AC2_i* and *AC3_i* are independent.

Example C.1 (*AC1, AC2, AC3, Not TC*)

Let $X = \{a, b, c\} \times \{d, e, f\}$; $x \succsim y$ iff $G(p_1(x_1, y_1), p_2(x_2, y_2)) \geq 0$ with

$$G(\alpha, \beta) = \begin{cases} \alpha + \beta & \text{if } |\alpha + \beta| > 2 \\ 0 & \text{otherwise} \end{cases}$$

and p_1 and p_2 given in the following tables

p_1	a	b	c
a	0	-2	-1
b	2	0	1
c	1	-1	0

p_2	d	e	f
d	0	0	-2
e	0	0	-2
f	2	2	0

G is odd and non-decreasing and p_1, p_2 are skew symmetric; hence \succsim is complete and satisfies *RC1, RC2*. Condition *TC* is violated since $(c, d) \succsim (a, f), (a, e) \succsim (c, d), (a, d) \succsim (b, e)$ but $\text{Not}[(a, d) \succsim (b, f)]$.

It is easily checked that *AC1, AC2* and *AC3* hold with $b \succ_1^\pm c \succ_1^\pm a$ and $f \succ_2^\pm [d \sim_2^\pm e]$. \diamond

Example C.2 (*Not AC1, AC2, AC3, TC*)

Let $X = \{a, b, c\} \times \{d, e, f\}$. Define \succsim letting $x \succsim y$ iff $G(p_1(x_1, y_1), p_2(x_2, y_2)) \geq 0$, with p_1 and p_2 given in the following tables

p_1	a	b	c
a	0	2	-1
b	-2	0	-1
c	1	1	0

p_2	d	e	f
d	0	2	2
e	-2	0	2
f	-2	-2	0

and G such that :

G	-2	-1	0	1	2
-2	-41	-30	-21	-10	0
-1	-31	-20	-9	0	10
0	-19	-11	0	11	19
1	-10	0	9	20	31
2	0	10	21	30	41

G is odd and increasing in its two arguments and p_1, p_2 are skew symmetric implying that \succsim is complete and satisfies TC .

It is easy to check that we have: $c \succ_1^- a \succ_1^- b, c \succ_1^+ b, a \succ_1^+ b, \text{Not}[c \succ_1^\pm a], \text{Not}[a \succ_1^\pm c], d \succ_2^\pm e \succ_2^\pm f$.

Hence $AC2$ and $AC3$ hold but $AC1_1$ is violated (while $AC1_2$ holds). One verifies indeed that we have $(c, f) \succ (c, f)$ and $(a, f) \succ (b, e)$ but neither $(a, f) \succ (c, f)$ nor $(c, f) \succ (b, e)$. \diamond

Example C.3 (AC1, Not AC2, AC3, TC)

This example is the same as example C.2 except that p_1 becomes $-p_1$. The effect of this modification is to interchange the roles of $AC1_1$ and $AC2_1$ since the value associated to the pair (y_1, x_1) is the value that was formerly associated to (x_1, y_1) in example C.2. Therefore \succsim is complete and verifies TC .

We have: $b \succ_1^+ a \succ_1^+ c, b \succ_1^- c, b \succ_1^- a, \text{Not}[a \succ_1^- c], \text{Not}[c \succ_1^- a], d \succ_2^\pm e \succ_2^\pm f$. Hence $AC1$ and $AC3$ hold but $AC2_1$ is violated (while $AC2_2$ holds). One verifies indeed that $(c, f) \succ (c, f)$ and $(b, e) \succ (a, d)$ but neither $(c, f) \succ (a, f)$ nor $(b, e) \succ (c, d)$. \diamond

Example C.4 (AC1, AC2, Not AC3, TC)

Let $X = \{a, b, c, d\} \times \{w, x, y, z\}$. Define \succsim as in example C.2 with the same table for G and p_1, p_2 given in the following tables

p_1	a	b	c	d
a	0	1	2	2
b	-1	0	1	0
c	-2	-1	0	-2
d	-2	0	2	0

p_2	w	x	y	z
w	0	2	2	2
x	-2	0	2	2
y	-2	-2	0	2
z	-2	-2	-2	0

Since G is odd and increasing and p_1, p_2 are skew symmetric, we know that \succsim is complete and verifies TC .

It can be checked that we have: $w \succ_2^\pm x \succ_2^\pm y \succ_2^\pm z, a \succ_1^+ d \succ_1^+ b \succ_1^+ c$ and $a \succ_1^- b \succ_1^- d \succ_1^- c$. Hence $AC1$ and $AC2$ hold but $AC3$ is violated since neither $d \succ_1^\pm b$ nor $b \succ_1^\pm d$. \diamond

Proofs

D Proof of lemma 6.5

We prove lemma 6.5 in the most constrained case, i.e. when f is skew symmetric. Since U is countable, the set S of values r for which there are

$a, b, c, d \in U$ such that either:

$$\begin{aligned} f(a, c) = f(b, c) = r & \quad \text{with } a > b \text{ or} \\ f(a, c) = f(a, d) = r & \quad \text{with } c > d, \end{aligned}$$

is countable. Consider separately the positive part S^+ ($r > 0$), the null part $S^=$ ($r = 0$) and the negative part S^- ($r < 0$) of S . We number arbitrarily the elements of S^+ and S^-

$$\begin{aligned} r_1^+, r_2^+, \dots & \text{ for the elements of } S^+ \\ r_1^-, r_2^-, \dots & \text{ for the elements of } S^-. \end{aligned}$$

For each u, v in $U \times U$, we define $g(u, v)$ letting:

$$g(u, v) = \begin{cases} u - v & \text{if } f(u, v) = 0, \\ f(u, v) + 1 + \sum_{k:r_k^+ < f(u,v)} (1/2)^k & \text{if } f(u, v) > 0 \text{ and } f(u, v) \notin S \\ r_i^+ + 1 + \sum_{k:r_k^+ < r_i^+} (1/2)^k + (1/2)^{i+1}(1 + u - v) & \text{if } f(u, v) = r_i^+ \\ f(u, v) - 1 - \sum_{k:r_k^- > f(u,v)} (1/2)^k & \text{if } f(u, v) < 0 \text{ and } f(u, v) \notin S \\ r_i^- - 1 - \sum_{k:r_k^- > r_i^-} (1/2)^k - (1/2)^{i+1}(1 - u + v) & \text{if } f(u, v) = r_i^- \end{cases}$$

The function g is now fully described. It is skew symmetric since f is. It is strictly monotonic on $f^{-1}(r)$, for all $r \in S$. By construction, we have:

$$[f(u, v) = f(u', v')] \Rightarrow [g(u, v) = g(u', v')] \text{ iff } f(u, v) \notin S \quad (79)$$

It is easy to check that g is increasing in its first argument, decreasing in its second argument and such that (78) holds.

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